浙江大学 2020 - 2021 学年 春夏 学期 《数据库系统》课程期末考试试卷参考答案和评分标准

课程号:	21121350 , 开课学院: <u>计算机学院</u>
考试试卷:	√A卷、B卷(请在选定项上打√)
考试形式:	√闭、开卷(请在选定项上打√),允许带 <u>一张 A4 纸笔记入场</u>
考试日期:	2021 年 7 月 2 日,考试时间: 120 分钟

诚信考试,沉着应考,杜绝违纪。

考生姓名:			学号:			所属院系:			
题序	_	<u> </u>	=	四	五	六	七	八	总 分
得分									
评卷人									

Problem 1: Relational Model and SQL (18 points)

Following are the relational schemas of a SRTP (Student Research Training Program) project database.

```
student (sId, sName, dId)
teacher (tId, tName, dId)
department (dId, dName)
project (pId, pName, tId, startTime, endTime)
participate (pId, sId, role)
```

The underlined attributes are primary keys, and foreign keys are listed as follows:

"dId" in "student" references "department";

"dId" in "teacher" references "department";

"tId" in "project" references "teacher";

"pId" and "sId" in "participate" reference "project" and "student", respectively. In "participate", only two different roles are permitted: "leader" and "member". Based on the above relational schemas, please answer the following questions:

- (1) Write a relational algebra expression to find the names of the projects that are instructed by a teacher from the department "Computer Science". (4 points)
- (2) Write SQL statements to create tables project and participate with all the necessary

- constraints (Note: Tables student, teacher, and department have already been created and can be referenced). (6 points)
- (3) Write a SQL statement to find the names of the teachers that instruct at least one project started in the year 2020. (4 points)
- (4) Write a SQL statement to find the names of the students participating more than 2 projects. (4 points)

Answers of Problem 1:

```
(1)
\prod_{pName}(project \bowtie teacher \bowtie (\sigma_{dName="Computer Science"}(department)))
评分细则:每个操作错扣1分,全对给4分
(2)
CREATE TABLE project
    (pId char(10),
    pName varchar(20),
    tId char(10),
    startTime date,
    endTime date,
    primary key (pId),
    foreign key (tId) references teacher);
评分细则: 写出 schema、primary、foreign key 给 2 分,类型全对给 3 分
CREATE TABLE participate
    (pId char(10),
    sId char(10),
    role varchar(20),
    primary key (pId, sId),
    foreign key (pId) references project,
    foreign key (sId) references student,
    check (role="leader" or role="member"));
评分细则: 写出 primary、foreign key、check 给 2 分,全对给 3 分
(3)
select distinct tName
from project, teacher
where project.tId=teacher.tId and startTime between '2020-01-01' and '2020-12-31'
评分细则:每个 where 条件错扣 1 分,全对给 4 分
(4)
select sName
from student
```

where sId in

(select sId

from participate

group by sId

having count(pId) > 2)

评分细则: 写出 having...且对给 1 分,内表写对给 2 分,全对给 4 分

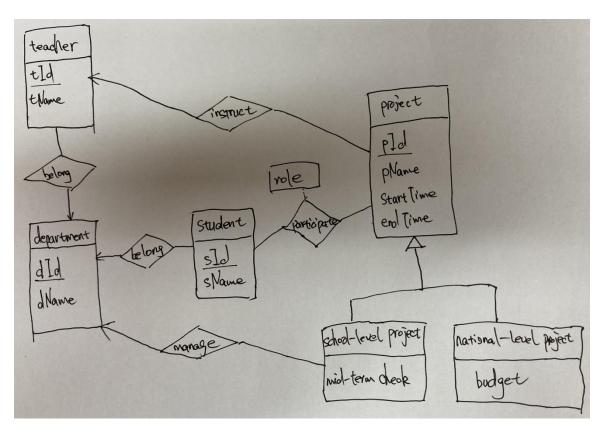
Problem 2: E-R Model (11 points)

Based on the SRTP project management scenario in Problem 1, some new requirements are added as follows:

- (1) There are two kinds of SRTP projects, i.e., school-level projects and national-level projects, and a project is either school-level or national-level.
- (2) National-level projects have budget information, and school-level projects have mid-term check information.
- (3) A school-level project is associated with exactly a department that is in charge of the management of the project.

