

# EE 709 : BDD Assignment

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## 1 Describing the 4 bit adder

- The variables for the inputs are created as `bdd` arrays `x[4]`, `y[4]` and for the adder outputs `sum[4]` and `cout`, created using `bdd_new_var_last()` function.
- Temporary `bdd` arrays `s[4]` and `c[4]` are used to represent the actual sum and carry in terms of the input bits. They are recursively defined as follows following **Ripple Carry Adder Architecture**

$$s[0] = x[0] \oplus y[0] \quad (1)$$

$$c[0] = x[0] \cdot y[0] \quad (2)$$

$$s[i] = (x[i] \oplus y[i]) \oplus c[i-1], 1 \leq i \leq 3 \quad (3)$$

$$c[i] = x[i] \cdot y[i] + c[i-1] \cdot (x[i] \oplus y[i]), 1 \leq i \leq 3 \quad (4)$$

## 2 Question 1

### 2.1 Method

- Let  $A$  be the subset of the domain and  $P(x, y)$  denote the function which generates this subset, i.e,  $P(x, y) = 1 \iff$  there are odd number of bits 1 in both  $x[:]$  and  $y[:]$ . Then  $P(x, y)$  is represented as

$$P(x, y) = (x_0 \oplus x_1 \oplus x_2 \oplus x_3) \cdot (y_0 \oplus y_1 \oplus y_2 \oplus y_3) \quad (5)$$

- Let the equations (1), (2), (3), (4) be captured by  $F(x, y)$
- The BDD of the image be represented as  $Q(s, c_{out})$ , which can be obtained using the existential quantification operation.

$$Q(s, c_{out}) = \exists_{x, y} P(x, y) \cdot (c_{out}, s \iff F(x, y)) \quad (6)$$

- Where  $c_{out}, s \iff F(x, y)$  denotes equivalence

$$(c_{out} \iff c[3]) \cdot (sum[i] \iff s[i]), 0 \leq i \leq 3 \quad (7)$$

The entire code can be found in **Appendix (5)**

## 2.2 Image BDD Obtained

bdd of image, where  $\text{var.8} = \text{cout}$

$\text{var.9} = \text{sum}[0]$ ,  $\text{var.10} = \text{sum}[1]$ ,  $\text{var.11} = \text{sum}[2]$ ,  $\text{var.12} = \text{sum}[3]$

```
if var.8
  if var.9
    if var.10
      !var.11
    else if !var.10
      0: if var.11
        !var.12
      else if !var.11
        1
      endif var.11
    endif var.10
  endif var.9
else if !var.8
  if var.9
    if var.10
      subformula 0
    else if !var.10
      1
    endif var.10
  endif var.9
else if !var.8
  if var.9
    if var.10
      if var.11
        var.12
      else if !var.11
        1
      endif var.11
    else if !var.10
      1: if var.11
        1
      else if !var.11
        var.12
      endif var.11
    endif var.10
  else if !var.9
    if var.10
      1
    else if !var.10
      subformula 1
    endif var.10
  endif var.9
endif var.8
```

## 3 Question 2

### 3.1 Method

- Let  $B$  the subset of the range and let  $B(s, c_{out})$  be the function which generates this function, i.e,  $B(s, c) = 1 \iff$  there are odd number of bits 1 in output 5 bits.  $B$  is represented as

$$B(s, c_{out}) = c_{out} \oplus s_3 \oplus s_2 \oplus s_1 \oplus s_0 \quad (8)$$

- Let the BDD of the pre-image be represented as  $S(x, y)$  which can be obtained by the existential quantification operation

$$S(x, y) = \exists_{s,c} B(s, c) \cdot (s, c \iff F(x, y)) \quad (9)$$

- Where  $c_{out}, s \iff F(x, y)$  denotes equivalence

$$(c_{out} \iff c[3]) \cdot (sum[i] \iff s[i]), 0 \leq i \leq 3 \quad (10)$$

The code for the implementation can be found in **Appendix** (5)

