

Patt-Ch7 Assembly Language

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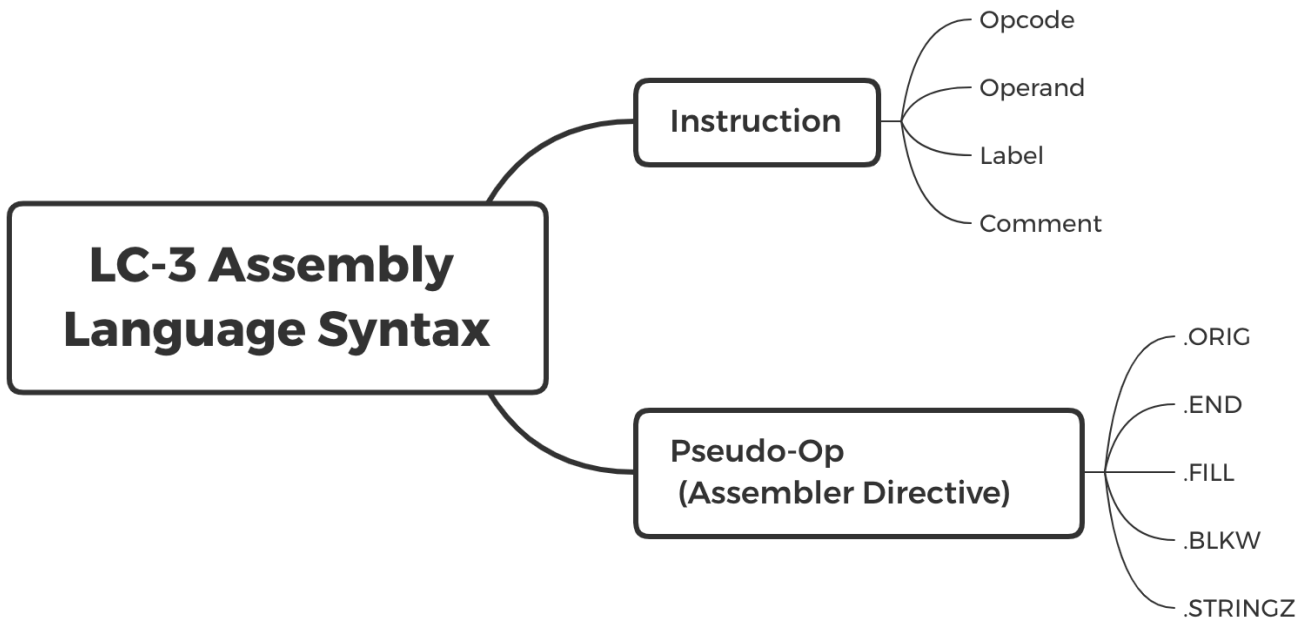
7.1 From Machine Code to Assembly Language

Assembler is a program that turns symbols into machine instructions.

Assembly Language

- Low-level language, ISA-specific
 - it is usually the case that **each ISA has only one assembly language**
- User-Friendly
 - mnemonics(助记符) for opcodes
 - **labels** for memory locations
- still provide detailed control over the instructions

7.2 LC-3 Assembly Language Syntax



Machine Code

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ADD*	0001				DR			SR1	0	00					SR2	
ADD*	0001				DR			SR1	1						imm5	
AND*	0101				DR			SR1	0	00					SR2	
AND*	0101				DR			SR1	1						imm5	
BR	0000		n		z		p									PCoffset9
JMP	1100			000				BaseR								000000
JSR	0100		1													PCoffset11
JSRR	0100		0		00			BaseR								000000
LD*	0010				DR											PCoffset9
LDI*	1010				DR											PCoffset9
LDR*	0110				DR			BaseR								offset6
LEA	1110				DR											PCoffset9
NOT*	1001				DR			SR								111111
RET	1100			000				111								000000
RTI	1000															000000000000
ST	0011				SR											PCoffset9
STI	1011				SR											PCoffset9
STR	0111				SR			BaseR								offset6
TRAP	1111			0000												trapvect8
reserved	1101															