Please draw an E-R diagram for the scenario.

Answers of Problem 2:



评分细则: 联系写成实体扣 1 分,没写联系扣 1 分,属性写错漏写扣 1 分,少写实体或联系扣 1 分,全对给 11 分

Problem 3: Relational Formalization (12 points)

For relation schema R (A, B, C, D, E, F) with functional dependencies set $F = \{A->B, A->C, B->C, D->E, D->F, EF->D\}$. Answer the following questions:

- (1) Find all the candidate keys. (3 points)
- (2) Find the canonical cover Fc. (3 points)
- (3) If R is not in BCNF, decompose it into BCNF schemas. (4 points) Is this decomposition dependency preserving? (2 points)

Answers of Problem 3:

(1)

AD AEF

评分细则: 少写一个扣1分, 多写扣1分, 全对给3分

(2)

A->B, B->C, D->EF, EF->D

评分细则:少写一个扣1分,全对给3分

(3)

There are different decomposition results and the following is just an example.

This decomposition is not dependency preserving (e.g., B->C is not preserved).

Following is another solution:

This decomposition is dependency preserving, because A->B can be checked on R21, B->C can be checked on R1, D->EF and EF->D can be checked on R221.

评分细则:分解部分少写一个扣1分,全对给4分,依赖保持判断错误扣1分,原 因错误扣1分,全对给2分

Problem 4: XML (8 points)

The following is a simplified DTD for the SRTP project database given in Problem 1:

```
<!ELEMENT SRTP(department+, teacher+, student+, project*)>
```

<!ELEMENT department (dname)>

```
<!ATTLIST
                 department dId ID #REQUIRED>
                 teacher (tname)>
   <!ELEMENT
   <!ATTLIST
                 teacher
          tId ID #REQUIRED
          dId IDREF #REQUIRED>
   <!ELEMENT
                 student (sname)>
   <!ATTLIST
                 student
          sId ID #REQUIRED
          dId IDREF #REQUIRED>
   <!ELEMENT
                 project (pname, starttime, endtime)>
   <!ATTLIST
                 project
          pId ID #REQUIRED
          tId IDREF #REQUIRED
          sIds IDREFS #REQUIRED >
                 dname (#PCDATA)>
   <!ELEMENT
   <!ELEMENT
                 tname (#PCDATA)>
   <!ELEMENT
                 sname (#PCDATA)>
   <!ELEMENT
                 pname(#PCDATA)>
                 starttime(#PCDATA)>
   <!ELEMENT
   <!ELEMENT
                 endtime(#PCDATA)>
1>
```

Please answer the following questions:

- (1) Give an XPath expression to return the names of all the teachers who supervise SRTP projects. (4 points)
- (2) Give an XQuery expression to return all the projects and their corresponding instructors, in the form of project_instructor elements that have a project subelement and a teacher subelement. (4 points)

Answers of Problem 4:

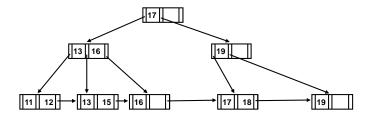
```
(1)
/SRTP/project/id(@tId)/tname/text()
评分细则:每个路径错误扣 1 分,全对给 4 分

(2)
for $p in /SRTP/project,
$t in /SRTP/teacher,
where $p/@tId= $t/@tId
return <project_instructor> { $p $t } </project_instructor>
评分细则:漏一个条件扣 1 分,逻辑错误扣 2 分,全对给 4 分
```

Problem 5: B+ -Tree and Query Processing (10 points)

Table student in Problem 1 is stored sequentially on sId. The following B+-tree is built for the table on attribute dId. Please answer the following questions:

- (1) Is the built index a primary index? Why? (2 points)
- (2) Draw the B+-tree after inserting entry 14. (4 points)
- (3) Draw the B+-tree after deleting entry 19 from the original B+-tree. (4 points)



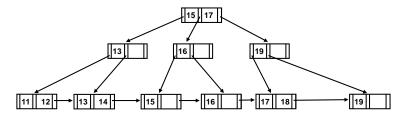
Answers of Problem 5:

(1)

The built index is not a primary index, as the search key of the index is not the search key of the sequentially ordered data file.

