CMSC 420, Spring 2019

Homework #1

First & Last Name:
UID (9 digits):
Problem 1: You suck at cooking (15 pts)
Figure 1: Yeah, you totally suck.
Please note: This is a real homework assignment, worth 15% of your homework credit.
Watch the recipes for Butternut Squash, Modified Ramen and the Sandwich of Justice from the YouTube channel You Suck At Cooking and answer the following questions:
Question (a): Butternut Squash (5 pts)
(i) Butternut squash is also known as nature's
(ii) What is the name of the family that butternut squash comes from?
(iii) If you want your butternut squash to be more buttery and nutty , how long must you marinate it for?
(iv) What should you do to counter-act the healthy effects of butternut squash?

(v) Name all 6 (six) variations of squash that you can try if you don't like butternut squash.
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Question (b): Modified Ramen (5 pts)
(i) According to whom do you need ramen to survive if you're in college?
(ii) To counter-act WABS , where should you place the flavor packet ?.
(iii) In the Spicy Peanut ramen variation, why should you use aggression ?
(iv) What kind of cement should you not confuse artery cement with?
(v) What product does the video creator's startup company mail to customers?
Question (c): Sandwich Of Justice (5 pts)
(i) What is the name of the egg that is obsessed with the possible homicide of a fellow egg?
(ii) What is the name of the dog of that particular egg ?
(iii) What does a perfectly timed egg create when placed inside the sandwich?

(iv) Which aspect of the sandwich is compromised if the bread slices aren't wide enough?
(v) After serving this sandwich to somebody, what should you say ?
Question (d): Your own kitchen nightmare(s) (5 pts extra credit)
Use the space below to describe a funny personal failure that involves food preparation. Perhaps you confused sugar with salt or cinnamon with paprika. ¹ Or maybe you thought it was a good idea to warm up olive oil in a pot with the lid on. ² It can be anything you have experienced or witnessed. The top 3 funniest responses will be shared on Piazza, anonymously. If you would rather we don't share your response, let us know in the response itself!

 $^{^{1}}$ Neither one of those things have happened to Jason. 2 Or this one.

Problem 2: Syllabus / Deadlines (15 pts)

To answer some of the questions of this problem, you will need to navigate our course's ELMS page and study our syllabus.

(a) What is the due date of your first homework ?	
(b) What is the due date of your first project ?	
(c) What is the date and time of your midterm?	
(d) What is the building & room where your midterm will take place?	
(e) What is the date and time of your final ?	

Problem 3: Git Basics (20 pts)

To pull the starter code for our projects, you will need to familiarize yourselves with the VCS **Git**, since the entire skeleton code will be hosted on a **GitHub repository** that we maintain for you. To answer this problem effectively, you will need to study up on Git. Please navigate to the page named "More Resources" on our **ELMS** website if you would like a starting point. Of course, Google is always your friend.

While a large subset of Git's functionality can be implemented using various graphical tools, there is **no substitute** for the modularity, ease and cross-platform compatibility of using simple command line instructions from the shell. In this class, you will need to familiarize yourselves with a **very basic** subset of available Git commands.

For every one of the following tasks, write the **full** command that would be necessary to type on a command line in order for you to achieve the relevant task. Your answer should include:

- The system program git.
- The command to be passed to the program git (e.g. add, checkout, push, etc).
- Any flags to be passed to the command, e.g -b, --cached, -u.
- Arguments that the command itself may need (e.g branch names, file paths, names of remotes, entire quoted strings, etc).

You are **quite encouraged** to try the commands out in your console before writing your responses below. In fact, the first four commands are a fine example of how you could start implementing your first project! We will provide you with the answer to the first question to help start you off. Note that this command does not have any **flags**, but it does, of course, have a URL argument.

(a) Clone the repository https://github.com/JasonFil/CMSC420-Spring2019.git into the current directory, using HTTPS.

git clone https://github.com/JasonFil/CMSC420-Spring2019.git

- (c) Rename the current branch to skeletoncode.

 (d) Switch from the current branch to a new branch called myimplem.
- (e) Stage a Java source file which you just created and which you named MyStack. java.
- (f) Make a **commit** from the current branch (myimplem) with message "Fixed stack popping issue.".

