Code

```
#include <bits/stdc++.h>
#include <utility>
using namespace std;
#ifndef ONLINE_JUDGE
#include "Debug.cpp"
#define FileI(fileName) freopen(fileName, "r", stdin);
#define FileO(fileName) freopen(fileName, "w", stdout);
#else
#define FileI(fileName)
#define FileO(fileName)
#define debug(...) 1
#define Time(i, x...) x
#endif
#define un unsigned
#define ld long double
#define LL long long
#define allc(x) begin(x), end(x)
#define rallc(x) rbegin(x), rend(x)
#define all(x,i,j) begin(x) + i, begin(x) + j
#define Test int TC, _(0); cin >> TC; while(_++ < TC)</pre>
template<typename T>
using HashedBoard = vector<pair<T, T>>;
string output, expected;
const string Dir[] = {"left", "up", "right", "down"};
struct Graph
```

```
{
    enum Direction {NONE = -1, LEFT, UP, RIGHT, DOWN};
    enum Algorithm {S_LOG_S, N_SQUARED};
    struct Node
    {
        HashedBoard<un short> ballBoardState;
        Direction parentDir = NONE;
        int sourceDistance = 0;
        Node* Parent = nullptr;
        Node() = default;
        Node(const HashedBoard<un short>& ballBoardState, i
nt sourceDistance = 0, Node* Parent = nullptr, Direction pa
rentDir = NONE)
        {
            this->ballBoardState = ballBoardState;
            this->sourceDistance = sourceDistance;
            this->Parent = Parent;
            this->parentDir = parentDir;
        }
    };
    unsigned moveTimer = 0; // counter for how many times w
e call move method
    Node* initialState, *finalState = nullptr;
    vector<vector<char>> globalBoard; // grid without balls
    vector<vector<pair<un short, un short>>> obstaclesRowPo
s, obstaclesColPos; // used for constant queries about the
nearest obstacle in 4 main directions
    vector<pair<un short, un>> ballsRowPos; // dynamically
updating the nearest ball (above or below) me for constant
queries
    pair<un short, un short> target; // coordinates of the
target
    const short toBase = UCHAR_MAX + 1; // toBase = 256
    vector<vector<string>> baseChar; // constant queries fo
r ContainerHashing function
    Algorithm choosenAlgo = N_SQUARED; // algorithm to appl
y on vertical moves
```

```
vector<vector<Direction>> adj = {{UP, DOWN}, {LEFT, RIG
HT}, {UP, DOWN}, {LEFT, RIGHT}}; // for every direction wha
t are the next directions
    Graph(const vector<vector<char>>& board, const pair<un</pre>
short, un short>& target) // O(max(slogs, n^2))
    {
        HashedBoard<un short> slidersPos;
        int n = board.size() - 2;
        baseChar = vector<vector<string>>(n + 2, vector<str</pre>
ing>(n + 2));
        for (int i = 0; i \le n + 1; i++)
            for (int j = 0; j \le n + 1; j++)
                ConvFromDec(i, baseChar[i][j]), ConvFromDec
(j, baseChar[i][j]);
        globalBoard = board;
        ballsRowPos = vector<pair<un short, un>>(n + 2, {-
1, moveTimer});
        obstaclesRowPos = obstaclesColPos = vector<vector<p</pre>
air<un short, un short>>>(n + 2, vector<pair<un short, un s
hort >> (n + 2));
        for (int i = 0; i \le n + 1; i++)
            for (int lj = 0, rj = n + 1, j = 0, jj = n + 1;
j <= n + 1; j++, jj--)
            {
                if (board[i][j] != '#')
                    obstaclesColPos[i][j].first = lj;
                else
                    1j = j;
                if (board[i][jj] != '#')
                    obstaclesColPos[i][jj].second = rj;
                else
                    rj = jj;
                if (board[i][j] == 'o')
                    slidersPos.emplace_back(i, j), globalBo
ard[i][j] = '.';
```

```
for (int j = 0; j \le n + 1; j++)
            for (int ui = 0, di = n + 1, i = 0, ii = n + 1;
i <= n + 1; i++, ii--)
            {
                if (board[i][j] != '#')
                     obstaclesRowPos[j][i].first = ui;
                else
                    ui = i;
                if (board[ii][j] != '#')
                     obstaclesRowPos[j][ii].second = di;
                else
                    di = ii;
            }
        initialState = new Node(slidersPos); // first state
have HashedBoard of balls sorted about rows then cols
        this->target = target;
        auto ci = chrono::high_resolution_clock::now();
        Move(slidersPos, UP);
        auto cf = chrono::high_resolution_clock::now();
        auto squareTime = chrono::duration_cast<chrono::mic</pre>
roseconds>(cf - ci).count();
        choosenAlgo = S_LOG_S;
        ci = chrono::high_resolution_clock::now();
        Move(slidersPos, UP);
        cf = chrono::high_resolution_clock::now();
        auto s lq s = chrono::duration cast<chrono::microse</pre>
conds>(cf - ci).count();
        if (squareTime <= s lq s)</pre>
            choosenAlgo = N SQUARED;
        debug(squareTime);
        debug(s_lg_s);
    }
    HashedBoard<un short> Move(const HashedBoard<un short>&
currentBoard, Direction dir) // O(\min(s*log(s), n^2)), \Omega(s)
```

```
{
        moveTimer++;
        un s = currentBoard.size(); // number of sliders
        switch (dir)
        {
            case RIGHT:
            {
                auto nextBoard = currentBoard;
                un short mnj;
                int it = s - 1;
                if (s)
                    nextBoard[it].second = obstaclesColPos
