

1. Consider the following grammar $G[S]$:

$$S \rightarrow aSbS \mid aS \mid c$$

Is this grammar ambiguous? if so, please give your reason with an example string and its parse trees.

2. Consider the following grammar $G[S]$:

$$S \rightarrow S0S \mid S1S \mid a$$

Is this grammar ambiguous? if so, please give your reason with an example string and its parse trees.

3. Consider the following grammar :

$$G[S]: S \rightarrow (L) \mid aS \mid a$$

$$L \rightarrow L,S \mid S$$

Please write the rightmost derivation for the sentential form '(S, (a))', and give the handle and the viable prefixes of this sentential form.

4. Consider the following grammar :

$$G[S]: S \rightarrow aAcB \mid Bd$$

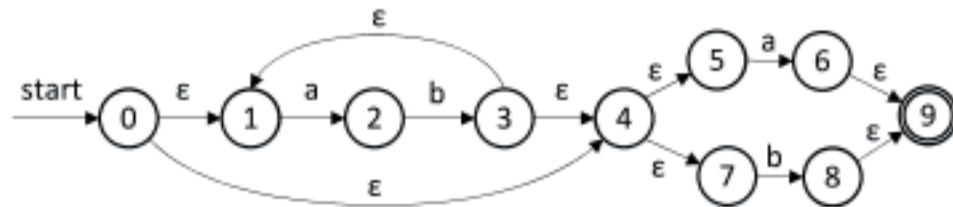
$$A \rightarrow AaB \mid c$$

$$B \rightarrow bScA \mid b$$

Please write the rightmost derivation for the sentential form 'aAcbBdcc', and give the handle and the viable prefixes of this sentential form.

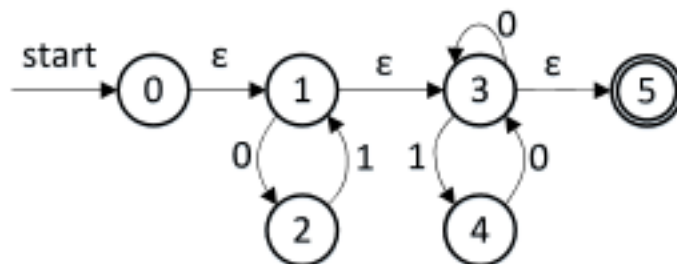
5. Construct the minimum-state DFA for the following NFA:

- 1) Convert this NFA into DFA by subset construction. Both the transition table and the transition graph of DFA are required.
- 2) Minimize the states of this DFA.



6. Construct the minimum-state DFA for the following NFA:

- 1) Convert this NFA into DFA by subset construction. Both the transition table and the transition graph of DFA are required.
- 2) Minimize the states of this DFA.



7. Consider the following grammar $G[S]$:

$$S \rightarrow k \mid (T)$$

$$T \rightarrow T * S \mid T / S \mid S$$

- (1) Please rewrite this grammar to eliminate left recursion.
- (2) Compute FIRST and FOLLOW for the grammar. Please explain the rewritten grammar is LL(1) grammar or not.
- (3) Construct the parsing table

8. Consider the grammar (decls, decl, type, varlist and varlist' are non-terminals):

$$\text{decls} \rightarrow \text{decl}; \text{decls} \mid \epsilon$$

$$\text{type} \rightarrow \text{int} \mid \text{bool}$$

$$\text{varlist}' \rightarrow , \text{varlist} \mid \epsilon$$

$$\text{decl} \rightarrow \text{type varlist}$$

$$\text{varlist} \rightarrow \text{id varlist}'$$

- 1) Construct First and Follow sets for the nonterminals.

9. Given the grammar $G[S]: S \rightarrow (S)A \mid aA, \quad A \rightarrow BA \mid \varepsilon, \quad B \rightarrow S \mid +S \mid *.$ (here, $A \rightarrow BA$ having higher priority on $A \rightarrow \varepsilon$), and the parsing table of $G[S]$ as follow.

	a	()	+	*	\$
S	$S \rightarrow aA$	$S \rightarrow (S)A$				
A	$A \rightarrow \epsilon$ $A \rightarrow BA$	$A \rightarrow \epsilon$ $A \rightarrow BA$	$A \rightarrow \epsilon$	$A \rightarrow \epsilon$ $A \rightarrow BA$	$A \rightarrow \epsilon$ $A \rightarrow BA$	$A \rightarrow \epsilon$
B	$B \rightarrow S$	$B \rightarrow S$		$B \rightarrow +S$	$B \rightarrow *$	

To achieve the predictive parsing, the rule " $A \rightarrow BA$ having higher priority on $A \rightarrow \epsilon$ when selecting a production of A for derivations" is added into the grammar $G[S]$. Please give the parsing process for the input string a^*a in the following table.

