

MPCS 51040 – C Programming

Lecture 4 – Pointers & Recursion

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Announcements

- ▶ Reminder: quiz next week.
- ▶ Possibly reschedule 11/14 lecture to Saturday 11/19. Please keep free.



Operator Precedence and Associativity

Precedence	Operator	Description	Associativity
1	++ --	Suffix/postfix increment and decrement	Left-to-right
	()	Function call	
	[]	Array subscripting	
	.	Structure and union member access	
	->	Structure and union member access through pointer	
	(type){list}	Compound literal(c99)	
2	++ --	Prefix increment and decrement	Right-to-left
	+ -	Unary plus and minus	
	! ~	Logical NOT and bitwise NOT	
	(type)	Type cast	
	*	Indirection (dereference)	
	&	Address-of	
	sizeof	Size-of	
	_Alignof	Alignment requirement(c11)	
3	* / %	Multiplication, division, and remainder	Left-to-right
4	+ -	Addition and subtraction	
5	<< >>	Bitwise left shift and right shift	
6	< <=	For relational operators < and ≤ respectively	
	> >=	For relational operators > and ≥ respectively	
7	== !=	For relational = and ≠ respectively	
8	&	Bitwise AND	
9	^	Bitwise XOR (exclusive or)	
10		Bitwise OR (inclusive or)	
11	&&	Logical AND	
12		Logical OR	
13	?:	Ternary conditional	Right-to-Left
14	=	Simple assignment	
	+= -=	Assignment by sum and difference	
	*= /= %=	Assignment by product, quotient, and remainder	
	<<= >>=	Assignment by bitwise left shift and right shift	
	&= ^= =	Assignment by bitwise AND, XOR, and OR	
15	,	Comma	Left-to-right

Associativity defines the order in which adjacent operators *with the same precedence level* are evaluated.

▶ $a - b - c \Rightarrow (a - b) - c$

▶ $q \ \&\& \ r \ || \ s \Rightarrow (q \ \&\& \ r) \ || \ s$



Does not say in which order the operator *operands* have to be evaluated!



precedence.c



Evaluation Order

```
1 // Unspecified function
2 // evaluation order;
3 // Associativity only
4 // defines addition order.
5 f1()+f2()+f3();
6
7 // bad; no sequence point
8 a[i]=b[i++];
9
10 // bad; no sequence point
11 // however, all side effects
12 // done *before* entering
13 // function
14 f(i++,i,j++);
15
16 // OK; , is seq point
17 i++,a=i;
```

Order of operations of the operands of *almost all* operators is not defined by the language. The compiler is free to evaluate operands in any order (and does not have to be consistent).

Only the sequential-evaluation (`,`), logical-AND (`&&`), logical-OR (`||`), conditional-expression (`? :`), and function-call operators constitute **sequence points** and therefore guarantee a particular order of evaluation for their operands.

(sequence point: all side effects of previous evaluations have occurred)

The function-call operator is the set of parentheses following the function identifier. The sequential-evaluation operator (`,`) is guaranteed to evaluate its operands from left to right. (Note that the comma operator in a function call is not the same as the sequential-evaluation operator and does not provide any such guarantee.)



Boolean Short-circuit Evaluation

```
1 // call to f never happens
2 0 && f();
3
4 // OK but avoid due to short circuit
5 // with side-effects;
6 // logical operator is seq point
7 q && r || s--;
```

Logical operators also guarantee evaluation of their operands from left to right. However, they evaluate the smallest number of operands needed to determine the result of the expression. This is called "short-circuit" evaluation. Thus, some operands of the expression may not be evaluated.



Expert topic; Better to avoid relying on specific sequencing whenever possible.



Pointers: recap

A pointer p is:

- ▶ A variable
- ▶ which has as **type** a pointer (to some type)
- ▶ which (like all other variables) has a *static* type (i.e. a pointer to a char will always be a pointer to a char)
- ▶ which holds as **value** a memory address
- ▶ which (like any variable) itself has a memory address
- ▶ which can be *dereferenced* (using $*$) to obtain an *lvalue* (i.e. something which can appear on the left of $=$)
- ▶ when dereferenced, an lvalue of type T (for pointers of type $T *$)
- ▶ when dereferenced, an lvalue for which, $\&(*p) == p$
- ▶ capable of arithmetic (unless function pointer): add, subtract integer, increment, decrement, subtract two pointers to same type.

