Correction to exercise 03:

For Riscode
$$k_1 = 3$$
, $n_1 = 6$

Hence $R = \frac{k_1}{n_1} = \frac{1}{2}$

The rate of a code belongs to the interval [0:1]. e) A rate of O means that for every information boit, an infinite number of redundancy bik is transmitted - this has no practiced interest. a) A rate of 1 means Rat no error correction was implemented.

Decreasing the add rate entails an increase in redundancy, hence, a better error probability (bower).

Question 4:

For this code, since le=3, there are $l^3=8$ possible to de words.

To check of a codeword is valid or not, one has to compute C. Ht and check if it is equal to O. (1) G = [101010].

$$C. H^{+} = (101010) \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix}$$
 [mod 2]

brewsbes a for ([000] + [110] =

Question 6:

- e) Code 3 des a repetition Coode.
- e) Code 2 is a permetation of the generalise matrix of code 4.
- She he bek & BLER curves are equal, one should not use the BER & BLER curves are experition while probability of error boy the fact that a repetition and uses 3 times more energy.

 Since the BER & BLER curves are equal, one should not use repetition and be course it divides the bit rate by 3, without an proving the BER & BLER.
- The codes E, and E have the same performance.

 This is due the fact that since the generator matrix

 Ga is a permutation of G, (swap between 1st and last)

 row

the obtained codebooks are equivalent, and hence the performances are the same.

Uncoded repetition coding out performs codeing for low Es, while coding with Con Co is better at high Es

Question 7:

For a BER close to 10 dB, I would suggest using either code & on &.

Now, code & has the advantage of being a systematic code, which allows a quick encoding, hence, & would be faster to implement than . Co.