Lecture 6

Projection methods

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Roadmap

- 1. What is projection
- 2. How we approximate functions

An arbitrary infinite horizon problem can be represented using a Bellman equation

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We'll focus on value functions here, but we can approximate policy functions as well

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Rearrange the Bellman equation and define a new functional $m{H}$ that maps the problem into a more general framework

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We can find some function V that solves H(V)=0

How do we do this?

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We solve this by specifying some linear combination of **basis functions** $\Psi_i(\mathbf{S})$

$$V^j(\mathbf{S}| heta) = \sum_{i=0}^j heta_i \Psi_i(\mathbf{S})$$

with coefficients $\theta_0, \ldots, \theta_j$

We then define a residual

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We still have some choices to make:

What basis do we select?

How do we project (select the coefficients)?

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Ordinary least squares linear regression

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Ordinary least squares linear regression

We can think of the problem as searching for some unknown conditional expectation E[Y|X], given outcome variable Y and regressors X

We don't know the true functional form, but we can approximate it using the first two monomials on X: 1 and X

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In OLS we use observed data, but in theory we use the operator H(V)

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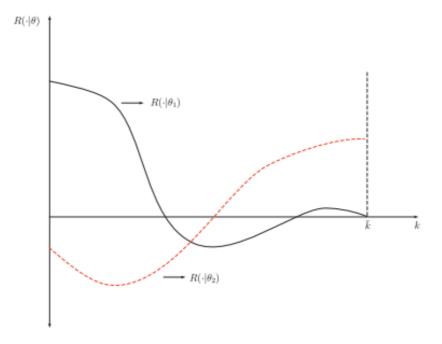
 ρ tells us how close our residual function is to zero over the domain of our state space

Example residuals given different projections

Example: The figure shows two different residuals on some capital domain of $[0,ar{k}]$

The residual based on the coefficient vector $heta_1$ is large for small values of capital but near-zero everywhere else

Figure 2: Residual Functions



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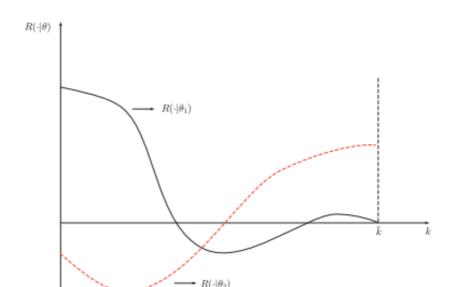


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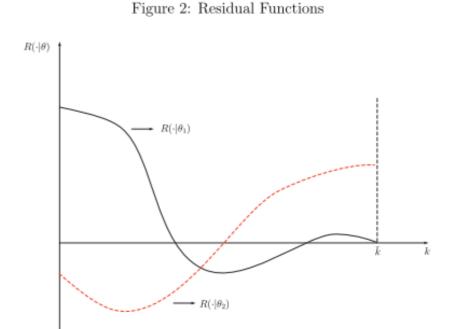
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Which is closer to zero over the interval? It will depend on our selection of ho



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The one-dimensional metric is defined as

$$ho(R\cdot| heta,0) = egin{cases} 0 & ext{if } \int_{\Omega}\phi_i(\mathbf{S})R(\cdot| heta)d\mathbf{S} = 0, i = 1,\ldots,j+1 \ 1 & ext{otherwise} \end{cases}$$

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Where we want to solve for $heta=rgmin
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First lets begin with a simple example before moving into the most commonly used weight functions

Suppose we selected the weight function to be

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Then we would be performing least squares! Why?

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The partial derivative weight function yields a metric function that solves the least squares problem!

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Here our weight function is

$$\phi_i(\mathbf{S}) = \delta(\mathbf{S} - \mathbf{S}_i)$$

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What does this weight function mean?

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Since we have a finite set of points we do not need to solve difficult integrals but only a system of equations

$$R(\mathbf{S}_i| heta)=0, i=1,\ldots,j+1$$

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But how do we implement collocation?

Recall we solve for coefficients θ by setting the residual to be zero at all of our collocation points

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We can solve this problem by *iterating* on the problem, continually setting the residuals equal to zero, recovering new θ s, and repeating

In any given iteration, we:

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- 5. Use these new maximized values to obtain updated coefficients solving the system of linear equations, and repeat the process until we have "converged"

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Do your answers to the previous two questions really matter?

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Do your answers to the previous two questions really matter?

Yes they are crucial

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Therefore we can simulate anything we want and recover optimal policy functions given many different sets of initial conditions or realizations of random variables

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We do so by selecting a specific finite number of points in our state space and use them to construct a *collocation grid* that spans the domain of our problem

Using our knowledge of how the value function behaves at the limited set of points on our grid, we can interpolate our value function approximant at all points off the grid points, but within the domain of our grid

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Note: The value function approximant is not very good outside the grid's domain since that would mean extrapolating beyond whatever information we have gained from analyzing our value function on the grid

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Each $\psi_j(x)$ is a basis function, and the coefficients c_j determine how they are combined at some point \bar{x} to yield our approximation $\hat{V}(\bar{x})$ to $V(\bar{x})$

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This is what happens we select our number of grid points in the state space to be equal to the number of coefficients (which induces a Dirac delta weighting function)

Solve a system of equations, linear in c_j that equates the function approximant at the grid points to the recovered values

$$\Psi \mathbf{c} = \mathbf{y}$$

where Ψ is the matrix of basis functions, c is a vector of coefficients, and y is a vector of the recovered values

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Why?

The Stone-Weierstrass Theorem which states (for one dimension)

Suppose f is a continuous real-valued function defined on the interval [a,b]. For every $\epsilon>0,\ \exists$ a polynomial p(x) such that for all $x\in [a,b]$ we have $||f(x)-p(x)||_{sup}\leq \epsilon$

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What does the SW theorem say in words?