Beware: arithmetic in base types, not in bytes!



Pointers and Arrays

```
1 char test[] = "test";  
2 somefunctions(test); // decay  
3 test[0]='a';          // decay  
4  
5 // no decay; is somefunc(5);  
6 somefunc(sizeof(test));  
7  
8 // decay into char *  
9 somefunc(test);
```

- ▶ Arrays are not really a first-class data type in C
- ▶ Arrays almost always 'decay' into pointers; Exceptions:
 - ▶ Argument of operator `&`
 - ▶ Argument of operator **sizeof**
 - ▶ Argument of operator `alignof` (C11)
 - ▶ Use as string literal

In particular, they decay when used in an expression (other than above) or when passed to a function.



Even the subscript operator (`[]`) is not what it seems:
 $a[2] \Rightarrow *(a+2)$



Function Pointers

Pointers to functions are possible:

```
1  int test(int i);  
2  int (*alstest)(int i) = test;  
3  
4  test(2);  
5  alstest(2);
```

- ▶ For functions, & is optional.
(clear from context)
- ▶ For function pointers, * is optional.
(clear from context)
- ▶ Pointer arithmetic is not allowed on
function pointers.
(some compilers allow as extension)
 - ▶ ++ will *not* move to the next
line of code!

In normal usage, operations on function pointers are restricted to setting to the address of a function, assigning to another function pointer, setting to NULL or dereferencing (calling).



sorting: qsort.c



Decoding Declarations

Right-left rule

► * → Pointer to

► [] → Array of

► () → function returning

Rule

1. Find identifier
2. Go right (until no more symbols or until parenthesis)
3. Go left (until no more symbols or until left parenthesis)

Example

```
1 // array of int
2 int a[10];
3 // function returning pointer
4 int * f();
5 // pointer to array of array of ints
6 int (*a)[][];
```

```
1 // function ret pointer to array of int
2 int (*f())[];
3 // typedef for function pointer
4 typedef int(*TypeAlias)(int);
5 // function returning pointer to function
6 // returning int
7 int (*f())();
```



What is: `void (*signal(int sig, void (*func)(int)))(int);`
(help? <http://cdecl.org>)



Scope and lifetime revisited

```
1  // global variable
2  // - file scope
3  // - program lifetime
4  // - program visibility
5  int global = 10;
6
7  // global variable
8  // internal linkage
9  // - file scope
10 // - program lifetime
11 // - file visibility
12 static int private = 2;
13
14 int func() {
15     // local variable
16     // (automatic storage duration)
17     // - block scope
18     // - block lifetime
19     int local = 2;
20 }
21
22 int func2() {
23     // local variable
24     // (static storage duration)
25     // - block scope
26     // - program lifetime
27     static int local = 2;
28 }
```

Scope and lifetime:

- ▶ Scope of a variable is where the name is accessible.
- ▶ Lifetime is the duration for which the variable has a storage location (in memory) associated with it, i.e. how long it can hold (remember) a value.

So far, we have not specified anything for scope or lifetime. We have used:

- ▶ *local variables*: scope local to the statement block.
- ▶ *global variables*: file scope (starts at point where declared until the end of the (preprocessed!) file).

Regarding lifetime:

- ▶ *local variables* by default have *automatic storage duration*: they are destroyed when they go out of scope (end of block) and recreated when re-entered. They *do not hold their value when destroyed*.
- ▶ *static storage duration* means that the lifetime is *until the program ends*. Global variables (outside of any function) have static storage.



For global variables, adding the **static** keyword

Dynamic Memory Allocation

Why do we need dynamic memory allocation

Pointers

Pointers have a number of benefits:

- ▶ Allow 'aliasing' of variables
- ▶ Allow easy traversing of arrays and other data structures
- ▶ Provide access to variables outside of the scope their normal scope.