[nextBoard[it].first][nextBoard[it].second].second - 1;
                for (it--; ~it; it--)
                {
                    auto& [x, y] = nextBoard[it];
                    mnj = obstaclesColPos[x][y].second;
                    if (nextBoard[it + 1].first == x)
                        mnj = min(mnj, nextBoard[it + 1].se
cond);
                    y = mnj - 1;
                }
                return nextBoard;
            }
            case LEFT:
                auto nextBoard = currentBoard;
                un short mxj;
                int it = 0;
                if (s)
                    nextBoard[it].second = obstaclesColPos
[nextBoard[it].first][nextBoard[it].second].first + 1;
                for (it++; it < s; it++)
                {
                    auto& [x, y] = nextBoard[it];
                    mxj = obstaclesColPos[x][y].first;
                    if (nextBoard[it - 1].first == x)
                        mxj = max(mxj, nextBoard[it - 1].se
```

```
cond);
                    y = mxj + 1;
                 }
                 return nextBoard;
            }
            case UP:
            {
                 switch (choosenAlgo)
                     case S_LOG_S:
                     {
                         HashedBoard<un short> nextBoard;
                         un short mxi;
                         for (int it = 0; it < s; it++)
                         {
                             auto &[x, y] = currentBoard[i
t];
                             mxi = obstaclesRowPos[y][x].fir
st;
                             if (ballsRowPos[y].second < mov</pre>
eTimer)
                                 ballsRowPos[y].second = mov
eTimer;
                             else
                                 mxi = max(mxi, ballsRowPos
[y].first);
                             ballsRowPos[y].first = mxi + 1;
                             nextBoard.emplace_back(mxi + 1,
y);
                         stable_sort(allc(nextBoard));
                         return nextBoard;
                     }
                     case N_SQUARED:
                     {
                         un n = globalBoard.size() - 2;
                         HashedBoard<un short> nextBoard;
                         un short mxi;
```

```
for (int it = 0; it < s; it++)
                         {
                             auto \&[x, y] = currentBoard[i]
t];
                             mxi = obstaclesRowPos[y][x].fir
st;
                             if (ballsRowPos[y].second < mov</pre>
eTimer)
                                  ballsRowPos[y].second = mov
eTimer;
                             else
                                  mxi = max(mxi, ballsRowPos
[y].first);
                             ballsRowPos[y].first = mxi + 1;
                             globalBoard[mxi + 1][y] = 'o';
                         }
                         for (int i = 1; i \le n; i++)
                             for (int j = 1; j <= n; j++)
                                  if (globalBoard[i][j] ==
'0')
                                      globalBoard[i][j] =
'.', nextBoard.emplace_back(i, j);
                         return nextBoard;
                     }
                 }
            }
            case DOWN:
            {
                 switch (choosenAlgo)
                 {
                     case S_LOG_S:
                     {
                         HashedBoard<un short> nextBoard;
                         un short mni;
                         for (int it = s - 1; ~it; it--)
                         {
                             auto& [x, y] = currentBoard[i
t];
```

```
mni = obstaclesRowPos[y][x].sec
ond;
                              if (ballsRowPos[y].second < mov</pre>
eTimer)
                                  ballsRowPos[y].second = mov
eTimer;
                              else
                                  mni = min(mni, ballsRowPos
[y].first);
                              ballsRowPos[y].first = mni - 1;
                              nextBoard.emplace_back(mni - 1,
y);
                         }
                         stable_sort(allc(nextBoard));
                         return nextBoard;
                     }
                     case N_SQUARED:
                     {
                         un n = globalBoard.size() - 2;
                         HashedBoard<un short> nextBoard;
                         un short mni;
                         for (int it = s - 1; ~it; it--)
                         {
                              auto\& [x, y] = currentBoard[i]
t];
                             mni = obstaclesRowPos[y][x].sec
ond;
                              if (ballsRowPos[y].second < mov</pre>
eTimer)
                                  ballsRowPos[y].second = mov
eTimer;
                             else
                                  mni = min(mni, ballsRowPos
[y].first);
                              ballsRowPos[y].first = mni - 1;
                              globalBoard[mni - 1][y] = 'o';
                         }
                         for (int i = 1; i \le n; i++)
```

```
for (int j = 1; j \le n; j++)
                                 if (globalBoard[i][j] ==
'0')
                                     globalBoard[i][j] =