### 3.2 Pre-Image BDD Obtained

bdd of pre-image, where var.0 - x[0]  
var.1 - x[1], var.2 - x[2], var.3 - x[3], var.4 - y[0]  
var.5 - y[1], var.6 - y[2], var.7 - y[3]

```
-----
if var.0
  if var.1
    if var.2
      if var.3
        if var.4
          0: if var.5
            1: if var.6
              !var.7
            else if !var.6
              var.7
            endif var.6
          else if !var.5
            !subformula 1
          endif var.5
        else if !var.4
          2: if var.5
            subformula 1
          else if !var.5
            3: if var.6
              !var.7
            else if !var.6
              0
            endif var.6
          endif var.6
        endif var.4
      endif var.3
    endif var.2
  endif var.1
endif var.0
```

```

        endif var.5
    endif var.4
else if !var.3
    if var.4
        4: if var.5
            var.6
        else if !var.5
            !var.6
        endif var.5
    else if !var.4
        5: if var.5
            var.6
        else if !var.5
            6: if var.6
                1
            else if !var.6
                !var.7
            endif var.6
        endif var.5
    endif var.4
endif var.3
else if !var.2
    if var.3
        if var.4
            7: if var.5
                !subformula 3
            else if !var.5
                subformula 3
            endif var.5
        else if !var.4
            8: if var.5
                !subformula 3
            else if !var.5
                !var.6
            endif var.5
        endif var.4
    else if !var.3
        if var.4
            9: if var.5
                !subformula 6
            else if !var.5
                subformula 6
            endif var.5
        else if !var.4
            10: if var.5
                !subformula 6
            else if !var.5
                subformula 1
            endif var.5
        endif var.4
    endif var.3
endif var.2
endif var.1

```

```

        endif var.5
    endif var.4
endif var.3
endif var.2
else if !var.1
    if var.2
        if var.3
            if var.4
                !subformula 2
            else if !var.4
                !subformula 7
            endif var.4
        else if !var.3
            if var.4
                !subformula 5
            else if !var.4
                !subformula 9
            endif var.4
        endif var.3
    else if !var.2
        if var.3
            if var.4
                !subformula 8
            else if !var.4
                !subformula 4
            endif var.4
        else if !var.3
            if var.4
                !subformula 10
            else if !var.4
                subformula 0
            endif var.4
        endif var.3
    endif var.2
endif var.1
else if !var.0
    if var.1
        if var.2
            if var.3
                if var.4
                    subformula 2
                else if !var.4
                    !subformula 2
                endif var.4
            else if !var.3
                if var.4
                    subformula 5
                else if !var.4

```

```

        !subformula 5
    endif var.4
endif var.3
else if !var.2
    if var.3
        if var.4
            subformula 8
        else if !var.4
            !subformula 8
        endif var.4
    else if !var.3
        if var.4
            subformula 10
        else if !var.4
            !subformula 10
        endif var.4
    endif var.3
endif var.2
else if !var.1
    if var.2
        if var.3
            if var.4
                !subformula 7
            else if !var.4
                subformula 7
            endif var.4
        else if !var.3
            if var.4
                !subformula 9
            else if !var.4
                subformula 9
            endif var.4
        endif var.3
    else if !var.2
        if var.3
            if var.4
                !subformula 4
            else if !var.4
                subformula 4
            endif var.4
        else if !var.3
            if var.4
                subformula 0
            else if !var.4
                !subformula 0
            endif var.4
        endif var.3
    endif var.2
endif var.2

```

```

endif var.1
endif var.0

```

## 4 Question 3

### 4.1 Method

Note: The claim *every even 4-bit number ( $n$ ) can be expressed as a sum of two prime numbers*, is true only if the number is greater than 2. **Assumption:**  $n > 2$

- Let  $s : (s_3, s_2, s_1, s_0)$  denote the four bit number, let it be the sum of two four bit numbers  $x : (x_3, x_2, x_1, x_0)$  and  $y : (y_3, y_2, y_1, y_0)$
- Let  $R$  the subset of the range such that  $s > 2$  is even and let  $R(s)$  be the function which generates this function.  $s \in \{4, 6, 8, 10, 12, 14\}$ . On solving K-Map we get

$$R(s) = \bar{s}_0 \cdot (s_2 + s_3) \quad (11)$$

- Let  $M$  be the subset of the domain such that  $x, y$  are both prime. Let  $M(x, y)$  be the generating function for this.  $x, y \in \{2, 3, 5, 7, 11, 13\}$ . On solving K-Map we get  $M(x, y)$  as

$$(x_0 \cdot x_1 \cdot \bar{x}_2 + \bar{x}_3 \cdot \bar{x}_2 \cdot x_1 + \bar{x}_3 \cdot x_2 \cdot x_0 + \bar{x}_1 \cdot x_0 \cdot x_2) \cdot (y_0 \cdot y_1 \cdot \bar{y}_2 + \bar{y}_3 \cdot \bar{y}_2 \cdot y_1 + \bar{y}_3 \cdot y_2 \cdot y_0 + \bar{y}_1 \cdot y_0 \cdot y_2) \quad (12)$$

