

## 《编译原理》期末试题试卷(A)参考答案

(考试形式：闭卷 考试时间：2 小时)



《中山大学授予学士学位工作细则》第六条

考试作弊不授予学士学位

方向：\_\_\_\_\_ 姓名：\_\_\_\_\_ 学号：\_\_\_\_\_

注意：答案一定要写在答卷中，写在本试题卷中不给分。本试卷要和答卷一起交回。

1. (8 points) Give a regular expression for each of the following languages over the alphabet  $\{a, b\}$ :

(1) (4 points) All nonempty strings that start and end with the same symbol.

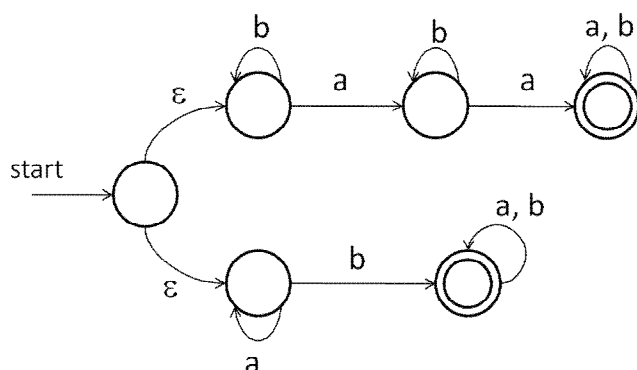
(2) (4 points) All strings that contain no repeated b's (including the empty string).

【参考答案】

(1)  $a(a|b)^*a \mid b(a|b)^*b \mid a \mid b$

(2)  $a^*(ba+)^*b?$

2. (12 points) Consider the following NFA:



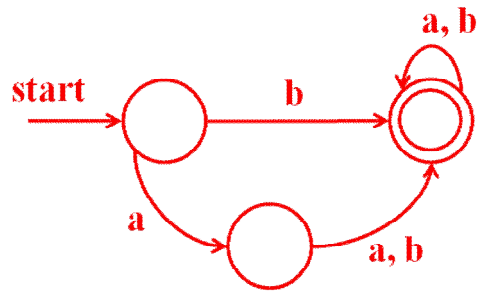
(1) (4 points) What language does the NFA accept? Please describe it in natural language.

(2) (8 points) Convert the NFA to an equivalent DFA. You may construct the DFA directly.

【参考答案】

(1) 所有至少包含 2 个 a 或 1 个 b 的由 a, b 组成的字符串.

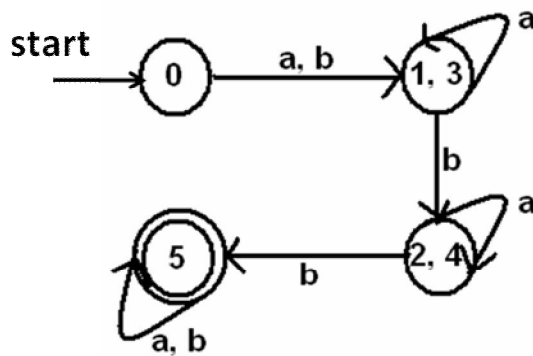
(2)



3. (7 points) Minimize the DFA represented by the following transition table, where state 0 is the start state and state 5 is the only accepting state.

	a	b
0	1	3
1	1	2
2	2	5
3	3	4
4	4	5
5	5	5

【参考答案】



4. (10 points) Give a context-free grammar (CFG) for each of the following languages over the alphabet  $\{a, b\}$ :

- (1) (5 points)  $L = \{a^i b^j \mid i \geq 0 \text{ and } 2 \cdot i \leq j \leq 3 \cdot i\}$ .
- (2) (5 points)  $L = \{w \mid w \text{ contains an odd number of symbols, and the symbol in the middle of } w \text{ is } a\}$ .

【参考答案】

- (1)  $S \rightarrow aSbb \mid aSbbb \mid \epsilon$
- (2)  $S \rightarrow aSa \mid aSb \mid bSa \mid bSb \mid a$

5. (8 points) Consider the following grammar over the alphabet  $\{a, b, c\}$ :

$S \rightarrow Xa$   
 $X \rightarrow bX$   
 $X \rightarrow Y$

$Y \rightarrow Zc$

$Z \rightarrow bZ$

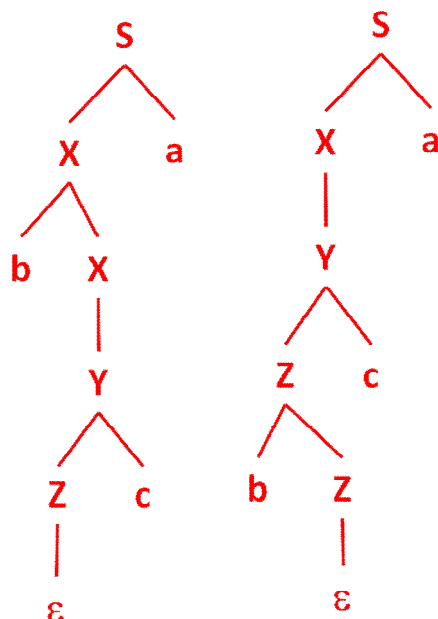
$Z \rightarrow \varepsilon$

(1) (5 points) Demonstrate that this grammar is ambiguous.

