

**CS 3050 Homework # 1.****Name : KEY****Submitted to Canvas, due at 11:59pm on Feb. 6, 2018.**

1. Use Insertion Sort on array A and show the numbers and their position in the Array for each Iteration; also show the “key” position in each iteration.

A = {3,1,4,1,5,9,2,6}

**Answer:**

3	1	4	1	5	9	2	6	<- Swap A[0] and A[1]
1	3	4	1	5	9	2	6	<- No swap needed
1	3	4	1	5	9	2	6	<- Swap A[2] & A[3] and insert A[2] into A[1] and shift elements
1	1	3	4	5	9	2	6	<- No swap needed
1	1	3	4	5	9	2	6	<- No swap needed
1	1	3	4	5	9	2	6	<- Swap A[5] and A[6] and insert A[6] into A[2] and shift elements
1	1	2	3	4	5	9	6	<- Swap A[6] and A[7], no insertion necessary.
1	1	2	3	4	5	6	9	<- The list is now sorted

2. Use Bubble Sort on the following letters to sort it alphabetically and give each outer step with i and j.

{S T R A I G H T}

Before	<i>i</i>	<i>j</i>	
S T R A I G H T	S	T	No swap
S T R A I G H T	T	R	Swap
S R T A I G H T	T	A	Swap
S R A T I G H T	T	I	Swap
S R A I T G H T	T	G	Swap
S R A I G T H T	T	H	Swap
S R A I G H T T	T	T	Swap
S R A I G H T T	S	R	Swap
R S A I G H T T	S	A	Swap
R A S I G H T T	S	I	Swap
R A I S G H T T	S	G	Swap
R A I G S H T T	S	H	Swap
R A I G H S T T	S	T	No swap
R A I G H S T T	R	A	Swap
A R I G H S T T	R	I	Swap
A I R G H S T T	R	G	Swap

A I G R H S T T	R	H	Swap
A I G H R S T T	R	S	No swap
A I G H R S T T	A	I	No swap
A I G H R S T T	I	G	Swap
A G I H R S T T	I	H	Swap
A G H I R S T T	I	R	No swap
A G H I R S T T	A	G	No swap
A G H I R S T T	G	H	No swap
A G H I R S T T	H	I	No swap
A G H I R S T T	I	R	No swap
A G H I R S T T	R	S	No swap
A G H I R S T T	S	T	No swap
A G H I R S T T	T	T	No swap
<b>A G H I R S T T</b>	-	-	Sorted

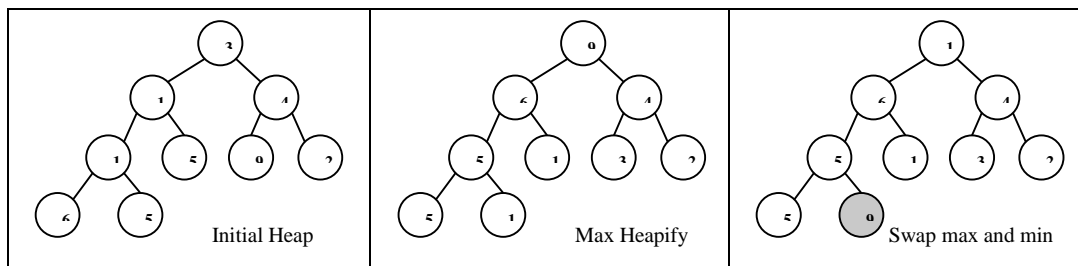
3. What is the number of basic steps executed by the following method (as a function of  $n$ )? What is the big-O of the method?

