* 继MC:←MC可理件为图上的概率随机游走 lumping

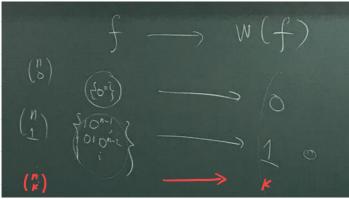
Ehrenfest urn:

$$\chi_t + \gamma_t$$
 $\chi_t + \gamma_t = n$

$$\begin{cases} \gamma_{t+1} = \gamma_{t-2}, \\ \gamma_{t+1} = \gamma_{t-2}, \end{cases}$$

Reversible MC:

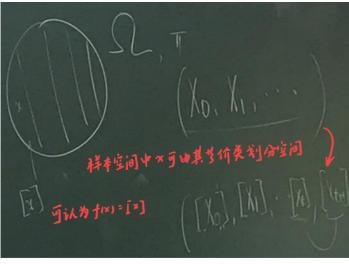




$$\begin{array}{c} X_i: S \ni V \\ \hline f: V \ni W \end{array} \right) \begin{array}{c} f(x_i) \\ S \ni W \end{array}$$

Mi 为MC, 给たナ、ナ(xi)是否为MC?

不一定是,给出反例以及为MC时所需条件



使[xiii]只与[xi]相关

P(x.[y]) = P(x', [y']) whenever [x] = [x'] $\Lambda^{\#} = \{[x] : x \in \Lambda\}$

P#(ix], [y]) = P(x,[y]) 若满足滋养件, fon 仍为MC

S={s.,...,sk}有限状态空间

si → sj : ∃n · (P");j >o i有概率多次转移后到达j

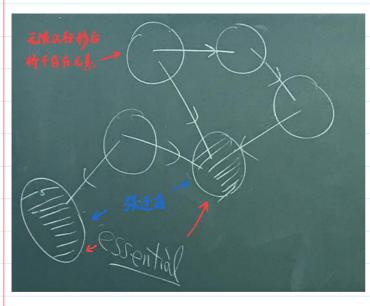
si → sj : ∃n·(P")ij >o i有概率多次转移后到达j

* 张连通分量:

A MC is irreducible if for all si.sjes 131=n, we have si > sj

> I + P + P + ... > 0

1+ P + P + ... + P -1 > 0



1 定理:

For any finite state irreduciable Markor chain, there exists

at least one stationary distribution - AP: A. A 30

≫ 希 P 不可约: ス。P。 = ス。

(00(120)) (100000 元月0) = (020) ←矩阵论

* 4n DP

←不动点、定理

 $f \in \Delta_{n-1}$, $f \Rightarrow f P$

pf: (xo. x.. xs. …) 时间转移矩阵

For $x \in \Omega$, hitting time for x, $T_x := min\{t>0; x_t=x\}$

Tx := min { t = 1; xt = x}

Lem: For any states x.y of an irreducible chain,

 $E_{\pi}(\tau_y^*) < \infty$

圣徒德堡悖论

 $pt: E_{x}(ty^{\dagger}): \sum_{t=1}^{\infty} t \cdot P_{x}(ty^{\dagger} = t)$ 太过国理 $\longrightarrow P(X = 2^{n}) = (\frac{1}{2})^{n+1}$ $\begin{cases} P(X < \infty) = 1 \\ EX = \infty \end{cases}$

There 3 8>0 & r>o s.t.: For Yz, west 3jer s.t. Pizzw)> E 全田强连通性疾易理中

Px { Ty + > kr } < (1-E) Px { Ty + > (k-1) r} > Px { Ty + > kr } < (1-E)*

 $E_x(\overline{\imath_y}^*):\sum_{t=1}^\infty t\cdot P_x(\overline{\imath_y}^*=t)=\sum_{t\geq 0}P_x\{\overline{\imath_y}^*>t\}$ ← 期望面积换维度计算

 $\leq \sum_{k\geq 0} r P_x \left\{ \overline{l}_y^+ > k_r \right\} < r \sum_{k\geq 0} \left(l - \varepsilon \right)^k = \frac{r}{\varepsilon} < \infty$



주可行: dim Ker(P-1):1 , Ker(P-1): {f: Pf:f}



EPai,j)faj)=fa) faj max

) ラ所有値均相な、ナカー作

下找该唯一任任: $\pi(x) = \frac{1}{E_x(\tau_x^*)}$ $\sum_{n \in \Omega} \pi(x) = 1$

 $Pick \ z \in \Omega$, $\widetilde{\pi}(y) = E_z(z)$ 过程中访问y次数) $\leq E_z T_z$

Claim: $\widetilde{\pi}P = \widetilde{\pi}$