[illegible]

10. Given the grammar $G[S]: S \rightarrow (S)A \mid aA, A \rightarrow BA \mid \varepsilon, B \rightarrow S \mid +S \mid *$. (here, $A \rightarrow BA$ having higher priority on $A \rightarrow \varepsilon$), and the parsing table of $G[S]$ as follow.

	a	()	+	*	\$
S	$S \rightarrow aA$	$S \rightarrow (S)A$				
A	$A \rightarrow \varepsilon$ $A \rightarrow BA$	$A \rightarrow \varepsilon$ $A \rightarrow BA$	$A \rightarrow \varepsilon$	$A \rightarrow \varepsilon$ $A \rightarrow BA$	$A \rightarrow \varepsilon$ $A \rightarrow BA$	$A \rightarrow \varepsilon$

11. Consider the following augmented grammar $G[S']$:

- (0) $S' \rightarrow S$ (1) $S \rightarrow Pa$ (2) $S \rightarrow Pb$ (3) $S \rightarrow c$
 (4) $P \rightarrow Pd$ (5) $P \rightarrow Se$ (6) $P \rightarrow f$

- 1) Construct the DFA of LR(0) items for this augmented grammar.
- 2) Is this grammar the LR(0) or SLR(1) grammar ? Give your reason.
- 3) Construct the SLR(1) parsing table.

12. Consider the following augmented grammar $G[S']$:

- (0) $S' \rightarrow S$ (1) $S \rightarrow iDeD$ (2) $S \rightarrow iD$ (3) $D \rightarrow Sb$ (4) $D \rightarrow \epsilon$

- 1) Construct the DFA of LR(0) items for this augmented grammar.
- 2) Is this grammar the LR(0) or SLR(1) grammar ? Give your reason.
- 3) Construct the SLR(1) parsing table.

13. Consider the following attribute grammar:

Grammar	Semantic Rules	
While-stmt \rightarrow while E do S	While-stmt.begin=newlabel; While-stmt.next=newlabel; E.true=newlabel; E.false= While-stmt.next; S.next= While-stmt.begin; While-stmt.code=Label While-stmt.begin E.code Label E.true S.code goto While-stmt.begin Label While-stmt.next	<pre> graph TD Start(()) --> Begin[While begin:] Begin --> ETrue[E.true] ETrue --> Box[E.code S.code] Box -- goto While.begin --> Begin EFalse[E.false] --> Exit(()) </pre>

$E \rightarrow E1 \text{ or } E2$	$E1.true = E.true;$ $E1.false = \text{newlabel};$ $E2.true = E.true;$ $E2.false = E.false;$ $E.code = E1.code \parallel \text{Label } E1.false \parallel E2.code$	<pre> graph LR E1_false[E1.false] --> E2_false[E2.false] E1_true[E1.true] --> E_true[E.true] E2_true[E2.true] --> E_true E2_false[E2.false] --> E_false[E.false] </pre>
$E \rightarrow id1 \text{ relop } id2$	$E.code = \text{if } id1.name \text{ relop } id2.name \text{ goto } E.true \parallel \text{goto } E.false$	
$S \rightarrow id = id + num$	$S.code = id.name = id.name + num.val$	

Given the source code: while a<b or c<d do t= t+a

- (1) Draw the Abstract Syntax Tree
- (2) According to the semantic rules, calculate the inherited attributes ‘true’, ‘false’ and ‘next’ on the corresponding nodes of the syntax tree, to form the semantic tree.
- (3) Consider the step (2) result and the synthetic attribute ‘code’ , translate the three address code in a bottom-up order, recursively.

14. Consider the following attribute grammar:

Grammar	Semantic Rules	
$\text{if-stmt} \rightarrow \text{if } E \text{ then } S1 \text{ else } S2$	$\text{If-stmt.next} = \text{newlabel};$ $E.true = \text{newlabel};$ $E.false = \text{newlabel};$ $S1.next = \text{if-stmt.next};$ $S2.next = \text{if-stmt.next}$ $\text{If-stmt.code} = E.code \parallel \text{Label } E.true \parallel S1.code \parallel \text{goto } S.next \parallel \text{Label } E.false \parallel S2.code \parallel \text{if-stmt.next}$	<pre> graph TD E_true[E.true] --> S1_code[S1.code] E_false[E.false] --> S2_code[S2.code] S1_code --> if_stmt_next[if-stmt.next] S2_code --> if_stmt_next if_stmt_next --> E_true </pre>

$E \rightarrow E1 \text{ or } E2$	$E1.true = E.true;$ $E1.false = \text{newlabel};$ $E2.true = E.true;$ $E2.false = E.false;$ $E.code = E1.code \parallel \text{Label } E1.false$ $\parallel E2.code$	
$E \rightarrow id1 \text{ relop } id2$	$E.code = \text{if } id1.name \text{ relop } id2.name \text{ goto } E.true \parallel \text{goto } E.false$	
$S \rightarrow id = num$	$S.code = id.name = num.val$	

Given the source code: if a<b or c<d then t= 5 else t=10

- (1) Draw the Abstract Syntax Tree
- (2) According to the semantic rules, calculate the inherited attributes ‘true’, ‘false’ and ‘next’ on the corresponding nodes of the syntax tree, to form the semantic tree。
- (3) Consider the step (2) result and the synthetic attribute ‘code’ , translate the three address code in a bottom- up order, recursively。

考试中有选择题型 20 分。这部分不出复习题了