



# Pointers cont'd

# Pointers can help save memory!

- In astronomy, we sometimes work with very large images
- Having multiple copies of those images “in memory” can crash your machine
  - good idea to avoid duplication where possible

# memory

- There is a way to find out how much memory IDL is using
- It's the `memory` function!
- `print, memory ( )` returns a 4-element array
  - the first element is the memory in use in *bytes*

# memory checking

```
IDL> print,memory()  
1803740      3078      946      1807690  
IDL> print,(memory())[0]/1024  
1761  
IDL> print,(memory())[0]/1024L^2  
1
```

- Memory given in # of bytes
- Divide by 1024 bytes/kilobyte
- Or 1024<sup>2</sup> bytes/megabyte

# Pointers & Memory

- Copying pointers is “cheap”, copying data is “expensive”
  - not going to quantify that though



# Free Pointers

```
x = ptr_new(5)
y = x
*x = 2
print,*x,*y
```

```
ptr_free,x
print,*y
```

```
% Invalid pointer: Y.
```

```
% Execution halted at: $MAIN$
```

- `y` points to the same data as `x`, but now it's gone



```
x = ptr_new(5)
y = x
*x = 2
print, *x, *y
```

What will print?

A) 5, 2

B) 2, 5

C) 2, 2

D) 5, 5

E) None of the above



# Arrays of Pointers

- Can declare an array of pointers!

```
pa = ptrarr(10)
```

- Defaults to all NULL pointers



Which line is a valid assignment operation?

pa = ptrarr(10)

- \*pa[1] = 5 ; A
- pa[2] = 5 ; B
- \*pa[3] = ptr\_new(5) ; C
- pa[4] = ptr\_new(5) ; D
- ; None of the above E

# Arrays of Pointers

- However, there is a keyword to `ptrarr` that “allocates memory on the heap”

## **ALLOCATE\_HEAP**

Normally, `PTRARR` sets every element of the result to the null pointer. If you wish IDL to allocate heap variables for every element of the array instead, set the `ALLOCATE_HEAP` keyword. In this case, every element of the array will be initialized to point at an undefined heap variable.

- All elements will start as `!NULL`



Which line is *not* a valid assignment operation?

```
pa = ptrarr(10,/allocate)
```

```
*pa[1] = 5 ; A
```

```
pa[2] = 5 ; B
```

```
*pa[3] = ptr_new(5) ; C
```

```
pa[4] = ptr_new(5) ; D
```

```
; None of the above E
```





What will be printed?

```
pa = ptrarr(10, /allocate)  
*pa[1] = 5  
print, *pa[1]
```

- A) <PtrHeapVar22>
- B) % Unable to convert variable to type pointer.
- C) 5
- D) 1
- E) None of the above



What will be printed?

```
*pa[3] = ptr_new(5)
pa[4] = ptr_new(5)
print,*pa[3],*pa[4]
```

- A) <PtrHeapVar22>                      5
- B) 5<PtrHeapVar22>
- C) % Unable to convert variable to type pointer.
- D) 5                      5
- E) None of the above

# Coding Example

- Problem: We want to create a sorted array, and then add elements into it. We want to keep the array sorted at all times, so we need to add each new element in the right place.



# Add a new element in order

- Could do this:

```
x = [2,3,1,5,6,9,12]
x = x[sort(x)]
print,x
x = [x,4]
x = x[sort(x)]
```

- But it's kind of a waste, since you can figure out pretty easily where that 1 new number (4) should go

# Inserting Elements

```
IDL> x = [2,3,1,5,6,9,12]
IDL> ; sort(x) returns indices
IDL> x = x[sort(x)]
IDL> print,x
      1      2      3      5      6      9      12
IDL> print,fix(sort(x))
      0      1      2      3      4      5      6
IDL> ; insert 4 - should be between 3 and 5
IDL> x = [x[0:2],4,x[3:]]
IDL> print,x
      1      2      3      4      5      6      9
```

# Inserting Generally

- You can insert any number this way
- How do we automate it?

```
IDL> print,x,fix(sort(x))
```

1	2	3	4	5	6	9	12
0	1	2	3	4	5	6	7

```
new_number = 7
```

```
ii = 0
```

```
while new_number gt x[ii] do ii++
```

```
print,ii
```



```
IDL> print,x,fix(sort(x))
```

1	2	3	4	5	6	9	12
0	1	2	3	4	5	6	7

```
new_number = 7
```

```
ii = 0
```

```
while new_number gt x[ii] do ii++
```

```
print,ii
```

What's ii?