They do not:

- ▶ change the lifetime of existing variables
- ▶ enable 'variable length' variables

Remaining issues; examples:

- ▶ Read a variable length line from a file
- ▶ Create a variable in a function and pass it to the caller (i.e. lifetime not bound to scope)



Dynamic Memory Allocation

malloc, free

```
1 // request memory
2 char * ptr = (char *)
3     malloc(100);
4
5 // Use ptr as usual
6 copyAndConvert("test",
7     ptr, 100);
8
9 // free memory
10 free(ptr);
```

- ▶ malloc does not know what you want to store in the requested memory (i.e. type) so it returns **void ***. You will need to cast.
- ▶ There is no garbage collection; You are responsible for freeing resources.
- ▶ There are no checks to ensure you stay within your allocated memory.
- ▶ You should not try to pass an invalid pointer to free or not try to free the pointer more than once. (undefined behaviour will result)
- ▶ The returned memory region is *not initialized*. It may contain random bytes. (if you need to clear it, use calloc)

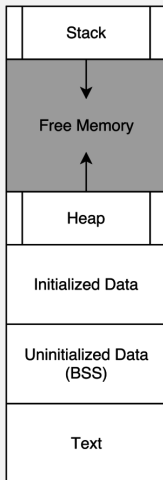


firstmalloc .c: malloc, free examples.



Memory Regions

(Implementation Detail)



Typical memory layout:

- ▶ automatic variables go on the stack.
- ▶ uninitialized global and static storage duration variables go into BSS.
- ▶ Initialized global and static storage duration variables go into data.
- ▶ malloc allocates memory from the heap.



This is an implementation detail, and not required by the standard. The standard only specifies variable lifetime etc., not how to map these concepts to memory.



Walk through stack changes on function call.



Recursion

```
1 void repeat(unsigned int i)
2 {
3     if (!i)
4         return;
5     // do something
6     puts("X");
7     repeat(i-1);
8 }
```

Recursion:

- ▶ Recursion happens when a function calls itself.
- ▶ Each invocation of the function receives its own copy of local (automatic) variables (variables go on the stack).
- ▶ Local variables exist until the function returns.
 - ▶ Watch out for memory (stack) usage!
- ▶ Recursion is a form of looping.

Recursion needs to end: *base case*.



Recursion

Example

Calculate fibonacci numbers: $F_n = F_{n-1} + F_{n-2}$ where $F_0 = 0, F_1 = 1$

Solution

The base case and recursion step are explicit in the mathematical definition.



Calculate fibonacci numbers



Recursion

Example

Recursion is very useful for divide-and-conquer type algorithms: split the problem in smaller problems and try to solve the smaller problem.

Example

Count the number of occurrences of a number in an array.

Solution

- ▶ *Base case (conquer): array of size 1*
- ▶ *Divide: split in two arrays, add counts.*

```
unsigned int count (int * array , unsigned int size );
```



Implement count().



Valgrind

Valgrind

Valgrind is an instrumentation framework for building dynamic analysis tools. There are Valgrind tools that can automatically detect many memory management and threading bugs, and profile your programs in detail. You can also use Valgrind to build new tools.

- ▶ Valgrind can help debug pointer issues (null pointer usage, accessing memory outside of allocated memory region, using memory that has been freed, memory leaks, ...)
- ▶ Valgrind will benefit from compiling with debug information enabled (more later).
- ▶ Homework 3 will use valgrind to validate your code.
- ▶ Valgrind is available on `linux.cs.uchicago.edu`.



demo memory leak, out-of-bounds, use-after-free



- ▶ General documentation: <http://valgrind.org/>
- ▶ Memory-debugging: <http://valgrind.org/docs/manual/quick-start.html#quick-start.mcrun>



Valgrind

Using valgrind to detect memory leaks

Example output



GDB

GDB

GDB is a command-line debugger (what's a debugger?)

- ▶ GDB is installed on `linux.cs.uchicago.edu`.
- ▶ GDB allows us to run the program step-by-step and to inspect the value of variables.
- ▶ GDB can do post-mortem analysis (i.e. on a coredump file).



