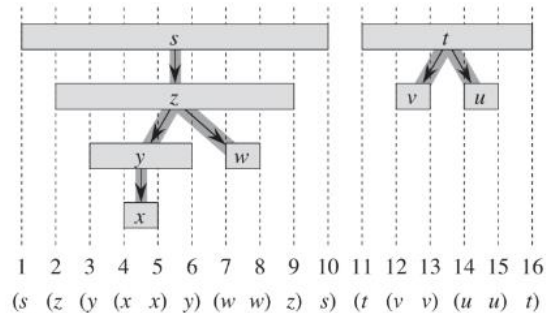
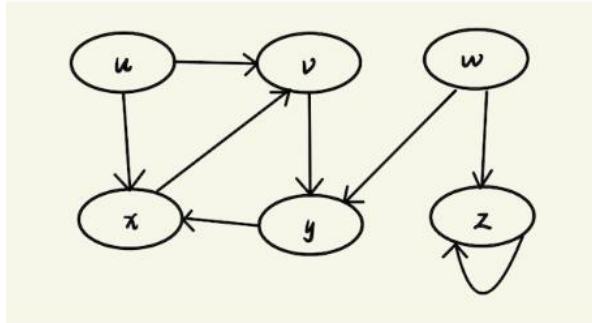


EL9343 Homework 8

Due: Nov. 10th 11:00 a.m.

- What is the running time of DFS if the graph $(G = (V, E))$, where there are $|V|$ nodes and $|E|$ edges) is given as an adjacency list and adjacency matrix? Justify your running time.
- Draw the parenthesis structure of the DFS of bottom left figure (start from u , assume that DFS considers vertices in alphabetical order) and see the example parenthesis structure as is shown in the bottom right.



- Bipartiteness.** Given an undirected graph $G = (V, E)$, it is *bipartite* if there exist U and W such that $U \cup W = V$, $U \cap W = \emptyset$, and every edge has one endpoint each in U and W .
 - Prove: G is bipartite only if G has no odd cycle. (Hint: proof by contradiction)
 - In fact, G is bipartite if and only if G has no odd cycle. Suppose this is given, consider the Algorithm 1 and briefly describe why it is correct.
 - Analyze the time complexity (worst-case, big-O) of the algorithm, in terms of $|V|$ and $|E|$. (Hint: can we check that there is no edge inside any layer in $O(|E|)$? Why?)

Algorithm 1 Testing bipartiteness of graphs

Do BFS starting at any node u

Let $L_0, L_1, L_2, \dots, L_k$ be layers in the resulting breadth-first tree ($L_0 = \{u\}$, $L_i, i = 1, 2, \dots, k$ contains the vertices at distance i from u)

if There is no edge inside any layer L_i **then**

 Declare G to be bipartite, and $U = L_0 \cup L_2 \cup L_4 \cup \dots, W = L_1 \cup L_3 \cup L_5 \cup \dots$ are the bipartition

else

 Declare G to be non-bipartite

end if

- You're helping a group of ethnographers analyze some oral history data they've collected by interviewing members of a village to learn about the lives of people who've lived there over the past two hundred years. From these interviews, they've learned about a set of n people (all of them now deceased), whom we'll denote P_1, P_2, \dots, P_n . They've also collected facts about when these people lived relative to one another. Each fact has one of the following two forms:

- For some i and j , person P_i died before person P_j was born; or
- for some i and j , the life spans of P_i and P_j overlapped at least partially.

Naturally, they're not sure that all these facts are correct; memories are not so good, and a lot of this was passed down by word of mouth. So what they'd like you to determine is whether the data they've collected is at least internally consistent, in the sense that there could have existed a set of people for which all the facts they've learned simultaneously hold.

The original problem asks you to give an efficient algorithm to either produce proposed dates of birth and death for each of the n people so that all the facts hold true, or report (correctly) that no such dates can exist — that is, the facts collected by the ethnographers are not internally consistent. Since topological sort will be taught next week, now let's do make it simple. **You only need to give an efficient algorithm to determine whether the data is internally consistent.**

Hints: Let's say there are totally m recorded facts, and your algorithm should run in $O(m+n)$. It should

be useful to turn the data into a graph. If each node represents an event (birth or death of someone), what is the total number of nodes? How to define the edges, directed or undirected? What makes the data inconsistent? What is the key structure to check in the graph?