

**TOAST**

**Programmer's Guide**

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## Table of Contents

Introduction to TOAST.....	4
Event Condition Action Rules.....	4
Rule Bindings.....	5
Built-in Conditions and Actions.....	5
Extending or Replacing the TOAST Language Built-ins.....	5
Agent Transformation.....	6
ECA Rule Engine.....	7
The TOAST Rule Language.....	9
Rule Events.....	9
Location Specifiers.....	10
Rule Bindings.....	12
Rule Expressions.....	13
Rule Conditions.....	13
Rule Actions.....	14
Built-In Calls.....	15
User-Defined Rule Helpers.....	16
TOAST Rule Language Standard Built-Ins.....	18
CountDowns.....	18
Waiters.....	18
Flags.....	19
Counters.....	20
Debugging and Tracing.....	20
Aborting Execution.....	21
Other Built-ins.....	22
Using TOAST.....	23
Obtaining the sources.....	23
Building TOAST.....	23
Using TOAST.....	23

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## Introduction to TOAST

TOAST is an execution scripting tool designed primarily to support automation of tests for multi-threaded and multi-JVM Java applications (TOAST stands for Thread Orchestration Automation Scripting Tool). It uses bytecode transformation to modify the runtime execution of the application without having to modify the source code. The possible modifications are essentially unlimited but the tool provides explicit support for test automation in three main areas:

- orchestrating the timing of activities performed by independent application threads
- tracing application execution
- subverting normal method execution

TOAST is actually more general than this description appears to suggest. The functionality in the TOAST script language which is specific to test automation and thread management is supplied via a simple POJO plug-in. It is trivial to modify the script language to support language operations which are specific to other application domains. Indeed this can be done on a piecemeal basis for individual components of TOAST scripts. The core engine which underlies TOAST is a general purpose code injection program based around Event Condition Action rules.

### ***Event Condition Action Rules***

TOAST primarily operates by introducing side-effects at specified points during execution according to an orchestration script. A script comprises a sequence of Event Condition Action (ECA) rules language which define how the application behaviour should be transformed at runtime. The three components of these rules, event, condition and action, are used, respectively, to define:

- *where* during application execution a side-effect should occur
- *whether* the side-effect should happen or not
- *what* the side effect should be

For example, in the following example rule the event specifies a trigger point in method `get()` of class `BoundedBuffer`. The precise location of the trigger point is just before method `get()` makes a call to method `Object.wait()` – the example assumes that `get()` suspends its caller by calling this method if the buffer is empty.

```
RULE throw on Nth empty get
CLASS org.my.BoundedBuffer
METHOD get()
AT INVOKE Object.wait()
BIND buffer = $0
IF countDown(buffer)
DO throw org.my.ClosedException(buffer)
ENDRULE
```

The event also establishes a binding for variable `buffer` assigning it with the value `$0` which refers to the recipient (`this`) argument of the `get()` call which triggered the rule. The condition invokes the standard TOAST built-in `countDown(Object)` which decrements a `CountDown` associated with the buffer – the example assumes some other rule has called `addCountDown(buffer, N)` to create this `CountDown` and initialise it with value `N`. The `countDown` built-in returns `true` when the value of the `CountDown`

decrements to zero.

So, in this example the condition will evaluate to `false` the first  $N-1$  times that a getter attempts to wait. At the  $N$ th triggering the condition will evaluate to `true` and the rule will fire, running the built-in action `throw`. This will cause the triggering thread to throw a `ClosedException` from the call to `get()`.

In order to use TOAST to test a Java application the JVM must be started with extra command line arguments. These arguments point the JVM at the TOAST scripting engine and identify a set of TOAST rule files specifying the side-effects to be introduced into the application under test. The engine applies the rules to any code which contains a location matching a rule trigger point. Rules may be supplied in a single script or as a set of scripts, the latter allowing reuse of rules across related tests.

## **Rule Bindings**

Rule conditions and actions may be parameterised by *bindings*, contextual data obtained from the code location targeted by the rule. This allows actions associated with one rule to be correlated with actions performed by other rules. It also means that actions can be specific to the particular case in hand.

For example, a rule might be attached to a database insertion method whose first parameter was the primary key of the record being created. The rule action could refer to this parameter (it would use the parameter reference `$1`) supplying it as the argument to an `addCountDown` call. Another rule attached to a database read method could employ a condition which passed this parameter to built-in `countDown`, decrementing any counter associated with this key. Its action might throw a read exception. With this example a separate countdown would be added for each newly inserted record. An exception would be thrown for each inserted record at the  $N$ th read. However, an exception would not be thrown for reads on existing records since built-in `countDown` returns `false` if no countdown exists with the supplied key.

## **Built-in Conditions and Actions**

TOAST provides a suite of built-in “orchestration” conditions and actions used to coordinate the activities of independent threads e.g. delays, waits and signals, countdowns, flag operations and so on. These are particularly useful for testing multi-threaded programs subject to arbitrary scheduling orders. Judicious insertion of orchestration actions can guarantee that thread interleavings in a given test run occur in a desired order, enabling test code to reliably exercise parallel execution paths which do not normally occur with synthetic workloads.