(2) (3 points) Please remove exactly one production from this grammar to obtain an unambiguous grammar generating the same language.

【参考答案】

(1)



(2) 删除  $X \rightarrow bX$  或  $Z \rightarrow bZ$  均可.

6. (12 points) Compute FIRST and FOLLOW for each nonterminal in the following grammar:

$S \rightarrow A$

$A \rightarrow BA'$

$A' \rightarrow iBA' \mid \varepsilon$

$B \rightarrow CB'$

$B' \rightarrow +CB' \mid \varepsilon$

$C \rightarrow )A^* \mid ($

【参考答案】

$\text{FIRST}(S) = \text{FIRST}(A) = \text{FIRST}(B) = \text{FIRST}(C) = \{ (, ) \}$

$\text{FIRST}(A') = \{ i, \varepsilon \}$

$\text{FIRST}(B') = \{ +, \varepsilon \}$

$\text{FOLLOW}(S) = \{ \$ \}$

$\text{FOLLOW}(A) = \text{FOLLOW}(A') = \{ \$, * \}$

$\text{FOLLOW}(B) = \text{FOLLOW}(B') = \{ \$, *, i \}$

$\text{FOLLOW}(C) = \{ \$, *, +, i \}$

7. (13 points) Consider the following grammar:

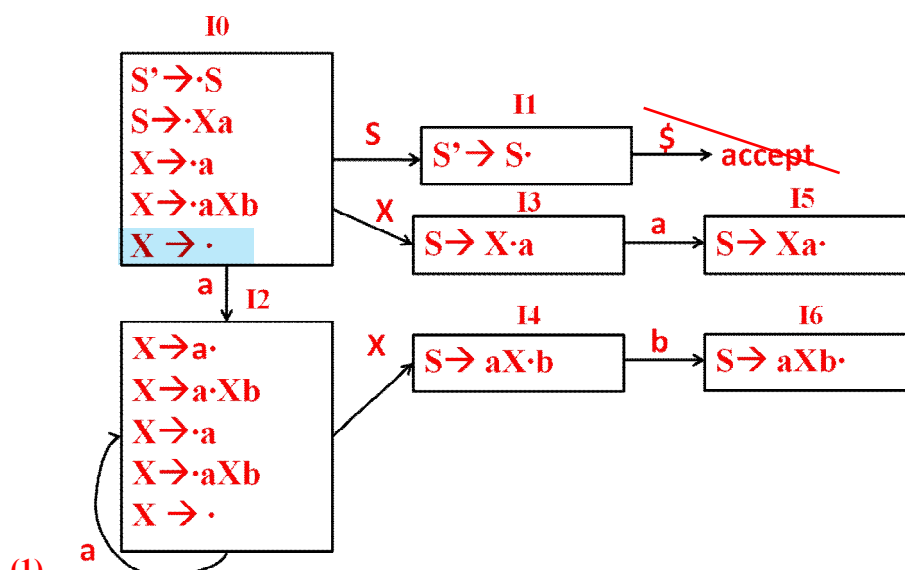
$S \rightarrow Xa$

$X \rightarrow a \mid aXb \mid \varepsilon$

(1) (9 points) Construct a DFA for viable prefixes of this grammar using LR(0) items.

(2) (4 points) Identify a shift-reduce conflict and a reduce-reduce conflict under SLR(1) parsing.

【参考答案】



- (1)  $a$
- (2) 由于  $a \in \text{FOLLOW}(X) = \{a, b\}$ , 因而在状态 I0 (或 I2), 遇到输入符号  $a$  时既可移进也可归约; 由于  $a \in \text{FOLLOW}(X) = \{a, b\}$ , 因而在状态 I2, 遇到输入符号  $a$  时既可用产生式  $X \rightarrow a$  归约, 也可用产生式  $X \rightarrow \varepsilon$  归约。

8. (10 points) The following grammar generates binary fractions.

$F \rightarrow 0.B$

$B \rightarrow 0B$

$\mid 1B$

$\mid 0$

$\mid 1$

Design a syntax-directed definition (SDD) for the above grammar such that the nonterminal  $F$  has an attribute  $Eval$  which keeps the decimal value of the binary fraction generated by  $F$ . Please use as few attributes as possible and do NOT modify the grammar.

