

Lecture 18
(16 October 2023)

Functions of Random Variables

$y = g(x)$, i.e., $y(\omega) = g(x(\omega))$.

For the discrete case

$$P_Y(y) = \sum_{x: g(x)=y} P_X(x),$$

For the continuous case we follow the two-step procedure outlined below.

1) calculate the CDF F_Y of y using

$$F_Y(y) = P(g(x) \leq y) = \int_{\{x: g(x) \leq y\}} f_X(x) dx.$$

2) Differentiate the CDF to obtain the PDF of y :

$$f_Y(y) = \frac{d}{dy} F_Y(y).$$

Linear Function of a RV:

$$Y = ax + b, \quad a \neq 0.$$

Discrete case -

$$P_Y(y) = P_X\left(\frac{y-b}{a}\right) \text{ for all } y,$$

continuous case -

$$Y = ax + b, \quad a \neq 0$$

$$\Rightarrow f_Y(y) = \frac{1}{|a|} f_X\left(\frac{y-b}{a}\right).$$

Proof. $a > 0$ (proved earlier)

Let $a > 0$,

$$P(ax+b \leq y) = P(x \geq \frac{y-b}{a})$$

$$= 1 - P(x < \frac{y-b}{a})$$

$$= 1 - F_X\left(\frac{y-b}{a}\right)$$

$$\Rightarrow f_Y(y) = \frac{d}{dy} F_Y(y) = \frac{1}{|a|} f_X\left(\frac{y-b}{a}\right).$$

Example. X is uniform on $[0, 1]$ and let $Y = 5X$.

For $0 \leq y \leq 5$,

$$F_Y(y) = P(\sqrt{x} \leq y)$$

$$= P(x \leq y^2)$$

$$= y^2$$

$$\Rightarrow f_Y(y) = 2y, \quad 0 \leq y \leq 1.$$

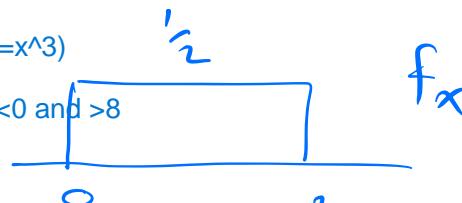
$$F_Y(y) = 0 \quad y \leq 0 \quad \text{and} \quad F_Y(y) = 1 \quad y \geq 1.$$

$\therefore f_Y(y) = 0$ for y outside $[0, 1]$.

Example. $y = x^3$ - $x \sim \text{uniform}[0, 2]$.

For $0 \leq y \leq 8$ if x can vary from 0 to 2, then realize that y can vary from 0 to 8 (because $y=x^3$)

Accordingly make the intervals. $[0, 8]$, <0 and >8

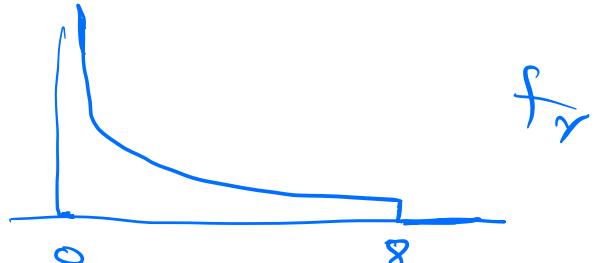


$$F_Y(y) = P(x^3 \leq y)$$

$$= P(x \leq y^{1/3})$$

$$= \frac{y^{1/3}}{2}$$

$$\Rightarrow f_Y(y) = \frac{1}{6} y^{-2/3} \quad 0 \leq y \leq 8.$$



Example. $y = \frac{10}{x}$ - $f_X(x) = \frac{1}{5}, \quad 5 \leq x \leq 10$,

For $1 \leq y \leq 2$

$$F_Y(y) = P\left(\frac{10}{x} \leq y\right) = P\left(x \geq \frac{10}{y}\right)$$

$$= 1 - P(X \leq \frac{10}{y})$$

$$= 1 - \frac{1}{5}(-5 + \frac{10}{y}) = -\frac{2}{y} + 2$$

$$\Rightarrow f_y(y) = \frac{2}{y^2}, \quad 1 \leq y \leq 2.$$

Monotonic Function:

Let x be a continuous RV and suppose that its range is contained in a certain interval I , in the sense that $f_x(x)=0$ for $x \notin I$. We consider $y=g(x)$ and assume that g is strictly monotonic over the interval I :

X is just a random variable.

$g(x) < g(x')$ for all $x, x' \in I$ s.t. $x < x'$

(monotonically increasing)

or

$g(x) > g(x')$ for all $x, x' \in I$ s.t. $x < x'$.

(monotonically decreasing)

For $y \in g^{-1}(I)$,

$P(g(x) \leq y)$

incr.g

$$= P(X \leq g^{-1}(y)) = F_x(g^{-1}(y))$$

$$\Rightarrow f_y(y) = f_x(g^{-1}(y)) (g^{-1})'(y).$$

For decreasing g ,

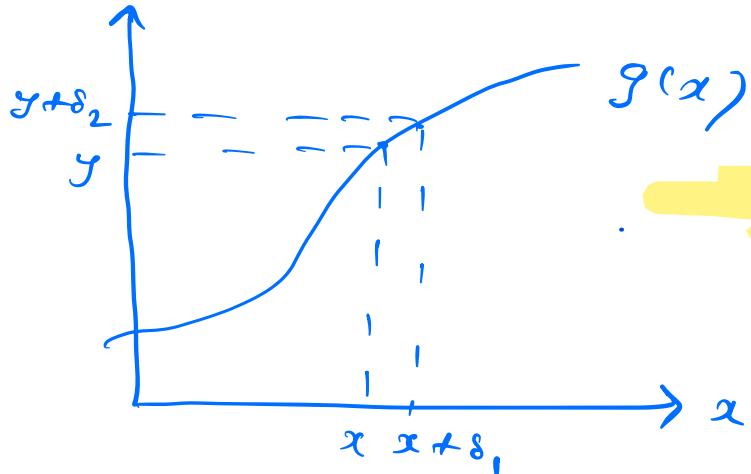
$$\begin{aligned} F_Y(y) &= P(g(x) \leq y) = P(x \geq g^{-1}(y)) \\ &= 1 - P(x \leq g^{-1}(y)) = 1 - F_X(g^{-1}(y)) \\ \Rightarrow f_Y(y) &= f_X(g^{-1}(y)) |(g^{-1})'(y)| \end{aligned}$$

$$\therefore f_Y(y) = f_X(g^{-1}(y)) |(g^{-1})'(y)|.$$

This can also be written as

$$f_Y(y) = \begin{cases} f_X(g^{-1}(y)) / |g'(g^{-1}(y))| & \text{if } g(x)=y \text{ for some } x \\ 0 & \text{if } g(x) \neq y \text{ for all } x \text{ with } f_X(x)>0 \end{cases}$$

Illustration:



$$\frac{\delta_2}{\delta_1} = g'(x)$$

why are both of them equal?

because $P(X=x) = P(Y=y)$ because there is only one point since one to one function. and small delta doesn't make a difference.

$$\begin{aligned} P(y \leq Y \leq y + \delta_2) &\approx f_Y(y) \delta_2 \\ &= P(x \leq X \leq x + \delta_1) = f_X(x) \delta_1, \end{aligned}$$

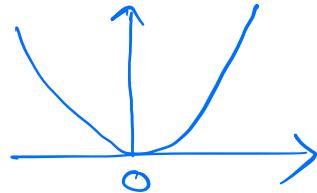
treat $f_Y(y)$ as a constant in this interval

$$f_y(y) = f_x(x) \frac{\delta_1}{\delta_2} = \frac{f_x(x)}{g'(x)} = \frac{f_x(g^{-1}(y))}{g'(g^{-1}(y))}$$

Non-monotonic Function

$$Y = g(x), \quad g(x) = x^2$$

For $y \geq 0$



$$\begin{aligned} P(g(x) \leq y) &= P(x^2 \leq y) \\ &= P(-\sqrt{y} \leq x \leq \sqrt{y}) \\ &= F_x(\sqrt{y}) - F_x(-\sqrt{y}) \end{aligned}$$

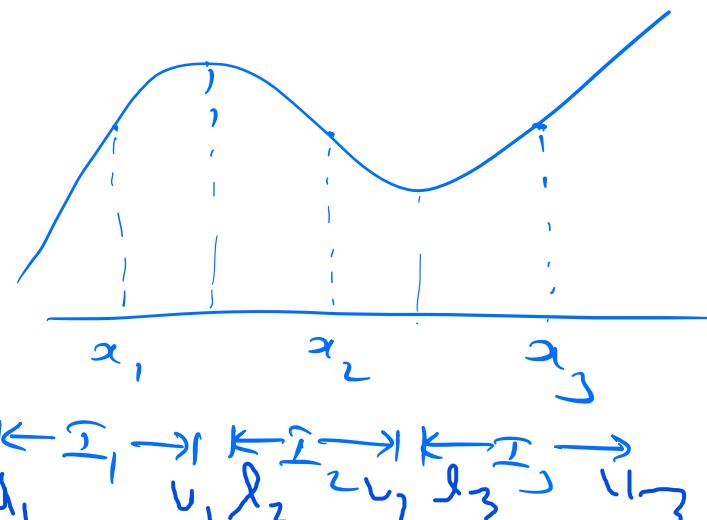
$$\Rightarrow f_y(y) = \frac{f_x(\sqrt{y})}{2\sqrt{y}} + \frac{f_x(-\sqrt{y})}{2\sqrt{y}}$$

Theorem. Consider a continuous RV x with PDF f_x , and let $y = g(x)$, suppose we partition the domain of f_x into a finite number of intervals such that $g(x)$ is strictly monotone and differentiable on each interval. Then the PDF of y is given by

$$f_y(y) = \sum_{i=1}^n \frac{f_x(x_i)}{|g'(x_i)|} - \text{where}$$

x_1, x_2, \dots, x_n are real solutions to $g(x)=y$ in the domain of f_x .

Proof.



$$\begin{aligned} I_1 &= [l_1, u_1], \\ I_2 &= [l_2, u_2], \\ I_3 &= [l_3, u_3], \\ I_4 &= [l_4, u_4]. \end{aligned}$$

For a given y let $g(x_i) = y$ $i \in \{1, 2, 3\}$.
 $g(x)$ is increasing at x_1 ,

decreasing at x_2 and
increasing at x_3 .

$$\begin{aligned} P(g(x) \leq y) &= P(\{l_1 \leq x \leq x_1\} \cup \{x_2 \leq x \leq u_2\} \\ &\quad \cup \{l_3 \leq x \leq x_3\}) \\ &= P(l_1 \leq x \leq x_1) + P(x_2 \leq x \leq u_2) + P(l_3 \leq x \leq x_3) \\ \Rightarrow f_y(y) &= f_x(x_1) \frac{dx_1}{dy} - f_x(x_2) \frac{dx_2}{dy} + f_x(x_3) \frac{dx_3}{dy} \end{aligned}$$

$$= \frac{f_x(x_1)}{|g'(x_1)|} + \frac{f_x(x_2)}{|g'(x_2)|} + \frac{f_x(x_3)}{|g'(x_3)|}.$$

→ For $y = x^2$, $x_1 = \sqrt{y}$, $x_2 = -\sqrt{y}$.

so $f_y(y) = \frac{f_x(\sqrt{y})}{2\sqrt{y}} + \frac{f_x(-\sqrt{y})}{2\sqrt{y}}$.

Lecture 19
(26 October 2023)

Functions of Two Random Variables

$$Z = g(X, Y),$$

$$\text{i.e., } Z(\omega) = g(X(\omega), Y(\omega)),$$

Example. $Z = \max(X, Y)$, X & Y are independent.

$$P(Z \leq z) = P(\max(X, Y) \leq z)$$

$$= P(X \leq z, Y \leq z)$$

$$= P(X \leq z) P(Y \leq z)$$

$$= F_X(z) F_Y(z)$$

$$\Rightarrow f_Z(z) = F'_X(z) f_Y(z) + f'_X(z) F_Y(z).$$

As a special case $X, Y \sim \text{Uniform}[0, 1]$,

$$f_Z(z) = \begin{cases} 2z, & 0 \leq z \leq 1 \\ 0 & \text{otherwise.} \end{cases}$$

Exercise: $Z = \min\{X, Y\}$ X and Y are indep. Find f_Z .