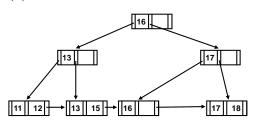
评分细则:结论正确给1分,解释正确给1分,结论错误不给分

(2)



评分细则:每个 entry 插入错扣 1 分,图中缺箭头或线扣 1 分,全对给 4 分

(3)



评分细则:每个entry删除错扣1分,图中缺箭头或线扣1分,全对给4分

Problem 6: Query Processing (14 points)

There are two relations r (100 blocks) and s (20 blocks), and hash-join algorithm is used to perform natural join between these two relations (memory size M=6 blocks). Please answer the following questions:

- (1) How many partitions can be constructed? Why? (3 points)
- (2) Which relation is best to choose as the build relation? Why? (3 points)
- (3) Is recursive partition needed? Why? (3 points)
- (4) Please compute the cost (numbers of seeks and block transfers) of the hash-join. (5 points)

Answers of Problem 6:

(1)

5 partitions, as the number of partitions is M-1.

评分细则:结论正确给1分,解释正确给2分,结论错误酌情给分

(2)

Relation s, as relation s is smaller than relation r.

评分细则:结论正确给1分,解释正确给2分,结论错误酌情给分

(3)

Recursive partition is not needed, as the size of the partitions of relation s (i.e., 4) is less than or equal to M-2 (i.e., 4).

评分细则:结论正确给1分,解释正确给2分,结论错误酌情给分

(4)

Number of block transfers: $3 \times (100+20)+4 \times 5$

Note: 4×5 is not necessary, which considers partially filled blocks.

Number of seeks: $2 \times (100+20)+2 \times 5$

评分细则: block transfers 的数量 3 分, seeks 的数量 2 分, 全对给 5 分

Problem 7: Concurrency Control (13 points)

Given the following schedule, please answer the following questions: (1) Draw the precedence graph for the read B schedule. (3 points) write C (2) Is the schedule conflict serializable? read A Why? (2 points) read C (3) Is it possible that the schedule is write A generated by the 2PL protocol with lock read A conversions? Explain. (5 points) write C (4) Which conditions should be satisfied if write B we want the schedule to be recoverable? read B (3 points)

Answers of Problem 7:

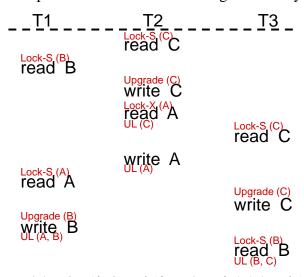
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评分细则:图中结点错误或连接错误,每错一处扣1分,全对给3分

(2)

The schedule is conflict serializable, as the precedence graph is acyclic. 评分细则:结论正确给 1 分,解释正确给 1 分,结论错误不给分

(3) It is possible that the schedule is generated by the 2PL protocol with lock conversions.



评分细则:结论正确给1分,流程图正确给4分,或者文字解释,表明题意也可给

分,结论错误不给分

(4)

T1 must commit before T3 does.

T2 must commit before T1 does.

T2 must commit before T3 does.

评分细则:每条执行顺序正确给1分,全对给3分

Problem 8: Recovery (14 points)

Given the following log file that supports logical undo, please answer the following questions:

- (1) The system crashes just after the last log record. What are the values of B and C in the database after system crash? (3 points)
- (2) Which transactions should redo and undo, respectively? (3 points)
- (3) What are the start and end points for redo and undo, respectively? (3 points)
- (4) What are the log records added during recovery? (5 points)

Answers of Problem 8:

14 $\langle T_2, \text{ commit} \rangle$

(1)

B = 2100

C= 300 or 400 or 700

评分细则: B 值正确给 1 分, C 值漏写扣 1 分, 三个值完整写出给 2 分

(2)

redo: T_0 and T_2 undo: T_1

评分细则: redo 完整写出给 2 分,漏写扣 1 分,undo 写对给 1 分

(3)

redo: 7-14 undo: 14-3

评分细则: redo、undo 顺序错写一个扣 1 分,redo 顺序写成 6-14 也可给分,全对给 3 分

(4)

<T₁, C, 600>

 $\langle T_1, O_1, operation-abort \rangle$

<T₁, B, 2050>

<T₁, abort>

评分细则: 错一个扣1分, 写出一个给1分, 全对给5分