Tracing actions are provided so that test deployment scripts can track progress of a test run and identify successful or unsuccessful test completion. Trace output can also be used to debug rule execution. The use of local context data and orchestration data in rule conditions allows trace output to be quite finely tuned. Trace actions can insert this data into message strings, allowing detailed scrutiny of test execution paths.

A few special built-in actions can be used to subvert the behaviour of application code by modifying execution paths. This is particularly important in a test environment where it is often necessary to force application methods to generate dummy results or simulate an error. A *return* action forces an early return from the code location targeted by the rule. If the method is non-void then the return action supplies a value to use as the method result. A *throw* action enables runtime exceptions (i.e. instances of `RuntimeException` or its subclasses) to be thrown from any target location, effectively either aborting a thread or, at

least, where a catch-all handler is employed, the portion of the thread's call tree between the rule location and the catch block. Other exceptions may be thrown so long as the method in which the rule is located declares the exception in its throws list. Finally, a *kill* action allows a machine crash to be simulated by configuring an immediate exit from the JVM.

It is worth noting that rules are not just restricted to using built-in operations. Application-specific side-effects can also be introduced by invoking public Java methods in rule events, conditions or actions. The obvious target for such methods is objects supplied from the triggering method using bindings. However, it is also possible to access static data and invoke static methods of any class accessible from the classloader of the triggering method. So, it is quite feasible to use TOAST rules to apply arbitrary modifications to the original program.

## ***Extending or Replacing the TOAST Language Built-ins***

Another option to bear in mind is that the set of built-in operations available to TOAST rules is not fixed. The rule engine works by mapping built-in operations which occur in a given rule to public instance methods of a helper class associated with the rule. By default, this helper class is `org.jboss.jbossts.orchestration.rule.helper.Helper`, which provides the standard set of built-ins designed to simplify management of threads in a multi-threaded application. However, it is possible to specify an alternative helper class for each individual rule.

Any non-abstract class may be specified as the helper. Its public instance methods automatically become available as built-in operations in the rule event, condition and action. For example, by specifying a helper class which extended the default class, `Helper`, a rule would be able to use any of the existing built-ins and/or also make rule-specific (or application-specific) built-in calls. So, although the default TOAST rule language is oriented towards orchestrating the behaviour of independent threads in multi-threaded tests, TOAST can easily be reconfigured to support a much wider range of application requirements.

## ***Agent Transformation***

The bytecode modifications performed by TOAST are implemented using a *Java agent* program. JVM class loaders provide agents with an opportunity to modify loaded bytecode just prior to compilation (see package `java.lang.Instrumentation` for details of how Java agents work). The TOAST agent reads the rule script at JVM bootstrap. It then monitors method code as it is loaded looking for *trigger points*, locations in the method bytecode which match the locations specified in rule events.

The agent inserts *trigger calls* into code at each point which matches a rule event. Trigger calls are calls to the rule execution engine which identify:

- the *trigger method*, i.e. the method which contains the trigger point
- the rule which has been matched
- the arguments to the trigger method

If several rules match the same trigger point then there will be a sequence of trigger calls, one for each matching rule, and rules will be triggered in the order they appear in their script(s).

When a trigger call occurs the rule execution engine locates the relevant rule and then executes it. The rule execution engine establishes bindings for variables mentioned in the

rule event and then *tests* the rule condition. If the condition evaluates to true it *fires* the rule, executing each of the rule actions in sequence.

Trigger calls pass the method recipient (this) and method arguments to the rule engine. These values may be referred to in the condition and action with a standard naming convention, \$0, \$1 etc. The event specification can introduce bindings for additional variables. Bindings for these variables may be initialized using literal data or by invoking methods or operations on the method parameters and/or static data. Variables bound in the event can simply be referred to by name in the condition or action. Bindings allow arbitrary data from the triggering context to be tested in the condition in order to decide whether to fire the rule and to be employed as a target or parameter for rule actions. Note that the agent will eventually be updated to pass local variables which are in scope at the trigger point as arguments to the trigger call, making them available as default bindings.

The agent also compiles exception handler code around the trigger calls in order to deal with exceptions which might arise during rule processing. This is not intended to handle errors detected during operation of the rule execution engine (they should all be caught and dealt with internally). Exceptions are thrown out of the execution engine to alter the flow of control through the triggering method. Normally, after returning from a trigger call the triggering thread continues to execute the original method code. However, a rule can use the return and throw built-in actions to specify that an early return or exception throw should be performed *from the trigger method*. The rule language implementation achieves this by throwing its own private, internal exceptions below the trigger call. The handler code compiled into the trigger method catches these internal exceptions and then either returns to the caller or recursively throws a runtime or application-specific exception. This avoids normal execution of the remaining code in the body of the triggering method. If there are other trigger calls pending at the trigger point then these are also bypassed when a return or throw action is executed.

## **ECA Rule Engine**

The TOAST rule execution engine consists of a rule parser, type checker and interpreter/compiler. The rule parser is invoked by the agent during bootstrap. This provides enough information to enable the agent to identify potential trigger points.

Rules cannot be type checked and compiled until the class and method bytecode they refer to has been loaded. This is because type checking requires identifying properties of the trigger class and, potentially, of classes it mentions using reflection. The type checker needs to identify properties of loaded classes such as the types and accessibility of fields, method signatures etc. So, in order to ensure that the trigger class and all its dependent classes have been loaded before the type checker tries to access them, rules are type checked and compiled the first time they are triggered. This also avoids the cost of checking and compiling rules included in the rule set which do not actually get called.

