COMPUTER ENGINEERING DEPARTMENT

BDA Assignment 2

COURSE: **B.E.** YEAR: **2020-2021** SEMESTER: **VII**

DEPT: Computer Engineering

SUBJECT CODE: CSDLO7032 DATE OF ASSIGNMENT: 08-10-2021

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CLASS: COMPS BE B DATE OF SUBMISSION: 08-10-2021

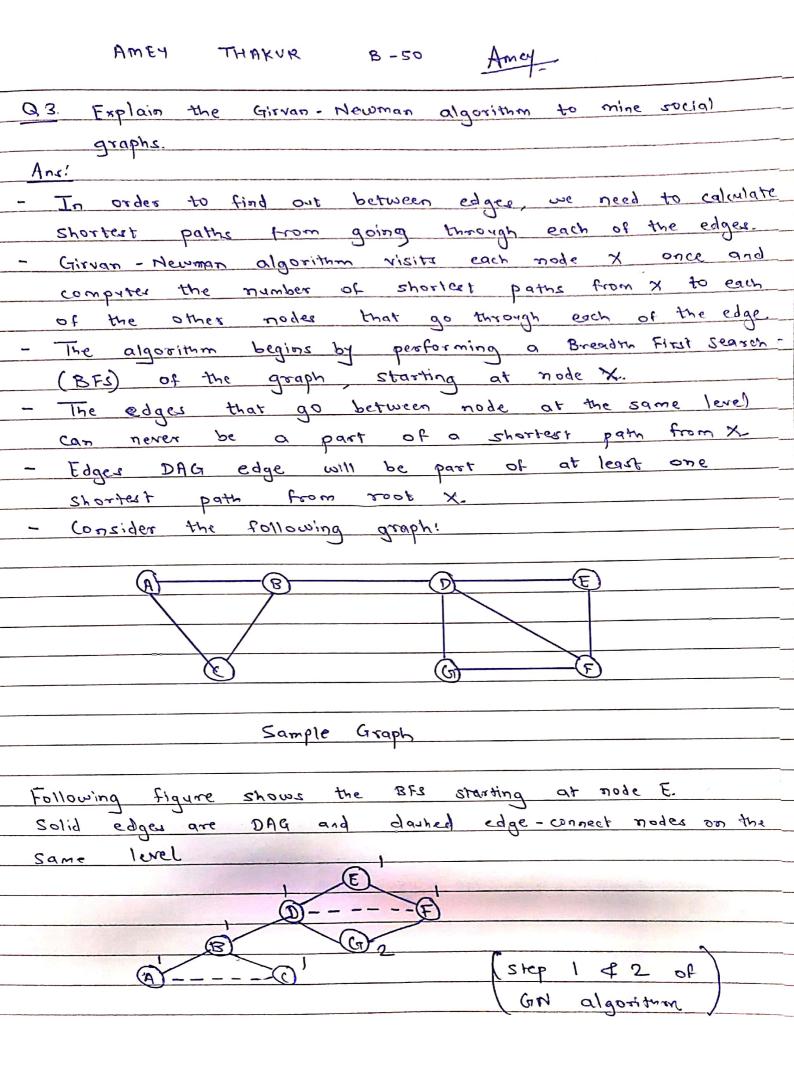
Sr. No.	Questions					
1	What is a Data Stream Management System? Explain with Block Diagram.					
2	Why is finding similar items important in Big Data? Illustrate using two example applications.					
3	Explain the Girvan-Newman algorithm to mine Social Graphs.					

Signature of Student

AMEY THAKUR	B - 50	Amey					
Q1. What is Data s	Stream Manag	ement systems					
Explain with Blo							
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Anu! - A Deme is a co	morter softw	are system	to manage				
continuous data st	TA	of relies si	DBMs, which				
is however design	and for st	atic data in	conventiona)				
	ned tol 22	5010					
- A DSMS also of	the flexible	QUEXY STO CES	ssing so that				
the information n	anded con	be expressed	using queries.				
It executes a conti	711274 QUEST	that is per	manently installed				
- Since most Dams	and data	- griven a	2404011000				
- Since most Dsing	s secules	so long as	new data				
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	Proces						
time	1						
	Limited	Archival					
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	Any number of streams can enter the system
-	Each stream can provide elements at its own schedule:
	they need not have the same data rates or data types
**	and the time between elements of one stream need not
	be uniform.
_	The fact that the rate of arrival of stream elements
	is not under the combrol of the system distinguisher -
	stream processing from the processing of the data -
	that goes on within a database management system,
-	The latter system controls the rate at which data is
	read from the disk and therefore never has to
	worry about data getting lost as it attempts to
	execute queries.
	Streams may be archieved in a large archival store,
	but we assume it is not possible to answer
	queries from the archival store.
-	It could be examined only under special circumstance -
	Using time- consuming retrieval processes
_	There is also a working store into which summaries -
	or barte of streams may be blaced and which can pe
	used for answering queries
-	The working store might be disk . or it might be
	the main memory, depending on how fast we need to
11	
_	But either way, it is of sufficiently limited capacity
	11
	that is cannot store all the data from all the streams.

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<u> </u>	Why	is	finding	similar.	Items	imbarga	ant in	_Big		
	Illus	strate	using	two	example	applica	itions.			
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AMEY THAKOR B-50 Amey The second step of GN algo is to label each node by the number of shortest paths that reach it from the root. Start by labeling the root 1, then, from the top down label each node y by the sum of the labels of its The labeling of nodes is shown in above figure. The final step is to calculate for each edge e the sum over nodes 4 of the fraction of shown paths from the root x to Y then go through e. Each node other than the root & given a credit 1. This credit may be divided among nodes and edges apone The rules for calculation are as follows: - Each node in DAG gets a coedit. - Each non-leaf node gets a credit equal to 1 plus the sum of the credits of the DAG edges from that node to the level A DAG edge e entering node & from the level above is given a share of the credit of 2 proportionals to the fraction of shortest paths from the root to Z that go through e. - After performing the credit calculate with each node as the root, we sum the credits for each edge. edge. - As each shortest path will have been discovered twice we must divide the result (credit for each edge by 2