(g)	Switch back to the branch skeletoncode.
(h)	Merge the branch myimplem with the current branch, skeletoncode.
(i)	Push your latest commit upstream.
(j)	List all your remotes.
(k)	Delete a remote called server2.
(l)	Assuming you have applied changes to 10,000 different files (e.g by having your IDE automatically create Scaladocs for the entire code base of Twitter), type a single command that will stage the files tracked by git and to which the changes have been made.
(m)	Tweak your global, system-wide Git configuration so that instead of typing git log -1 HEAD to show your last commit's information, you could instead type the shorthand git last .

Problem 4: Sparse Matrix Representations (25 pts)

Whether stored in column or row - major order, two-dimensional matrices can be a big bottleneck in terms of memory if they are **sparse**, that is, they have lots of **null** cells. For example, the matrix shown in Table 1 is quite sparse; only about 14% of it is occupied!

	0	1	2	3	4	5	6	7
0	11	X	Χ	X	X	X	X	X
1	X	X	X	X	X	3	X	X
2	X	-3	4	X	X	X	X	X
3	X	X	X	X	X	X	X	X
4	X	X	11	6	X	X	X	e^2
5	-1	X	X	X	X	X	$\begin{array}{c} X \\ X \\ X \\ X \\ X \\ X \\ \end{array}$	X
6	X	X	X	X	X	X	$\sqrt{3}$	X
7	X	X	X	X	X	X	X	X

Table 1: A sparse matrix. X= null for brevity.

Programming languages such as MATLAB offer built-in functionality (methods or operators included in the "standard library" of the language) to sparsify a matrix. Sparsifying a matrix consists of transforming it into a linked list which contains information only for those cells of the matrix that have been filled in. Every node of this list contains the (i, j) coordinates of the relevant cell, as well as the data originally held in the cell.

To answer some of the following questions, you will have to review **asymptotic notation**. We have a resource page on ELMS, with CMSC250 slides linked, where you can review "big-Oh" $(\mathcal{O}(\cdot))$, "big-Theta" $(\mathcal{O}(\cdot))$ and "big-Omega" $(\Omega(\cdot))$.

For example, if you believe that a certain operation has time complexity at most proportional to M (e.g, 2M, 2.5M, 4M - 7), you should answer with $\mathcal{O}(M)$. If instead you believe that it has time complexity at most proportional to the **square** of M multipled by N (e.g, $10M^2N + 2N$), you should answer with $\mathcal{O}(M^2N)$. If you believe that the operation runs in time **exactly** proportional to $M^2 \cdot N$, you should answer with $\mathcal{O}(M^2 \cdot N)$.

Questions on opposite page...

(i)	Make a drawing of the sparsification of matrix shown in Table 1, assuming row-major order. Recall what your nodes should contain.					
	BEGIN YOUR ANSWER BELOW THIS LINE					
(ii)	Perform the same task, this time assuming column-major order.					
	BEGIN YOUR ANSWER BELOW THIS LINE					
(iii)	Suppose that our original matrix has dimensions $M \times N$. As a function of those two variables, and using "big-Theta" $(\Theta()\cdot)$ notation, what is the space that the sparse matrix occupies in memory in its original representation?					
(iv)	Using "big-Oh" $(\mathcal{O}()\cdot)$ notation, what is the space that the matrix occupies in its sparsified format?					
(v)	Suppose that $0 \le i < M$ and $0 \le j < N$. In its <i>original</i> representation, and using whicheven asymptotic notation you consider appropriate, what is the time complexity of assigning the cell $((i,j))$ a value of, say, 3 in the matrix's original representation?					
(vi)	What is the time complexity of the same operation in the sparsified representation (again, in the most appropriate asymptotic notation)?					

Problem 5: Graph representations (25 pts)

Observe the graph of figure 2.

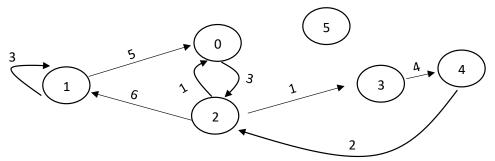


Figure 2: A conceptual image of a graph. Nodes are enumerated from 0 (zero) for consistency with later drawings.