A) 0

B) 7

C) 6

D) 5

E) None of the above



What's the best (i.e., most efficient) way to put 7  
in the right place?  
(note that all of these \*should\* work)

```
x = [x,7] & x = x[sort(x)] ; A
x = [7,x] & x = x[sort(x)] ; B
x = [x[0:ii],7,[ii+1]] ; C
x = [x[where(x lt 7)],7,x[where(x gt 7)]] ; D
; None of the above E
```

IDL> print,x,fix(sort(x))

1	2	3	4	5	6	9	12
0	1	2	3	4	5	6	7

We want to add a new number, 0. Will all of the same approaches work?

```
x = [x,0] & x = x[sort(x)] ; 1
x = [0,x] & x = x[sort(x)] ; 2
x = [x[0:ii-1],0,x[ii:~]] ; 3
x = [x[where(x lt 0)],0,x[where(x gt 0)]]; 4
```

A) Yes

B) No, 1 & 2 will fail

C) No, 3 & 4 will fail

D) No, just 3 will fail

E) No, just 4 will fail



# Building up our algorithm...

- We need an “if” statement to check the first number:

```
if new_number < x[0] then begin
    x = [new_number, x]
endif else begin
    ii = 0
    while new_number > x[ii] do ii++
    x = [x[0:ii-1], new_number, x[ii: *]]
endelse
```





```
IDL> print,x,fix(sort(x))
```

1	2	3	4	5	6	9	12
0	1	2	3	4	5	6	7

What if new\_number is greater than every element in x? (new\_number=13)

```
ii = 0  
while new_number gt x[ii] do ii++  
x = [x[0:ii-1],new_number,x[ii:*]]
```

- A) No problem, it will stick on the end
- B) We'll get an infinite loop
- C) We'll end up in the wrong place
- D) We'll get an indexing error
- E) None of the above

# Building up our algorithm...

- Need another “if” statement

```
if new_number < x[0] then begin
    x = [new_number, x]
endif else if new_number > x[-1] then begin
    x = [x, new_number]
endif else begin
    ii = 0
    while new_number > x[ii] do ii++
    x = [x[0:ii-1], new_number, x[ii: *]]
endelse
```

```
; insert a number into a sorted array
function insert,number,srtd_arr

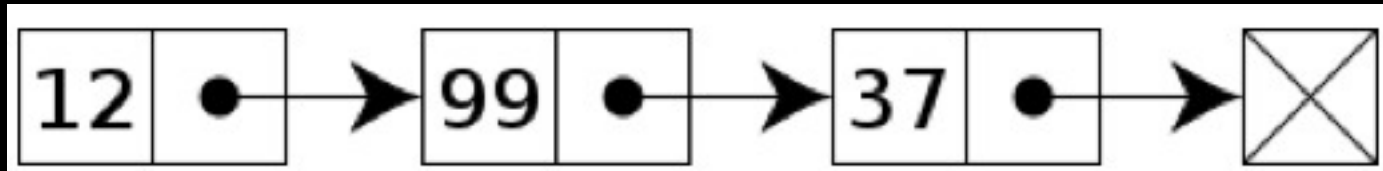
    if number < srtd_arr[0] then begin
        srtd_arr = [number,srtd_arr]
    endif else if number > srtd_arr[-1] then begin
        srtd_arr = [srtd_arr,number]
    endif else begin
        ii = 0
        while number > srtd_arr[ii] do ii++
        srtd_arr = [srtd_arr[0:ii-1],$
                    number,srtd_arr[ii:*]]
    endelse

    return,srtd_arr
end ; insert
```

# Linked Lists

- Linked Lists are another data type
- They're very efficient for inserting new elements (adding new elements in order)
- As far as I know, they're (almost) never used directly, but they serve as the conceptual basis for other things

# Linked Lists



- Made of “nodes”
- Each node has two things:
  1. data
  2. a pointer to the next node