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Note that the SW theorem $does\ not\ say\ what\ kind\ of\ polynomial\ can\ approximate\ f$ arbitrarily well, just that some polynomial exists

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Practice

code up a function project_monomial(f, n, lb, ub) that takes in some function f, degree of approximation n, lower bound lb and upper bound ub, and constructs a monomial approximation on an evenly spaced grid via collocation

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monomial approximation on an evenly spaced grid via collocation

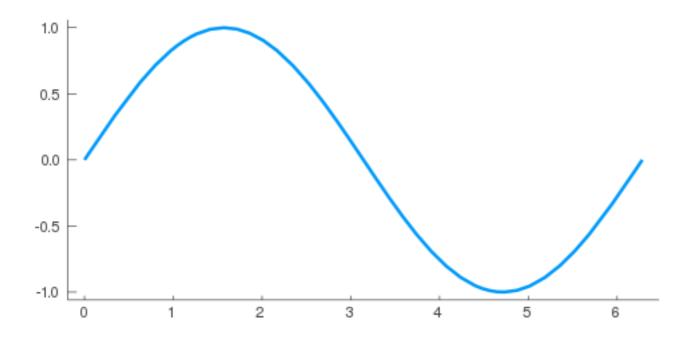
We will be plotting stuff, see http://docs.juliaplots.org/latest/generated/gr/ for example code using the GR backend

The monomial basis

Let's approximate sin(x)

```
using Plots
gr();
f(x) = sin.(x);

Plots.plot(f, 0, 2pi, line = 4, grid = false, legend = false, size = (500, 250))
```



```
function project monomial(f, n, lb, ub)
    # solves \Psi c = v \rightarrow c = \Psi \backslash v
    \# \Psi = matrix \ of \ monomial \ basis \ functions \ evaluted \ on \ the \ grid
    coll points = range(lb, ub, length = n)
                                                                        # collocation points
                                                                        # function values on the grid
    v values = f(coll points)
    basis functions = [coll points.^degree for degree = 0:n-1] # vector of basis functions
    basis matrix = hcat(basis functions...)
                                                                        # basis matrix
    coefficients = basis matrix\y values
                                                                        \# c = \Psi \backslash V
    return coefficients
end;
coefficients 4 = project monomial(f, 4, 0, 2pi);
coefficients 5 = project monomial(f, 5, 0, 2pi);
coefficients 10 = project monomial(f, 10, 0, 2pi)
```

```
## 10-element Array{Float64,1}:
## 0.0
## 0.9990725797458863
## 0.004015857153649684
## -0.1738437387373486
## 0.007075663351630060
```

Now we need to construct a function f_approx(coefficients, plot_points) that takes in the coefficients vector, and an arbitrary vector of points to evaluate the approximant at (for plotting)

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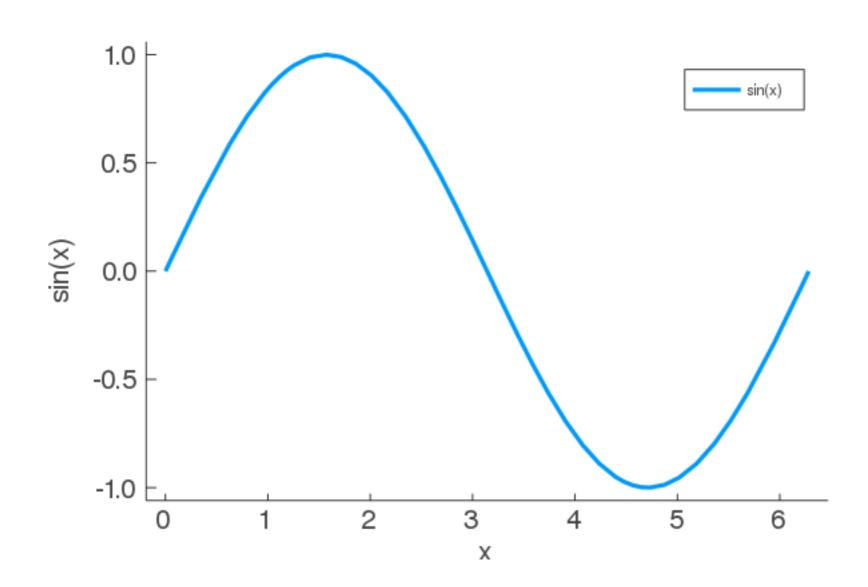
```
function f_approx(coefficients, points)
    n = length(coefficients) - 1
    basis_functions = [coefficients[degree + 1] * points.^degree for degree = 0:n] # evaluate basis functions
    basis_matrix = hcat(basis_functions...) # transform into matrix
    function_values = sum(basis_matrix, dims = 2) # sum up into function value
    return function_values
end;
```