Example. X and Y are independent RVS that are uniformly distributed on $[0, 1]$. What is the PDF of $Z = Y/X$?

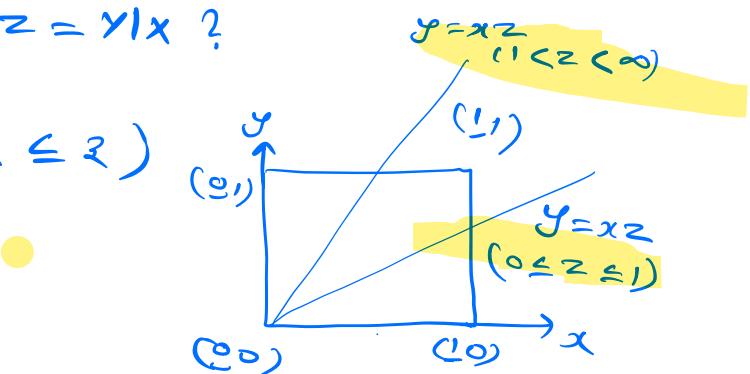
$$P(Z \leq z) = P(Y/X \leq z)$$

z can be from 0 to infinity

For $0 \leq z \leq 1$

$$P(Y \leq xz)$$

$$= \int_{x=0}^{zx} \int_{y=0}^1 1 dx dy = z/2.$$



For $1 < z < \infty$

$$P(Y \leq xz) = 1 - \frac{1}{2z}.$$

$$\Rightarrow f_Z(z) = \begin{cases} z/2 & \text{if } 0 \leq z \leq 1 \\ 1/2z & \text{if } z > 1 \\ 0 & \text{otherwise.} \end{cases}$$

$$f_Z(z) = \begin{cases} \frac{1}{2} & \text{if } 0 \leq z \leq 1 \\ \frac{1}{2z} & \text{if } z > 1 \\ 0 & \text{otherwise.} \end{cases}$$

Sum of Independent Random Variables

Let $Z = X+Y$ where X and Y are independent random variables,

Discrete case:

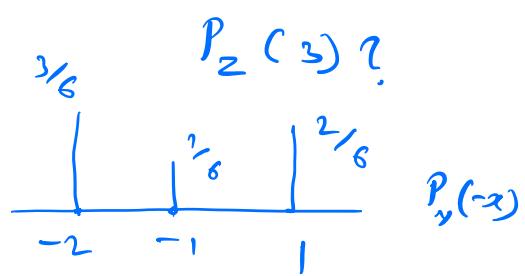
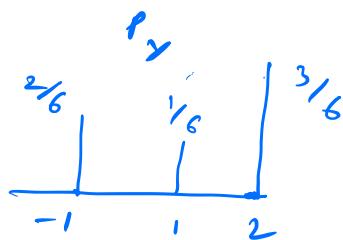
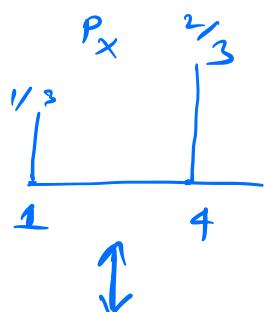
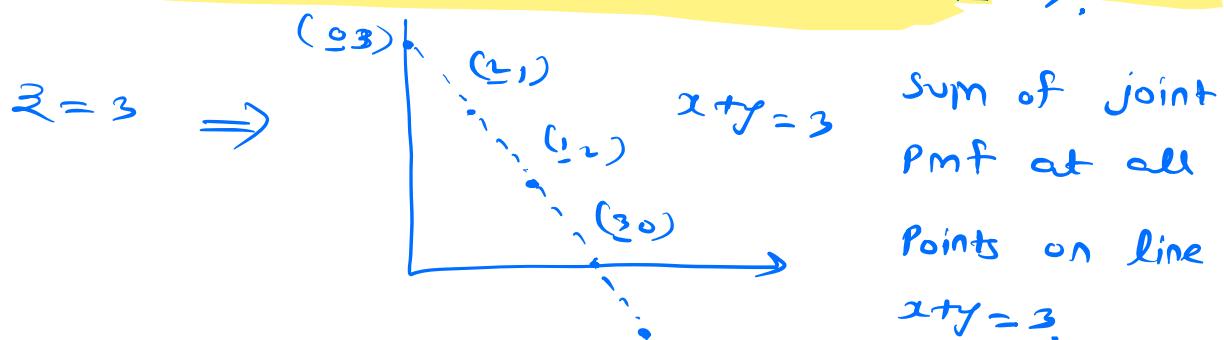
$$P_Z(z) = P(X+Y=z)$$

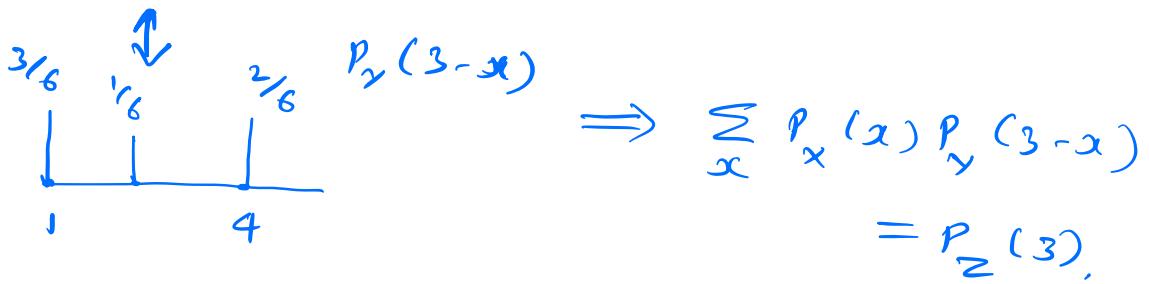
$$= \sum_{(x,y): x+y=z} P_{XY}(x,y)$$

$$= \sum_x P_{XY}(x, z-x) = \sum_x P_X(x) P_Y(z-x),$$

because X and Y are independent RVs

The resulting PMF P_Z is called the convolution of the PMFs of X and Y .





Continuous case:

$$\begin{aligned}
 P(X+Y \leq z) &= \int_{x=-\infty}^{\infty} \int_{y=-\infty}^{z-x} f_{xy}(x,y) dy dx \\
 &= \int_{x=-\infty}^{\infty} f_x(x) \int_{y=-\infty}^{z-x} f_y(y) dy dx \\
 &= \int_{x=-\infty}^{\infty} f_x(x) F_y(z-x) dx
 \end{aligned}$$

$d(F_y(z-x))$

$$\Rightarrow f_z(z) = \frac{d}{dz} F_z(z)$$

$\frac{d}{dz} F_z(z) = \frac{d}{dy} F_y(z-y) \cdot \frac{dy}{dz}$

$$\begin{aligned}
 &= \int_{-\infty}^{\infty} f_x(x) \frac{d}{dz} F_y(z-x) dx \\
 &= \int_{x=-\infty}^{\infty} f_x(x) f_y(z-x) dx
 \end{aligned}$$

$y + n = z$

$f_y(z-n) \text{ ch. } \rightarrow \frac{dy}{dz} = 1$

Theorem. If $x \sim N(\mu_1, \sigma_1^2)$, $y \sim N(\mu_2, \sigma_2^2)$
 are independent then
 $x+y \sim N(\mu_1 + \mu_2, \sigma_1^2 + \sigma_2^2)$.

Proof. $z = x+y$,

$$\begin{aligned}
 f_z(z) &= \int_{-\infty}^{\infty} f_x(x) f_y(z-x) dx \\
 &= \int_{-\infty}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}} \cdot \frac{1}{\sqrt{2\pi}} e^{-\frac{(z-x)^2}{2}} dx \\
 &= \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-\frac{-x^2 - z^2 + 2xz}{2}} dx \\
 &= \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2 - z^2 + 2xz}{2}} dx \\
 &= \frac{1}{\sqrt{4\pi}} \cdot e^{-\frac{z^2}{4}} \cdot \int_{-\infty}^{\infty} \frac{1}{\sqrt{\pi}} e^{-\frac{(x-z)^2}{2}} dx \\
 &= \frac{1}{\sqrt{4\pi}} \cdot e^{-\frac{z^2}{4}} \cdot 1
 \end{aligned}$$

Exercise. $x \sim N(\mu_x, \sigma_x^2)$, $y \sim N(\mu_y, \sigma_y^2)$ indep $\Rightarrow x+y \sim N(\mu_x + \mu_y, \sigma_x^2 + \sigma_y^2)$.

Sum of a discrete RV and a continuous RV:

$$Z = X + Y$$

X-discrete, Y-continuous, X & Y are independent

$$P(Z \leq z) = P(X+Y \leq z)$$

$$= \sum_x P(X+Y \leq z | X=x) P_X(x)$$

$$= \sum_x P(Y \leq z-x | X=x) P_X(x)$$

$$= \sum_x F_Y(z-x) P_X(x)$$

$$= \sum_x \int_{-\infty}^{z-x} f_Y(y) dy P_X(x) \quad y = t-x$$

$$= \sum_x \int_{-\infty}^z f_Y(t-x) dt P_X(x)$$

$$= \int_{-\infty}^z \left(\sum_x f_Y(t-x) P_X(x) \right) dt$$

$$\therefore f_Z(z) = \sum_x f_Y(z-x) P_X(x).$$

Functions of Two Random Variables

$(X, Y) \sim f_{XY}$.

$Z = g_1(X, Y)$, $W = g_2(X, Y)$.

$$P(Z \leq z, W \leq w) = P(g_1(X, Y) \leq z, g_2(X, Y) \leq w)$$

$$= \iint_{\substack{xy \\ g_1(x,y) \leq z \\ g_2(x,y) \leq w}} f_{XY}(x, y) dx dy$$

$$g_1(x, y) \leq z$$

$$g_2(x, y) \leq w$$

Sometimes it may be easier to find out the region under the above integral and compute the integral. Then we can find f_{ZW} by taking double derivative:

$$f_{ZW}(z, w) = \frac{\partial^2 F_{ZW}(z, w)}{\partial z \partial w}.$$

Exercise. Suppose x and y are independent uniformly distributed RVs in $[0, 1]$,

$$z = \min\{x, y\} \quad w = \max\{x, y\}, \text{ Find } f_{zw}.$$

- In some cases we may be able to obtain a closed-form expression for f_{zw} in terms of f_{xy} .

Assume that g_1 & g_2 satisfy;

$$(1) \quad z = g_1(x, y), \quad w = g_2(x, y)$$

$$\Leftrightarrow x = h_1(z, w), \quad y = h_2(z, w).$$

$$\text{Example: } z = xy, \quad w = x - y,$$

(2) The functions g_1 & g_2 have continuous partial derivatives at all points (x, y) and are such that

$$J(x, y) = \begin{vmatrix} \frac{\partial g_1}{\partial x} & \frac{\partial g_1}{\partial y} \\ \frac{\partial g_2}{\partial x} & \frac{\partial g_2}{\partial y} \end{vmatrix} = \frac{\partial g_1}{\partial x} \frac{\partial g_2}{\partial y} - \frac{\partial g_1}{\partial y} \frac{\partial g_2}{\partial x} \neq 0$$

at all points (x, y) .

Theorem. In the setting above,

$$f_{\geq \omega}(\geq \omega) = f_{xy}(x,y) | J(x,y)|^{-1}.$$

This is analogous to

$$f_y(y) = \frac{f_x(g^{-1}(y))}{|g'(g^{-1}(y))|},$$

when $y=g(x)$ for a monotonic g .

Lecture 20

(30 October 2023)

Two Functions of Two Random Variables

$(x, y) \sim f_{xy}$.

Let $z = g_1(x, y)$, $w = g_2(x, y)$.

