

编译原理第五次理论作业

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Exercise 5.1

Given the following grammar

$$S \rightarrow (L) \mid a$$

$$L \rightarrow L, S \mid S$$

- Construct an LL(1) parsing table for the grammar
 - Note: you must eliminate the left recursion first.
- Draw the detailed process of the parsing of the sentence **(a, (a, a))**, follow the style in the previous slides.

- 消除左递归如下:

$$\begin{aligned} S &\rightarrow (L) \mid a \\ L &\rightarrow SL' \\ L' &\rightarrow , SL' \mid \epsilon \end{aligned}$$

然后可求FIRST和FOLLOW: (因为文法中包含逗号, 故将文法中的逗号用单引号括起来以示区分)

$$\begin{aligned} FIRST(S) &= \{ (, a \} & FOLLOW(S) &= \{ \$,) \} \\ FIRST(L) &= \{ (, a \} & FOLLOW(L) &= \{ \} \} \\ FIRST(L') &= \{ ', ', \epsilon \} & FOLLOW(L') &= \{ \} \} \end{aligned}$$

故可以得出下表：

	()	a	,	\$
S	$S \rightarrow (L)$		$S \rightarrow a$		
L	$L \rightarrow SL'$		$L \rightarrow SL'$		
L'		$L' \rightarrow \epsilon$		$L' \rightarrow, SL'$	

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Matched	Stack	Input	Action
	$S\$$	$(a, (a, a))\$$	
	$(L)\$$	$(a, (a, a))\$$	output $S \rightarrow (L)$
$($	$L)\$$	$a, (a, a))\$$	
$($	$SL')\$$	$a, (a, a))\$$	output $L \rightarrow SL'$
$($	$aL')\$$	$a, (a, a))\$$	output $S \rightarrow a$
$(a$	$L')\$$	$, (a, a))\$$	
$(a$	$, SL')\$$	$, (a, a))\$$	output $L' \rightarrow, SL'$
$(a,$	$SL')\$$	$(a, a))\$$	
$(a,$	$(L)L')\$$	$(a, a))\$$	output $S \rightarrow (L)$
$(a, ($	$L)L')\$$	$a, a))\$$	
$(a, ($	$SL')L')\$$	$a, a))\$$	output $L \rightarrow SL'$
$(a, ($	$aL')L')\$$	$a, a))\$$	output $S \rightarrow a$
$(a, (a$	$L')L')\$$	$, a))\$$	
$(a, (a$	$, SL')L')\$$	$, a))\$$	output $L' \rightarrow, SL'$
$(a, (a,$	$SL')L')\$$	$a))\$$	
$(a, (a,$	$aL')L')\$$	$a))\$$	output $S \rightarrow a$
$(a, (a, a$	$L')L')\$$	$)\$$	
$(a, (a, a$	$)L')\$$	$)\$$	output $L' \rightarrow \epsilon$
$(a, (a, a)$	$L')\$$	$)\$$	
$(a, (a, a)$	$)\$$	$)\$$	output $L' \rightarrow \epsilon$
$(a, (a, a))$	$\$$	$\$$	

Exercise 5.2

Given the following grammar

$$A \rightarrow B \mid BC$$

$$B \rightarrow aB \mid \epsilon$$

$$C \rightarrow ab$$

- Left factor the grammar.
- After left factoring, is the grammar an LL(1) grammar? or is it an LL(k) grammar? and why?
 - Note: you may try the input string **ab**.

- 提取左公因子：

$$A \rightarrow BA'$$

$$B \rightarrow aB \mid \epsilon$$

$$A' \rightarrow C \mid \epsilon$$

$$C \rightarrow ab$$

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首先，这个语法不是LL(1),因为

$$A \rightarrow BC \rightarrow \epsilon C \rightarrow C \rightarrow ab$$

$$A \rightarrow B \rightarrow aB$$

第一个都是a，故不是LL(1)文法

此外，这个语法是LL(k)的。