- The claim can be expressed as the following logical expression

$$\forall_s (R(s) \rightarrow \exists_{x,y} M(x, y) \cdot (s \iff F(x, y))) \quad (13)$$

- To prove the claim, we need to show that the above expression is equivalent to a tautology.
- Where  $s \iff F(x, y)$  denotes equivalence

$$(sum[i] \iff s[i]), 0 \leq i \leq 3 \quad (14)$$

### 4.2 Implementation

This **if-else** block checks if the final BDD for the logical expression equals to the bdd for one.

```

if(final == bdd_one(bddm)){
    printf("The claim is correct, Hence Proved!\n");
}
else{
    printf("The claim is incorrect, Hence Disproved!\n");
}

```

The entire code can be found in **Appendix (5)**

### 4.3 Output Obtained

The claim is correct, Hence Proved!

## 5 Appendix for Code

### 5.1 fourbitadder.c

```
#include <stdlib.h>
#include <stdio.h>
#include <bdduser.h>

// number of bits in adder

void visualize_bdd(bdd_manager bddm, bdd x){
    printf("-----\n");
    bdd_print_bdd(bddm,x,NULL, NULL,NULL, stdout); //prints bdd
    printf("-----\n");
}

int main (int argc, char* argv[])
{
    // create the universe

    bdd_manager bddm = bdd_init();

    // make inputs of adder

    bdd x[4];
    for(int i=0;i<4;i++){
        x[i] = bdd_new_var_last(bddm);
    }

    bdd y[4];
    for(int i=0;i<4;i++){
        y[i] = bdd_new_var_last(bddm);
    }

    // sum and carry of adder

    bdd s[4];
    bdd c[4];

    // first bit of adder

    s[0] = bdd_xor(bddm, x[0], y[0]);
    c[0] = bdd_and(bddm, x[0], y[0]);

    // subsequent rippling of adder bits

    for(int i=1; i<4; i++){
        s[i] = bdd_xor(bddm, x[i], y[i]);
        bdd temp = bdd_and(bddm, s[i], c[i-1]);
        s[i] = bdd_xor(bddm, s[i], c[i-1]);
    }
}
```



```

        c[i] = bdd_and(bddm, x[i], y[i]);
        c[i] = bdd_or(bddm, c[i], temp);
    }

    bdd cout = bdd_new_var_last(bddm);
    bdd sum[4];

    for(int i=0; i<4; i++){
        sum[i] = bdd_new_var_last(bddm);
    }

    // approach to question 1

    bdd x_has_odd_high_bits = bdd_xor(bddm, x[0], x[1]);

    for(int i=2; i<4; i++){
        x_has_odd_high_bits = bdd_xor(bddm, x[i], x_has_odd_high_bits);
    }

    bdd y_has_odd_high_bits = bdd_xor(bddm, y[0], y[1]);

    for(int i=2; i<4; i++){
        y_has_odd_high_bits = bdd_xor(bddm, y[i], y_has_odd_high_bits);
    }

    // function P(x, y) for subset A

    bdd has_odd_high_bits = bdd_and(bddm, x_has_odd_high_bits, y_has_odd_high_bits);

    // function equivalence check c, s <-> F(x, y)

    bdd equivalence = bdd_xnor(bddm, cout, c[4 - 1]);

    for(int i=0; i<4; i++){
        bdd curr_eq = bdd_xnor(bddm, s[i], sum[i]);
        equivalence = bdd_and(bddm, equivalence, curr_eq);
    }

    bdd condition = bdd_and(bddm, equivalence, has_odd_high_bits);

    // To get image apply existence quantification
    // Q(c, s) = Exists x,y P(x, y).(c, s <-> F(x))

    bdd quantified_vars[9];

    for(int i=0; i<4; i++){
        quantified_vars[i] = x[i];
    }

    for(int i=4; i<8; i++){
        quantified_vars[i] = y[i - 4];
    }