Suppose g_1 & g_2 are such that

(i) \exists functions h_1, h_2 satisfying

$$x = h_1(z, w), \quad y = h_2(z, w).$$

(ii)

$$\begin{aligned} J(x, y) &:= \begin{vmatrix} \frac{\partial g_1(x, y)}{\partial x} & \frac{\partial g_1(x, y)}{\partial y} \\ \frac{\partial g_2(x, y)}{\partial x} & \frac{\partial g_2(x, y)}{\partial y} \end{vmatrix} \\ &= \frac{\partial g_1}{\partial x} \frac{\partial g_2}{\partial y} - \frac{\partial g_2}{\partial x} \frac{\partial g_1}{\partial y} \neq 0 \end{aligned}$$

at all x, y .

Theorem. In the setting above,

$$f_{zw}(z \leq w) = f_{xy}(x \leq y) |J(x,y)|^{-1},$$

where $x = h_1(z \leq w)$, $y = h_2(z \leq w)$.

This is analogous to

$$f_y(y) = \frac{f_x(g^{-1}(y))}{|g'(g^{-1}(y))|},$$

when $y = g(x)$ for a monotonic g .

Proof. Consider

$$P(z \leq Z \leq z + \Delta z, w \leq W \leq w + \Delta w)$$

$$= f_{zw}(z \leq w) \Delta z \Delta w.$$

Alternately,

$$P(z \leq Z \leq z + \Delta z, w \leq W \leq w + \Delta w)$$

$$= P(x \leq X \leq x + \Delta x, y \leq Y \leq y + \Delta y) = f_{xy}(x \leq y) \Delta x \Delta y$$

where $x = h_1(z \leq w)$, $y = h_2(z \leq w)$.

$$\Rightarrow f_{zw}(z, \omega) = f_{xy}(x, y) \frac{\Delta x \Delta y}{\Delta z \Delta w}$$

$$= \frac{f_{xy}(x, y)}{\left(\frac{\Delta z \Delta w}{\Delta x \Delta y} \right)} = \frac{f_{xy}(x, y)}{|J(x, y)|},$$

where

$$J(x, y) = \frac{\partial g_1}{\partial x} \frac{\partial g_2}{\partial y} - \frac{\partial g_1}{\partial y} \frac{\partial g_2}{\partial x}.$$

[See chapter 6
from Papoulis &
Pillai's Book for

Analogy: $g = g(x) \Rightarrow \frac{dg}{dx} = g'(x)$, more details
and a formal
justification]

Similarly $z = g_1(x, y)$ $w = g_2(x, y)$

$$\Rightarrow \frac{\Delta z \Delta w}{\Delta x \Delta y} = \left| \det \begin{bmatrix} \frac{\partial g_1}{\partial x} & \frac{\partial g_1}{\partial y} \\ \frac{\partial g_2}{\partial x} & \frac{\partial g_2}{\partial y} \end{bmatrix} \right|.$$

- Note that Jacobian is essentially doing the following change of variables:

$$\begin{aligned} P((z, \omega) \in B) &= P((x, y) \in A) = \iint_A f_{xy}(x, y) dx dy \\ &= \iint_B f_{xy}(h_1(z, \omega), h_2(z, \omega)) \cdot \frac{1}{|J(h_1(z, \omega), h_2(z, \omega))|} dz d\omega, \end{aligned}$$

→ more generally, for a given point (z, w) ,
 $f_1(x, y) = z$, $f_2(x, y) = w$ can have many solutions.
 Let $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)$ represent these multiple solutions such that

$$f_1(x_i, y_i) = z, \quad f_2(x_i, y_i) = w, \quad i \in [1:n].$$

Then

$$f_{zw}(z, w) = \sum_{i=1}^n f_{xy}(x_i, y_i) |J(x_i, y_i)|^{-1}.$$

Example. $\begin{bmatrix} z \\ w \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix}$.

$$J(x, y) = \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} = a_{11}a_{22} - a_{12}a_{21} \neq 0,$$

$$f_{zw}(z, w) = \frac{f_{xy}(x, y)}{|a_{11}a_{22} - a_{12}a_{21}|}, \quad \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}^{-1} \begin{bmatrix} z \\ w \end{bmatrix}.$$

Moment Generating Functions

- A transform associated with a random variable and provides an alternative representation of a probability law.
- Particularly useful for certain types of mathematical manipulations
- Moment generating function (MGF) of a RV X is a function $M_X(s)$ defined as
$$M_X(s) = E[e^{sx}], \text{ for a scalar } s.$$

Discrete case:

$$M_X(s) = \sum_x e^{sx} P_X(x).$$

Continuous case:

$$M_X(s) = \int_{-\infty}^{\infty} e^{sx} f_X(x) dx,$$

$$\rightarrow M_X(0) = 1.$$

Example. Let $P_X(1) = \frac{1}{2}$, $P_X(2) = \frac{1}{4}$, $P_X(3) = \frac{1}{4}$,

$$M_X(s) = \frac{1}{2} e^s + \frac{1}{4} e^{2s} + \frac{1}{4} e^{3s}.$$

Example. Exponential RV with parameter λ .

$$f_X(x) = \lambda e^{-\lambda x}, x \geq 0,$$

$$M_X(s) = \int_0^\infty e^{sx} \lambda e^{-\lambda x} dx$$

$$= \lambda \int_0^\infty e^{(s-\lambda)x} dx$$

$$= \lambda \left[\frac{e^{(s-\lambda)x}}{(s-\lambda)} \right]_0^\infty$$

$$= \frac{\lambda}{\lambda-s} \quad \text{if } s < \lambda,$$

otherwise the integral is infinite.

— $M_X(s)$ is only defined for those values of s for which $E[e^{sx}]$ is finite.

Exercise.

1) If x takes only non-negative integer values then show that

$$\lim_{s \rightarrow -\infty} m_x(s) = P(x=0).$$

2) Let x be a Poisson RV with parameter λ . Find $m_x(s)$.

3) Prove that $m_{ax+b}(s) = e^{sb} m_x(as)$.

Example. $x \sim N(0,1)$.

$$m_x(s) = \int_{-\infty}^{\infty} e^{sx} \frac{1}{\sqrt{2\pi}} e^{-x^2/2} dx$$

$$= \int_{-\infty}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2} + sx} dx$$

$$= \int_{-\infty}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-\frac{(x-s)^2}{2}} \cdot e^{sx} dx$$

$$= e^{s^2/2} \cdot \int_{-\infty}^{\infty} \frac{1}{\sqrt{2\pi}} e^{-\frac{(x-s)^2}{2}} dx$$

$$= e^{s^2/2}$$

$$Y \sim N(\mu - \sigma^2) \Rightarrow M_Y(s) = e^{sb} m_X(s)$$

$\left[Y = \sigma X + \mu \right]$

$$= e^{sb} \cdot e^{\frac{-\sigma^2 s^2}{2}}$$

$$= e^{\mu s + \frac{-\sigma^2 s^2}{2}}.$$

- MGFs are useful mainly for two reasons

(1) MGFs of x gives us all moments of x .

$$M_X(s) = E[e^{sx}]$$

$$= \int_{-\infty}^{\infty} e^{sx} f_X(x) dx$$

$$\frac{d}{ds} M_X(s) = \int_{-\infty}^{\infty} x e^{sx} f_X(x) dx$$

$$\frac{d}{ds} M_X(s) \Big|_{s=0} = \int_{-\infty}^{\infty} x f_X(x) dx$$

Similarly $\frac{d^n}{ds^n} M_X(s) \Big|_{s=0} = \int_{-\infty}^{\infty} x^n f_X(x) dx.$

(2) MGF (if it exists) uniquely determines

the CDF of random variable.

Consider two random variables x and y , suppose that there exists a positive constant c such that MGFs of x and y are finite and identical for all values of s in $[-c, c]$. Then,

Uniqueness

$$F_x(t) = F_y(t) \text{ for all } t \in \mathbb{R}. \quad \text{Property}$$

Proof follows from connection with Laplace transform.

This can also be seen through characteristic functions.

$$\Phi_x(t) = E[e^{itx}] \quad i = \sqrt{-1}.$$

- characteristic functions always exist.

- Related to Fourier transform.

Fact:

$\int_{-\infty}^{\infty} g(x) dx$ exists if $\int_{-\infty}^{\infty} |g(x)| dx$ exists.

Since $|e^{itx}| = |\Phi_x|$ always exists,

- If x and y are independent RVs

$$M_{x+y}(s) = M_x(s)M_y(s).$$

Exercise . (i) $X \sim N(\mu_x, \sigma_x^2)$, $Y \sim N(\mu_y, \sigma_y^2)$

$$\Rightarrow X+Y \sim N(\mu_x + \mu_y, \sigma_x^2 + \sigma_y^2).$$

(ii) Let x and y be independent Poisson RVs with parameters λ_1 & λ_2 , respectively.

Then $X+Y$ is Poisson with parameter $\lambda_1 + \lambda_2$

Lecture 21

(2 November 2023)

Module 5 : Probability Bounds & Limit Theorems

— Suppose we want to compute $P(x \geq a)$. In some scenarios it may be sufficient to have bounds on this probability instead of its exact value, e.g., when the distribution of x is unavailable or hard to compute. In such scenarios if we have exact values or bounds for the mean and variance of x , we can obtain meaningful bounds on the quantity of interest.

Markov's Inequality

If x is a non-negative random variable

then $P(x \geq a) \leq \frac{E[x]}{a}$, for all $a > 0$.

Interpretation:

"If $x \geq 0$ and $E[x]$ is small, then the probability that x takes a large value must be small".



$$E[x] = \sum_x x P_x(x), \quad P_x(1000) \rightarrow \text{small}$$

Proof. we first prove that if $x \geq y$, then $E[x] \geq E[y]$.

$z = x - y \geq 0$, i.e., $x(\omega) - y(\omega) \geq 0$ whenever,

$$E[z] \geq 0 \Rightarrow E[x] \geq E[y].$$

Let $y = a \mathbf{1}\{x \geq a\}$, we have

$$x \geq y.$$

$$\Rightarrow E[x] \geq E[y]$$

$$\Rightarrow E[x] \geq E[a \mathbf{1}\{x \geq a\}]$$

$$= a P(x \geq a)$$

$$\therefore P(x \geq a) \leq \frac{E[x]}{a}.$$

Example. Let $X \sim \text{Binomial}(n, p)$. Using Markov's inequality find an upper bound on $P(X \geq \alpha n)$, where $0 < \alpha < 1$. Evaluate the bound for $p = \frac{1}{2}$ and $\alpha = \frac{3}{4}$.

$$P(X \geq \alpha n) \leq \frac{E[X]}{\alpha n} = \frac{np}{n\alpha} = \frac{p}{\alpha}$$

$$= \frac{1}{2} \cdot \frac{4}{3} = \frac{2}{3}.$$

$$\therefore P(X \geq \alpha n) \leq \frac{2}{3}.$$

Example. $X \sim \text{Uniform}[-4, 4]$.

$$P(X \geq 3) \leq P(|X| \geq 3)$$

$$\leq \frac{E[|X|]}{3}$$

$$= \frac{2}{3}.$$

Chebychev's Inequality

If x is a random variable with mean μ and variance σ^2 then

$$P(|x-\mu| \geq c) \leq \frac{\sigma^2}{c^2} \text{ - for all } c > 0,$$

Proof. Let $y = |x-\mu|$.

$$P(y \geq c) = P(y^2 \geq c^2)$$

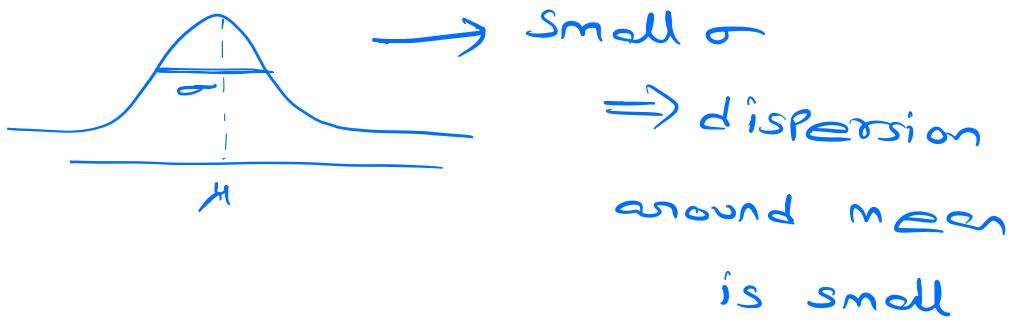
$$\begin{aligned} &\leq \frac{E[y^2]}{c^2} = \frac{E[|x-\mu|^2]}{c^2} \\ &= \frac{\sigma_x^2}{c^2}. \end{aligned}$$

Interpretation:

"If a random variable has small variance then the probability that it takes a value far from its mean is also small."

