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FLA Worksheet-1,2

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R2 ; CSE-IT

QD. $S \rightarrow NPVP$
 $NP \rightarrow \text{Det } N \mid \text{'Aajun'}$
 $VP \rightarrow VNP \mid VP PP$
 $PP \rightarrow PNP$
 $\text{Det} \rightarrow \text{'a' } \mid \text{'the'}$
 $N \rightarrow \text{'tiger' } \mid \text{'forest'}$
 $V \rightarrow \text{'hunte'}$
 $P \rightarrow \text{'in'}$

Left recursion:

it occurs when a non-terminal A can derive itself directly or indirectly.

$\rightarrow S \rightarrow NPVP$
 $NP \rightarrow \text{Det } N NP' \mid \text{'Aajun' } NP'$
 $NP' \rightarrow \epsilon \mid NNP'$
 $VP \rightarrow VNPVP' \mid VP'PP$
 $VP' \rightarrow \epsilon \mid NPVP'$
 $PP \rightarrow PNP$
 $\text{Det} \rightarrow \text{'a' } \mid \text{'the'}$
 $N \rightarrow \text{'tiger' } \mid \text{'forest'}$
 $V \rightarrow \text{'hunte'}$
 $P \rightarrow \text{'in'}$

- b) No null productions.
- c) No unit productions.

Q2) CNF:-

$S \rightarrow AB$

$A \rightarrow DN/a$

$B \rightarrow VA/BX$

$X \rightarrow PA$

$D \rightarrow b/c$

$N \rightarrow d/e$

$V \rightarrow f$

$P \rightarrow g$

$NP \rightarrow A, VP \rightarrow B$

$Det \rightarrow D, N \rightarrow N$

${}^6Aajun^d \rightarrow a$

$V \rightarrow V, PP \rightarrow P$

${}^6a^d - b, {}^6the^d - c$

${}^6tiger^d - d, {}^6forest^d - e$

${}^6hunte^d - f, {}^6in^d - g$

• Convert to GNF:-

$A \rightarrow a^d$

$A \rightarrow$ non terminal, $\alpha \rightarrow$ string of non terminal.

$S \rightarrow a_1 a_2$

$NP \rightarrow a_3 a_4 / {}^6Aajun^d$

$VP \rightarrow a_5 a_6 / a_5 a_7$

$PP \rightarrow a_8 a_4$

$Det \rightarrow {}^6a^d / {}^6the^d$

$N \rightarrow {}^6tiger^d / {}^6forest^d$

$V \rightarrow {}^6hunte^d$

$P \rightarrow {}^6in^d$

a) Leftmost Derivation

"while $a < b$ do $c = d$ "

stmt \Rightarrow while cond do stmt,
lm

\Rightarrow while rd relop rd do stmt.
lm

\Rightarrow while $a < b$ do stmt

lm \Rightarrow while $a < b$ do rd = Expr

lm \Rightarrow while $a < b$ do $c = \text{Expr}$.

