

# MATH/CSCI 4116: CRYPTOGRAPHY, WINTER 2019

## Handout 3: Problems for Homeworks 3 and 4

**Problem 1.** Consider the following  $S$ -box with 3 input bits and 3 output bits (a  $3 \times 3$   $S$ -box):

Input:	000	001	010	011	100	101	110	111
Output:	110	101	001	000	011	010	111	100

- Make a linear approximation table for this  $S$ -box, analogous to Table 4 in Heys.
- What is the highest probability bias (either positive or negative) of any input sum and output sum?

**Problem 2.** Consider the simple substitution permutation network shown in Figure 1 at the end of this handout. Assume that the  $S$ -box used is that from Problem 1. Find the encryption of the plaintext “011010”, using the key  $(K_1, K_2, K_3, K_4) = (010101, 001011, 111000, 111110)$ . For convenience, also show the intermediate results (i.e., the rows  $A$ ,  $B$ ,  $D$ ,  $E$ ,  $F$ ,  $G$ ,  $H$ , and  $J$  from Figure 1).

**Problem 3.** For the substitution permutation network from Problem 2:

- Find a linear approximation, analogous to Figure 3 in Heys, which relates the plaintext bits  $P_1$ ,  $P_2$ ,  $P_4$ , and  $P_5$  to a suitable subset of the inputs to the last round of  $S$ -boxes, i.e., a subset of the bits  $H_1, \dots, H_6$ .
- What is the total bias of the linear approximation you found in part (a)? What does this mean?
- Suppose you are given the following known plaintext/ciphertext pairs

for this cipher, all encrypted with the same (unknown) key:

Plaintext	Ciphertext
100111	100100
000111	110010
001100	111001
011000	011101
001000	001101
011010	101001

Using the linear approximation from part (a), determine the first and third bits of the subkey  $K_4$ . (Bonus question: why is this information insufficient to determine the second bit of the subkey  $K_4$ ?)

**NOTE:** this problem has been specifically constructed so that a very small number of plaintexts and ciphertexts is sufficient to determine two subkey bits. Unlike the general method of Heys, a probabilistic analysis is not necessary in this problem — the bits in question can be determined with certainty.

**Problem 4.** Make a difference distribution table, analogous to Table 7 in Heys, for the  $S$ -box from Problem 1.

**Problem 5.** For the substitution permutation network from Problem 2: Consider two plaintexts  $P$  and  $P'$  such that  $\Delta P = P \oplus P' = 000001$ .

- Using your difference distribution table from the previous problem, what are the possible values for  $\Delta H$ ? Remark: note that there are only six possible values.
- Suppose you are given the following additional chosen plaintext, and corresponding ciphertext:

Plaintext	Ciphertext
100110	111110
000110	110110
001101	100000
011001	011111
001001	000011
011011	101000

This is in addition to the plaintext/ciphertext pairs from Problem 3. Use your answer to part (a) to determine the last three bits of the subkey  $K_4$ .

**NOTE:** this problem has been specifically constructed so that a very small number of plaintexts and ciphertexts is sufficient to determine three subkey bits. Unlike the general method of Heys, a probabilistic analysis is not necessary in this problem — the bits in question can be determined with certainty.

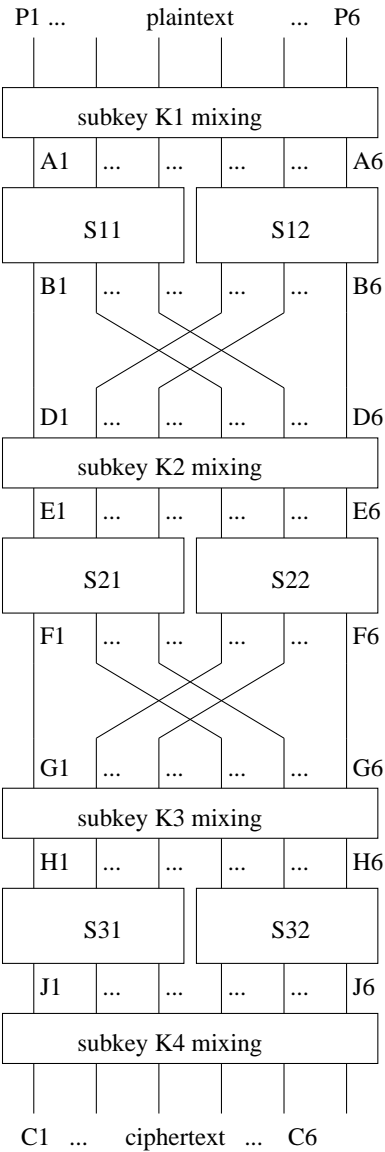


Figure 1: A very simple SPN network