- Recall that variance measures the spread

of Rv X around its mean.



- An alternative form of chebyshев's inequality:

$$P(|X-\mu| \geq k\sigma) \leq \frac{1}{k^2}.$$

Example. $X \sim \text{Binomial}(n, p)$. Using chebyshev's inequality find an upper bound on $P(X \geq \alpha n)$ where $p < \alpha < 1$. Evaluate for $p = \frac{1}{2}$, $\alpha = \frac{3}{4}$.

$$P(X \geq \alpha n) = P(X - np \geq n(\alpha - p))$$

$$\leq P(|X - np| \geq n(\alpha - p))$$

$$\begin{aligned} &\leq \frac{\text{Var}(X)}{n^2(\alpha - p)^2} = \frac{np(1-p)}{n^2(\alpha - p)^2} \\ &= \frac{4}{n}. \end{aligned}$$

Chernoff Bounds

If X is a random variable then for any $a \in \mathbb{R}$, we can write

now apply Markov Inequality here

$$P(X \geq a) = P(e^{sx} \geq e^{sa}), \text{ for } s > 0$$

$$P(X \leq a) = P(e^{sx} \geq e^{sa}), \text{ for } s < 0.$$

Note that e^{sx} is always a positive random variable for all $s \in \mathbb{R}$. Thus we can apply Markov's inequality.

$$\text{For } s > 0, P(X \geq a) = P(e^{sx} \geq e^{sa})$$

$$\leq \frac{E[e^{sx}]}{e^{sa}}$$

$$\leq \min_{s > 0} E[e^{sx}] \cdot e^{-sa}$$

Similarly for $s < 0$

$$P(X \leq a) \leq \min_{s < 0} E[e^{sx}] e^{-sa}$$

Example, $X \sim \text{Binomial}(n, p)$. Using Chernoff bounds bound $P(X \geq \alpha n)$ where $p < \alpha < 1$. Evaluate the bound for $p = \frac{1}{2}$ and $\alpha = \frac{3}{4}$.

$$X = \sum_{i=1}^n X_i, \quad P_{X_i}(1) = p = 1 - P_{X_i}(0),$$

X_i are i.i.d.

$$\begin{aligned} E[e^{sx}] &= m_X(s) \\ &= \prod_{i=1}^n m_{X_i}(s) \\ &= (pe^s + 1-p)^n. \end{aligned}$$

$$P(X \geq \alpha n) \leq \min_{s > 0} e^{-\alpha ns} (pe^s + 1-p)^n$$

$$\frac{d}{ds} (e^{-\alpha ns} (pe^s + 1-p)^n) =$$

find the s that minimizes the function

$$\Rightarrow e^s = \frac{\alpha(1-p)}{p(1-\alpha)}$$

$$\Rightarrow s = \log \frac{\alpha(1-p)}{p(1-\alpha)} > 0$$

Since

$$\frac{1-p}{p} \cdot \frac{\alpha}{1-\alpha} > 1,$$

Also check double derivative ≥ 0 .

We get

$$\begin{aligned} P(X \geq \alpha n) &\leq \left(\frac{\alpha(1-p)}{p(1-\alpha)} \right)^{-\alpha n} \left(\frac{p\alpha(1-p)}{p(1-\alpha)} + 1-p \right)^n \\ &= \left(\frac{\alpha(1-p)}{p(1-\alpha)} \right)^{-\alpha n} \left(\frac{1-p}{1-\alpha} \right)^n \\ &= \left(\frac{1-p}{1-\alpha} \right)^{(1-\alpha)n} \left(\frac{p}{1-\alpha} \right)^{\alpha n} \\ &= \left(\frac{1-\frac{1}{2}}{1-\frac{3}{4}} \right)^{\frac{n}{4}} \cdot \left(\frac{2}{3} \right)^{\frac{3n}{4}} \\ &= \frac{2^{\frac{n}{4}} \cdot 2^{\frac{3n}{4}}}{27^{\frac{n}{4}}} = \left(\frac{16}{27} \right)^{\frac{n}{4}}. \end{aligned}$$

Comparison between Markov, Chebyshev
and Chernoff bounds:

$$P(X \geq \alpha n) \leq \frac{2}{3} \quad [\text{Markov}]$$

$$P(X \geq \alpha n) \leq 4/n \quad [\text{Chebyshev}]$$

$$P(X \geq \alpha n) \leq \left(\frac{16}{27} \right)^{\frac{n}{4}} \quad [\text{Chernoff}]$$

- The bound given by Markov is the weakest bound. It is a constant (does not depend on n).
- Chebyshev's bound is stronger than Markov's. In particular note that $4/n \rightarrow 0$ as $n \rightarrow \infty$.
- Chernoff bound is the strongest bound. It goes to zero exponentially fast.

Example. Suppose X is a RV taking values in $[a, b]$. Obtain a bound on $P(|X - \mu| \geq c)$ using Chebyshev's inequality.

$$\begin{aligned} \text{claim: } X \in [a, b] \Rightarrow \sigma_x^2 &\leq \frac{(b-a)^2}{4}, \\ E[(X-\mu)^2] &= E[X^2] - 2E[X]\mu + \mu^2 \\ \sigma_x^2 &\leq E[(X - \frac{a+b}{2})^2] = E[((X-a) + (X-b))^2]/4 \\ &= \frac{1}{4} E \left[((X-a) - (X-b))^2 + 4(X-a)(X-b) \right] \\ &\leq \frac{(b-a)^2}{4}. \end{aligned}$$

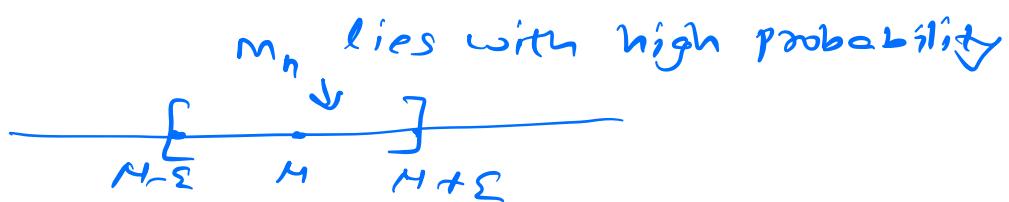
$$P(|X - \mu| \geq c) \leq \frac{\sigma_x^2}{c^2} \leq \frac{(b-a)^2}{4c^2}.$$

Weak Law of Large Numbers (WLLN)

Let x_1, x_2, \dots be a sequence of independent and identically distributed (i.i.d.) random variables with mean μ . For every $\varepsilon > 0$, we have

$$P\left(\left|\frac{\sum_{i=1}^n x_i}{n} - \mu\right| \geq \varepsilon\right) \rightarrow 0 \text{ as } n \rightarrow \infty.$$

Interpretation: Let $M_n = \frac{\sum_{i=1}^n x_i}{n}$. WLLN asserts that the sample mean of a large number of i.i.d. RVS is very close to the true mean.



Proof. $E[M_n] = \mu$, $\text{Var}(M_n) = \frac{\sigma^2}{n}$.

$$P(|M_n - E[M_n]| \geq \epsilon)$$

$$= P(|M_n - M| \geq \epsilon) \leq \frac{\text{Var}(M_n)}{\epsilon^2}$$

$$= \frac{\sigma^2}{n\epsilon^2} \rightarrow 0 \text{ as } n \rightarrow \infty,$$

Lecture 22

(6 November 2023)

WLLN. Let x_1, x_2, \dots be a sequence of independent and identically distributed RVs with mean μ . For every $\varepsilon > 0$ we have

$$P\left(\left|\frac{\sum_{i=1}^n x_i}{n} - \mu\right| \geq \varepsilon\right) \rightarrow 0 \text{ as } n \rightarrow \infty.$$

Remark. Note that the proof we have seen assumes finite variance. It turns out that this law remains true even if the x_i have infinite variance, but a much more elaborate argument is needed.

Another Interpretation:

Let a_1, a_2, \dots, a_n be the realizations of a sequence of i.i.d. random variables $\sim p_x$.

$$a_i \in X = \{x_1, x_2, \dots, x_k\}.$$

$$\sum_{i=1}^n a_i/n = \frac{\sum_{i=1}^K x_i k_i}{n} \rightarrow \sum_{i=1}^K k_i = n$$

k_i = no. of times x_i occurs

If we interpret probabilities as relative frequency

$$\sum_{i=1}^K \left(\frac{k_i}{n}\right) x_i \rightarrow \sum_{i=1}^K P_x(x_i) x_i = E[x]$$

Example (Polling). Each voter in a population selects a particular candidate A with probability p and chooses another with probability 1-p.

$$x_i = \begin{cases} 1 & \text{ith voter chooses A} \\ 0 & \text{o.w.} \end{cases}$$

$$P_{x_i}(1) = p = 1 - P_{x_i}(0)$$

x_i are independent and identically distributed. We are interested in knowing the value of p.

$$M_n = \sum_{i=1}^n x_i / n,$$

the mean here is p.

$$P(|M_n - p| \geq 0.1) \leq \frac{p(1-p)}{n(0.1)^2} \leq \frac{1}{400(0.1)^2} = 0.25 \quad \text{for } n=100$$

With a sample size of 100, the probability that our estimate is incorrect by more than 0.1 is smaller than 0.25.

$$P(|M_n - \mu| \geq 0.01) \leq \frac{1}{4n(0.01)^2}$$

$$n \geq 50000 \Rightarrow P(|M_n - \mu| \geq 0.01) \leq 0.05.$$

with a sample size of 100, the probability that our estimate is incorrect by more than 0.01 is smaller than 0.05

Convergence in Probability

We can interpret WLLN as stating that

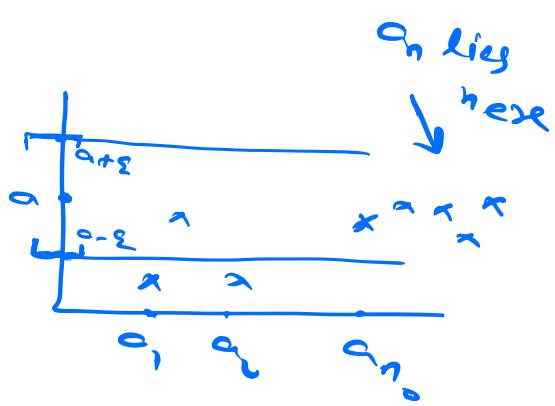
" $M_n = \frac{1}{n} \sum_{i=1}^n X_i$ converges to mean μ " in the

sense of "convergence in probability".

Recall the convergence of a deterministic sequence.

Let a_1, a_2, \dots be a sequence of real numbers and let a be another real. We say that a_n converges to a , or $\lim_{n \rightarrow \infty} a_n = a$, if for every $\epsilon > 0$ there exists some n_0 s.t.

$|Y_n - a| \leq \varepsilon$ for all $n \geq n_0$.



Convergence in Probability

Let Y_1, Y_2, \dots be a sequence of random variables, and let a be a real number. We say that Y_n converges to a in probability if for every $\varepsilon > 0$ we have

$$\lim_{n \rightarrow \infty} P(|Y_n - a| \geq \varepsilon) = 0.$$

- If RVs Y_1, Y_2, \dots have a PMF or a PDF and converge in probability to a , then "almost all" of the PMF or PDF of Y_n is concentrated within ε of a for large values of n .
- The above condition of convergence in probability can be equivalently written as follows: for every $\varepsilon > 0$ and for every $\delta > 0$ there exists some n_0 such that

$$P(|Y_n - a| \geq \varepsilon) \leq \delta \text{ for all } n \geq n_0.$$

$\varepsilon \rightarrow$ accuracy level, $\delta \rightarrow$ confidence level.

