

A Recursive Strategy for Symbolic Execution Expressed in Coq

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1 Data Types

- Object E to represent set of concrete states.
- Object of type n to represent nodes on the symbolic execution tree.
 - $\text{get_s}(n)$ returns s , the symbolic state.
 - $\text{get_pc}(n)$ returns π , the path constraint.
- Object of type \mathcal{E} to represent a symbolic execution tree made of objects of type n .
 - $\text{is_leaf}(\mathcal{E}, n)$ returns true if n is a leaf in \mathcal{E} .
 - $\text{is_root}(\mathcal{E}, n)$ returns true if n is the root in \mathcal{E} .
 - $\text{get_root}(\mathcal{E})$ returns object of type n that is the root of the tree.

Shorthand:

- $\bar{s}_{r,m} = \text{get_s}(\text{get_root}(\mathcal{E}_m))$
- $\pi_{r,m} = \text{get_pc}(\text{get_root}(\mathcal{E}_m))$
- $\bar{s}_{l,m} = \text{get_s}(n_{l,m})$, where $n_{l,m} \in \mathcal{E}_m$.
- $\pi_{l,m} = \text{get_pc}(n_{l,m})$, where $n_{l,m} \in \mathcal{E}_m$.

2 Circle Operations

These use the definitions defined in Arusoae et al.'s paper [ALR13].

Definition 1 $\text{circle_op_1} = \text{the set } \varphi \in \llbracket \beta \rrbracket \ \forall \ \varphi' \in \llbracket \beta' \rrbracket \text{ of a given } \beta' \text{ such that } [\varphi]_{\sim} \Rightarrow^S [\varphi']_{\sim} \text{ where } \beta \Rightarrow \beta'.$

Definition 2 $\text{circle_op_2} = \text{the set } \varphi' \in \llbracket \beta' \rrbracket \ \forall \ \varphi \in \llbracket \beta \rrbracket \text{ of a given } \beta \text{ such that } [\varphi]_{\sim} \Rightarrow^S [\varphi']_{\sim} \text{ where } \beta \Rightarrow \beta'.$

circle_op_1 represents all concrete states that will take us down exactly one path in the symbolic execution tree. circle_op_2 represents all concrete output states of a given path in the symbolic execution tree.

3 Properties

For a given $E, X = \text{a sequence } \mathcal{E}_0, \dots, \mathcal{E}_m \text{ such that } \forall \mathcal{E}_x, \exists n_{l,x} \text{ such that } \text{is_leaf}(\mathcal{E}_x, n_{l,x}) = \text{true} \rightarrow \text{the conjunction of the following:}$

1. $s_0 \in \text{circle_op_1}(\bar{s}_{r,0}, \pi_{l,0})$
2. $E \cap \text{circle_op_2}(\bar{s}_{l,m}, \pi_{l,m}) \neq \{\}$
3. for $j = \{0, \dots, m-1\}$, $\text{circle_op_2}(\bar{s}_{l,j}, \pi_{l,j}) \subseteq \text{circle_op_1}(\bar{s}_{r,j+1}, \pi_{l,j+1})$

4 Next Step

Currently, I am working on unifying the circle operation definitions and the tree representation.

References

- [ALR13] Andrei Arusoaie, Dorel Lucanu, and Vlad Rusu. A generic framework for symbolic execution. In *International Conference on Software Language Engineering*, pages 281–301. Springer, 2013.