A single rule may be associated with more than one trigger point. Firstly, depending upon how precisely the rule specifies its event, it may apply to more than one class or more than one method within a class. But secondly, even if a rule specifies a class and method unambiguously the same bytecode file may be loaded by different class loaders. So, the rule has to be type checked and compiled for each applicable trigger point.

If a type check or compile operation fails the rule engine prints an error and disables execution of the trigger call. Note that in cases where the event specification is ambiguous a rule may type check successfully against one trigger point but not against another. Rule execution is only disabled for cases where the type check fails.

In the basic operating mode, trigger calls execute a rule by interpreting the rule parse tree.

Execution is performed with the help of an auxiliary class generated at runtime by the TOAST agent called a *helper adapter*. This class is actually a subclass of the helper class associated with the rule. It inherits from the helper class so that it knows how to execute built-in operations defined by the helper class. A subclass is used to add extra functionality required by the rule system, most notably an execute method to be called at the trigger point.

A helper adapter class includes a field referencing the rule it is derived from. This allows the execute method to locate the triggered rule and, from there, look up the rule's event, condition and action. The basic execute method traverses the parse trees of these three rule components evaluating each expression in turn. Another generated field holds bindings for the method parameters and for variables introduced by the rule event. These bindings can be looked up or assigned during rule execution when they are referred to from within the rule event, condition or action. When the execute method encounters a call to a built-in it can execute this call by invoking one of the methods inherited from its helper superclass.

When a rule is triggered the rule engine creates an instance of the rule's helper adapter class to provide a context for the trigger call. It uses setter methods generated by the TOAST agent to initialise the rule and bindings fields and then it calls the adapter instance's execute method. Since each rule triggering is handled by its own adapter instance this ensures that concurrent triggers of the same rule from different threads do not interfere with each other.

Work is in progress to support compilation of rules to bytecode. When compilation of rules is enabled the TOAST agent generates an execute method which contains inline bytecode derived from the rule event condition and action. Although this takes slightly more time to generate it should provide a performance pay off where the trigger method gets called many times. Ideally, compilation should be selectable per rule or across the board for all rules in a rule set.



## The TOAST Rule Language

Rules are defined in scripts which consists of a sequence of rule definitions interleaved with comment lines. Comments may occur within the body of a rule definition as well as preceding or following a definition but must be on separate lines from the rule text.

Comments are lines which begin with a # character:

```
#####  
# Example Rule Set  
#  
# a single rule definition  
RULE example rule  
# comment line in rule body  
.  
.  
.  
ENDRULE
```

### Rule Events

Rule event specifications identify a specific *location* in a *target* method associated with a *target* class. Target methods can be either static or instance methods or constructors. If no detailed location is specified the default location is entry to the target method. So, the basic schema for a single rule is as follows:

```
# rule skeleton  
RULE <rule name>  
CLASS <class name>  
METHOD <method name>  
BIND <bindings>  
IF <condition>  
DO <actions>  
ENDRULE
```

The name of the rule following the RULE keyword can be any free form text with the restriction that it must include at least one non-white space character. Rule names do not have to be unique but it obviously helps when debugging rule scripts if they clearly identify the rule. The rule name is printed whenever an error is encountered during parsing, type checking, compilation or execution.

The class and method names following the CLASS and METHOD keywords must be on the same line. The class name can identify a class either with or without the package qualification. The method name can identify a method with or without an argument list or return type. A constructor method is identified using the special name <init>. For example,

```
# class and method example  
RULE any commit on any coordinator engine  
CLASS CoordinatorEngine  
METHOD commit  
.  
.  
.  
ENDRULE
```

matches the rule with any class whose name is CoordinatorEngine, irrespective of the package it belongs to. When any class with this name is loaded then the agent will insert a trigger point at the beginning of any method named commit. If there are several occurrences of this method, with different signatures then each method will have a trigger

point inserted.

More precise matches can be guaranteed by adding more detail. For example,

```
# class and method example 2
RULE commit with no arguments on wst11 coordinator engine
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD State commit()
AT LINE 324
. . .
ENDRULE
```

This rule will only match the `CoordinatorEngine` class in package `com.arjuna.wst11.messaging.engines` and only match a method `commit` with no arguments and with a return type whose name is `State`. Note that the package for class `State` is left unspecified. The type checker infers this information from the matched method.

The previous example also employs the location specifier `AT LINE`. The text following the line keyword must be able to be parsed to derive an integer line number. This directs the agent to insert the trigger call at the start of a particular line in the source code.

Note:

The TOAST agent *will not* transform any classes in package `java.lang` nor classes in package `org.jboss.jbossts.orchestration` (the orchestration package itself).