Example. Let x_1, x_2, \dots be a sequence of independent and uniformly distributed rvs in $[0, 1]$, and let $y_n = \min\{x_1, \dots, x_n\}$.

Note that $y_1 \geq y_2 \geq \dots$. In fact,

$y_n \rightarrow 0$ in probability.

$$\begin{aligned} P(|y_n - 0| \geq \varepsilon) &= P(x_1 \geq \varepsilon, \dots, x_n \geq \varepsilon) \\ &= (1 - \varepsilon)^n \rightarrow 0 \text{ as } n \rightarrow \infty \end{aligned}$$

for every $\varepsilon > 0$.

- The following example shows that if $y_n \rightarrow a$, then $E[y_n]$ may not converge to a .

Example. $P(y_n = y) = \begin{cases} 1/n, & \text{for } y = 0 \\ 1/n, & \text{for } y = n \\ 0, & \text{elsewhere} \end{cases}$

For every $\varepsilon \geq 0$ $P(|y_n - 0| \geq \varepsilon) = \frac{1}{n} \rightarrow 0$ as $n \rightarrow \infty$.

$E[y_n] = n \rightarrow \infty$ as $n \rightarrow \infty$,

i.e., $E[y_n] \not\rightarrow a$.

Central Limit Theorem (CLT)

Let x_1, x_2, \dots be a sequence of independent and identically distributed RVs with common mean μ and variance σ^2 .

Define $Z_n = \frac{\sum_{i=1}^n x_i - n\mu}{\sqrt{n}\sigma}$. Then

$$\begin{aligned} \lim_{n \rightarrow \infty} P(Z_n \leq z) &= \Phi(z) \\ &= \frac{1}{\sqrt{2\pi}} \int_{-\infty}^z e^{-x^2/2} dx. \end{aligned}$$

$$- S_n = \sum_{i=1}^n x_i \quad - \text{Var}(S_n) = n\sigma^2 \quad \rightarrow \infty \text{ as } n \rightarrow \infty$$

$$M_n = \sum_{i=1}^n x_i / n \quad - \text{Var}(M_n) = \frac{\sigma^2}{n} \quad \rightarrow 0 \text{ as } n \rightarrow \infty$$

$$Z_n = \sum_{i=1}^n x_i / \sqrt{n} \quad - \text{Var}(Z_n) = 1 \text{ independent of } n \quad (\text{constant})$$

Proof of CLT:

Fact, If $M_{Z_n}(s) \rightarrow M_Z(s)$ then

$$F_{Z_n}(z) \rightarrow F_Z(z) \text{ for all } z.$$

Proof of the fact is related to the continuity of inverse Fourier transform,

We assume that $M_X(s)$ is finite when $-d < s < d$. Let x_i has zero mean & variance 1.

$$Z_n = \frac{1}{\sqrt{n}} \sum_{i=1}^n x_i / \sqrt{n}.$$

$$M_{Z_n}(s) = E \left[e^{s \sum_{i=1}^n x_i / \sqrt{n}} \right]$$

$$= \prod_{i=1}^n E \left[e^{s x_i / \sqrt{n}} \right] = M_X(s/\sqrt{n})^n.$$

$$\text{Let } \angle(t) = \log M_X(t).$$

$$\angle(0) = 0, \quad \angle'(0) = \frac{m_x'(0)}{m_x(0)} = E[x] = 0,$$

$$\angle''(0) = \frac{m_x(0)m_x''(0) - m_x'(0)^2}{m_x(0)^2} = E[x^2] = 1,$$

we first show that

$$m_x(s/\sqrt{n}) \xrightarrow{\text{MGF of } N(0,1)} e^{t^2/2}$$

or equivalently that

$$n\angle(s/\sqrt{n}) \rightarrow t^2/2 \text{ as } n \rightarrow \infty.$$

$$\lim_{n \rightarrow \infty} \frac{\angle(s/\sqrt{n})}{n^{-1}} = \lim_{n \rightarrow \infty} \frac{-\angle'(s/\sqrt{n}) \cdot s}{-\frac{2}{n^2}} \quad (\text{by L'Hopital Rule})$$

$$= \lim_{n \rightarrow \infty} \frac{\angle'(s/\sqrt{n}) \cdot s}{2n^{-3/2}}$$

$$= \lim_{n \rightarrow \infty} \frac{-\angle''(s/\sqrt{n}) s^2}{-2n^{-3/2}} \quad (\text{by L'Hopital Rule})$$

$$= \lim_{n \rightarrow \infty} \frac{\angle''(s/\sqrt{n}) s^2}{2} = \frac{s^2}{2}, \quad (\because \angle''(0) = 1)$$

If x_i 's are of mean μ & variance σ^2

$$\sum_{i=1}^n \frac{x_i - \mu}{\sqrt{n}\sigma} = \frac{\sum_{i=1}^n \frac{x_i - \mu}{\sigma}}{\sqrt{n}} \rightarrow \begin{matrix} \text{mean } 0 \\ \text{var} = 1 \end{matrix}$$

Fact. If $m_{Z_n}(t) \rightarrow m_Z(t)$ then

$$F_{Z_n}(z) \rightarrow F_Z(z) \text{ for all } z.$$

Proof is related inverse Fourier transform
and its continuity.

This completes the proof of the central limit theorem.

Lecture 23

(9 November 2023)

Central Limit Theorem.

Let x_1, x_2, \dots be a sequence of i.i.d. random variables with common mean μ and variance σ^2 . Define

$$Z_n = \frac{\sum_{i=1}^n x_i - n\mu}{\sigma\sqrt{n}}. \quad \text{Then, the CDF of}$$

Z_n converges to the standard normal CDF in the sense that

$$\lim_{n \rightarrow \infty} P(Z_n \leq z) = \Phi(z)$$

$$= \frac{1}{\sqrt{2\pi}} \int_{-\infty}^z e^{-x^2/2} dx,$$

for every z .

Proof. we prove this for $\mu=0$ & $\sigma^2=1$,
i.e., $Z_n = \frac{\sum_{i=1}^n x_i}{\sqrt{n}}$.

$$\begin{aligned}
 M_{Z_n}(s) &= E[e^{sZ_n}] \\
 &= E\left[e^{s\sum_{i=1}^n X_i/\sqrt{n}}\right] \\
 &= \prod_{i=1}^n M_X(s/\sqrt{n}) = (M_X(s/\sqrt{n}))^n.
 \end{aligned}$$

Let $\angle(s) = \log M_X(s)$.

$$\log M_{Z_n}(s) = n \log M_X(s/\sqrt{n}) = n \angle(s/\sqrt{n}).$$

we have $\angle(0) = \angle'(0) = 0$ & $\angle''(0) = 1$,

$$\text{Consider } \lim_{n \rightarrow \infty} n \log M_X(s/\sqrt{n})$$

$$= \lim_{n \rightarrow \infty} \frac{\log M_X(s/\sqrt{n})}{n^{-1}}$$

$$= \lim_{n \rightarrow \infty} \frac{M_X'(s/\sqrt{n}) s n^{-3/2}}{2 M_X(s/\sqrt{n}) n^{-2}}$$

$$= \lim_{n \rightarrow \infty} \frac{\angle'(s/\sqrt{n}) s}{2 n^{-1/2}}$$

$$= \lim_{n \rightarrow \infty} \frac{\angle''(s/\sqrt{n}) s^2 n^{-3/2}}{2 n^{-3/2}} = s^2/2.$$

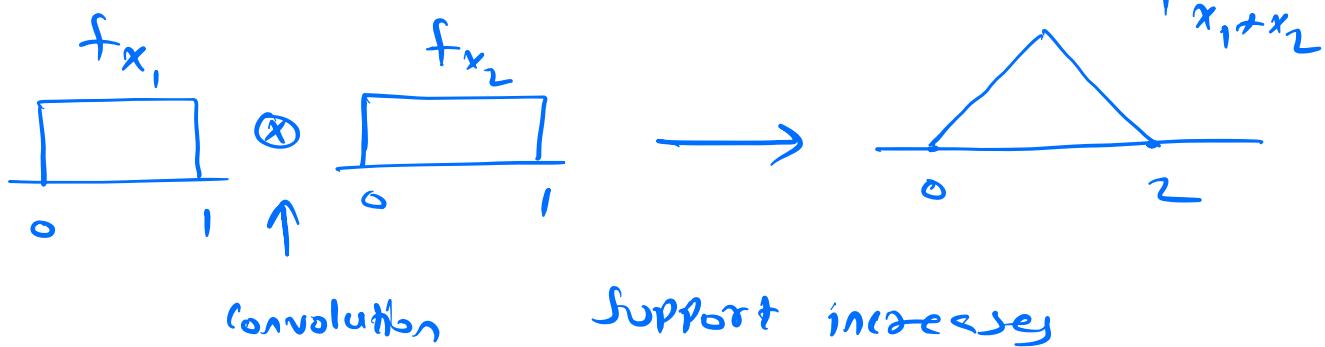
Examples. (i) $X_i \sim N(\mu, \sigma^2)$

$$\frac{\sum_{i=1}^n X_i - n\mu}{\sqrt{n}\sigma} \sim N(0, 1) \text{ for all } n \in \mathbb{N},$$

$$\frac{X_1 - \mu}{\sigma}, \frac{X_1 + X_2 - 2\mu}{\sqrt{2}\sigma}, \dots \sim N(0, 1),$$

$$\lim_{n \rightarrow \infty} P(Z_n \leq z) = P(Z \leq z), \quad Z \sim N(0, 1).$$

(ii) $X_i \sim \text{uniform}[0, 1]$.



Intuitively as $n \rightarrow \infty$, support will be real line
and $\sum_{i=1}^n X_i \rightarrow \text{Gaussian RV.}$

Normal Approximation Based on CLT:

Let $S_n = \sum_{i=1}^n x_i$, where x_i are i.i.d.

with mean μ and variance σ^2 . We are interested in computing $P(S_n \leq c)$. If n is large, it can be approximated in the following way.

$$P(S_n \leq c) = P\left(\frac{S_n - n\mu}{\sqrt{n}\sigma} \leq \frac{c - n\mu}{\sqrt{n}\sigma}\right)$$

$$\approx \Phi(z) \text{ where } z = \frac{c - n\mu}{\sqrt{n}\sigma}.$$

Example, X_i , $i \in [1:100]$ are i.i.d. and uniformly distributed in the interval $[5, 50]$. Find $P(S_{100} > 3000)$, $S_{100} = \sum_{i=1}^{100} X_i$. It is not easy to calculate CDF and the desired probability exactly, but an approximate answer can be quickly obtained using CLT.

$$\mu = \frac{5+50}{2} = 27.5, \quad \sigma^2 = \frac{(50-5)^2}{12} = 168.75.$$

$$P(S_{100} > 3000) = 1 - P(S_{100} \leq 3000)$$

$$= 1 - P\left(\frac{S_{100} - n\mu}{\sqrt{n}\sigma} \leq \frac{3000 - 275}{10(168.75)}\right)$$

$$= 1 - P\left(\frac{S_{100} - n\mu}{\sqrt{n}\sigma} \leq 1.92\right)$$

$$= 1 - \Phi(1.92)$$

$$\approx 1 - 0.9726$$

(using standard Normal table)

$$= 0.0274.$$

Example (Polling). $x_i = \begin{cases} 1 & \text{if } i\text{th voter votes for A} \\ 0 & \text{otherwise} \end{cases}$

$$E[x_i] = p \quad M_n = \sum_{i=1}^n x_i / n.$$

WLLN (or Chebyshev's inequality):

For $n=100$, we got that

$$P(|M_n - p| \geq 0.1) \leq 0.25.$$

We will see now if CLT can improve this bound.

$$P\left(\left|\frac{\sum_{i=1}^n x_i}{n} - p\right| \geq 0.1\right)$$

$$= P\left(\left|\frac{\sum_{i=1}^n x_i - np}{n}\right| \geq 0.1n\right)$$

$$= P\left(\left|\frac{\sum_{i=1}^n x_i - np}{\sqrt{n}\sigma}\right| \geq \frac{0.1\sqrt{n}}{\sigma}\right)$$

$$= P\left(\frac{\sum_{i=1}^n x_i - np}{\sqrt{n}\sigma} \geq \frac{0.1\sqrt{n}}{\sigma}\right) + P\left(\frac{\sum_{i=1}^n x_i - np}{\sqrt{n}\sigma} \leq -\frac{0.1\sqrt{n}}{\sigma}\right)$$

$$\approx 2P\left(Z \geq \frac{0.1\sqrt{n}}{\sigma}\right) \sim N(0,1)$$

$$\leq 2P\left(Z \geq 0.2\sqrt{n}\right) \text{ as } \sigma \leq \frac{1}{2}$$

$$= 2 - 2P(Z \leq 0.2\sqrt{n})$$

$$= 2 - 2\Phi(-0.2\sqrt{100}) = 2 - 2\Phi(-2) = 2 - 2 \times 0.977 = 0.046.$$

This is a much better estimate than 0.25 from WLLN.