# Linked Lists

- Nodes aren't arrays (they have two different data types)
- They COULD be lists
- But it makes more sense to make each node a struct

```
node = {Node, data: 0, next: ptr_new()}
```

# LL notation

- Here's a linked list:

`[1|next] -> [2|next] -> [3|next] -> !null`

- Each `[ | ]` is a node.
  - The left side is data
  - The right side is a pointer to the next node



# Make a linked list

- Let's start with a linked list with 3 elements: 1, 4, 2
  - We're going to keep it sorted
  - We'll add each new value in turn

# First Node

```
node1 = {Node, data: 1, next: ptr_new()}
```

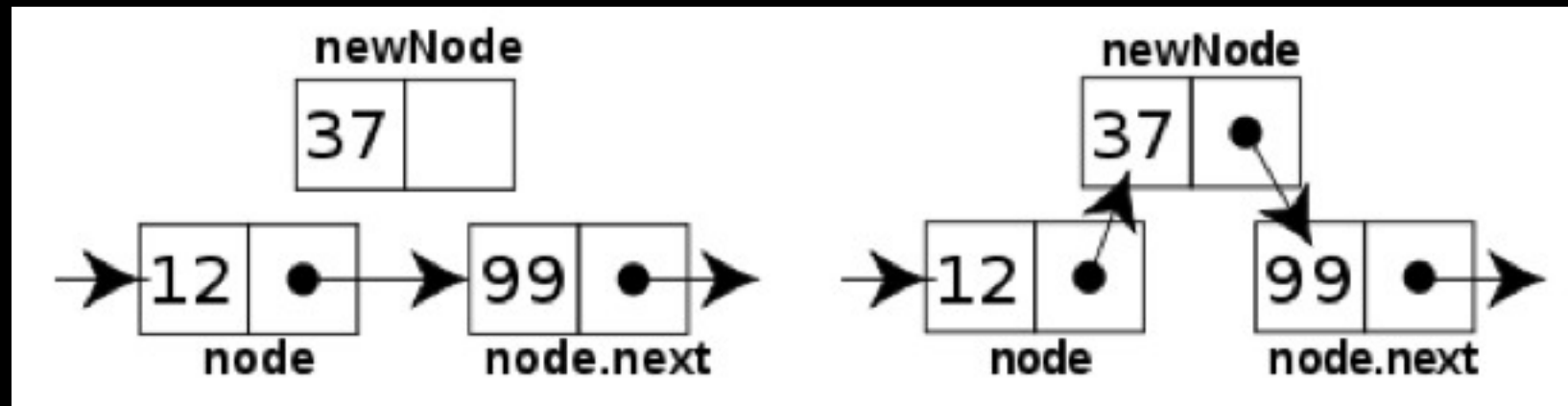
- [ 1 | !null ]
- We'll add a second node, with data=4, after the first

# Second Node

```
node1 = {Node, data: 1, next: ptr_new()}  
node1.next = ptr_new({Node, data:4, next:ptr_new()})
```

- [1 | node2] -> [4 | !null]
- For the next node, we'll loop through all nodes until we find the right place
- How do we know if we've hit the end?
  - if node.next eq !null, there's no next node