'.', nextBoard.emplace_back(i, j);
                         return nextBoard;
                     }
                }
            }
            case NONE:
                return currentBoard;
        }
        return currentBoard;
    }
    vector<vector<char>> Debug(HashedBoard<un short> c)
    {
        vector<vector<char>> tmp = globalBoard;
        for (auto\& [x, y] : c)
            tmp[x][y] = 'o';
        return tmp;
    }
    void PrintAnswer(int Difficulty)
    {
        Time(Algorithm, ShortestPath_BFS();)
        if (!finalState)
            return cerr << (output += "Unsolvable\n"), void</pre>
();
        auto curr = finalState;
        deque<Direction> path;
        while (curr)
        {
            path.push_front(curr->parentDir);
            curr = curr->Parent;
        }
```

```
int k = path.size();
        output += "Solvable\nMin number of moves: " + to_st
ring(k - 1) + "\nSequence of moves: ";
        for (int i = 1; i < k; i++)
            output += Dir[path[i]] + ", ";
        cerr << (output += '\n');</pre>
        if (!Difficulty)
            return;
        auto curBallBoardState = initialState->ballBoardSta
te;
        cerr << "Initial\n";</pre>
        debug(Debug(curBallBoardState));
        for (int i = 1; i < k; i++)
            cerr << Dir[path[i]] << "\n", debug(Debug(curBa</pre>
1lBoardState = Move(curBallBoardState, path[i])));
    }
    inline void ConvFromDec(un short n, string& ret) // O(1
og256(n))
    {
```

```
while (n)
            ret += (n % toBase), n /= toBase;
        ret += ' ';
    }
    string ContainerHashing(const HashedBoard<un short>& co
ntainer) // O(s*log256(n))
    {
        string ret;
        un s = container.size();
        for (int it = 0; it < s; it++)
            ret += baseChar[container[it].first][container
[it].second];
        ret.pop_back();
        return ret;
    }
    void ShortestPath_BFS() // O(2^min(k, d) * min(s log
(s), n^2) : (k) is the depth of the answer if solvable,
(d) is the max depth of the graph, (d) is O(\log((n^2 - b))
C(s)))
    {
        queue<Node*> nxt;
        unordered set<string> vis;
        vis.insert(ContainerHashing(initialState->ballBoard
State));
        if (binary_search(allc(initialState->ballBoardStat
e), target))
            return finalState = initialState, void();
        for (auto dir : {LEFT, UP, RIGHT, DOWN})
        {
            auto newBoardState = move(Move(initialState->ba
llBoardState, dir));
            if (vis.insert(ContainerHashing(newBoardStat
e)).second)
            {
                auto child = new Node(newBoardState, initia
```

```
lState->sourceDistance + 1, initialState, dir);
                if (binary_search(allc(newBoardState), targ
et))
                    return finalState = child, void();
                nxt.push(child);
            }
        }
        while (!nxt.empty())
        {
            auto cur = nxt.front(); nxt.pop();
            for (auto dir : adj[cur->parentDir])
            {
                auto newBoardState = move(Move(cur->ballBoa
rdState, dir));
                if (vis.insert(ContainerHashing(newBoardSta
te)).second)
                {
                    auto child = new Node(newBoardState, cu
r->sourceDistance + 1, cur, dir);
                    if (binary_search(allc(newBoardState),
target))
                        return finalState = child, void();
                    nxt.push(child);
                }
            }
        }
        return void();
    }
};
void Solve()
{
    short chooseDifficulty; cout << "Choose Test Difficulty"
1 : sample , 2 : Hard .\n"; cin >> chooseDifficulty;
    string FileName, outFileName;
    if (--chooseDifficulty)
    {
        FileName += "Complete Tests/";
```

```
cout << "\nChoose Case size: ";</pre>
        short sz; cin >> sz;
        if (sz == 1)
            FileName += "1 small/";
        else if (sz == 2)
            FileName += "2 medium/";
        else if (sz == 3)
            FileName += "3 large/";
        short inputCase; cout << "\nChoose Test Case File:</pre>
";
        cin >> inputCase;
        FileName += "Case " + to_string(inputCase) + "/Cas
e" + to_string(inputCase);
    }
    else
    {
        short inputCase; cout << "\nChoose Test Case File:</pre>
";
        cin >> inputCase;
        FileName = "Sample Tests/Case" + to_string(inputCas
e);
    }
    outFileName = FileName + "-output.txt";
    FileName += ".txt";
    FileI(FileName.c_str());
    int n; cin >> n;
    vector<vector<char>> grid(n + 2, vector<char>(n + 2,
'#'));
    char cell;
    for (int i = 1; i \le n; i++)
        for (int j = 1; j <= n; j++)
        {
            while (cin >> cell, cell != '#' && cell != '.'
&& cell != 'o');
            grid[i][j] = cell;
```

```
}
   pair<un short, un short> target;
   cin >> target.second, cin.ignore(), cin >> target.firs
t;
   target.first++;
   target.second++;
   Graph g(grid, target);
   g.PrintAnswer(chooseDifficulty - 1);
   fclose(stdin);
   FileI(outFileName.c_str())
   string tmp;
   while (getline(cin, tmp) && tmp != "Initial")
      expected += tmp + '\n';
   cerr << "\t\t\t\t" << (output == expected ? "ACCEPTE</pre>
}
signed main()
   Solve(), cout << '\n';
   return 0;
}
```