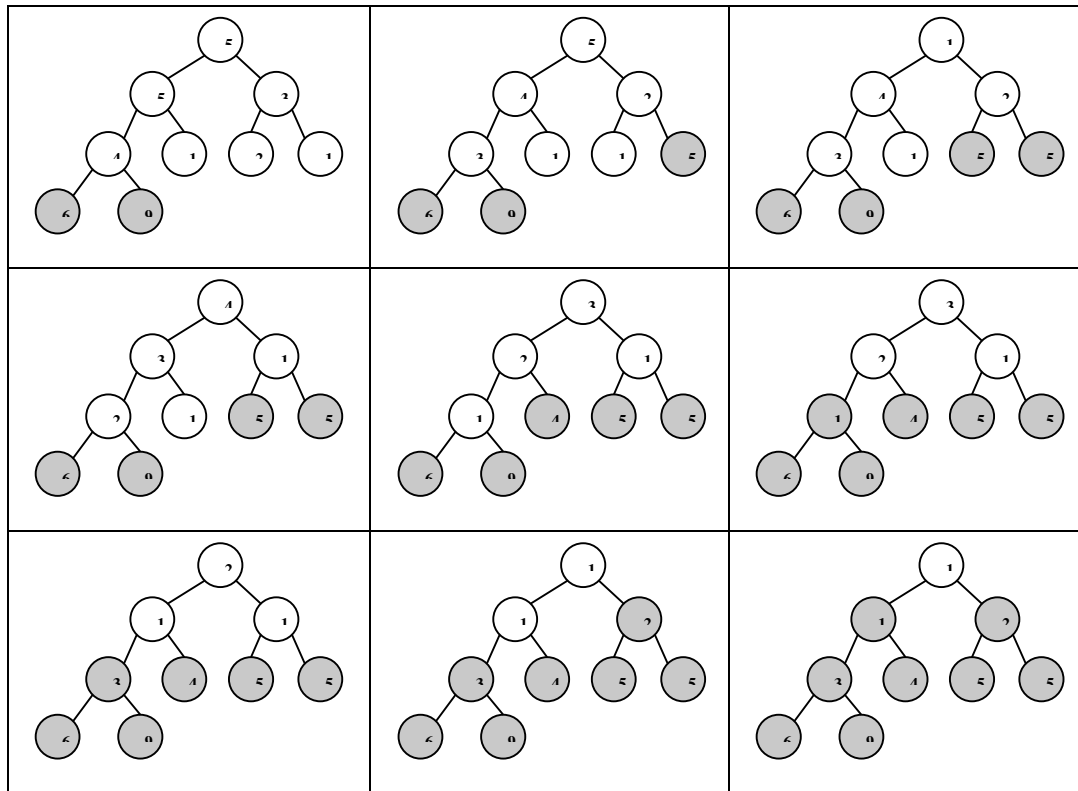
```

Public int howLongA(int n) {
    int k = 0, kk = 0;
    for(int i = 0; i < n/2 ; i++) {
        k++;
        for(int j = 0; j < n/2; j++)
            kk++;
    }
    return k*kk;
}

```

4. Using Figure 6.4 in the textbook as a model, illustrate the operation of HEAPSORT on the array:  
 $A = \{3, 1, 4, 1, 5, 9, 2, 6, 5\}$





**5. Analysis of tertiary heaps:** A ternary heap is like a binary heap, non-leaf nodes have 3 children instead of 2 children.

a. How would you represent a ternary heap in an array?

Solution #1(Index starts a 1):

A binary heap is a special case of a d-ary heap where  $d=2$ . We can extend the properties of d-ary heap to any arbitrary number. In this case where  $d=3$ , we can then use the following formulae for representing the ternary heap in an array:

- Root =  $\text{Floor}[(k+1)/3]$
- Left Child =  $3*k-1$
- Center Child =  $3*k$
- Right Child =  $3*k+1$

Solution #2(Index starts at 0):

- Root =  $((i-2)/3) + 1$
- Left Child =  $3(i-1)+2$
- Center Child =  $3(i-1)+3$
- Right Child =  $3(i-1)+4$

a. What is the height of a ternary heap of  $n$  elements in terms of  $n$ ?

**Answer:**

$$h = \log_3 n$$

b. What is the height of a ternary heap of  $n$  elements in terms of  $n$ ?

c. Give an efficient implementation (in pseudo code) of EXTRACT-MAX in a tertiary max-heap. Analyze its running time in terms of  $n$ .

**Answer:**

<b>HeapExtract-Max(A)</b>	
if A.heap-size < 1	$O(1)$
<b>error</b> "heap underflow"	$O(1)$
max = A[1]	$O(1)$
A[1] = A[A.heap-size]	$O(1)$
A.heap-size = A.heap-size-1	$O(1)$
Max-Heapify(A,1)	$O(\log_3 n)$
<b>return</b> max	Total $O(\log_3 n)$

d. Give an efficient implementation (in pseudo code) of INSERT in a ternary max-heap. Analyze its running time in terms of  $n$ .

**Answer:**

<b>Heap-Insert(A, key)</b>	
A.heap-size++	$O(1)$
i = A.heap-size(A)	$O(1)$
A[i] = $-\infty$	$O(1)$
HeapChangeKey(A, i, key)	$O(\log_3 n)$
	Total $O(\log_3 n)$

e. Give an efficient implementation (in pseudo code) of INCREASE-KEY(A,i,k), which flags an error if  $k < A[i]$ , but otherwise sets  $A[i] = k$  and then updates the ternary maxheap structure appropriately. Analyze its running time in terms of  $n$ .

**Answer:**

<b>HeapIncrease-Key(A, i, key)</b>	
A[i] = key	$O(1)$
while (i > 1 and A[parent(i)] < A[i])	$O(1)$
swap(A[i], A[parent(i)])	$O(\log_3 n)$
	Total $O(\log_3 n)$