## ***Location Specifiers***

It is easy and convenient to use line numbers to specify locations in code which is not subject to change. However, this is less useful for automated testing because modifications to the code base can shift the line numbers of unmodified lines invalidating test scripts. Luckily there are several other ways of specifying where a trigger point should be inserted into a target method which relate to the structure of the code. For example,

```
# location specifier example
RULE countdown at commit
CLASS CoordinatorEngine
METHOD commit
AT READ state
. . .
ENDRULE
```

In this rule the trigger point will be inserted just before the first location in the bytecode where a `getField` operation is performed on field called `state`. This is effectively the same as saying that the trigger point will occur at the first point in the source code of the method where field `state` is accessed. By contracts, the following rule would locate the trigger point after the first write to field `recovered`:

```
# location specifier example 2
RULE add countdown at recreate
CLASS CoordinatorEngine
METHOD <init>
AFTER WRITE CoordinatorEngine.recovered
. . .
ENDRULE
```

Note that in the last example the field type is qualified to ensure that the write is to the field

belonging to an instance of class `CoordinatorEngine`.

The full set of location specifiers is as follows

```
AT ENTRY
AT EXIT
AT LINE number
AT READ [type .] field [count]
AFTER READ [ type .] field [count]
AT WRITE [ type .] field [count]
AFTER WRITE [ type .] field [count]
AT INVOKE [ type .] method [ ( argtypes ) ] [count]
AFTER INVOKE [ type .] method [ ( argtypes ) ][count]
AT SYNCHRONIZE [count]
AFTER SYNCHRONIZE [count]
AT THROW [typename] [count]
```

If a location specifier is provided it must immediately follow the `METHOD` specifier. If no location specifier is provided it defaults to `AT ENTRY`.

An `AT ENTRY` specifier normally locates the trigger point before the first executable instruction in the trigger method. An exception to this occurs in the case of a constructor method in which case the trigger point is located before the first instruction following the call to the super constructor or redirection call to an alternative constructor. This is necessary to ensure that rules do not attempt to bind and operate on the instance before it is constructed

An `AT EXIT` specifier locates a trigger point at each location in the trigger method where a normal return of control occurs (i.e. wherever there is an implicit or explicit return but not where a throw exits the method).

An `AT LINE` specifier locates the trigger point before the first executable bytecode instruction in the trigger method whose source line number is greater than or equal to the line number supplied as argument to the specifier. If there is no executable code at (or following) the specified line number the agent will not insert a trigger point (note that it does not currently print an error in such cases because this may merely indicate that the rule does not apply to this particular class or method – perhaps this behaviour needs revising?).

An `AT READ` specifier locates the trigger point before the first mention of an object field whose name matches the supplied field name i.e. it corresponds to the first occurred of a corresponding `getField` instruction in the bytecode. If a type is specified then the `getField` instruction will only be matched if the named field is declared by a class whose name matches the supplied type. If a count *N* is supplied then the *N*th matching `getField` will be used as the trigger point. Note that the count identifies to the *N*th textual occurrence of the field access, not the *N*th field access in a particular execution path at runtime.

An `AFTER READ` specification is identical to an `AT READ` specification except that it locates the trigger point after the `getField` bytecode.

`AT WRITE` and `AFTER WRITE` specifiers are the same as the corresponding `READ` specifiers except that they correspond to assignments to the named field in the source code i.e. they identify `putField` instructions.

`AT INVOKE` and `AFTER INVOKE` specifiers are like `READ` and `WRITE` specifiers except that they identify invocations of methods or constructors within the trigger method as the

trigger point. The method may be identified using a bare method name or the name may be qualified by a, possibly package-qualified, type or by a descriptor. A descriptor consists of a comma-separated list of type names within brackets. The type names identify the types of the method parameters and may be prefixed with package qualifiers and employ array bracket pairs as suffixes.

AT SYNCHRONIZE and AFTER SYNCHRONIZE specifiers identify synchronization blocks in the target method, i.e. they correspond to MONITORENTER instructions in the bytecode. Note that AFTER SYNCHRONIZE identifies the point immediately after entry to the synchronized block rather than the point immediately after exit from the block.

An AT THROW specifier identifies a throw operation within the trigger method as the trigger point. The throw operation may be qualified by a, possibly package-qualified, typename identifying the lexical type of the thrown exception. If a count *N* is supplied then the location specifies the *N*th textual occurrence of a throw. n.b. the exception typename will be ignored in the current release.

n.b. for hysterical reasons CALL may be used as a synonym for INVOKE, RETURN may be used as a synonym for EXIT and the AT in an AT LINE specifier is optional.

## ***Rule Bindings***

The event specification includes a binding specification which computes values for variables which can subsequently be referenced in the rule body. These values will be computed each time the rule is triggered before testing the rule condition. For example,

```
# binding example
RULE countdown at commit
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
AT READ state
BIND engine:CoordinatorEngine = $0,
    recovered:boolean = engine.isRecovered(),
    identifier:String = engine.getId()
. . .
ENDRULE
```

creates a variable called `engine`. This variable is bound to the recipient of the `commit` method call which triggered the rule, identified by the parameter reference `$0` (if `commit` was a static method then reference to `$0` would result in a type check exception). Arguments to the trigger method can be identified using parameter references with successive indices, `$1`, `$2` etc. The declaration of `engine` specifies its type as being `CoordinatorEngine` though this is not strictly necessary since it can be inferred from the type of `$0`.

Similarly, variables `recovered` and `identifier` are bound by evaluating the expressions on the right of the `=` operator. Note that the binding for `engine` has been established before these variables are bound so it can be referenced in the evaluated expression. Once again, type specifications are provided but they could be inferred.