```

```

quantified_vars[8] = 0;

int assoc = bdd_new_assoc(bddm, quantified_vars, 0);
bdd_assoc(bddm, assoc);

// bdd for the image of the set

bdd image = bdd_exists(bddm, condition);

printf("bdd of image, where var.8 - cout\n");
printf("var.9 - sum[0], var.10 - sum[1], var.11 - sum[2], var.12 - sum[3]\n");
visualize_bdd(bddm, image);

bdd_free_assoc(bddm, assoc);

// approach for question 2

// bdd for B(s, c)

bdd output_odd_high_bits = bdd_xor(bddm, sum[0], cout);

for(int i=1; i<4; i++){
    output_odd_high_bits = bdd_xor(bddm, sum[i], output_odd_high_bits);
}

// to find the preimage of set B

condition = bdd_and(bddm, equivalence, output_odd_high_bits);

// To get pre-image apply existence quantification
// S(x, y) = Exists s,c B(s, c).(s, c <-> F(x, y))

for(int i=0; i<4;i++){
    quantified_vars[i] = sum[i];
}

quantified_vars[4] = cout;

quantified_vars[5] = 0;

assoc = bdd_new_assoc(bddm, quantified_vars, 0);
bdd_assoc(bddm, assoc);

// bdd for the pre-image of the set

bdd preimage = bdd_exists(bddm, condition);

printf("bdd of pre-image, where var.0 - x[0]\n");
printf("var.1 - x[1], var.2 - x[2], var.3 - x[3], var.4 - y[0]\n");
printf("var.5 - y[1], var.6 - y[2], var.7 - y[3]\n");
visualize_bdd(bddm, preimage);

```

```

bdd_free_assoc(bddm, assoc);

// approach for question 3

bdd even_four_bit_num = bdd_and(bddm, bdd_not(bddm, sum[0]), bdd_or(bddm, sum[2],
↪ sum[3])); // R(s)

bdd x_is_prime = bdd_or(bddm, bdd_and(bddm, x[0], bdd_xor(bddm, x[1], x[2])),
↪ bdd_and(bddm, bdd_not(bddm, x[3]), bdd_or(bddm, bdd_and(bddm, x[0], x[2]),
↪ bdd_and(bddm, x[1], bdd_not(bddm, x[2]))));
bdd y_is_prime = bdd_or(bddm, bdd_and(bddm, y[0], bdd_xor(bddm, y[1], y[2])),
↪ bdd_and(bddm, bdd_not(bddm, y[3]), bdd_or(bddm, bdd_and(bddm, y[0], y[2]),
↪ bdd_and(bddm, y[1], bdd_not(bddm, y[2]))));
bdd x_is_prime_and_y_is_prime = bdd_and(bddm, x_is_prime, y_is_prime); // M(x, y)

// s <-> F(x, y)
equivalence = bdd_xnor(bddm, s[0], sum[0]);

for(int i=1; i<4; i++){
    equivalence = bdd_and(bddm, equivalence, bdd_xnor(bddm, s[i], sum[i]));
}

for(int i=0; i<4; i++){
    quantified_vars[i] = x[i];
}

for(int i=4; i<8; i++){
    quantified_vars[i] = y[i - 4];
}

quantified_vars[8] = 0;

assoc = bdd_new_assoc(bddm, quantified_vars, 0);
bdd_assoc(bddm, assoc);

// M(x,y).(s <-> F(x, y))

condition = bdd_and(bddm, equivalence, x_is_prime_and_y_is_prime);

// Exists x,y M(x,y).(s <-> F(x, y))
bdd exists = bdd_exists(bddm, condition);

// R(s) -> Exists x,y M(x,y).(s <-> F(x, y))
bdd implication = bdd_or(bddm, exists, bdd_not(bddm, even_four_bit_num));

bdd_free_assoc(bddm, assoc);

// to prove for all outputs
for(int i=0; i<4; i++){
    quantified_vars[i] = sum[i];
}

```

```

quantified_vars[4] = 0;

assoc = bdd_new_assoc(bddm, quantified_vars, 0);
bdd_assoc(bddm, assoc);

// Forall s [ R(s) -> Exists x,y M(x,y).(s <-> F(x, y)) ]
bdd final = bdd_forall(bddm, implication);

if(final == bdd_one(bddm)){
    printf("The claim is correct, Hence Proved!\n");
}
else{
    printf("The claim is incorrect, Hence Disproved!\n");
}

bdd_free_assoc(bddm, assoc);
bdd_quit(bddm);
return(0);
}

```

## 5.2 compile.sh

```

gcc -o fourbitadder fourbitadder.c -I ../cmu_bdd/include -L ../cmu_bdd/lib -lbdd -lmem
./fourbitadder > outputs.txt

```

## 6 Submission

The submission contains the .pdf report, .c source file, fourbitadder executable file and the outputs file outputs.txt

- In this case, the fourbitadder.c is kept in the following directory structure.

```

-
|_ cmu_bdd/
|_ folder/
    |_ compile.sh
    |_ fourbitadder.c

```

- run the bash script as ./compile.sh to give final directory structure

```

-
|_ cmu_bdd/
|_ folder/
    |_ compile.sh
    |_ fourbitadder.c
    |_ fourbitadder
    |_ output.txt

```