【参考答案】

产生式	语义规则
$F \rightarrow 0.B$	$F.val = B.val$
$B \rightarrow 0B_1$	$B.val = B_1.val * 0.5$
$B \rightarrow 1B_1$	$B.val = 0.5 + B_1.val * 0.5$

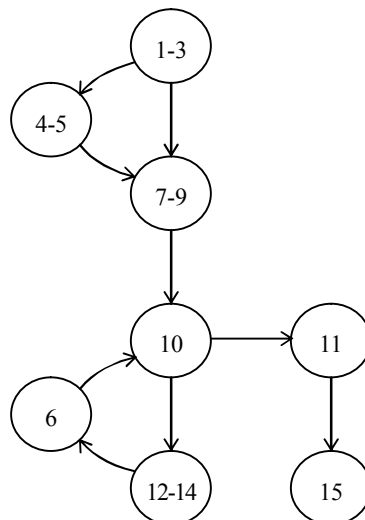
<b>B → 0</b>	<b>B.val = 0</b>
<b>B → 1</b>	<b>B.val = 0.5</b>

9. (8 points) Consider the following fragment of three-address instructions:

- |      |             |                            |
|------|-------------|----------------------------|
| (1)  |             | <b>b := 1</b>              |
| (2)  |             | <b>b := 2</b>              |
| (3)  |             | <b>if w &lt;= x goto B</b> |
| (4)  |             | <b>e := b</b>              |
| (5)  |             | <b>jump B</b>              |
| (6)  | <b>A:</b>   | <b>jump D</b>              |
| (7)  | <b>B:</b>   | <b>c := 3</b>              |
| (8)  |             | <b>b := 4</b>              |
| (9)  |             | <b>c := 6</b>              |
| (10) | <b>D:</b>   | <b>if y &lt;= z goto E</b> |
| (11) |             | <b>jump End</b>            |
| (12) | <b>E:</b>   | <b>g := g + 1</b>          |
| (13) |             | <b>h := 8</b>              |
| (14) |             | <b>jump A</b>              |
| (15) | <b>End:</b> | <b>h := 9</b>              |

Please partition these three-address instructions into basic blocks, and draw the control flow graph. You may draw the resulting graph directly, but you must mark each node by number n~m indicating that the corresponding basic block consists of instructions n through m, inclusive.

【参考答案】



10. (12 points) Consider the following basic blocks:

- |                |                |
|----------------|----------------|
| (1) T0 := 3.14 | (7) B := A     |
| (2) T1 := 2*T0 | (8) T5 := 2*T0 |
| (3) T2 := R+r  | (9) T6 := R+r  |

(4)  $T3 := R - r$                       (10)  $T7 := T3 - T5$

(5)  $T4 := T3 - T1$                     (11)  $B := A - T7$

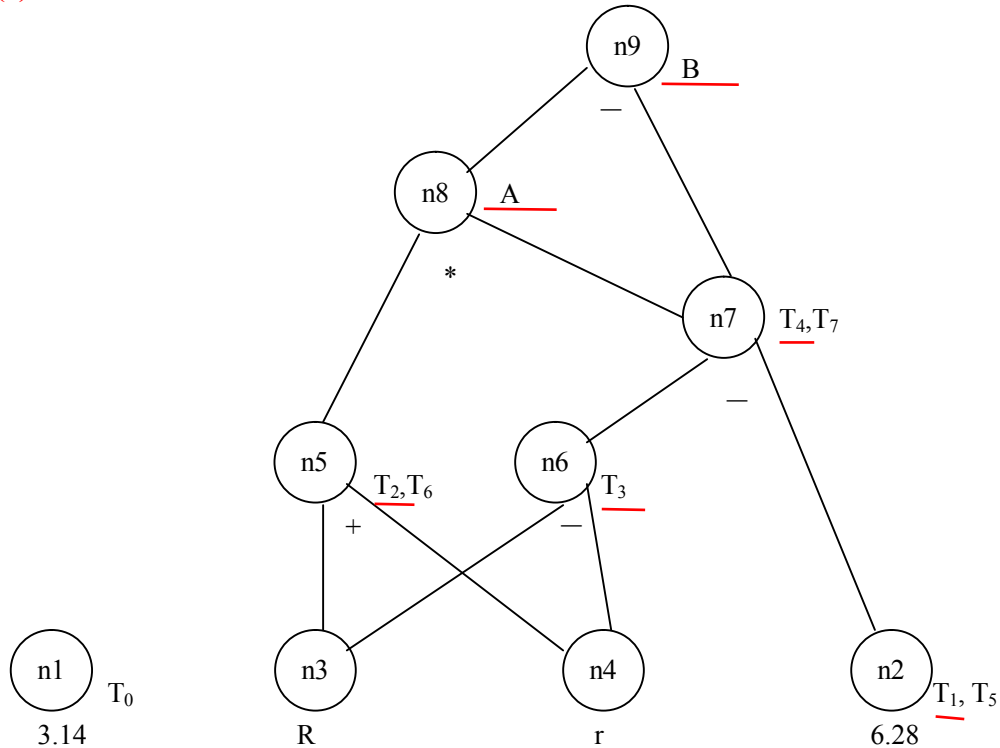
(6)  $A := T2 * T4$

(1) (7 points) Construct a DAG for this basic block.

(2) (5 points) Assuming that only A and B are live on exit from this basic block, simplify the three-address code.

【参考答案】

(1)



(2)

$S1 := R + r$

$S2 := R - r$

$S3 := S2 - 6.28$

$A := S1 * S3$

$B := A - S3$

注意本题中的临时变量 S1~S3 可以替换为其它名称.