

EECS 336 Fall 2015
Homework Problem 6.3

Define $s_i = d_i - f_i$

Step One: $OPT(i, j)$ = "set of all possible sums using j elements from $\{s_1, \dots, s_i\}$ "

Step Two: Recurrence. $OPT(i + 1, j + 1) = OPT(i, j + 1) \cup [OPT(i, j) + s_{i+1}]$

Step Three: Base case

$$OPT(i, i) = \sum_{l=1}^i s_l$$

$$OPT(i, 1) = \{s_1, \dots, s_i\}$$

Step Four: Iterative DP

Algorithm 1 Iterative DP to find all possible sums of subset

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1: procedure FINDSUMS( $d, f$ )
2:    $n = |d|$ 
3:   Calculate  $s_i = d_i - f_i$ 
4:   Calculate total sum  $S = \sum_{i=1}^n s_i$ 
5:   Initialize  $memo[j][k] = \emptyset$  for  $j, k$  in  $\{1, \dots, n\}$ 
6:   Set  $memo[i][i] = \sum_{l=1}^i s_l$ 
7:   Set  $memo[i][1] = \{s_1, \dots, s_i\}$ 
8:   for  $i = 2$  to  $n$  do
9:     for  $j = 1$  to  $i$  do
10:       $tempSet = \emptyset$ 
11:      for  $e$  in  $memo[i-1][j]$  do
12:         $tempSet = tempSet \cup \{e + s_i\}$ 
13:      end for
14:       $memo[i][j] = memo[i-1][j] \cup tempSet$ 
15:    end for
16:  end for
17:  for each  $sum$  in  $memo[n][n/2]$  do
18:    if  $0 < sum < S$  then
19:      return True
20:    end if
21:  end for return False
22:
23: end procedure
```

Correctness:

By induction on recurrence. If $OPT(i, j)$ correctly captures all possible sums using j elements from first $\{s_1, \dots, s_i\}$, then from the recurrence, $OPT(i, j + 1)$ denotes all possible sums taking one more $(j + 1)$ elements in subset $\{s_1, \dots, s_i\}$, and $[OPT(i, j) + s_{i+1}]$ denotes all possible sums in $OPT(i, j)$ added by s_{i+1} . The two components ensures that all possible sums with $j + 1$ elements are added from $\{s_1, \dots, s_{i+1}\}$, without and with s_{i+1} . Thus, the recurrence is correct.

Runtime Analysis:

Each iteration of inner two for loops takes $O(n)$. Assuming union operation takes constant time for quick-union, or $O(\log(n))$ for standard union as in union-find. Initialization of summation and memo table set up takes $O(n^2)$. Thus, the algorithm takes $O(n^2)$ in total using quick-union operations, or $O(n^2 \log(n))$ for standard union as in union-find.