

Embedding Generalized Parsing in Haskell

Jaro Reinders

Why parser combinators?

For the parser generator crowd

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- Directly parse into a semantic value

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numbers = (+) <$> number <*> char ' ' <*> numbers <|> pure 0
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> parse numbers "2 31 9"  
Just 42
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maybe p = Just <$> p <|> pure Nothing
```

- Monadic parsers enable data-dependent disambiguation (1)

```
ndots 0 = pure ()  
ndots n = char '.' *> ndots (n - 1)
```

```
> parse (number >=> ndots) "5....."  
Just ()
```

Why generalized parsing?

For the parser combinator crowd

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- Left-recursion

$$D ::= 0 \mid 1 \mid \dots \mid 9$$
$$N ::= N D \mid D$$

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- Left-recursion

$$\begin{aligned} D &::= 0 \mid 1 \mid \dots \mid 9 \\ N &::= N D \mid D \end{aligned}$$

- Compositionality

“... it can be quite difficult to determine what language is defined by a TDPL program.” ~ Aho and Ullman (2, p466)

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$$\begin{aligned} D &::= 0 \mid 1 \mid \dots \mid 9 \\ N &::= N D \mid D \end{aligned}$$

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- Disambiguation through annotation rather than deformation

Aside: eliminating left-recursion

“Can’t we just...”

- $N ::= N D \mid D$

becomes

$N ::= D N'$

$N' ::= D N' \mid \varepsilon$

- Complicated for hidden left recursion and semantic values
- Grammar size can grow exponentially

GLL (3)

- Slots, Extended Packed Nodes, Descriptors, Commencements, Continuations
- Essentially building up big set of intermediate results
- $O(n^3)$ time and space

Generalized Parser Combinators

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Generalized Parser Combinators

- Partial normalization up front (Free MonadPlus)
- Then simple driver
- Stack of continuations
- Actions
 - Descend (Push)
 - Loop (Append)
 - Continue (Read)
 - Ascend (Pop)

Stack and continuations

Stack and continuations

- Stack entries consist of

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 - nonterminal name and pivot

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 - Enter the right hand side

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 - Capture a slice up to that occurrence

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 - Append that slice along and current continuation to the loop continuations

Descend & Loop

- When a nonterminal is encountered
- If it is **not** on the stack at the current pivot: Descend
 - Push it to the stack with an empty list of loop continuations and the current continuation
 - Enter the right hand side
- If it is on the stack at the current pivot: Loop
 - Capture a slice up to that occurrence
 - Append that slice along and current continuation to the loop continuations
 - Bail out

Continue & Ascend

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- Continue
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 - Continue with its parser
- Ascend
 - Pop the stack

Continue & Ascend

- When a nonterminal is fully parsed
- Do either of
- Continue
 - Peek at the list of loop continuations
 - Choose one
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 - Continue with its parser
- Ascend
 - Pop the stack
 - Continue with the final continuation

Parsing $5 + 3 + 7$ with $X ::= X + X \mid \mathbb{N}$

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$X_0[+ X]; X_2[+ X]; X_4$	$X + X \mid \mathbb{N}$	7	fail*

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$X_0[+ X]; X_2[+ X]; X_4$	\mathbb{N}	7	parse \mathbb{N}
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$X_0[+ X]; X_2[+ X]$	X	7	descend X
$X_0[+ X]; X_2[+ X]; X_4$	$X + X \mid \mathbb{N}$	7	fail*
$X_0[+ X]; X_2[+ X]; X_4$	\mathbb{N}	7	parse \mathbb{N}
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$X_0[+ X]; X_2[+ X]$	$+ X$	$+ 7$	parse $+$
$X_0[+ X]; X_2[+ X]$	X	7	descend X
$X_0[+ X]; X_2[+ X]; X_4$	$X + X \mid \mathbb{N}$	7	fail*
$X_0[+ X]; X_2[+ X]; X_4$	\mathbb{N}	7	parse \mathbb{N}
$X_0[+ X]; X_2[+ X]; X_4$	ϵ	ϵ	ascend
$X_0[+ X]; X_2[+ X]$	ϵ	ϵ	ascend
$X_0[+ X]$	ϵ	ϵ	ascend
ϵ	ϵ	ϵ	done

Conflict-free nonterminal names

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- Template Haskell (4) quotes

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```
number :: Parser Int
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```
number = 'number
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```
  ::= (\x y → 10 * x + y) <$> number <*> digit  
  <|> digit
```

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- Alternative: GADTs (5)

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data Number a where
  Number :: Number Int
```

Conflict-free nonterminal names

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number :: Parser Int
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```

- Alternative: GADTs (5)

```
data Number a where
  Number :: Number Int
```

Combined with Data Types à la Carte (6)

Conclusion

- We can combine
 - lightweight
 - embedded
 - generalized
 - monadic
- Parser combinators

Future work

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- Disambiguation (Layout, Precedence, Fixity)

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- Memoization

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- Memoization
- Higher-order combinators

Future work

- Disambiguation (Layout, Precedence, Fixity)
- Memoization
- Higher-order combinators
- Actually start writing parsers

References

- (1) Trevor Jim, Yitzhak Mandelbaum, & David Walker. (2010, January). *Semantics and algorithms for data-dependent grammars*. In Proceedings of the 37th annual ACM SIGPLAN-SIGACT symposium on *Principles of programming languages* (pp. 417-430).
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