Running Time For Hard Tests

```
test:1,1 \rightarrow 45.5s
test:1,2 \rightarrow 32.4s
test:1,3 \rightarrow 9.6s
test:2,1 \rightarrow 10.8s
```

```
test:2,2 \rightarrow 0.012s
test:2,3 \rightarrow 29s
test3,1 \rightarrow 0.018s
test3,2 \rightarrow 4.5s
```

test:3,3 → 0.102s

Analysis

ConvFromDec(n)

given:

- 1) log(base, number) is the way to represent the log function
- 2) to Base = $UCHAR_MAX + 1 = 256$

this function takes the number and keeps dividing this number by toBase

- → Body of the loop consists of arithmetic operations so it takes Constant steps
- \rightarrow $\Theta(1)$
- → the loop it self takes exactly log(toBase, n) steps
- \rightarrow this is $\Theta(\log(n))$

ContainerHashing(container)

given:

1) n = container size

this function takes the container and keeps looping on its size

- → Body of the loop consists of arithmetic operations so it takes Constant steps
- $\rightarrow \Theta(1)$
- → the loop it self takes exactly n steps
- \rightarrow this is $\Theta(n)$

Debug(hashedBoard)

given:

1) n = hashed board size

this function takes the hashed board and keeps looping on its size

- ightarrow Body of the loop consists of assignment operation so it takes Constant steps
- → Θ(1)
- → the loop it self takes exactly n steps
- \rightarrow this is $\Theta(n)$

Move(board, direction)

code changed not comp

given:

- 1) s = number of sliders in the given grid
- 2) n*n is the size of the grid
- 3) this function has 2 modes that can work based on deciding what is minimum (s*log(s) and n*n) or s in the other case

let's assume that we will move 'UP'

 \rightarrow if s*log(s) > n*n \rightarrow this code take 2*s + n*n so it takes $\Theta(s + n*n) = \Theta(n^2)$