# Linked List Insertion



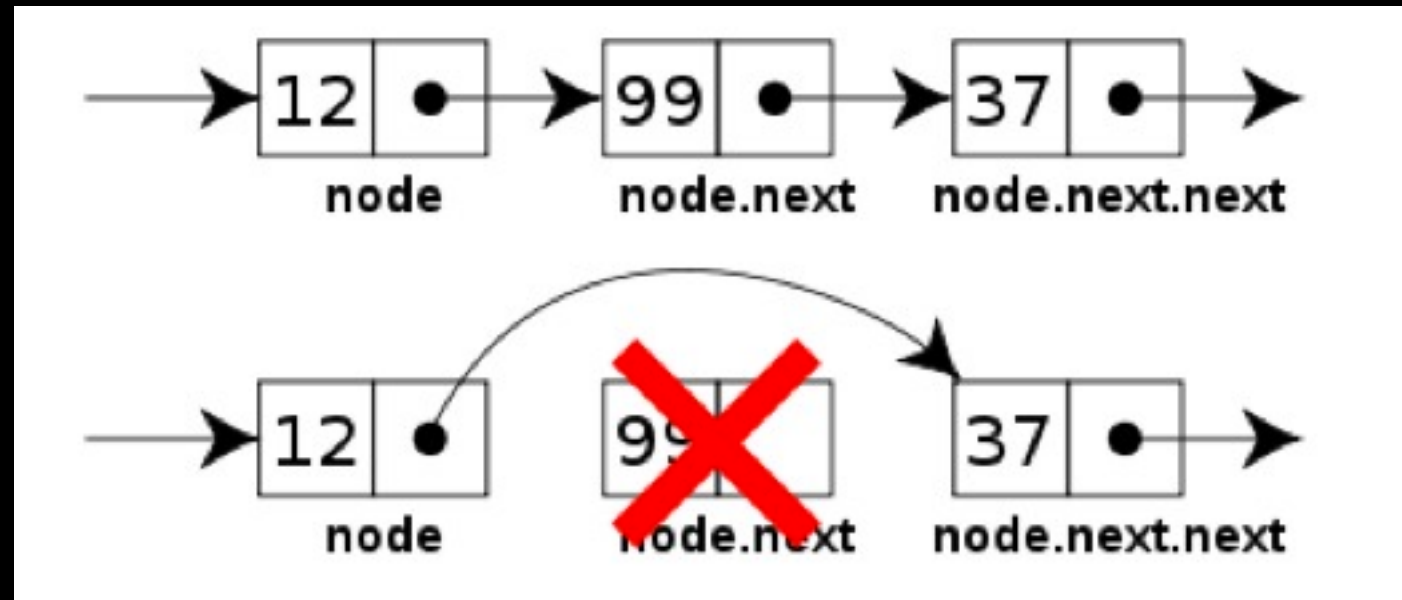
- Actually, really quite difficult to do.
- You need to loop through the linked list, while keeping track of both the previous and next element of the linked list

# Linked List Insertion

- If you have some new data “number” you want to insert into a linked list, you loop through the list until you have the previous & next node in hand...
- Then, make a new node, and stick it in

```
new_node = ptr_new({Node, $  
    data: number, $  
    next: node_ptr})  
(*prev).next = new_node
```

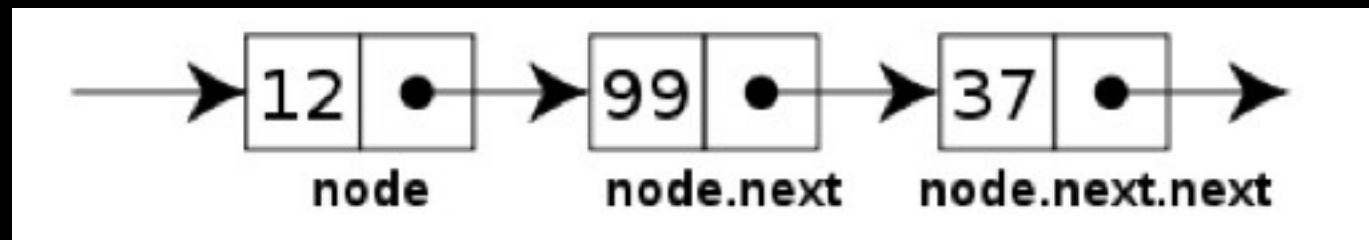
# Linked List Removal



- What if you want to get rid of an element in a linked list?
- You need to “know about” 3 different nodes! You have to have the first point to the third

# Dereferreferreferencing

- Not necessarily too hard to have 3 nodes in hand...
- `(* (* (*node).next).next).data` is pretty darned hard to read, though





# Third Node

```
node1 = {Node, 1, ptr_new()}  
node1.next = ptr_new({Node, 4, ptr_new()})  
node2 = *node1.next  
node3 = ptr_new({Node, 2, ptr_new(node2)})  
node1.next = node3
```

- [ 1 | node3 ] -> [ 2 | node2 ] -> [ 4 | !null ]

# Did it really work?

- [ 1 | node3 ] -> [ 2 | node2 ] -> [ 4 | !null ]

```
IDL> print,node1
{
    1<PtrHeapVar10>}
IDL> print,*node1.next
{
    2<PtrHeapVar9>}
IDL> print,*(*node1.next).next
{
    4<NullPointer>}
```