Add `—ggdb` to your compile commandline to instruct the compiler to include more information. Adding `—O0` is also recommended.



setting breakpoints, inspecting variables, use a coredump.



- ▶ Easy tutorial: <http://www.techbeamers.com/how-to-use-gdb-top-debugging-tips/>
- ▶ <https://www.gnu.org/software/gdb/documentation>



Value Types

Value Types and Reference Types

In general, there are two kinds of types (not meaning 'C datatype' here but logical type): *value* type and *reference* type.

```
1 // Example value type
2 int a = 10;
3 b = a;
4 ++b; // b != a
5
6 // Example reference type
7 FILE * f = fopen (...);
8 FILE * f2 = f;
9 fgetc(f);
10 // f2 is modified as well!
```

- ▶ value types (example: POD types in C++) are fully contained in the memory the underlying (C) type holds. They do not refer to outside data.
- ▶ Because of this, value types can be created, copied and 'destroyed' without special considerations. (Compare with FILE * which needs to be properly closed)
- ▶ The underlying C type for a reference type does not hold (or only partially holds) the state for the logical type. It cannot be copied by assignment (since this would not copy the external state).



Modular Code

Example

Good code can easily be reused. An often overlooked property is that ideally, changes in the implementation details of the code should not affect users of that code.

```
1 // Library for employee management
2 // header
3 struct Person
4 {
5     int age;
6 };
7
8 // User of library
9 struct Person p;
10 // bad!
11 printf ("%i", p.age);
```

This is not good reusable code.

- ▶ Can never change struct Person (for example store age elsewhere)
- ▶ (Example) Can not add access control or logging.



Person: Improved

```
1 // in header:
2 struct Person { int age; }
3 void setAge (Person * p, int age);
4
5 // in main
6 struct Person p;
7 setAge(&p, 21);
```

Possible issues:

- ▶ Can look in header and access age directly.
- ▶ Can assume Person to be value type.



Modular Code

Better: Encapsulation

```
1  // In java:
2  //   p.setAge(10);
3
4  // Update age
5  void setAge (struct Person * p, int a);
6
7  // Better
8  void setAge (Person * p, int age);
9
10 // Even better
11 void setAge (Person p, int age);
```

- ▶ Object-oriented principle: encapsulation
 - ▶ Information hiding
 - ▶ Only provide information needed (can't be 'abused')
- ▶ C does not support C++/Java syntax
 - ▶ but we can emulate it



Homework 3 - First Impressions

Some questions/remarks:

- ▶ Why do we need maxchar?
 - ▶ Think about somebody else using your code, and given only your header file.
 - ▶ Explicit checks/safety is better than documentation/assumptions.
 - ▶ Same reason why gets almost impossible to use safely.
- ▶ Be very careful with array lengths/indices!

```
1  ...
2  while (i<maxchar && ...)
3  {
4  }
5  dest[i]=0;
6  ...
```

- ▶ Avoid specifying explicit sizes!
- ▶ Think 0-based!

- ▶ Maximum line length. Those days are over! You should be able to handle any input provided the system has enough memory.
- ▶ Array decay and consequences for operator =

```
1  char a[] = "1234";
2  char * b;
3  b=a; // does *not* make a copy of a!
```



Assignment: Homework 4 & Reading

Homework Assignment

See <https://mit.cs.uchicago.edu/mpcs51040-aut-16/mpcs51040-aut-16/raw/master/homework/hw4/hw4.pdf>

Reading Assignment

O'Reilly Mastering Algorithms in C:

Required Chapter 4, 5

Recommended Chapter 1-3 (should mostly be refresh)



Discuss 'handle' types.



Quiz next week



Topics:

- ▶ Lectures 1–4
- ▶ Homework 1–4

Quiz Format

- ▶ Duration: max 1.5h
- ▶ Closed book / no laptop, phone, ...
- ▶ Mostly practical focus:
 - ▶ What is the output of this code?
 - ▶ Is this code correct?
 - ▶ Write small functions/programs.
 - ▶ Some terms and definitions.



Homework is excellent preparation!