 \rightarrow if s*log(s) < n*n \rightarrow this code takes s + s*log(s) so it takes $\Theta(s + s*log(s)) = \Theta(s*log(s))$

```
case S_LOG_S:
{
      HashedBoard<un short> nextBoard; /// 0(1)
      un short mxi; /// \Theta(1)
      for (int it = 0; it < s; it++) /// 0(s)
      {
               auto &[x, y] = currentBoard[it]; /// 0(1)
               mxi = obstaclesRowPos[y][x].first; /// 0(1)
               if (ballsRowPos[y].second < moveTimer) /// 0</pre>
(1)
                      ballsRowPos[y].second = moveTimer; ///
\Theta(1)
               else /// 0(1)
                      mxi = max(mxi, ballsRowPos[y].first);
/// O(1)
               ballsRowPos[y].first = mxi + 1; /// \theta(1)
               nextBoard.emplace_back(mxi + 1, y); /// 0(1)
       }
       stable_sort(allc(nextBoard)); /// O(slog(s))
```

```
return nextBoard;/// 0(s)
}
```

NOTE: the 'DOWN' runs in the same time complexity with little to no difference in constant times so we can consider that \rightarrow 'DOWN' = 'UP' = Θ (min(s*log(s), n*n))

let's assume that we will move 'RIGHT'

 \rightarrow we have 1 choice here and it is $\Theta(s)$

NOTE: the 'LEFT' runs in the same time complexity with little to no difference in constant times so we can consider that \rightarrow 'RIGHT' = 'LEFT' = $\Theta(s)$

In conclusion, this function depends on the direction of the move if it is right or left it runs in $\Theta(s)$ but if it is up or down it runs in $\Theta(min(s*log(s), n*n))$

 \rightarrow so the over all function complexity is $\Omega(s)$ and $\Omega(min(s*log(s), n*n))$

shortestPath_BFS()

given:

- 1) s = number of sliders in the given grid
- 2) n*n is the size of the grid
- 3) k = optimal path size (number of steps to reach the goal if it exists)
- 4) $d = max depth of tree (where <math>k \le d$)

we can divide this function to three main blocks so let's analyze them

1- first block complexity:

Θ(s*log(s))

```
queue<Node*> nxt; /// 0(1)
unordered_set<string> vis; /// 0(1)
vis.insert(ContainerHashing(initialState->ballBoardState));
/// 0(1)

if (binary_search(allc(initialState->ballBoardState), targe
t)) /// 0(slog(s))
    return finalState = initialState, void(); /// 0(1)
```

2- second block complexity:

because of observations in the third block we can conclude that calling the MOVE() will result in exact complexity of $\Theta(s + min(s*logs, n^2))$

```
for (auto dir : {LEFT, UP, RIGHT, DOWN}) /// 0(4*MOVE())
{
   auto newBoardState = move(Move(initialState->ballBoardState, dir)); /// 0(MOVE())
   if (vis.insert(ContainerHashing(newBoardState)).second)
/// 0(s)
   {
```

```
auto child = new Node(newBoardState, initialState->s
ourceDistance + 1, initialState, dir); /// 0(1)

if (binary_search(allc(newBoardState), target)) /// 0(sl
og(s))

    return finalState = child, void(); /// 0(1)
    nxt.push(child); /// 0(1)
}
```

3- third block is where we have the main logic of the BFS

first let's consider some facts that will be useful for the analysis:

- max nodes that the BFS will touch is 4 * (2^(d+1) 1) such that d is the depth of the tree
- (2^(k) 1) can become (2^(d+1) 1) if we build the whole tree
- sometimes we will break of this loop if the answer was found so this will be less than (2^(k) - 1)

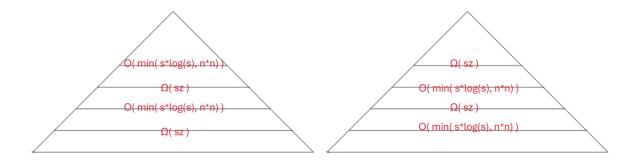
an interesting observation:

Move() can be called 4 times at the beginning however as we go we can call it 2 times as we go.. and another thing is that if we move 'LEFT' or 'RIGHT' we can move 'UP' or 'DOWN' after that and vice versa.. when we consider the 2 facts we conclude that 2 calls from the initial 4 calls will make Move() run in Ω (s) then O(min(s*log(s), n*n)) and the other calls will run O(min(s*log(s), n*n)) then Ω (s).. so this means half of the calls for MOVE() will be Ω (s) and the other half will be

O(min(s*log(s), n*n)).. so although MOVE() doesn't have an exact complexity we know that we will call it from BFS in O($2*(2^{(d+1)} - 1)*min(s*log(s), n*n)$).