There are two common ways in which graphs, directed or undirected, are represented in computer memory: the **adjacency matrix** (Table 2) and the **adjacency list** (Figure 3).

	0	1	2	3	4	5
0	0	0	3	0	0	0
1	5	3	0	0	0	0
2	1	6	0	1	0	0
3	0	0	0	0	4	0
4	0	0	2	0	0	0
5	0	0	0	0	0	0

Table 2: The **adjacency matrix** representation of the graph of Figure 2.

Let V be the number of **nodes** (or vertices) in a given graph G and E be the number of its **edges**. The adjacency matrix is a $V \times V$ matrix M where:

$$M(i,j) = \begin{cases} w, & \text{where } w > 0 \text{ is the weight of the edge } i \to j \\ 0, & \text{if there is no edge } i \to j \end{cases}$$

In **unweighted** graphs, w = 1 for all edges.

Undirected graphs, on the other hand, have the interesting property that their adjacency matrix is **symmetric**:

$$(\forall i, j \in \{1, 2, \dots, V\})[M(i, j) = M(j, i)]$$

The **adjacency list** representation of this graph is a single-dimensional array A of length V, defined as follows:

$$A(i) = \{v, \text{ if vertex } v \text{ is connected to vertex } i \text{ in } G\}$$

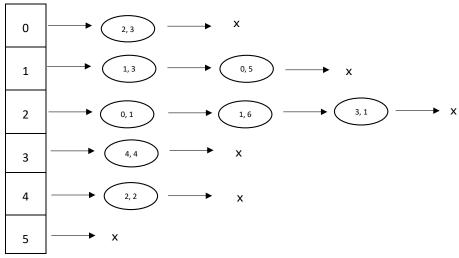


Figure 3: The **adjacency list** representation of the graph of Figure 2. Every node in the list contains a tuple (**target_node_id**, **weight**).

So A(i) points to the **set** of all vertices v on which i is incident. These sets are typically organized as **linked lists**, as implied by Figure 3. These lists do **not have to be sorted in any way**. For this question, we will assume that:

- The heads of the sets are contained in a classic, one-dimensional contiguous storage array (**not** a "dynamic" array like Java's ArrayList and C++'s Vector). So the entire adjacency list representation will consist of an array of lists (see Figure 3).
- The lists themselves are organized as **classic**, **bare-bones linked lists**. This means that we will **not** endow these lists with a pointer to their last element, and we will also **not** have a link pointing to the previous node from every node (i.e the linked lists are **not** doubly-connected) (see Figure 3).

Furthermore, we add a little twist to these representations, based on information from problem 4. We will consider a **sparse adjacency matrix** representation of this graph (linked list of only non-null elements). Figure 4 shows how the graph of Figure 2 is represented in such a format.

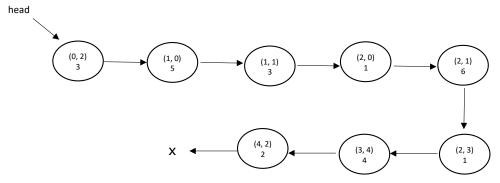


Figure 4: The **sparse adjacency matrix** representation of the graph of Figure 2. The matrix has been scanned in row-major order (but does it really matter which order we use?)

Your task is to fill-in Table 3 using appropriate complexity notation, coupled with the parameters V and E. For example, if you believe that a certain operation will take time *proportional* to the **cube** of the number of nodes V, you should answer with $\mathcal{O}(V^3)$. Or, if you believe that the operation will take time proportional to the product of vertices multiplied by the logarithm base 2 of the number of edges E, you should answer

 $\mathcal{O}(V \cdot \log_2 E)$. Finally, also consider the fact that $0 \le E \le V^2$, since the "worst case" for the number of edges in a graph is V^2 , for a V-clique.

Operation	Adjacency Matrix	Adjacency List	Sparse Adjacency Matrix
Is node i connected to node j ?			
Add a new node k .			
Add the edge $i \to j$.			
Remove a node ℓ .			
Remove the edge $i \to j$.			
Spatial Complexity			

Table 3: Fill in this table with appropriate asymptotic notation.