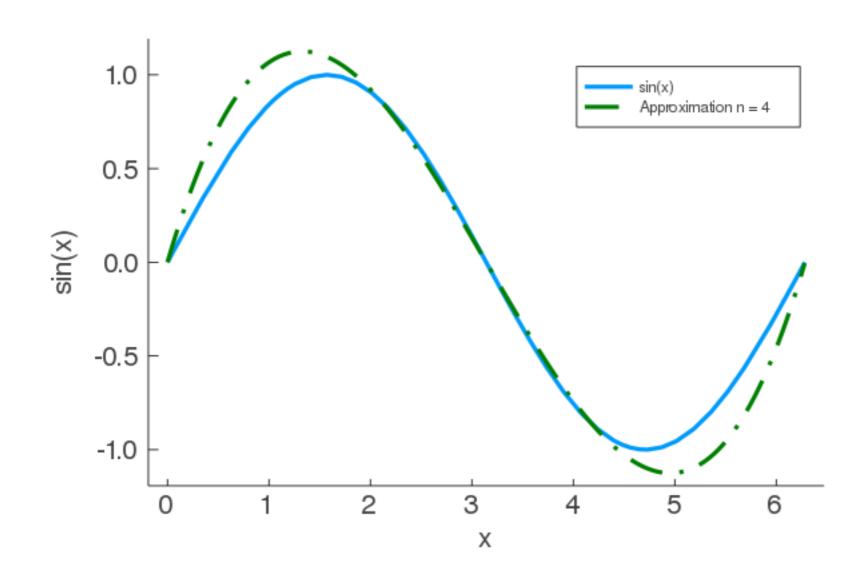
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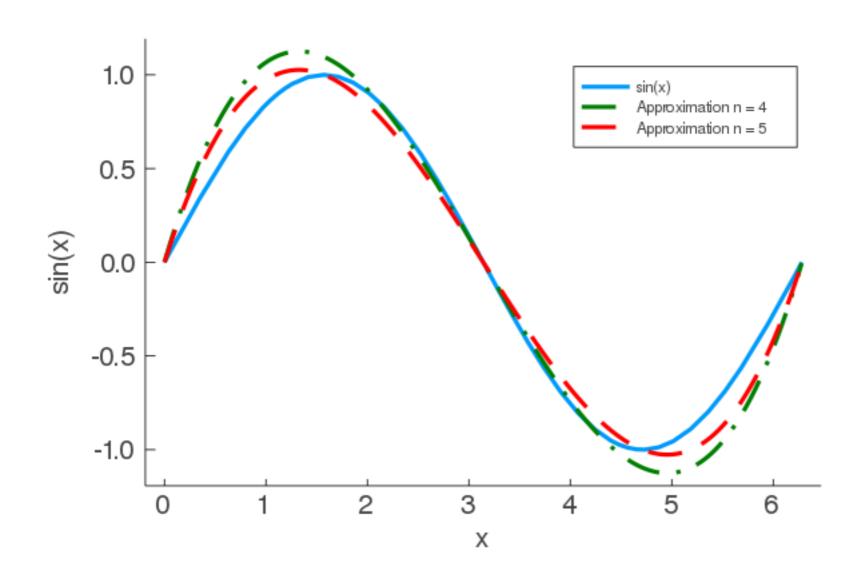
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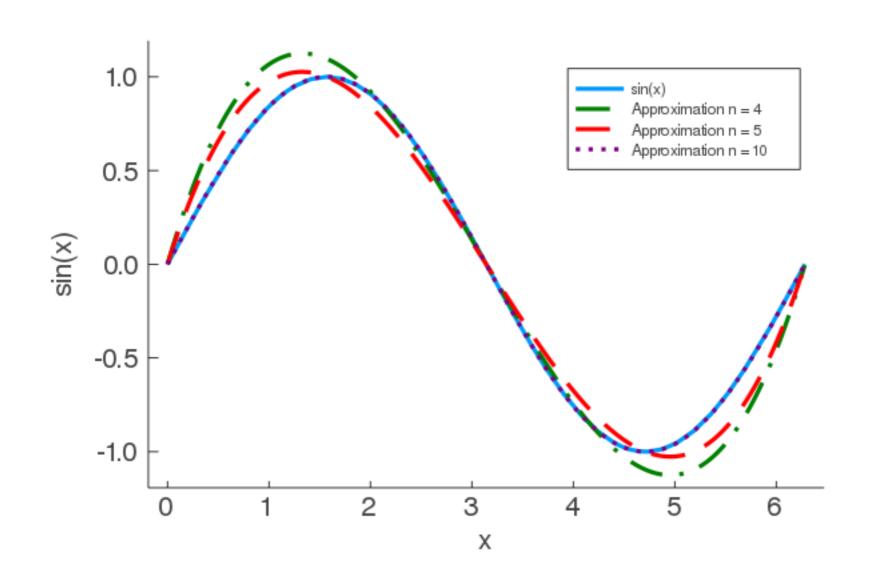
```
plot_points = 0:.01:2pi;
f_values_4 = f_approx(coefficients_4, plot_points);
f_values_5 = f_approx(coefficients_5, plot_points);
f_values_10 = f_approx(coefficients_10, plot_points)
```

```
## 629×1 Array{Float64,2}:
## 0.0
## 0.099999953610597868
```









We just wrote some code that exploits Stone-Weierstrauss and allows us to (potentially) approximate any continuous function arbitrarily well as n goes to infinity

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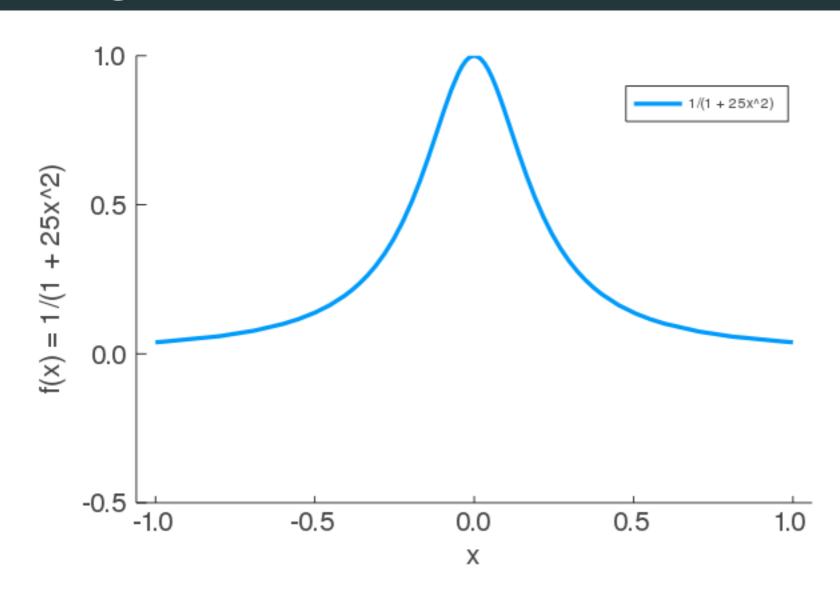
Why?

