

# 1 the whole problem

$$\begin{aligned}
& \min_{F, \Omega, \theta, \lambda^+, \lambda^-} F_{n+1} + \Omega \tag{1} \\
s.t. \quad & \Omega \geq \sum_{(v,w)=a \in V^2} (F_v - F_w + t_w) x_a^k + \Gamma \theta^k + \sum_{v \in V} (\lambda_v^{k,+} + \lambda_v^{k,-}) \quad \forall 1 \leq k \leq K, \\
& \theta^k + \lambda_v^{k,+} \geq \sum_{w \in V, (w,v) \in A} u_v \cdot x_{(w,v)}^k \quad \forall v \in V, \\
& \theta^k + \lambda_v^{k,-} \geq - \sum_{w \in V, (w,v) \in A} u_v \cdot x_{(w,v)}^k \quad \forall v \in V, \\
& z_{ij} = 1 \quad \forall (i, j) \in E, \\
& z_{ij} + z_{ji} \leq 1 \quad \forall i, j \in V, i < j, \\
& z_{ij} \geq z_{ip} + z_{pj} - 1 \quad \forall i, j, p \in V, i \neq j \neq p, \\
& f_{ij}^k \leq \min(\tilde{r}_i^c, \tilde{r}_j^c) z_{ij} \quad \forall i, j \in V, i \neq j, i \neq n+1, j \neq 0, \forall k \in \{1, 2, \dots, K\}, \forall c \in C, \\
& \sum_{j \in V \setminus \{i, 0\}} f_{ij}^k = \tilde{r}_i^c \quad \forall i \in V \setminus \{n+1\}, \forall k \in \{1, 2, \dots, K\}, \forall c \in C, \\
& \sum_{i \in V \setminus \{j, n+1\}} f_{ij}^k = \tilde{r}_j^c \quad \forall j \in V \setminus \{0\}, \forall k \in \{1, 2, \dots, K\}, \forall c \in C, \\
& f_{ij}^k \geq 0 \quad \forall i, j \in V, i \neq j, i \neq n+1, j \neq 0, \forall k \in \{1, 2, \dots, K\}, \\
& F_i \geq 0 \quad \forall i \in V \\
& x_a \leq M z_a \quad \forall a \in E, \\
& z_{ij} \in \{0, 1\} \quad \forall i, j \in V, i \neq j, \\
& \tilde{r}_i^c = R^c, \text{falls } i = 0, n+1, \forall k \in \{1, 2, \dots, K\}, \forall c \in C, \\
& \tilde{r}_i^c = r_i^c, \text{otherwise}, \forall k \in \{1, 2, \dots, K\}, \forall c \in C.
\end{aligned}$$

## 2 subproblem

$\mathcal{E}^t := \{(i, j) \in V^2 \mid z_{ij}^* > 0\}$ ,  $Z^t$  is accordingly set.

$$\begin{aligned}
M^{*t} = \max_{x, \xi^+, \xi^-} & \sum_{a \in A} (F_v - F_w + t_w) x_a + q_a^+ u_w - q_a^- u_w \\
s.t. & \sum_{(\cdot, v) \in \mathcal{E}^t} x_a - \sum_{(v, \cdot) \in \mathcal{E}^t} x_a \leq d_v^+ \quad \forall v \in V \setminus \{v_0\}, \\
& - \sum_{(\cdot, v) \in \mathcal{E}^t} x_a + \sum_{(v, \cdot) \in \mathcal{E}^t} x_a \leq d_v^+ \quad \forall v \in V \setminus \{v_0\}, \\
& \sum_{v \in V} (\xi_v^+ + \xi_v^-) \leq \Gamma, \\
& q_a^+, q_a^- \leq x_a \quad \forall a \in \mathcal{E}^t, \\
& x_a \leq M z_a \quad \forall a \in \mathcal{E}^t, \\
& q_{(v, w)}^+ \leq M \cdot \xi_w^+ \quad \forall (v, w) \in \mathcal{E}^t, \\
& q_{(v, w)}^- \leq M \cdot \xi_w^- \quad \forall (v, w) \in \mathcal{E}^t, \\
& q_{(v, w)}^+ \geq M \cdot \xi_w^+ + x_{(v, w)} - M \quad \forall (v, w) \in \mathcal{E}^t, \\
& q_{(v, w)}^- \geq M \cdot \xi_w^- + x_{(v, w)} - M \quad \forall (v, w) \in \mathcal{E}^t, \\
& x_a, q_a^+, q_a^- \geq 0 \quad \forall a \in \mathcal{E}^t, \\
& \xi_v^+, \xi_v^- \in \{0, 1\} \quad \forall v \in V
\end{aligned} \tag{2}$$

## 3 cut and algorithm

Cut:

$$\Omega \geq (M^{*t} - \lambda^t(1)) \left( \sum_{(i, j) \in Z^t} z_{ij} - |Z^t| \right) + M^{*t}$$

$N^t(1) = \{z \in \{0, 1\}^{|Z^t|} \mid \sum_{(i, j) \in Z^t} z_{ij} = 1\}$ ,  $\lambda^t(1)$  is the lower bound of  $\min_{N^t(1)} M^{*t}$ .

Starting with a (possibly empty) subset  $\hat{X} \subset \tilde{X}$  in each iteration, the LP (1) with only constraints corresponding to the nodes from  $\hat{X}$  is solved optimally. For the optimal solution  $F^*, \Omega^*$ , a new vertex  $x^*$  is computed by solving (2). If the objective is greater than  $\Omega^*$ , a violated constraints is found and we add  $x^*$  to  $\hat{X}$ . Else, the solution  $F^*, \Omega^*$  is optimal.

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**Algorithm 1** Exact Algorithm

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**Input:** Graph  $G = (V, A)$ , uncertainty parameter  $\Gamma \in \mathbb{N}_{<|V|}$ , execution times  $(t_v)_{v \in V}$ , uncertainties  $(u_v)_{v \in V}$ , costs per time unit for finishing earlier than scheduled  $(d_v^-)_{v \in V}$ , costs per time unit for finishing later than scheduled  $(d_v^+)_{v \in V}$ , flowvalue  $(f_{ij}^k)_{(i,j) \in V^2}$ , resource requirement  $(\tilde{r}_i^c)_{i \in V}$ , initial vertex set  $\hat{X} \subset \tilde{X}$ , sufficiently large constant  $M$ ,  $t = 0$

**Output:** the finish time

**while** true **do**

*schedule\_obj*,  $F^*$ ,  $\Omega^*$  optimal solution of (1) with constraints only for  $\hat{X}$

*subproblem\_obj*,  $x^*$  = optimal solution of (2)

**if**  $(M^{*t} - \lambda^t(1))(\sum_{(i,j) \in Z^t} z_{ij} - |Z^t|) + M^{*t} > \Omega$  **then**

$\hat{X} = \hat{X} \cup \{x\}$ ,  $t = t + 1$

**else**

        break;

**end if**

**end while**

**return** *schedule\_obj*,  $(F_v)_{v \in V}$

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