# Accessing Linked Lists

- How do you store a linked list?
  - For normal variables, you'd do something like:

```
x = fltarr(5)
```

- For linked lists, you store a single pointer to the beginning of the list

# The “Head” of a LL

`[1 | next] -> [2 | next] -> [3 | next] -> !null`

- To store this LL, we need to point to its start

`head = ptr_new()`

- `head` is not a node, but `*head` is

`head->[1 | next]->[2 | next]->[3 | next]->!null`

# Building a LL from scratch

```
pro node__DEFINE
    node = {Node, data: 0, next: ptr_new()}
end
```

```
IDL> first_node = {Node,data:1,next:ptr_new()}
IDL> head = ptr_new(first_node)
IDL> print,head
<PtrHeapVar5>
IDL> print,*head
{
    1<NullPointer>}
IDL> print,first_node
{
    1<NullPointer>}
```



```
IDL> first_node = {Node,data:1,next:ptr_new()}  
print,first_node.next
```

What will print?

- A. % Unable to dereference NULL pointer:  
    <POINTER (<NullPointer>)>.
- B. {           1<NullPointer>}
- C. <NullPointer>
- D. <PtrHeapVar6>
- E. None of the above





What will this do?

```
first_node.next = ptr_new({Node, 7, ptr_new()})
```

- A) Crash
- B) Extend the linked list with a new Node
- C) Make a new node, then lose it
- D) Make a new pointer to nothing
- E) None of the above



# Our linked list now...

```
head->[ 1 | next ]->[ 7 | next ]->!null
```

- Note that `*head` and `first_node` are both the same thing
  - `head` is a pointer to a node
  - `first_node` is a variable of type node

# Looping through Nodes

```
new_data = 2
while (node.data < new_data) and $
    (node.next != null) do $
    node = *(node.next)
```

- Now “node” is our loop variable
- The while loop will quit if either:
  - A. new\_data is greater than or equal to node.data
  - B. we’ve reached the end of the linked list

# Traversing a Linked List

head->[1|next]->[7|next]->!null

node = \*head

head->[1|next]->[7|next]->!null  
node is this LL item

node = \*(node.next)

head->[1|next]->[7|next]->!null  
now node is this LL item

# Inserting [ 2 | ... ]

head->[ 1 | next ]->[ 7 | next ]->!null

node = \*head

head->[ 1 | next ]->[ 7 | next ]->!null

node is this LL item

node.data = 1, so node.data != 2

node = \*(node.next)

head->[ 1 | next ]->[ 7 | next ]->!null

now node is this LL item

node.data = 7, so node.data != 2 is false!

# Inserting [ 2 | ... ]

head->[ 1 | next ]->[ 7 | next ]->!null

node is this LL item

node.data = 7, so node.data != 2 is false!

The while loop quits now.

We can make a new node:

```
new_node = {Node, 2, ptr_new(node)}
```

But where does this leave us?

# Inserting [ 2 | ... ]

[ 2 | next ]  
new\_node →

head → [ 1 | next ] → [ 7 | next ] → !null  
node is this LL item

How do we get the previous node to point to our `new_node`? We should have kept track of where we were before.

# Traversing a Linked List

```
head->[ 1 | next ]->[ 7 | next ]->!null  
node = *head  
next_node = *(node.next)
```

```
head->[ 1 | next ]->[ 7 | next ]->!null  
      node         next_node
```

Now we can do

```
new_node = {Node, 2, ptr_new(next_node)}
```

```
node.next = ptr_new(new_node)
```

to get

```
head->[ 1 | next ]->[ 2 | next ]->[ 7 | next ]->!null  
      node         new_node    next_node
```



# Oh no... that's all wrong still.

- Turns out, if you do  
`node = *head`  
it creates a *copy* of the node head points to
- This is bad - you actually have to use  
`current_ptr = head`  
to copy the *pointer* instead of the *node*
- We'll see this in tutorial....

# Working with nodes

- In tutorial, you'll write a toolkit for linked lists
- Be sure to TEST your results! Make sure you get out what you expect!