The special syntax `BIND NOTHING` is available for cases where the rule does not need to employ any bindings.

## ***Rule Expressions***

Expressions which occur on the right hand side of the = operator in event bindings can be simple expressions i.e.

- references to previously bound variables
- static data references
- primitive literals
- field accesses
- method invocations
- built-in operation invocations

n.b. built-in operations are explained in more detail below.

Expressions can also be complex expressions composed from other expressions using the usual Java operators: +, -, \*, /, %, &, |, ^, &&, ||, !, ==, !=, <, <=, >, >=, etc. The ternary conditional expression operator, ? :, can also be employed. The type checker does its best to identify the types of simple and complex expressions wherever possible. So, for example, if it knows the type of bound variable `engine` then it will be able to employ reflection to infer the type of a field access `engine.recovered`, a method invocation `engine.isRecovered()`, etc.

Note:

- The assignment operator is not available for use in expressions. It can only be employed at the top level in binding specifications.
- Use of the new operator is not currently allowed in expressions.
- throw and return operations are only allowed as the last action in a sequence of rule actions (see below).
- Expressions should obey the normal rules regarding associativity and precedence.
- It should eventually be possible to allow expressions to make references to local variables which are in scope at the trigger point as well as to method arguments. So, for example, if the synchronization block of method `commit` was preceded by a declaration for an `int` variable with name `idx` then the example rule should be able to include references to this variable in expressions by employing a local variable reference of the form `$idx`. This reference would evaluate to the value of the local variable when the trigger point was reached.

## ***Rule Conditions***

Rule conditions are nothing more than rule expressions with boolean type. For example,

```
# condition example
RULE countdown at commit
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
AT READ state
BIND engine:CoordinatorEngine = $0,
    recovered:boolean = engine.isRecovered(),
    identifier:String = engine.getId()
IF    recovered
. . .
ENDRULE
```

merely tests the value of bound variable recovered. The same effect could be achieved by using the following condition

```
# condition example 2
RULE countdown at commit
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
AT READ state
BIND engine:CoordinatorEngine = $0,
. . .
IF    engine.isRecovered()
. . .
ENDRULE
```

Alternatively, if, say, the instance employed a public field, recovered, to store the boolean value returned by method isRecovered then the same effect would be achieved by the following condition.

```
# condition example 3
RULE countdown at commit
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
AT READ state
BIND engine:CoordinatorEngine = $0,
. . .
IF    engine.recovered
. . .
ENDRULE
```

Note that the boolean literal true is available for use in expressions so a rule which should always fire can use this as the condition expression.

## **Rule Actions**

Rule actions are either a rule expression or a return or throw expression or a comma-separated sequence of rule expressions, possibly ending with a return or throw expression. Rule expressions occurring in an action list may have arbitrary type, including void type.

A return expressions is the return keyword possibly followed by a rule expression which is used to compute a return value. A return expression causes a return from the triggering method so it may omit a return value if and only if the method is void. If a return value is employed then the type checker will ensure that it's type is assignable to the return type of the trigger method. So, for example, the following use of return is legitimate assuming

method `commit` has return type `boolean`:

```
# return example
RULE countdown at commit
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
AT READ state

    . . .
DO debug("returning early with failure"),
    return false
ENDRULE
```

A `throw` expression is the `throw` keyword followed by an exception constructor expression. An exception constructor expression is the class name of the exception which is to be thrown followed by an argument list. The argument list may be empty i.e. it may consist of an open and close bracket pair. Alternatively, the brackets may include a single rule expression or a sequence of rule expressions separated by commas. If no arguments are supplied the exception type must implement an empty constructor. If arguments are supplied then the exception type must implement a constructor whose signature is type-compatible.

A `throw` expression causes an exception of the type named in the exception constructor to be created and thrown from the triggering method. In order for this to be valid the expression type must either be assignable to `java.lang.RuntimeException` or be explicitly declared in the triggering method's `throws` list. The type checker will throw a type exception if either of these conditions is not met. So, for example, the following use of `throw` is legitimate assuming method `commit` includes `WrongStateException` in its `throws` list.

```
# throw example
RULE countdown at commit
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
AT READ state

    . . .
DO debug("throwing wrong state"),
    throw WrongStateException()
ENDRULE
```

An empty action list may be specified using the keyword `NOTHING`.

## ***Built-In Calls***

Built-in calls are written without a recipient as though they were invocations of a method on `this`. The rule engine identifies calls in this format and translates them to runtime invocations of helper class instance methods. So, referring back to the last few examples, it is apparent that the helper class implements a debugging method with signature

```
boolean debug(String message)
```

This method prints the supplied string to `System.out` and always returns `true`. It can be used in a rule action to display a trace message, for example:

```
DO    debug("killing JVM"), killJVM()
```

When the `debug` built-in is executed the rule engine calls the corresponding method of the current helper instance passing it the string `"killing JVM"`. Method `killJVM` is

another built-in implemented by an instance method of the default helper class `Helper`.

Note that method `debug` has a boolean signature so that tracing can also be performed in rule conditions. This would normally occur in combination with a test of some bound variable or method parameter, for example:

```
IF    debug("checking for recovered participant")
AND
    participant.isRecovered()
AND
    debug("recovered participant " + participant.getId())
```

n.b. `AND` is an alternative token for the Java `&&` operator.