Pseudo-Ops: .ORIG x3000
 .FILL xFFFF
 .END

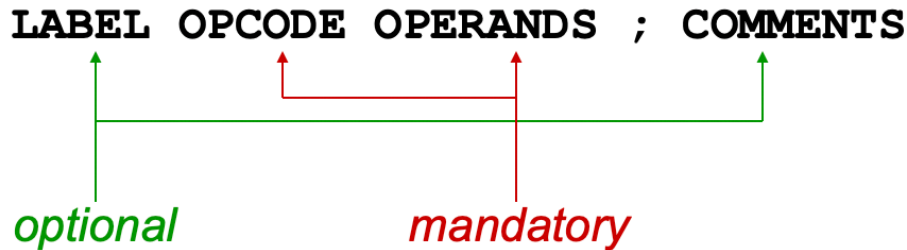
Assembly Language

ADD DR, SR1, SR2 ADD r1, r2, r3
 ADD DR, SR1, imm5 ADD r1, r2, #10 +进制
 AND DR, SR1, SR2 [E16,15] b1010 = 进制
 AND DR, SR1, imm5 xA +入进制
 BRn LABEL BR/BRnzp, BRz, BRp,
 BRnz, BRnp, BRzp
 JMP BaseR
 JSR LABEL
 JSRR BaseR
 LD DR, LABEL
 LDI DR, LABEL
 LDR DR, BaseR, offset6
 LEA DR, LABEL
 NOT DR, SR
 RET
 RTI
 ST SR, LABEL
 STI SR, LABEL
 STR SR, BaseR, offset6
 TRAP trapvect8 TRAP x25

.BLKW 2
 .STRINGZ "Hello World"
 .EXTERNAL START of FILE*

LDR	R4, R2, #-5	; R4 <- mem[R2-5]
LD	R4, VALUE	; R4 <- mem[VALUE]
LDI	R4, VALUE	; R4 <- mem[mem[VALUE]]
LEA	R4, TARGET	; R4 <- the address of TARGET
ST	R4, HERE	; mem[HERE] <- R4
STI	R4, THERE	; mem[mem[THERE]] <- R4
STR	R4, R2, #-5	; mem[R2-5] <- R4

7.2.1 Instructions



7.2.1.1 Opcodes & Operands

Opcodes

- *reserved words*
- Symbolic names

Operands

Operands

- **registers** -- specified by Rn, where n is the register number
- **numbers** -- indicated by # (decimal) or x (hex) or b(binary)
- **label** -- symbolic name of memory location
- **separated by ,**
- **number, order, and type correspond to instruction format**

➤ **ex:**

```
ADD R1 , R1 , R3
ADD R1 , R1 , #3
LD  R6 , NUMBER
BRz LOOP
```

7.2.1.2 Labels

- **placed at the beginning of the line**
- **assigns a symbolic name to the address corresponding to line**

➤ **ex:**

```
LOOP  ADD  R1 , R1 , #-1
      BRp  LOOP
```

2 reasons for using Labels:

- The location is the target of BR
- The location contains a value that is loaded or stored.

7.2.1.3 Comments

- **anything after a semicolon is a comment**
- **ignored by assembler**
- **used by humans to document/understand programs**
- **tips for useful comments:**
 - **avoid restating the obvious, as “decrement R1”**
 - **provide additional insight, as in “accumulate product in R6”**
 - **use comments to separate pieces of program**

7.2.2 Pseudo-Ops (Assembler Directives)

Pseudo-, /'sudu/, 假的

```
.ORIG x3000                ; The start address of program

.END                        ; The end of program

.FILL x0030                ; allocate one word, initialize with value n

.BLKW n(decimal)           ; BLocK of Words, allocate n words of storage

.STRINGZ "ABC"              ; allocate n+1 addresses, initialize n words
                           ; with Zero-EXT ASCII codes, and NULL terminated

.EXTERNAL                  ; send a message to LC-3 assembler that the
                           ; absence of the label is in another module.
                           ; Used for multi-file programming.
```

Note : .END does not stop the program & it does not even exist at the time of execution.