Strong Law of Large Numbers

- SLLN also deals with the convergence of the sample mean to the true mean. However, SLLN refers to a different type of convergence.

WLLN

$$M_n = \frac{1}{n} \sum_{i=1}^n x_i \rightarrow E[x] \text{ in probability}$$

SLLN

$$M_n = \frac{1}{n} \sum_{i=1}^n x_i \rightarrow E[x] \text{ with probability 1,}$$

(or almost sure
convergence)

Almost Sure convergence:

Let x_1, x_2, \dots be a sequence of RVS and c be a real number. Then we say

$$x_n \xrightarrow{\text{a.s.}} c \quad \text{if } P(\{\omega \in \Omega : \lim_{n \rightarrow \infty} x_n(\omega) = c\}) = 1.$$

Example. Let $\Omega = [0, 1]$, consider a probability law defined by $P([a, b]) = b - a$ for all $0 \leq a \leq b \leq 1$.

Define $X_n(\omega) = \omega^n$ for $n \in \mathbb{N}$,

Note that $\lim_{n \rightarrow \infty} X_n(\omega) = \begin{cases} 0 & \text{if } 0 \leq \omega < 1 \\ 1 & \text{if } \omega = 1 \end{cases}$
 always find this first

$$P\left(\{\omega : \lim_{n \rightarrow \infty} X_n(\omega) = 0\}\right) = P([0, 1)) = 1$$

Since the singleton set $\{1\}$ has zero probability,

$$\therefore X_n \xrightarrow{\text{a.s.}} 0.$$

SLLN. Let x_1, x_2, \dots be a sequence of i.i.d. RVs with mean μ . Then,

$$M_n = \frac{\sum_{i=1}^n x_i}{n} \xrightarrow{\text{a.s.}} \mu, \quad \text{i.e.,}$$

$$P\left(\{\omega : \lim_{n \rightarrow \infty} M_n(\omega) = \mu\}\right) = 1.$$

Proof. Assume that $E[x_i^4] = k < \infty$.

$$S_n = \sum_{i=1}^n x_i. \quad E[S_n^4] = E\left[\left(\sum_{i=1}^n x_i\right)^4\right].$$

This will have terms of the form

$$x_i^4, x_i^3 x_j, x_i^2 x_j^2, x_i^2 x_j x_k, x_i x_j x_k x_l$$

where i, j, k, l are different.

Assume $n=0$. Then because of independence it follows that

$$E[x_i^3 x_j] = E[x_i^3] E[x_j] = 0$$

$$E[x_i^2 x_j x_k] = E[x_i^2] E[x_j] E[x_k] = 0$$

$$E[x_i x_j x_k x_l] = E[x_i] E[x_j] E[x_k] E[x_l] = 0.$$

$$\begin{aligned} \text{so } E[S_n^4] &= n E[x_i^4] + 6 \binom{n}{2} E[x_i^2 x_j^2] \\ &= n E[x_i^4] + 3n(n-1) E[x_i^2] E[x_j^2] \\ &\qquad\qquad\qquad \underbrace{\qquad\qquad\qquad}_{= (E[x_i^2])^2 \leq E[x_i^4]} \\ &\leq nk + 3n(n-1)k \stackrel{\text{as } \text{var}(x_i^2) \geq 0}{\leq} 3n^2 k \\ \Rightarrow E[S_n^4/n^4] &\leq 3k/n^2 \end{aligned}$$

$$\Rightarrow E\left[\sum_{n=1}^{\infty} \frac{s_n^4}{n^4}\right] = \sum_{n=1}^{\infty} E[s_n^4/n^4]$$

$$= \sum_{n=1}^{\infty} \frac{1}{n^2} \cdot (3K) < \infty.$$

This implies that

$$P\left(\sum_{n=1}^{\infty} \frac{s_n^4}{n^4} < \infty\right) = 1$$

$$\sum_{n=1}^{\infty} \frac{s_n^4}{n^4} < \infty \Rightarrow \lim_{n \rightarrow \infty} \frac{s_n^4}{n^4} = 0$$

$$\Rightarrow 1 = P\left(\sum_{n=1}^{\infty} \frac{s_n^4}{n^4} < \infty\right) \leq P\left(\lim_{n \rightarrow \infty} \frac{s_n^4}{n^4} = 0\right)$$

$$\Rightarrow P\left(\lim_{n \rightarrow \infty} \frac{s_n}{n} = 0\right) = 1.$$

In the above proof we have used

$$E\left[\sum_{i=1}^{\infty} z_i\right] = \sum_{i=1}^{\infty} E[z_i], \text{ This is not}$$

necessarily true always for any z_i 's.

However this holds true when all z_i are non-negative random variables. This is because of monotone convergence theorem (not covered in this course).

Lecture 24

(16 November 2023)

Random Processes

- A random or stochastic process is a mathematical model of a probabilistic experiment that evolves in time and generates a sequence of numerical values.
- Each numerical value in the sequence is modeled by a random variable so a random process is simply a (finite or infinite) sequence of random variables.

Recall that a RV $X : \Omega \rightarrow \mathbb{R}$.

Formally, a random process is a family of random variables $(X_t : t \in T)$, all on the same probabilistic model (Ω, P) .

In many applications, T is a set of times. If $T = \mathbb{N}$ then $(X_t : t \in T)$ is called a discrete-time process. If $T = \mathbb{R}$, then continuous-time process.

- For each $t \in T$, x_t is a RV.

$$\omega \mapsto x_1(\omega) \ x_2(\omega) \ \dots$$

Discrete-time

$$\omega \mapsto x_t(\omega) \quad t \in T.$$

continuous-time

For a fixed $\omega \in \Omega$, $(x_t(\omega), t \in T)$ is called the sample path at ω .

Mean Function:

$$M_x(t) = E[x_t] \text{ for } t \in T,$$

Covariance Function:

$$C_x(t_1, t_2) = \text{Cov}(x_{t_1}, x_{t_2})$$

$$= E[x_{t_1} x_{t_2}] - E[x_{t_1}] E[x_{t_2}]$$

Correlation function:

$$R_x(t_1, t_2) = E[x_{t_1} x_{t_2}].$$

- A random process is statistically specified by its complete set of n th order distribution function for all $n \in \mathbb{N}$.

$$F_x(x_1, \dots, x_n; t_1, t_2, \dots, t_n) = P(X_{t_1} \leq x_1, \dots, X_{t_n} \leq x_n).$$

For the continuous case we can obtain the PDF as

$$f(x_1, x_2; t_1, t_2) = \frac{\partial^2 F(x_1, x_2; t_1, t_2)}{\partial x_1 \partial x_2},$$

Arrival - Type Processes we are interested in occurrences that have the character of an arrival such as message receptions at a receiver, customer purchases at a store etc. we will focus on models in which the interarrival times (the times between successive arrivals) are independent random variables.

Bernoulli process - case where arrivals occur in discrete time and the interarrival times are geometrically distributed

Poisson process - case where arrivals occur in continuous time and the interarrival times are exponentially distributed.

The Bernoulli process

Consider a sequence of independent coin tosses with the probability of heady p in the range $0 < p < 1$.

- A Bernoulli process is a sequence x_1, x_2, \dots of independent Bernoulli Rvs x_i with
 - $P(x_i = 1) = P(\text{success at the } i\text{th trial})$
 - $P(x_i = 0) = P(\text{failure at the } i\text{th trial})$
- No. of arrivals within a certain time period or no. of successes in n independent trials:

$$P_s(k) = \binom{n}{k} p^k (1-p)^{n-k} \quad k=0, 1, \dots, n$$

$$E[S] = np, \quad \text{var}(S) = np(1-p).$$

This is Binomial with parameters n & p .

- The time until the first arrival or no. of trials up to (and including) the first success :

$$P_T(t) = (1-p)^{t-1} p, \quad t=1, 2, \dots$$

$$E[T] = \frac{1}{p}, \quad \text{var}(T) = \frac{1-p}{p^2}.$$

Independence & Memorylessness

Due to independence property in the Bernoulli process — whatever has happened in the past trials provides no information on the outcome of future trials.

Fresh-start property :

For any given n , consider

$x_1, x_2, \dots, x_n, x_{n+1}, x_{n+2}, \dots$, we notice that

$x_i = x_{n+i}$ are independent Bernoulli trials and therefore form a Bernoulli process,

Memolessness property: T = time until k^{th} success

$$P(T-n=t | T>n)$$

$$= \frac{P(T-n=t - T>n)}{P(T>n)}$$

$$= \frac{P(T=n+t)}{P(T>n)} = \frac{(1-p)^{n+t-1} \cdot p}{(1-p)^n}$$

$$= (1-p)^{t-1} p = P(T=t),$$

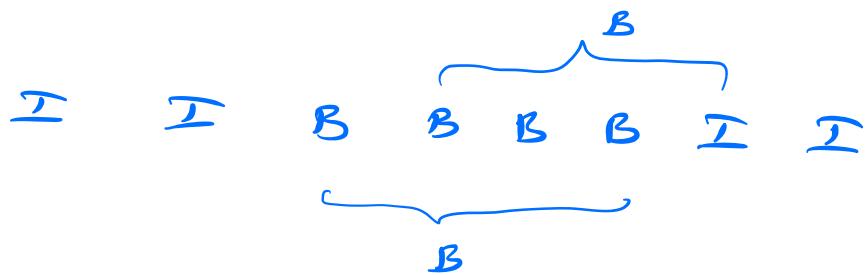
- we call a slot *busy* (B) if there is an arrival, and otherwise let us call it *idle*, we call a string of idle slots flanked by busy slots, an 'idle period'. Similarly, we can define 'busy period'.

Let us derive the PMF, mean, and variance of the following random variables.

T = the time index of the first idle slot

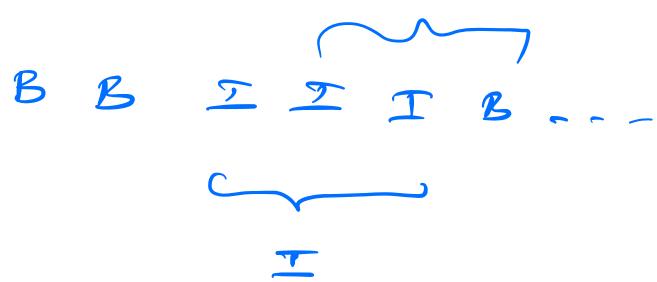
$$P_T(k) = p^{k-1} (1-p) \quad k=1, 2, \dots$$

\angle = the length (no. of slots) of the first busy period



$$P_B(k) = p^{k-1} (1-p) \quad k=1, 2, \dots$$

Σ = the length of the first idle period



$$P_I(k) = (1-p)^{k-1} p \quad k=1, 2, \dots$$

Example. Let N be the first time that we have a success immediately following a previous success. That is N is the first time i for which $x_{i-1} = x_i = 1$. What is the probability that there are no successes in the two trials that follow, i.e., $P(x_{N+1} = x_{N+2} = 0)$?

$$P(X_{N+1} = x_{N+2} = 0) \\ = \sum_{n=1}^{\infty} P(X_{N+1} = x_{N+2} = 0 \mid N=n) P(N=n)$$

$$P(X_{N+1} = x_{N+2} = 0 \mid N=n) \\ = P(X_{n+1} = x_{n+2} = 0) = (1-p)^2$$

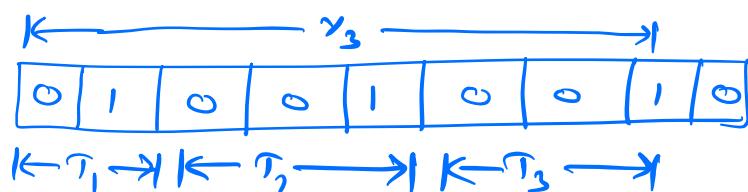
Equality because $N=n$ is determined by some function of (x_1, x_2, \dots, x_n) .