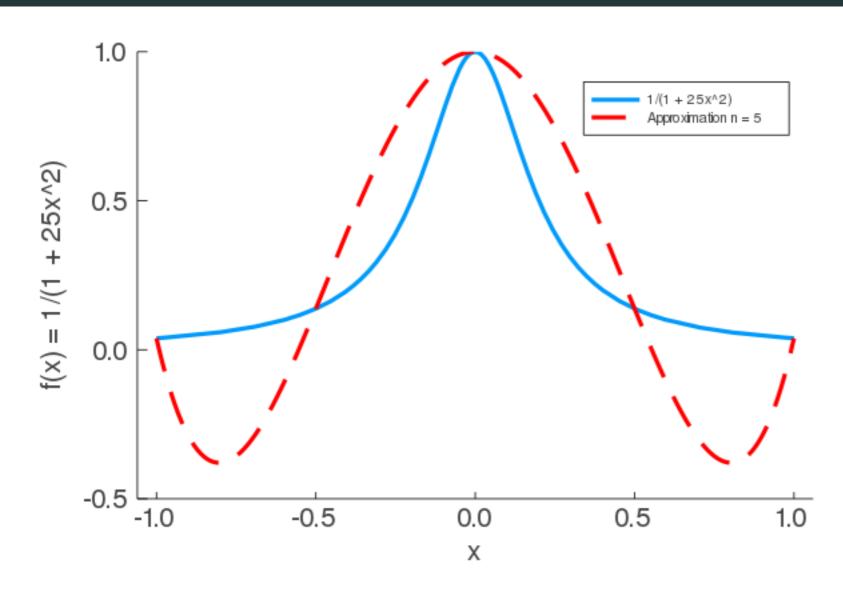
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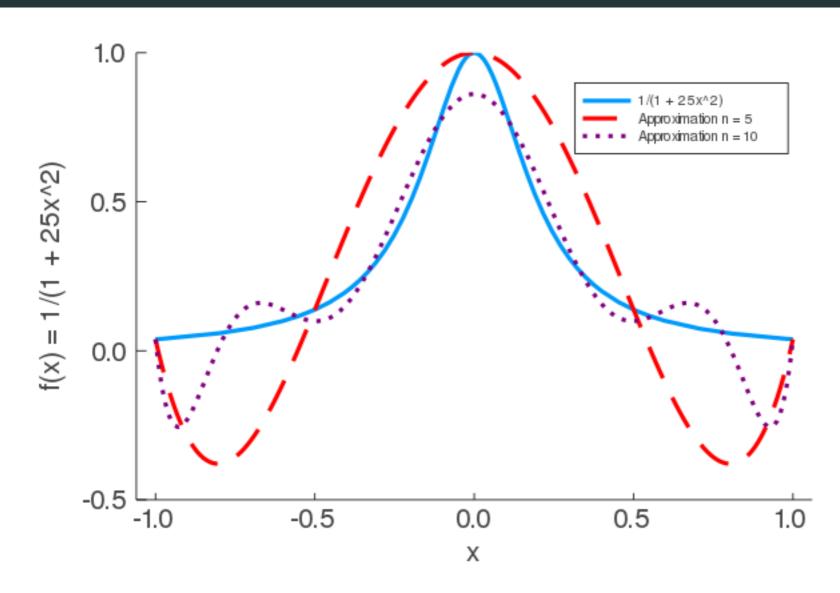
Why?

Try approximating **Runge's function**: $f(x) = 1/(1 + 25x^2)$

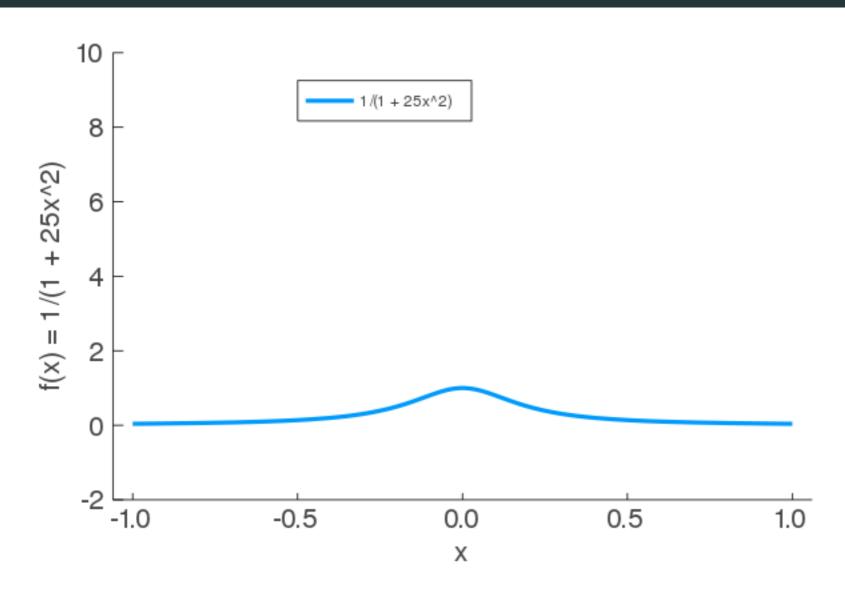
```
runge(x) = 1 ./ (1 .+ 25x.^2);
coefficients_5 = project_monomial(runge, 5, -1, 1);
coefficients_10 = project_monomial(runge, 10, -1, 1);
plot_points = -1:.01:1;
runge_values_5 = f_approx(coefficients_5, plot_points);
runge_values_10 = f_approx(coefficients_10, plot_points);
```



