

# Solving Alphametics

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# 1 Problem Definition

We want to find a set of operations to solve an alphametic puzzle. An example of such a puzzle is this:

$$\begin{array}{rcccccc} & & S & E & N & D \\ + & & M & O & R & E \\ \hline = & M & O & N & E & Y \end{array}$$

Where

- $E, N, O, R, Y \in \mathbb{N}_0$  and  $M, S \in \mathbb{N}$
- $E \neq N \neq O \neq R \neq Y \neq M \neq S \neq E$
- 

$$\begin{aligned} Y &= D + E & (\text{mod } 10) \\ E &= N + R + \lfloor \frac{D+E}{10} \rfloor & (\text{mod } 10) \\ N &= E + O + \lfloor \frac{N+R + \lfloor \frac{D+E}{10} \rfloor}{10} \rfloor & (\text{mod } 10) \\ &\vdots \end{aligned}$$

We can introduce a term  $c_n$  for the overflow where  $c \in \{0, 1\}$  in this case.

$$\begin{aligned} Y &= D + E & (\text{mod } 10) \\ E &= N + R + c_1 & (\text{mod } 10) \\ N &= E + O + c_2 & (\text{mod } 10) \\ O &= S + M + c_3 & (\text{mod } 10) \\ M &= c_4 & (\text{mod } 10) \end{aligned}$$

$c_{max}$  can be generalized as it is linked to the number of addends where  $c_{max} = \max\{n_{addends} - 1, 9\}$

## 2 A Simpler Example

Let's experiment by solving a simple example

$$\begin{array}{r} \phantom{+} \phantom{=} \phantom{O} \phantom{U} \phantom{T} \\ \phantom{+} \phantom{=} \phantom{O} \phantom{U} \phantom{T} \\ + \phantom{=} \phantom{O} \phantom{U} \phantom{T} \\ \hline = \phantom{O} \phantom{U} \phantom{T} \end{array}$$

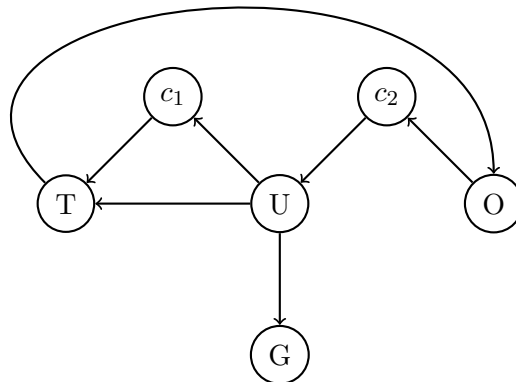
This alphametic results in the following equations

$$T = O + O \pmod{10} \quad (1)$$

$$U = T + G + c_1 \quad (\text{mod } 10) \quad (2)$$

$$O = c_2 \pmod{10} \quad (3)$$

We can construct a graph like structure to see what variables depend on each other:



### 3 Solving the more complicated case

the alphametic from the introduction above gives us the following graph