Interarrival Times

Let y_k denote the time of the k^{th} success (or arrival).

$$T_1 = y_1 \quad T_k = y_k - y_{k-1}, \quad k=2, 3, \dots$$

T_k represents the k^{th} interarrival time i.e., no. of trials following the $(k-1)^{\text{st}}$ success until the next success.



We have $y_k = \sum_{i=1}^k T_i$.

Interarrival times T_1, T_2, T_3, \dots are independent and all have the same geometric distribution.

Properties of the k^{th} Arrival Time

$$y_k = T_1 + T_2 + \dots + T_k$$

Each T_i is a geometric RV with parameter p and all T_i 's are independent.

$$E[y_k] = \sum_{i=1}^k E[T_i] = k/p,$$

$$\text{Var}(y_k) = \sum_{i=1}^k \text{Var}(T_i) = \frac{k(1-p)}{p^2},$$

To find the pmf of y_k

$$P(y_k = t) = P(A \cap B), \text{ where}$$

$$A = \left\{ \text{trial } t \text{ is a success} \right\}$$

$$B = \left\{ \text{exactly } k-1 \text{ successes in } t-1 \text{ trials} \right\}$$

$$P_{Y_k}(t) = P(A), P(B)$$

$$= p \cdot \binom{t-1}{k-1} p^{k-1} (1-p)^{t-k}$$

$$= \binom{t-1}{k-1} p^k (1-p)^{t-k},$$



Known as Pascal PMF of order k.

Poisson Process

- A counting process counts the number of arrivals, e.g., no. of customers who arrive at a restaurant — $X(t) = \text{no. of arrivals in time } t$.
- A counting process has independent and stationary increments. $[X(0)=0]$

Independent increments:

For all $0 \leq t_1 < t_2 < t_3 < \dots < t_n$ the RVs

$X(t_2) - X(t_1)$, $X(t_3) - X(t_2)$, ... are independent.

Note that $X(t_i) - X(t_{i-1})$ represents the no. of arrivals in the interval $[t_{i-1}, t_i]$.

Stationary increments:

For all $t_2 > t_1 \geq 0$ & $T > 0$, the RVs

$X(t_2) - X(t_1)$ and $X(t_2 + T) - X(t_1 + T)$ have the same distribution.

- Thus a counting process has independent increments if the no. of arrivals in non-overlapping intervals are independent.

Also it has stationary increments if, for all $t_2 > t_1 \geq 0$, $X(t_2) - X(t_1)$ has the same distribution as $X(t_2 - t_1)$.

- Before we define Poisson process we recall Poisson RV.

$X \sim \text{Poisson}(\lambda)$ means $P_X(k) = \frac{\lambda^k e^{-\lambda}}{k!}$ - $k=0, 1, 2, \dots$

- A counting process $(N(t), t \in [0, \infty))$ is called a Poisson process with rate λ if:

- (1) $N(0) = 0$,
- (2) $N(t)$ has independent increments,
- (3) the no. of arrivals in any interval of length $T > 0$ has Poisson (λT) distribution.

$$E[N(t)] = \lambda t, \quad \text{var}(N(t)) = \lambda t.$$

$$\begin{aligned}
 R_X(t_1, t_2) &= E[N(t_1)N(t_2)] \quad (t_1 < t_2) \\
 &= E[N(t_1)(N(t_2) - N(t_1) + N(t_1))] \\
 &= E[N(t_1)(N(t_2) - N(t_1))] + E[N^2(t_1)] \\
 &= E[N(t_1)] E[N(t_2 - t_1)] + E[N^2(t_1)] \\
 &= \lambda t_1 \cdot \lambda(t_2 - t_1) + \lambda t_1 + \lambda^2 t_1^2 \\
 &= \lambda^2 t_1 t_2 - \cancel{\lambda^2 t_1^2} + \lambda t_1 - \cancel{\lambda^2 t_1^2} \\
 &= \lambda t_1 + \lambda^2 t_1 t_2.
 \end{aligned}$$

$$C(t_1, t_2) = \lambda \min\{t_1, t_2\},$$

Lecture 25

(18 November 2023)

- A random process $X = (X_t; t \in T)$ is specified by n^{th} -order distribution function,

$$F_{X_{t_1}, \dots, X_{t_n}}(x_1, x_2, \dots, x_n)$$

$$= P(X_{t_1} \leq x_1, \dots, X_{t_n} \leq x_n)$$

for all $n \in \mathbb{N}$,

Example. Consider the random process

$$X(t) = A \sin(\omega_0 t + \phi)$$

where A and ϕ are independent and ϕ is uniformly distributed over $[-\pi, \pi]$.

$$M_x(t) = E[A \sin(\omega_0 t + \phi)]$$

$$= M_A \cdot E[\sin(\omega_0 t + \phi)]$$

$$= M_A \cdot \int_{-\pi}^{\pi} \frac{1}{2\pi} \cdot \sin(\omega_0 t + \phi) d\phi$$

$$\begin{aligned}
 R_x(t_1, t_2) &= E[x(t_1)x(t_2)] \\
 &= E[A^2 \sin(\omega_0 t_1 + \phi) \sin(\omega_0 t_2 + \phi)] \\
 &= E[A^2] E[\sin(\omega_0 t_1 + \phi) \sin(\omega_0 t_2 + \phi)] \\
 &= E[A^2] \frac{1}{2} \left(E[\cos(\omega_0(t_1 - t_2))] - E[\cos(\omega_0(t_1 + t_2) + 2\phi)] \right) \\
 &= \frac{1}{2} E[A^2] \underbrace{\left[\cos(\omega_0(t_1 - t_2)) \right]}_{=0}.
 \end{aligned}$$

Stationary Process

A random process $(X(t), t \in \mathbb{R})$ is stationary if, for all $t_1, t_2, \dots, t_n, T \in \mathbb{R}$,

$$F_{\underline{x(t_1)} \underline{x(t_2)} \dots \underline{x(t_n)}}^{(\underline{x_1} \underline{x_2} \dots \underline{x_n})} = F_{\underline{x(t_1+T)} \dots \underline{x(t_n+T)}}^{(\underline{x_1} \underline{x_2} \dots \underline{x_n})}$$

$\forall n \in \mathbb{N},$

Exercise. Write down a similar definition for discrete-time random process

Example. Consider i.i.d. discrete-time random process x_1, x_2, \dots , Is this stationary?
Yes!

$(x_{n_1}, x_{n_2}, \dots, x_{n_8})$ has same distribution as
 $(x_{n_1+T}, \dots, x_{n_8+T})$.

— If a process is stationary, the analysis is usually simpler as the probabilistic properties do not change by time.

However, it turns out that not many real-life processes are stationary.
Even if a process is stationary, it might be difficult to prove it. Fortunately, often a weaker notion of stationarity suffices.

Wide-Sense stationary (WSS)

A random process $x(t)$ is called WSS if for all $t, t_1, t_2 \in \mathbb{R}, T \in \mathbb{R}$,

$$\mu_x(t) = \mu_x(t+T)$$

$$R_x(t_1, t_2) = R_x(t_1+T, t_2+T).$$

so $\mu_x(t)$ is a constant and $R_x(t_1, t_2)$ is a function of $t_2 - t_1$.

Example. $x(t) = A \sin(\omega_0 t + \theta)$ — A, θ are independent & \sim uniform $[-\pi, \pi]$.

$x(t)$ is WSS because

$$\mu_x(t) = 0 \text{ (a constant)}$$

$$\begin{aligned} R_x(t_1, t_2) &= \frac{1}{2} E[A^2] \cos(\omega_0(t_2 - t_1)) \\ &= R_x(t_1, -t_2). \end{aligned}$$

- For WSS random process it suffices to denote autocorrelation by $R_x(T)$.

Exercise. Prove that a stationary process is WSS.

Properties of $R_x(\tau)$

- $R_x(0) = E[x(t)x(t)] = E[x^2(t)] \geq 0$,
 $E[x^2(t)]$ is called expected (or average) power in $x(t)$ at time t . For WSS this is not a function of time.
- $R_x(-\tau) = E[x(t+\tau)x(t)]$
= $E[x(t)x(t+\tau)]$
= $R_x(\tau)$.
- $|R_x(\tau)| \leq R_x(0)$ for all $\tau \in \mathbb{R}$.
This follows from Cauchy-Schwarz's inequality.

$$|E[xy]| \leq \sqrt{E[x^2]E[y^2]}.$$

Consider $E[(x-\alpha y)^2] \geq 0$

$$\Rightarrow E[x^2 + \alpha^2 y^2 - 2\alpha xy] \geq 0$$
$$\Rightarrow \alpha^2 E[y^2] - 2\alpha E[xy] + E[x^2] \geq 0$$

So discriminant ≤ 0

$$\Rightarrow 4E^2[xy] \leq 4E[x^2]E[y^2]$$

$$\Rightarrow |E[xx^T]| \leq \sqrt{E[x^2]E[g^2]}.$$

$$|E[x(t)x(t+\tau)]| \leq \sqrt{E[\tilde{x}(t)]E[\tilde{x}(t+\tau)]}$$

$$\Rightarrow |R_x(\tau)| \leq \sqrt{R_x(0) \cdot R_x(0)} = R_x(0).$$

Power Spectral Density (PSD)

Power spectral density is the Fourier transform of $R_x(\tau)$, mathematically

$$S_x(f) = \int_{-\infty}^{\infty} R_x(\tau) e^{-j2\pi f \tau} d\tau, \quad j = \sqrt{-1}$$

is PSD of $x(t)$.

It can be shown that the inverse Fourier transform of $S(f)$ is $R_x(\tau)$.

$$\int_{-\infty}^{\infty} S_x(f) e^{j2\pi f \tau} df = R_x(\tau).$$

Properties of PSD

- $S_x(-f) = S_x(f)$.

This is because $R_x(\tau)$ is even.

- $E[x(t)^2] = R_x(0) = \int_{-\infty}^{\infty} S_x(f) df$.

Inverse Fourier Transform

so we get expected or average power in $x(t)$ by integrating the PSD of $x(t)$.
This is why $S_x(f)$ is called power spectral density.

- $S_x(f) \geq 0$.