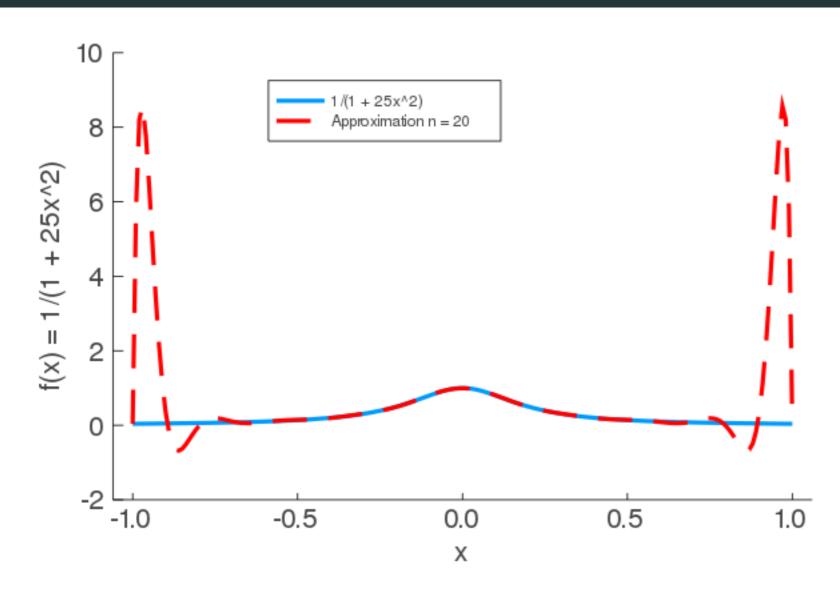




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Second, they can vary dramatically in size, which leads to scaling/truncation errors

```
runge(x) = 1 ./ (1 .+ 25x.^2);
coefficients_10 = project_monomial(runge, 10, -1, 1);
points = rand(10);
n = length(coefficients_10) - 1;
basis_functions = [coefficients_10[degree + 1] * points.^degree for degree = 0:n];
basis_matrix = hcat(basis_functions ...);
println("Condition number: $(cond(basis_matrix))")
```

Condition number: 1.7447981077501208e20

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Ideally we want an **orthogonal basis**: when we add another element of the basis, it has sufficiently different behavior than the elements before it so it can capture features of the unknown function that the previous elements couldn't

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The Chebyshev polynomial closely approximates the minimax polynomial: the polynomial, given degree d, that minimizes any approximation error to the true function

Chebyshev polynomials are defined by a recurrence relation,

$$egin{aligned} T_0(x) &= 1 \ T_1(x) &= x \ T_{n+1} &= 2xT_n(x) - T_{n-1}(x) \end{aligned}$$

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Chebyshev polynomials look similar to monomials but have better properties that are visually distinctive

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If you can't get the recurrence relation to work, you can use an alternative:

$$T_n(x) = cos(n \operatorname{arccos}(x))$$

The two basis functions

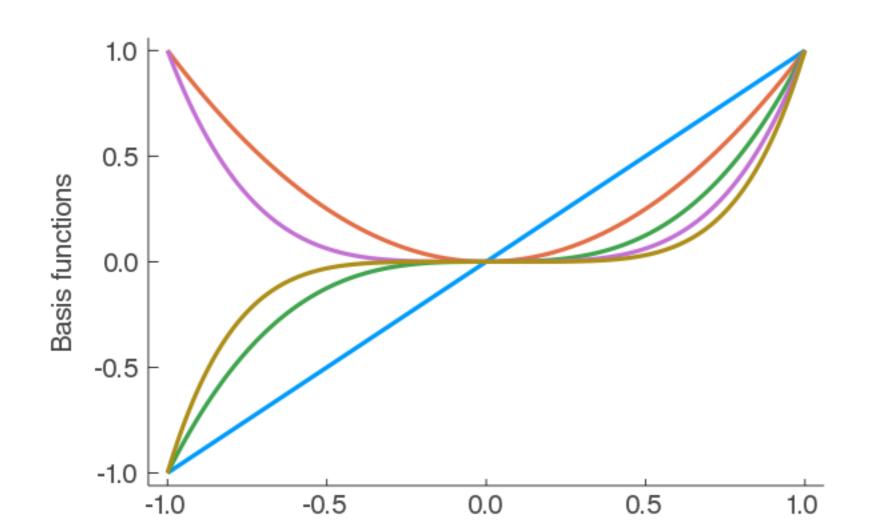
```
# Chebyshev polynomial function
function cheb_polys(x, n)
    if n = 0
                                   \# T O(x) = 1
        return x./x
    elseif n = 1
                                     \# T 1(x) = x
        return X
    else
        cheb_recursion(x, n) =
            2x.*cheb polys.(x, n - 1) \cdot - cheb polys.(x, n - 2)
        return cheb recursion(x, n) # T n(x) = 2xT \{n-1\}(x) - T \{n-2\}(x)
    end
end;
# Monomial function
monomials(x, n) = x.^n;
```

The plotting function

```
function plot function(basis function, x, n)
    for i = 1:n-1
        f data = basis function(x, i)
        if i = 1
            plot(x, f data, linewidth = 4.0, xlabel = "x", ylabel = "Basis functions", label = "",
                 tickfontsize = 14, guidefontsize = 14, grid = false);
        else
            plot!(x, f data, linewidth = 4.0, label = "");
        end
    end
    f data = basis function(x, n)
    plot!(x, f data, linewidth = 4.0, label = "")
end;
x = -1:.01:1;
```