The rule language implementation automatically exposes all public instance methods of class `Helper` as built-in operations. So when the rule type checker encounters an invocation of `debug` with no recipient supplied it identifies that `debug` is a method of class `Helper` and automatically type checks the call against this method. At execution time the call is executed by invoking the implementation of `debug` on the helper instance created under the rule trigger call.

This feature allows additional built-ins to be added to the rule engine simply by adding new helper implementations. No changes are required to the parser, type checker and compiler in order for this to work.

## ***User-Defined Rule Helpers***

A rule can specify its own helper class if it wants to extend, override or replace the set of built-in calls available for use in its event, condition or action. For example, in the following rule, class `FailureInjector` is used as the helper class. Its boolean instance method `injectWrongState(CoordinatorEngine)` is called from the condition to decide whether or not to throw a `WrongStateException`.

```
# helper example
RULE help yourself
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
HELPER com.arjuna.wst11.messaging.engines.FailureInjector
AT EXIT
BIND NOTHING
IF injectWrongState($0)
DO throw WrongStateException()
ENDRULE
```

A helper class does not need to implement any special interface or inherit from any pre-defined class. It merely needs to provide instance methods to resolve the built-in calls which occur in the rule. By sub-classing the default helper it is possible to extend or override the default set of methods. For example, the following rule employs a helper which adds emphasis to the debug messages printed by the rule.



```

# helper example 2
RULE help yourself but rely on others
CLASS com.arjuna.wst11.messaging.engines.CoordinatorEngine
METHOD commit
HELPER HelperSub
AT ENTRY
BIND NOTHING
IF NOT flagged($0)
DO debug("throwing wrong state"),
    flag($0)
    throw WrongStateException()
ENDRULE

```

```

class HelperSub extends Helper
{
    public boolean debug(String message)
    {
        super("!!! IMPORTANT EVENT !!! " + message);
    }
}

```

The rule is still able to employ the built-in methods `flag` and `flagged` defined by the default helper class.

## TOAST Rule Language Standard Built-Ins

The default helper class provides the following standard suite of built-in calls for use in rule expressions. These are primarily intended for use in condition and action expressions but they may also be called in event bindings. They provide features which are designed to make it easy to perform complex tests, in particular to coordinate the actions of threads in multi-threaded applications.

### **CountDowns**

The rule engine provides CountDowns which can be used to ensure that firing of some given rule will only occur after other rules have been triggered or fired a certain number of times. The API defined by the helper class is

```
public boolean addCountDown(Object identifier, int count)
public boolean getCountDown(Object identifier)
public boolean countDown(Object identifier)
```

CountDowns are identified by an arbitrary object, allowing successive calls to the countdown API to apply to the same or different cases. This identification can be made across different rule and helper instances. For example, one rule might include action `addCountDown($0, 1)` and another rule might include condition `countDown($0)`. A CountDown created by the first rule would only be decremented if the second rule was triggered from a method call with the same value for `this`. CountDowns created by invocations with distinct values for `this` would match up accordingly. However, if the CountDown was identified using a common String literal (i.e. action and condition were `addCountDown("counter", 1)` and `countDown("counter")`, respectively), then the CountDown created by the first rule would be decremented by the next firing of the second rule irrespective of whether the trigger method calls were on related instances.

`addCountDown` is used to create a CountDown. `count` specifies how many times the CountDown will be decremented before a decrement operation fails i.e. if `count` is 1 then the CountDown will decrement once and then fail at the next decrement. If `count` is supplied with a value less than 1 it will be replaced with value 1. `addCountDown` would normally be employed in a rule action. However, it is defined to return `true` if a new CountDown is created and `false` if there is already a CountDown associated with the identifier. This allows it to be used in rule conditions where several rules may be racing to create a CountDown.

`getCountDown` is for use in a rule condition to test whether a CountDown associated with a given identifier is present, returning `true` if so otherwise `false`.

`countDown` is for use in a rule condition to decrement a CountDown. It returns `false` if the decrement succeeds or if there is no CountDown associated with identifier. It returns `true` if the CountDown fails i.e. it has count 0. In the latter case the association between the identifier and the CountDown is removed, allowing a new CountDown to be started using the same identifier. Note that this behaviour ensures that a race between multiple threads to decrement a counter from one or more rule conditions can only have one winner.

### **Waiters**

The rule engine provides Waiters used to suspend threads during rule execution and then have other threads wake them up. The wakeup can simply allow the suspended thread to

resume execution of the rule which suspended it. Alternatively, it can force the waiting thread to exit from the triggering method with an exception. The API defined by the helper class is

```
public void waitFor(Object identifier)
public void waitFor(Object identifier, long millisecsWait)
public boolean waiting(Object identifier)
public boolean signalWake(Object identifier)
public boolean signalKill(Object identifier)
```

As with CountDowns, Waiters are identified by an arbitrary object. Note that the wait operation is not performed by invoking `Object.wait` on `identifier`. Doing so might interfere with locking and synchronization operations performed by the triggering method or its callers. The identifier is merely used by the rule engine to associate wait and signal operations. The Helper class employs it's own private Waiter object to manage the synchronization activity.