We can't conclude that this block have an exact complexity because as stated earlier sometimes we won't enter this loop at all so this is $\Omega(1)$

these pyramids represents the levels of the tree



```
while (!nxt.empty())
/// the touches of nodes is analyzed above
    auto cur = nxt.front(); nxt.pop();
    ///adj.size() = 2 except one time where it equals 4
    for (auto dir : adj[cur->parentDir])
    {
        auto newBoardState = move(Move(cur->ballBoardState,
dir)); ///0(Move())
        if (vis.insert(ContainerHashing(newBoardState)).sec
ond)
        {
            auto child = new Node(newBoardState, cur->sourc
eDistance + 1, cur, dir);
            if (binary_search(allc(newBoardState), target))
/// 0(slogs)
                return finalState = child, void(); /// 0
(1)
            nxt.push(child); /// 0(1)
        }
    }
return void();
```

so let's combine the three blocks of code together:

```
Ω(s + min(slogs, n^2) + 1)

→ this becomes Ω(min(slogs, n^2))
```

```
O(s + k + min(slogs, n^2) + k + (2<sup>(d+1)</sup> - 1) * min(s*log(s), n^n)

\rightarrow this becomes O(k + (2<sup>(d+1)</sup> - 1) * min(s*log(s), n^n)
```

Graph()

given:

1) n*n is the size of the grid

let's break this constructor function into 2 main blocks

1- first block: we have n*n nested loop that calls ConvFromDec() each call runs in logarithmic time so this runs in $\Sigma[\log(i)]$ such that (i is from 0 to n+1)

so this block runs in O(n*n*(log(n)*2)) = O(n*n*log(n))

```
vector<pair<un short, un short>> slidersPos; /// \theta(1) int n = board.size() - 2; /// \theta(1) globalBoard = board; /// \theta(n*n) ballsRowPos = vector<pair<un short, unsigned>>(n + 2, {-1, moveTimer}); /// \theta(n) baseChar = vector<vector<string>>(n + 2, vector<string>(n + 2)); /// \theta(n*n) for (int i = 0; i <= n + 1; i++) for (int j = 0; j <= n + 1; j++) ConvFromDec(i, baseChar[i][j]), ConvFromDec(j, baseChar[i][j]);
```

2-second block: this blocks performs n*n moves then it calls MOVE() so it becomes $\Omega(n*n + s + k)$ or O(n*n + min(s*log(s), n*n) + k)

```
/// O(n*n)
obstaclesRowPos = obstaclesColPos = vector<vector<pair<un s
hort, un short>>>(n + 2, vector<pair<un short, un short>>(n
+ 2));
/// O(n*n)
for (int i = 0; i \le n + 1; i++)
    for (int lj = 0, rj = n + 1, j = 0, jj = n + 1; j <= n
+ 1; j++, jj--)
    {
        if (board[i][j] != '#')/// 0(1)
            obstaclesColPos[i][j].first = lj; /// 0(1)
        else
            1j = j; /// \Theta(1)
        if (board[i][jj] != '#')/// 0(1)
            obstaclesColPos[i][jj].second = rj; /// 0(1)
        else
             rj = jj; /// \Theta(1)
        if (board[i][j] == 'o') /// \Theta(1)
             slidersPos.emplace_back(i, j), globalBoard[i]
[j] = '.'; /// \Theta(1)
    }
/// \Theta(n*n)
for (int j = 0; j \le n + 1; j++)
    for (int ui = 0, di = n + 1, i = 0, ii = n + 1; i <= n
+ 1; i++, ii--)
    {
        if (board[i][j] != '#') /// \Theta(1)
             obstaclesRowPos[j][i].first = ui; /// 0(1)
        else
            ui = i; /// \Theta(1)
        if (board[ii][j] != '#')/// 0(1)
             obstaclesRowPos[j][ii].second = di; /// 0(1)
        else
            di = ii; /// \Theta(1)
    }
```

this function runs in O(n*n + n*n*log(n) + min(s*log(s), n*n) + k) \rightarrow which can be simplified to O(n*n*log(n) + min(s*log(s), n*n) + k)

PrintAnswer()

given:

- 1) s = number of sliders in the given grid
- 2) n*n is the size of the grid
- 3) k = optimal path size (number of steps to reach the goal if it exists)
- 4) $d = max depth of tree (where <math>k \le d$)

we have several loops that needs k steps and the body of these loops take constant times then we have a loop that takes k steps and calls MOVE().. however, we calculated the exact for the third block in the BFS() function we can do the same here and get the exact complexity as the 2 trees balances each others

so this is $\Theta(k/2 * (min(slogs, n*n) + s)) = \Theta(k * min(s*log(s), n*n))$