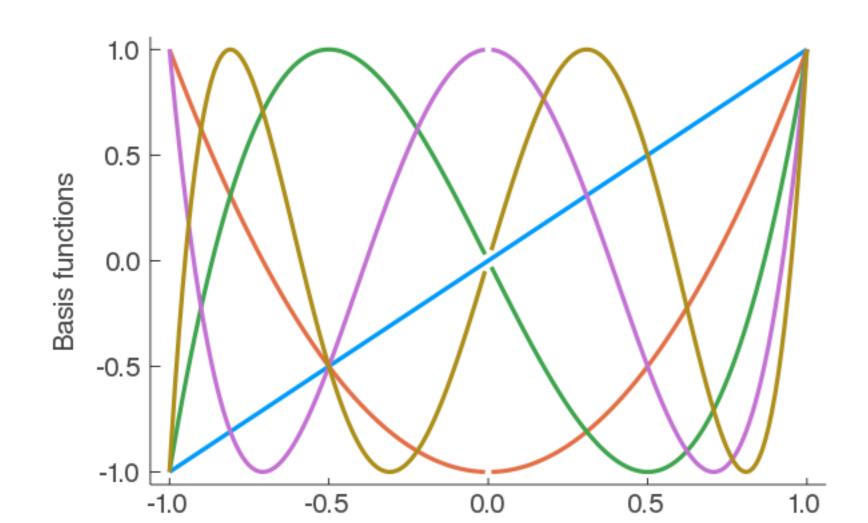
Monomials up to degree 5

```
plot_function(monomials, x, 5)
```



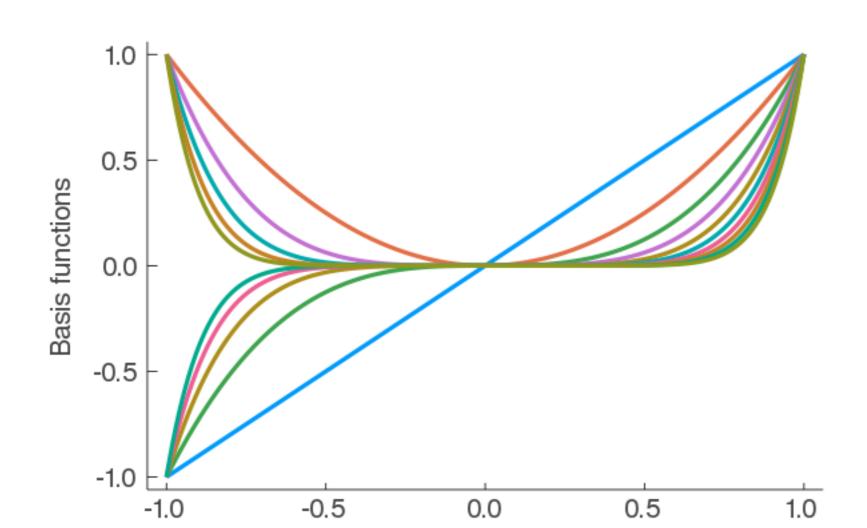
Chebyshev polynomials up to degree 5

plot_function(cheb_polys, x, 5)



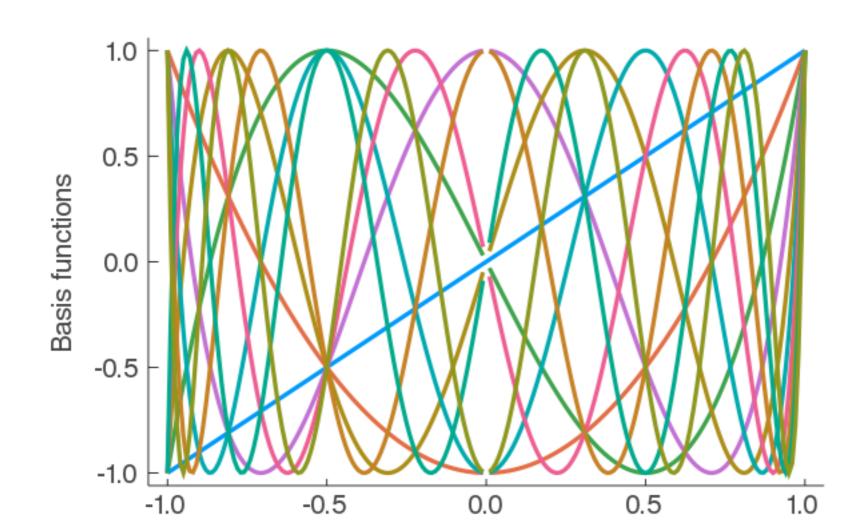
Monomials up to degree 10

```
plot_function(monomials, x, 10)
```



Chebyshev polynomials up to degree 10

plot_function(cheb_polys, x, 10)



Chebyshev polynomials span the space

Chebyshev polynomials span the space

Monomials clump together

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It also guarantees that Φ has full rank and is invertible

Chebyshev zeros and alternative rep

Also note that the Chebyshev polynomial of order n has n zeros given by

$$x_k = cos\left(rac{2k-1}{2n}\pi
ight), \;\; k=1,\ldots,n$$

which tend to cluster quadratically towards the edges of the domain

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You can think about this as projecting sequentially finer uniform grids from a hemicircle onto the x-axis

There are two important theorems to know about Chebyshev polynomials

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Chebyshev interpolation theorem: If $f(x) \in \mathbb{C}[a,b]$, if $\{\psi_i(x), i=0,\dots\}$ is a system of polynomials (where $\psi_i(x)$ is of exact degree i) orthogonal with respect to $\phi(x)$ on [a,b] and if $p_j = \sum_{i=0}^j \theta_i \psi_i(x)$ interpolates f(x) in the zeros of $\psi_{n+1}(x)$, then:

$$\lim_{j o\infty}ig(||f-p_j||_2ig)^2=\lim_{j o\infty}\int_a^b\phi(x)ig(f(x)-p_jig)^2dx=0$$

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What does this say?