`waitFor` is intended for use in a rule action. It suspends the current thread on the Waiter associated with the identifier until either a `signalWake` or a `signalKill` is called with the same identifier. In the former case the thread will continue processing any subsequent actions and then return from the trigger call. In the latter case the thread will throw a runtime exception from the triggering method call frame. The version without a wait parameter will never time out. The version which employs a wait parameter will time out after the specified number of milliseconds.

`waiting` is intended for use in rule conditions. it will return true if any threads are waiting on the relevant Waiter for a signal. It returns false if there are no threads waiting.

`signalWake` is intended for use in rule conditions or actions. If there are threads waiting on the Waiter associated with `identifier` it wakes them and returns true. If not it returns false. Note this behaviour ensures that a race between multiple threads to signal waiting threads from a rule condition can only have one winner.

`signalKill` is identical to `signalWake` except that it does not just wake any waiting threads. It also causes them to throw a runtime exception of type `ExecuteException` from their triggering method call frame when they wake up.

## **Flags**

The rule engine provides a simple mechanism for setting, testing and clearing global flags. The API defined by the helper class is

```
public boolean flag(Object identifier)
public boolean flagged(Object identifier)
public boolean clear(Object identifier)
```

As before, Flags are identified by an arbitrary object. All three methods are designed to be used either in conditions or actions.

`flag` can be called to ensure that the Flag identified by `identifier` is set. It returns true if the Flag was previously clear otherwise false. Note that the API is designed to ensure that race conditions between multiple threads trying to set a Flag from rule conditions can only have one winner.

`flagged` tests whether the Flag identified by `identifier` is set. It returns true if the Flag is set otherwise false.

`clear` can be called to ensure that the Flag identified by `identifier` is clear. It returns

true if the Flag was previously set otherwise false. Note that the API is designed to ensure that race conditions between multiple threads trying to clear a Flag from rule conditions can only have one winner.

## **Counters**

The rule engine provides Counters which maintain global counts across independent rule triggerings. They can be created and initialised, read, incremented and decremented in order track and respond to the number of times various triggerings or firings have happened. Note that unlike CountDowns there are no special semantics associated with decrementing a Counter to zero. They may even have negative values. The API defined by the helper class is

```
public boolean createCounter(Object o)
public boolean createCounter(Object o, int count)
public boolean deleteCounter(Object o)
public boolean readCounter(Object o)
public boolean incrementCounter(Object o)
public boolean decrementCounter(Object o)
```

As before, Counters are identified by an arbitrary object. All methods are designed to be used in rule conditions or actions.

createCounter can be called to create a new Counter associated with identifier. If argument count is not supplied then the value of the new Counter defaults to 0.

createCounter returns true if a new Counter was created and false if a Counter associated with identifier already exists. Note that the API is designed to ensure that race conditions between multiple threads trying to create a Counter from rule conditions can only have one winner.

deleteCounter can be called to delete any existing Counter associated with identifier. It returns true if the Counter was deleted and false if no Counter was associated with identifier. Note that the API is designed to ensure that race conditions between multiple threads trying to delete a Counter from rule conditions can only have one winner.

incrementCounter can be called to increment the Counter associated with identifier. If no such Counter exists it will create one with value 0 before incrementing it. It returns the new value of the Counter.

decrementCounter can be called to decrement the Counter associated with identifier. If no such Counter exists it will create one with value 0 before decrementing it. It returns the new value of the Counter.

readCounter can be called to read the value of the Counter associated with identifier. If no such Counter exists it will create one with value 0.

## **Debugging and Tracing**

The rule engine provides a simple built-in debug method to support conditional display of messages during rule execution. The API defined by the helper class is

```
public boolean debug(String message)
```

debug prints the supplied message to System.out, prefixed with the name of the rule being executed. It always returns true, allowing debug messages to be used in conditions

by ANDing them with other boolean expressions.

Generation of debug messages can be switched on by setting the following system property on the JVM command line:

```
org.jboss.jbossts.orchestration.debug
```

The rule engine provides a set of built-in methods to support logging of trace messages during execution. Messages may be logged to `System.out`, `System.err` or to a named file. The API defined by the helper class is

```
public boolean traceOpen(Object identifier, String filename)
public boolean traceOpen(Object identifier)
public boolean traceClose(Object identifier)
public boolean trace(Object identifier, String message)
public boolean traceLn(Object identifier, String message)
```

`traceOpen` opens the file identified by `fileName` and associates it with `identifier`, returning `true`. `filename` can be either a relative or absolute path. Relative file names are located relative to the current working directory of the JVM. If there is already a file associated with `identifier` then `traceOpen` immediately returns `false`. If a file with the given name already exists it is opened in append mode. If `filename` is omitted then a unique name is generated for the file which is guaranteed not to match any existing trace file in the current working directory.

`traceClose` closes the file associated with `identifier` and removes the association, returning `true`. If no open file is associated with `identifier` it returns `false`.

`trace` prints `message` to file associated with `identifier`, returning `true`. If no open file is associated with `identifier` then a file will be opened and associated with `identifier` as if a call to `trace` had been made with no file name supplied.

`traceLn` prints `message` to file associated with `identifier` and appends a newline to the file, returning `true`. If no open file is associated with `identifier` then a file will be opened and associated with `identifier` as if a call to `trace` had been made with no file name supplied.