(Proof omitted)

Lecture 26

(20 November 2023)

Q) (a) Is an event A independent of itself? No in general.

$$P(A \cap A) = P(A)^2 \Rightarrow P(A) = 0 \text{ or } 1.$$

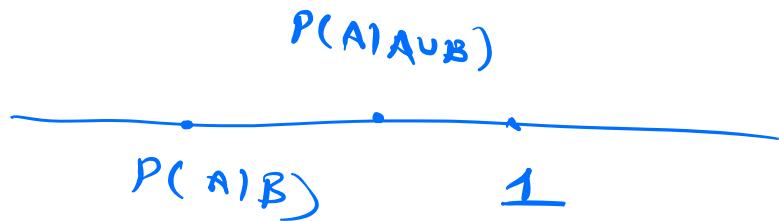
(b) Find the probability that exactly one of the events A or B occurs in terms of $P(A)$, $P(B)$, $P(A \cap B)$.

$$\begin{aligned} P((A \cap B^c) \cup (A^c \cap B)) &= P(A \cap B^c) + P(A^c \cap B) \\ &= P(A) + P(B) - 2P(A \cap B). \end{aligned}$$

(c) Is it true that $P(A|A \cup B) \geq P(A|B)$?

$$\begin{aligned} P(A|A \cup B) &= P(A|(A \cup B) \cap B) P(B|A \cup B) + \\ &\quad P(A|(A \cup B) \cap B^c) P(B^c|A \cup B) \\ &= P(A|B) P(B|A \cup B) + P(A|A \cap B^c) P(B^c|A \cup B) \\ &\quad \underbrace{\qquad\qquad}_{=1} \end{aligned}$$

$$= \neg P(A|B) + (1-\gamma),$$



$$\therefore P(A|A \cup B) \geq P(A|B),$$

Q) If X is a positive integer valued RV that satisfies memorylessness property

$$P(X > m+n | X > m) = P(X > n), \text{ for any } m, n \in \mathbb{N}.$$

Then prove that X is a geometric random variable.

Proof. Let $P(X > n) = a_n$ for $n \in \mathbb{N}$,

The memorylessness property gives

$$\frac{a_{m+n}}{a_m} = a_n$$

$$\Rightarrow a_{m+n} = a_m a_n, \forall m, n \in \mathbb{N}$$

$$\Rightarrow a_{m+1} = a_m a_1 = a_1^{m+1}$$

where $a_1 = P(X > 1) = 1 - P(X = 1) \triangleq 1 - p$,
 $P(X = n) = a_{n+1} - a_n = (1-p)^{n+1} - (1-p)^n = (1-p)^{n-1}p$.

Q) Let X be a discrete RV, Y be a continuous RV, and I is a binary RV s.t. X is independent of I , Y is independent of I .

Define

$$Z = \begin{cases} X & \text{if } I=1 \\ Y & \text{if } I=0 \end{cases}$$

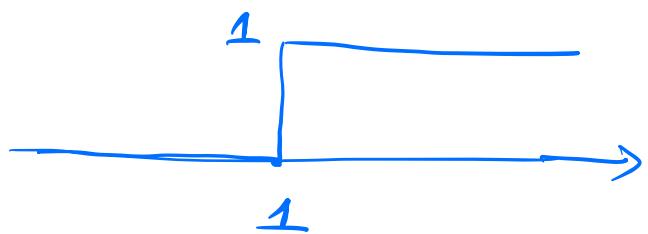
$$Z = \begin{cases} 1 & \text{if } I=1 \\ \text{Uniform } [0,2] & \text{if } I=0 \end{cases}$$

Z is neither a discrete nor a continuous RV.

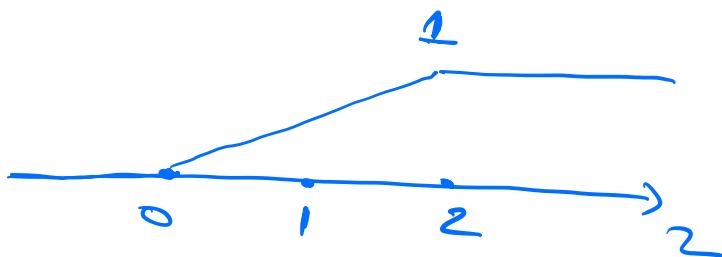
$$\begin{aligned} F_Z(z) &= P(Z \leq z) \\ &= P(Z \leq z | I=1)P(I=1) + P(Z \leq z | I=0) \\ &\quad P(I=0) \\ &= P F_X(z) + (1-p) F_Y(z) \\ &\quad (P = P(I=1)) \end{aligned}$$

Z is a mixed random variable,

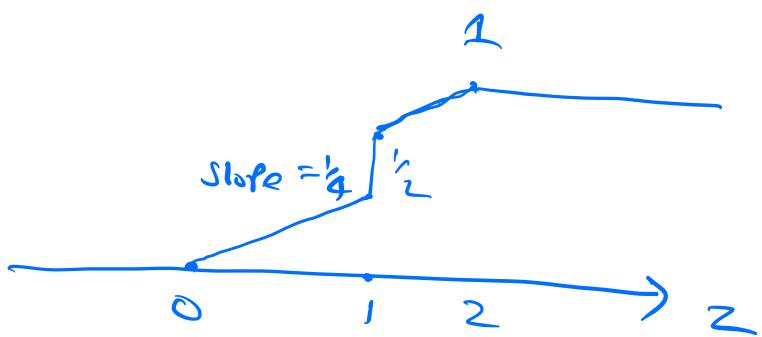
$$F_x(z)$$



$$F_y(z)$$



$$\text{For } p = \frac{1}{2} \quad f_z(2) =$$



Application of Probability in Information

Leakage,

The "information content" of a message depends on the degree to which the content of the message is surprising. If a highly likely event occurs, the message carries very little information. If a highly

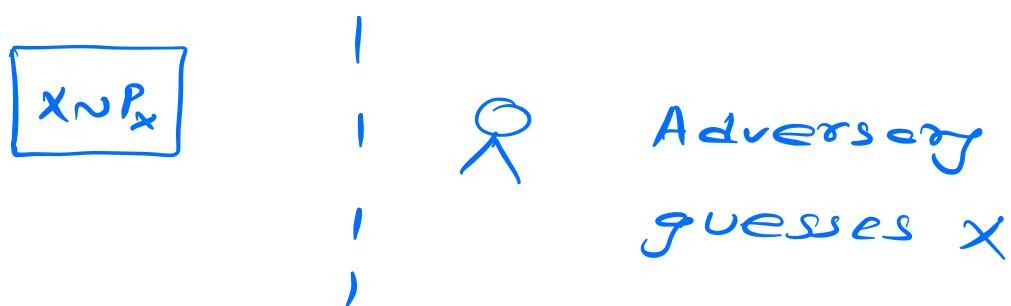
unlikely event occurs, the message is more informative.

The information content in an event with probability $p = \log \frac{1}{p}$.

[all logarithms are to the base 2]

Vulnerability: Suppose $x \sim p_x$. The vulnerability of x is given by

$$v(x) = \max_{x \in X} p_x(x).$$



$v(x)$ is the worst-case probability that an adversary could guess the value of x correctly.

$\log \frac{1}{v(x)}$ can be viewed as information measure

$$\text{min-entropy of } x - H_\infty(x) = \log \frac{1}{v(x)}$$

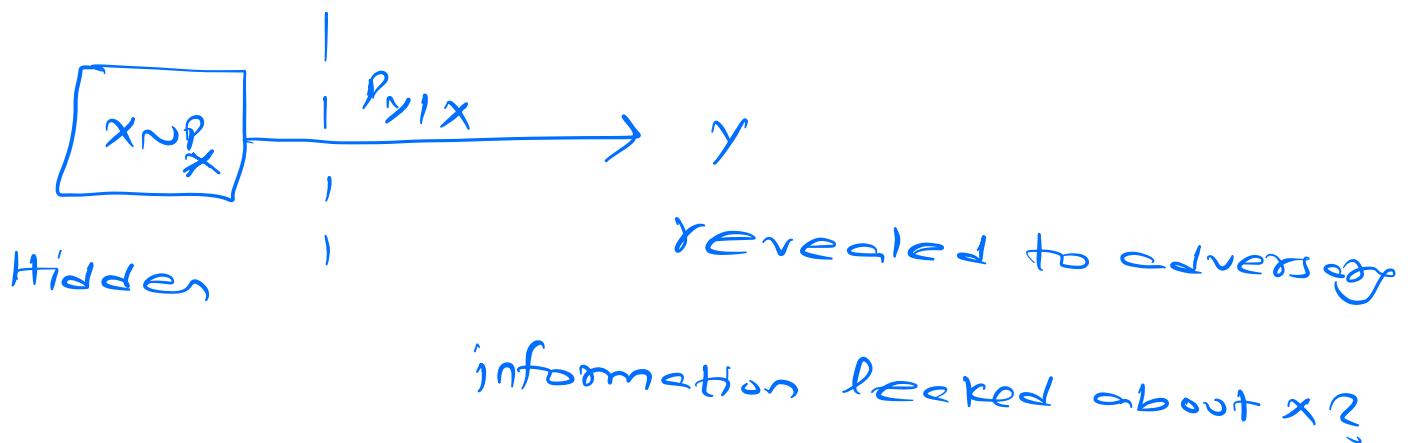
$$= \log \frac{1}{\max_x p_x(x)}.$$

- Given two RVS, $(x, y) \sim P_{xy}$, the conditional vulnerability

$$v(x|y) = \sum_{y \in Y} p_y(y) v(x|y=y)$$

$$= \sum_{y \in Y} p_y(y) \max_x p_{x|y}(x|y).$$

$$H_\infty(x|y) = \log \frac{1}{v(x|y)}.$$



Information leaked about x from y

= initial uncertainty in x

-

remaining uncertainty in x after observing y

$$\text{Leakage } \angle_{P_X}(x \rightarrow y) = H_\infty(x) - H_\infty(x|y)$$

$$= \log \frac{v(x|y)}{v(x)}.$$

If P_X is unknown, the worst-case leakage is of interest.

Theorem. $\max_{P_X} \angle_{P_X}(x \rightarrow y) = \angle_{P_U}(x \rightarrow y)$

$$= \sum_y \max_x P_{Y|X}(y|x).$$

(sum of maximum of the columns of $P_{Y|X}$ matrix)

Proof. $\angle_{P_X}(x \rightarrow y)$

$$= \frac{v(x|y)}{v(x)}$$

$$= \frac{\sum_y P_Y(y) \max_x P_{X|Y}(x|y)}{\max_{x'} P_X(x')}$$

$$= \sum_y P_Y(y) \max_x \frac{P_{Y|X}(y|x) P_X(x)}{P_Y(y)} / \max_{x'} P_X(x')$$

$$\begin{aligned}
 &= \sum_{\gamma} \frac{\max_x P_{Y|X}(\gamma|x) P_X(x)}{\max_{x'} P_X(x')} \\
 &\leq \sum_{\gamma} \max_x P_{Y|X}(\gamma|x) \cdot \frac{\max_{x'} P_X(x')}{\max_{x'} P_X(x')} \\
 &= \sum_{\gamma} L_{P_0}(x \rightarrow \gamma),
 \end{aligned}$$

where $P_0(x) = \frac{1}{|X|}$, $\forall x \in X$.

Remark. This is an operationally motivated leakage measure and satisfies all the axioms of a leakage measure. Interestingly, mutual information $I(X; Y) = H(X) - H(X|Y)$ does not satisfy all.

Q) Let (X, Y) have the joint PDF

$$f_{XY}(x, y) = \begin{cases} xy & (x, y) \in [0, 1]^2 \\ 0 & \text{otherwise,} \end{cases}$$

Let $Z = x(1+y)$, $W = x^2$.

$$g_1(x,y) = x(1+y) \quad g_2(x,y) = x^2.$$

$$x = h_1(z \geq \omega) = \sqrt{\omega}, \quad y = h_2(z \geq \omega) = \frac{z}{\sqrt{\omega}} - 1.$$

$$\frac{\partial g_1}{\partial x} = 1+y \quad - \frac{\partial g_1}{\partial y} = x$$

$$\frac{\partial g_2}{\partial x} = 2x \quad - \quad \frac{\partial g_2}{\partial y} = 0.$$

$$J = \begin{vmatrix} 1+y & x \\ 2x & 0 \end{vmatrix} = 2x^2 = 2\omega$$

$$f_{Z \geq \omega}(z \geq \omega) = \frac{f_{XY}(\sqrt{\omega}, \frac{z}{\sqrt{\omega}} - 1)}{2\omega}$$

$$= \frac{\sqrt{\omega} + \frac{z}{\sqrt{\omega}} - 1}{2\omega},$$

for $(z \geq \omega)$ s.t. $0 \leq \omega \leq 1$, $\sqrt{\omega} \leq z \leq 2\sqrt{\omega}$.

0 otherwise,