If we have an approximation set of basis functions that are exact at the roots of the n^{th} order polynomials, then as n goes to infinity the approximation error becomes arbitrarily small and converges at a quadratic rate

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Unfortunately we cant store an infinite number of polynomials in our computer, we would like to know how big our error is after truncating our sequence of polynomials

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Since Chebyshev polynomials are bounded between [-1,1], the sum of the omitted terms is bounded by the sum of the magnitude of the coefficients

So the error in the approximation is as well!

We often also have that Chebyshev approximations geometrically converge which give us the following convenient property:

$$|f(x) - f^j(x| heta)| \sim O(heta_j)$$

The truncation error by stopping at polynomial j is of the same order as the magnitude of the coefficient θ_j on the last polynomial

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Thus in many situations we can simply check the size of the last polynomial to gauge how accurate our approximation is

Boyd's moral principle

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John Boyd summarizes decades of experience with function approximation with his moral principle:

- When in doubt, use Chebyshev polynomials unless the solution is spatially periodic, in which case an ordinary fourier series is better
- Unless you are sure another set of basis functions is better, use Chebyshev polynomials
- Unless you are really, really sure another set of basis functions is better use Chebyshev polynomials

Grid point selection

We construct the basis function approximant by evaluating the function on a predefined grid in the domain of \emph{V}

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If we have precisely n nodes, x_i , we then have

$$\sum_{j=1}^n c_j \phi_j(x_i) = V(x_i) \,\, orall i = 1, 2, \ldots, n \,\,\,\, ext{(interpolation conditions)}$$

Grid point selection

We can write this problem more compactly as

$$\Phi c = y$$

(interpolation equation)

where

- ullet y is the column vector of $V(x_i)$
- c is the column vector of coefficients c_i
- ullet Φ is an n imes n matrix of the n basis functions evaluated at the n points

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If we recover a set of values at our interpolation nodes, $V^*(x_i)$, we can then simply invert Φ and right multiply it by y to recover our coefficients

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How do we select our set of nodes x_i ?

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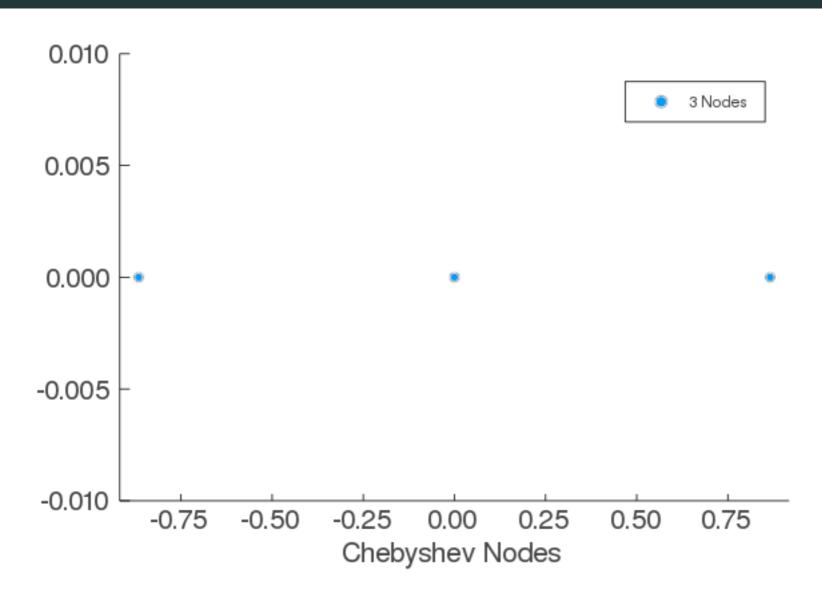
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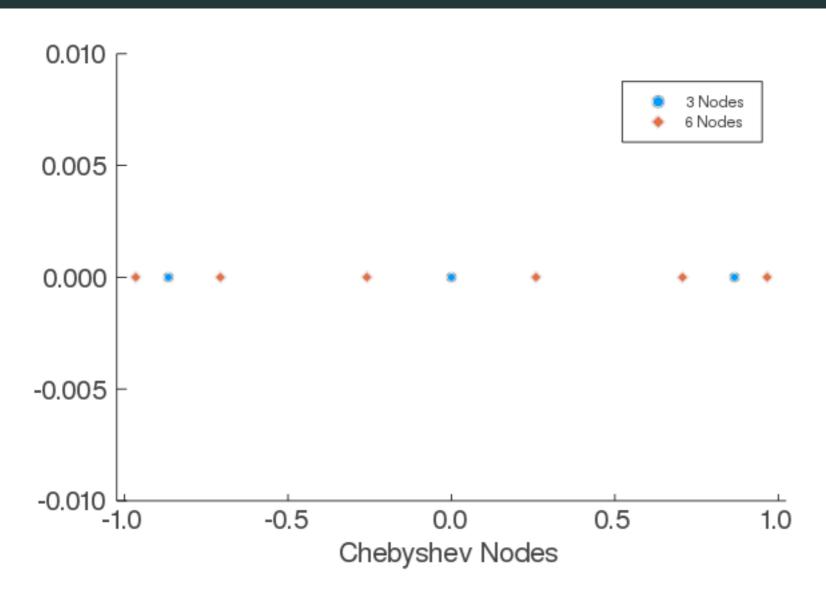
We can gain intuition by looking at a graph of where Chebyshev nodes are located, plot them yourself!