A caveat applies to the above descriptions for three special cases. If `identifier` is `null` or the string `"out"`, then `trace` and `traceLn` write to `System.out`. If `identifier` is the string `"err"`, then `trace` and `traceLn` write to `System.err`. `traceOpen` and `traceClose` always return `false` immediately if `identifier` has any of these values.

## ***Aborting Execution***

The rule engine provides two built-ins for use in rule actions which allow execution of the triggering method to be aborted. The API defined by the helper class is

```
public void killThread()
public void killJVM()
public void killJVM(int exitCode)
```

`killThread` causes a runtime exception of type `ExecuteException` to be thrown from the triggering method call frame. This will effectively kill the thread unless a catch-all exception handler is installed somewhere up the call stack.

`killJVM` results in a call to `java.lang.Runtime.getRuntime().halt()`. This effectively kills the JVM without any opportunity for any registered exit handlers to run, simulating a JVM crash. If `exitCode` is not supplied it is defaulted to `-1`.

## ***Other Built-ins***

It is intended to implement a rendezvous API along the following lines:

```
public boolean rendezvous(Object identifier, int count)
public int rendezvousCount(Object identifier)
public int rendezvousLimit(Object identifier)
```

`rendezvous` will cause the current thread to suspend until `count` other threads have arrived at the rendezvous identified by `identifier`, at which point it will return `true`. If `count` is less than 1 or if other threads are waiting for the rendezvous under a different count then it will return `false`. Note that it is legitimate (although pathological) to supply a count of 1. In this latter case the rendezvous call will return `true` without waiting.

`getRendezvous` will return the number of threads waiting at the rendezvous identified by `identifier` or 0 if no threads are currently waiting.

`rendezvousLimit` will return the count associated with the rendezvous identified by `identifier` or 0 if no threads are currently waiting.

# Using TOAST

## ***Obtaining the sources***

TOAST sources are available from the JBossTS SVN repository located under directory workspace/adinn/orchestration (they will have to move to trunk at some point). The source tree includes an ext directory containing the external javacup, JFlex and ObjectWeb asm jars needed by the agent and rule code.

The sources include some sample rule scripts located in directory dd/scripts.

## ***Building TOAST***

TOAST builds to produce a single jar in build/lib/orchestration.jar. This jar contains the Java agent and rule engine code. This jar currently also bundles in the contents of the external Jflex, javacup runtime and ObjectWeb asm package libraries for ease of use. These could be unbundled so long as they are available in the classpath of the JVM being used to run the test application and the rule code.

The top level directory contains a file build.xml with default target 'jar' which builds the orchestration jar. Other useful targets include 'parser' which rebuilds the javacup/JFlex parser from the grammar rules in dd/grammar and 'TestScript', which compiles and runs program TestScript.java. The latter is an offline type checker which parses and then type checks the rules defined in file handler.txt in the top level directory. Note that the latter target requires any classes mentioned in the rules to be in the class path of the TestScript program so it may be necessary to edit the rules in build.xml to get this to work.

## ***Using TOAST***

Using TOAST is refreshingly simple in that it only requires pointing the JVM at the agent code in the jar and at the script files containing the orchestration rules. This is specified using the java command flag

```
-javaagent:agentlib(=script:agentscript)+
```

This is a standard option for JDK 1.5 and upwards.

*agentlib* is a path to the orchestration jar. The build process inserts a metadata file in the jar which allows the JVM to identify the agent program entry point so eberthing else si shrink-wrapped.

*agentscript* is a path to a rule script to be used during the JVM run. Multiple scripts may be provided (the brackets and + sign are regexp syntax and do not actually appear in the command flag). The name of each script file is separated from the agentlib file or preceding script file names by an =script:separator string. For example, setting

```
export JAVA_OPTS="-javaagent:${HOME}/jboss/workspace/adinn/
orchestration/build/lib/orchestration.jar=script:${HOME}/j
boss/workspace/adinn/orchestration/dd/scripts/HeuristicSave
AndRecover.txt"
```

will cause the JVM (and indeed JBoss AS) to pick up the jar file from the build directory and use the script provided in the dd directory.

The transformations performed by the agent can be observed by setting several environment variables which cause the transformed bytecode to be dumped to disk.

If system property

```
org.jboss.jbossts.orchestration.dump.generated.classes
```

is set (with *any* value) the agent transformer code will dump a class file containing the bytecode of any class it has modified. The class file is dumped in a directory hierarchy derived from the package of the transformed class. So, for example, class `com.arjuna.Foo` will be dumped to file `com/arjuna/Foo.class`.

If system property

```
org.jboss.jbossts.orchestration.dump.generated.classes.directory
```

is set to the name of a directory writeable by the JVM then class files will be dumped in a package directory hierarchy below this directory. For example, if this property is set with value `/tmp/dump` then class `com.arjuna.Foo` will be dumped to file `/tmp/dump/com/arjuna/Foo.class`. If this property is unset or does not identify a writeable directory then class files will be dumped below the current working directory of the JVM.

If system property

```
org.jboss.jbossts.orchestration.verbose
```

is set then the rule execution engine will display a variety of trace messages as it parses, typechecks, compiles and executes rules.

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