7.2.3 Good Programming Style*

Style Guidelines

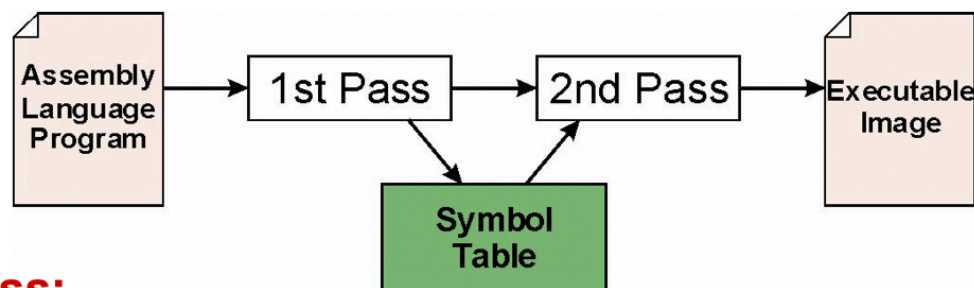
Use the following style guidelines to improve the readability and understandability of your programs:

1. Provide a program header, with author's name, date, etc., and purpose of program.
2. Start labels, opcode, operands, and comments in same column for each line. (Unless entire line is a comment.)
3. Use comments to explain what each register does.
4. Give explanatory comment for most instructions.
5. Use meaningful symbolic names.
 - Mixed upper and lower case for readability.
 - ASCIItoBinary, InputRoutine, SaveR1
6. Provide comments between program sections.
7. Each line must fit on the page -- no wraparound or truncations.
 - Long statements split in aesthetically pleasing manner.

7.3 The Assembly Process

Assembly Process

Convert assembly language file (.asm) into an executable file (.obj) for the LC-3 simulator.



First Pass:

- scan program file
- find all labels and calculate the corresponding addresses; this is called the symbol table

Second Pass:

- convert instructions to machine language, using information from symbol table

7.3.1 First Pass: Constructing the Symbol Table

- 1. Find the .ORIG statement, which tells us the address of the first instruction.**
 - Initialize location counter (LC), which keeps track of the current instruction.
- 2. For each non-empty line in the program:**
 - a) If line contains a label, add label and LC to symbol table.
 - b) Increment LC.
 - NOTE: If statement is .BLKW or .STRINGZ, increment LC by the number of words allocated.
- 3. Stop when .END statement is reached.**

NOTE: A line that contains only a comment is considered an empty line.

```

                .ORIG    x301C
                ST      R3, SAVE3
                ST      R2, SAVE2
                AND     R2, R2, #0
TEST           IN
                BRz     TEST
                ADD     R1, R0, #-10
                BRn     FINISH
                ADD     R1, R0, #-15
                NOT     R1, R1
                BRn     FINISH
                HALT
FINISH         ADD     R2, R2, #1
                HALT
SAVE3          .FILL   x0000
SAVE2          .FILL   x0000
                .END

```

简答题 (2 分)

Symbol	Address
TEST	x301F
FINISH	x3027
SAVE3	x3029
SAVE2	x302A

7.3.2 Second Pass: Generating Machine Language

Potential problems:

- Improper number or type of arguments
 - ex: NOT R1, #7
ADD R1, R2
ADD R3, R3, NUMBER
- Immediate argument too large
 - ex: ADD R1, R2, #20
- Address (associated with label) more than 256 from instruction
 - can't use PC-relative addressing mode

7.4 Multiple Object Files (Recommend Reading)

Multiple Object Files

An object file is not necessarily a complete program.

- system-provided library routines
- code blocks written by multiple developers

**For LC-3 simulator,
can load multiple object files into memory,
then start executing at a desired address.**

- system routines, such as keyboard input, are loaded automatically
 - loaded into “system memory,” below x3000
 - user code should be loaded between x3000 and xFDFF
- each object file includes a starting address
- be careful not to load overlapping object files

7.4.1 The Executable Image

When a computer begins execution of a program, the entity being executed is called an *executable image*. The executable image is created from modules often created independently by several different programmers. Each module is translated separately into an object file. We have just gone through the process of performing that translation ourselves by mimicking the LC-3 assembler. Other modules, some written in C perhaps, are translated by the C compiler. Some modules are written by users, and some modules are supplied as library routines by the operating system. Each object file consists of instructions in the ISA of the computer being used, along with its associated data. The final step is to combine (i.e., *link*) all the object modules together into one executable image. During execution of the program, the FETCH, DECODE, ... instruction cycle is applied to instructions in the executable image.

Linking and Loading

Loading is the process of copying an executable image into memory.

- more sophisticated loaders are able to relocate images to fit into available memory
- must readjust branch targets, load/store addresses

Linking is the process of resolving symbols between independent object files.

- suppose we define a symbol in one module, and want to use it in another
- some notation, such as `.EXTERNAL`, is used to tell assembler that a symbol is defined in another module
- linker will search symbol tables of other modules to resolve symbols and complete code generation before loading