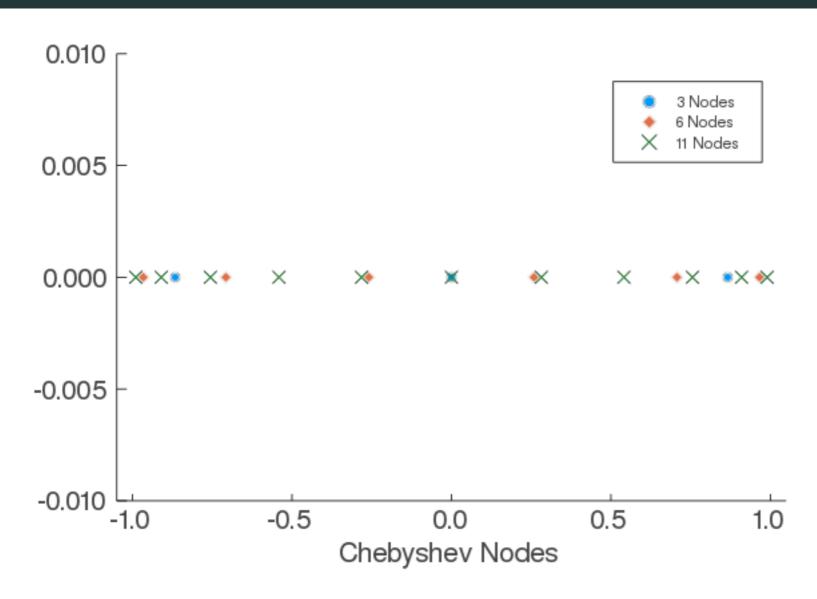
Chebyshev node function

```
cheb_nodes(n) = cos.(pi * (2*(1:n) •- 1)./(2n))
```

cheb_nodes (generic function with 1 method)







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Therefore, it's best to put more nodes in these areas to shore up this informational deficit and get good approximation quality near the edges of our domain

Discrete states

How do we handle a discrete state S_d when trying to approximate V?

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Just like you might expect, we effectively have a different function approximant over the continuous states for each value of S_d

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Construct a vector of polynomials $[\phi_{1,1}, \phi_{1,2}, \phi_{1,3}]$ for dimensions 1

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We can then solve for the 9 coefficients on these two dimensional polynomials

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We will cover smart ways to deal with this later

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For example, we can use **splines** (piecewise polynomials) over segments of our domains where they are spliced together at prespecified breakpoints, which are called knots

The higher the order the polynomial we use, the higher the order of derivatives that we can preserve continuity at the knots

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If we have a quadratic spline, we can also preserve the first derivative's continuity at the knots, but the second derivative will be a discontinuous step function

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These are what ensure continuity of higher order derivatives at the knots as the degree of the spline grows

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In numerical models we typically don't use linear splines because we often care about the quality of approximation of higher order derivatives, cubic splines are more common

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We can obtain 3(N-1) equations by ensuring that the approximant is continuous at all interior knots, and its first and second derivatives are continuous at all interior knots [3 imes (N+1-1-1)]

This means that the value of the left cubic polynomial equals the value of the right cubic polynomial at each interior knot

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What is often used is that the approximant's first or second derivative matches that of the function at the end points

If the derivative is of interest for optimization, or to recover some variable of economic meaning, then we may need to have these derivatives preserved well at the knots

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Useful spline packages out there: Dierckx, Interpolations, QuantEcon

Code it up!

Let's code up our own linear spline approximation function linear_spline_approx(f, knots), where f is the function we are approximating and knots are the knots

Have it return a function a function spline_eval that takes in evaluation_points as an argument where evaluation_points are the points we want to evaluate the spline approximant at

Hint: Linear splines are pretty easy, given two points (x_{i+1}, y_{i+1}) and (x_i, y_i) , the spline in between these points is given by

$$y(x) = y_i + rac{y_{i+1} - y_i}{x_{i+1} - x_i}(x - x_i)$$

Spline approximator

```
function linear spline approx(f, knots)
   function spline eval(evaluation points)
        prev knot = knots[1] # initialize previous knot
        if !(typeof(evaluation points) <: Number) # if using multiple points</pre>
            v eval = similar(evaluation points)
            y index = 1
            for knot in knots[2:end]
                current points = evaluation points[prev knot . ≤ evaluation points . < knot]
                y eval[y index:y index + length(current points) - 1] =
                    f(prev knot) .+ (f(knot) - f(prev knot))/(knot - prev knot)*(current points - prev knot)
                prev knot = knot
                v index += length(current points)
            end
        else # if using just a single point
            for knot in knots[2:end]
                if prev knot . ≤ evaluation points . < knot
                    y eval = f(prev knot) + (f(knot) - f(prev knot))/(knot - prev knot)*(evaluation point - prev
                end
                prev knot = knot
            end
        end
        return v eval
   end
```

```
f(x) = \sin(x)
## f (generic function with 1 method)
 knots coarse = 0:pi/2:2pi;
 spline_func_coarse = linear_spline_approx(f, knots_coarse);
 knots fine = 0:pi/4:2pi;
 spline func fine = linear spline approx(f, knots fine);
 knots superfine = 0:pi/12:2pi;
 spline func superfine = linear spline approx(f, knots superfine);
x vals =0:.05:2pi;
y_vals_coarse = spline_func_coarse(x_vals);
y vals fine = spline func fine(x vals);
y_vals_superfine = spline_func_superfine(x_vals);
```