lm \Rightarrow while $a < b$ do $c = d$

lm \Rightarrow stmt $\xRightarrow{*}$ while ($a < b$ do $c = d$.)
lm

b) Rightmost derivation:-

if $a > b$ then $c = d * e$

stmt \Rightarrow if cond then stmt.
rm

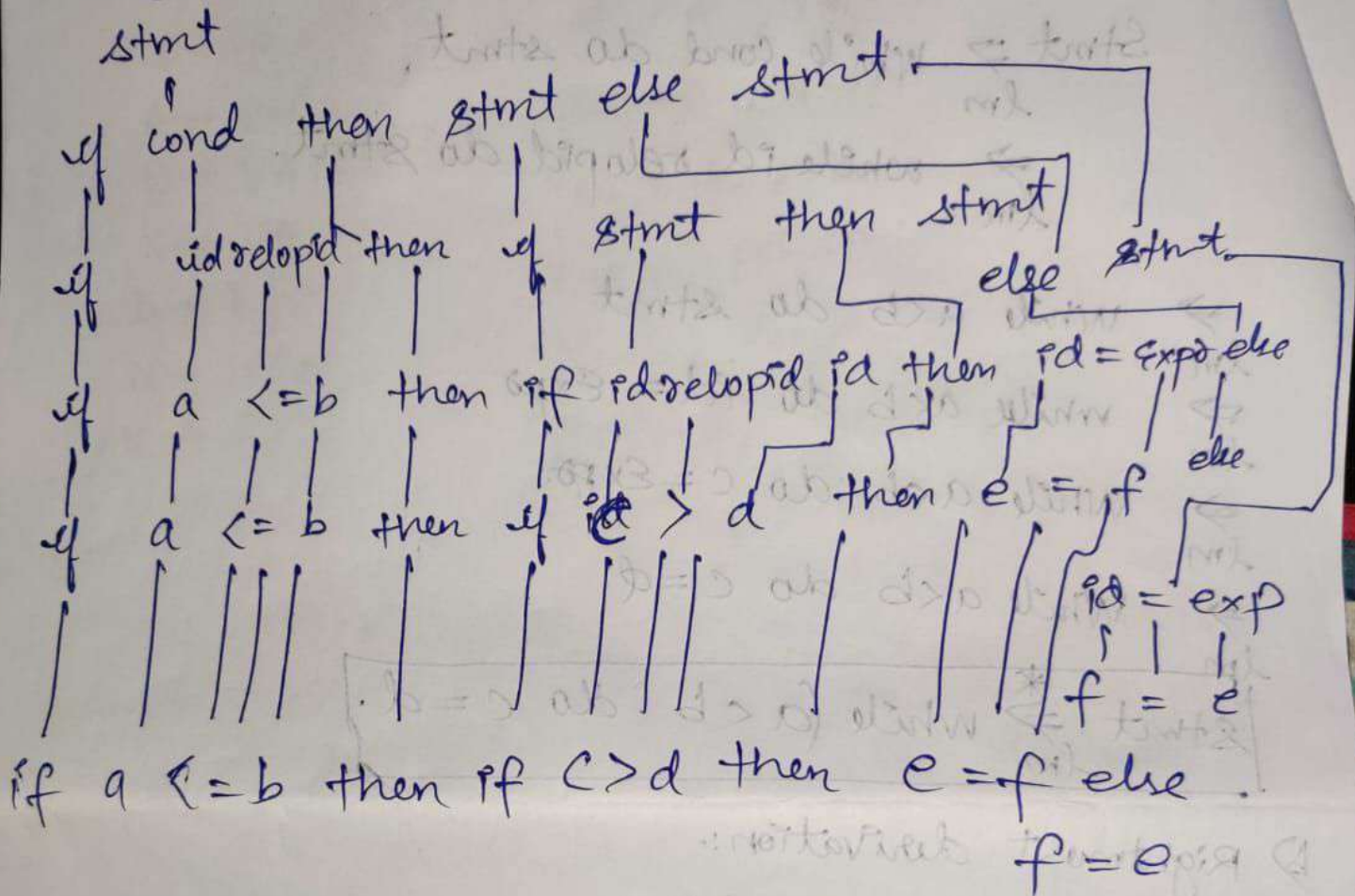
\Rightarrow if cond then rd = Expr.
rm

\Rightarrow if cond then rd = $d * e$
rm

\Rightarrow if cond then $c = d * e$
rm

\Rightarrow stmt $\xRightarrow{*}$ if $a > b$ then $c = d * e$
rm

c) By parse tree.



If two diff. form exist its ambiguous.

Q3)

A) All of these grammars generate languages containing strings composed of \underline{w} & \underline{v} the patterns in these lang. involve combⁿs of \underline{w} & \underline{v} .

$$L = \{wv, wvv, wwvv, wwwvv, \dots\}$$

B) pair of grammars generating the same lang. are:

Grammar 1 $(S \rightarrow wvS/wv)$

Grammar 7 $(S \rightarrow wv; v \rightarrow vS/v)$

Grammar 1 $\rightarrow (S \rightarrow wvS/wv)$
 $wv, wwvv, wvzv, \dots$

Grammar 7 $\rightarrow (S \rightarrow wv; v \rightarrow vS/v)$
 $wv, wwvv, wvzv, \dots$

C) Convert grammar 1 to CNF

$$A \rightarrow BC$$

$$A \rightarrow a$$

Grammar 1 to CNF.

$$S \rightarrow wvS/wv$$

1) $S \rightarrow AS/wv$

2) $A \rightarrow wv$

3) $w \rightarrow w$

4) $v \rightarrow v$

D) Grammar 1 to GNF.

$A \rightarrow a\alpha$.

$S \rightarrow wvS | wv$

1) $S \rightarrow wvS | wv$

2) $S \rightarrow wvB$

3) $B \rightarrow vS | \epsilon$

1) $S \rightarrow wvB$
2) $B \rightarrow vS | \epsilon$

Q4) $S \rightarrow NPVP$

$NP \rightarrow ADNO | DENO | DEADNO | NO$

$VP \rightarrow VVNP$

$DE \rightarrow the | in$

$AD \rightarrow short | tall | green.$

$VV \rightarrow lives$

$NO \rightarrow girl | boy | house.$

production.

$S \rightarrow AB$

$A \rightarrow cD | eD | eCD | A$

$B \rightarrow FA$

$E \rightarrow a | b$

$C \rightarrow e | d | e$

$F \rightarrow f$

$D \rightarrow g | h | i$

1) Terminal $\rightarrow \{a, b, c, d, e, f, g, h, i\}$.

2) Non-terminal $\rightarrow \{S, A, B, C, D, E, F\}$.

3) no the grammar cannot generate
 "The tall boy hides in the house."

4) $S \rightarrow AB$
 $A \rightarrow CD|ED|E|D$
 $B \rightarrow FA$
 $E \rightarrow a|b$
 $C \rightarrow c|d|e$
 $F \rightarrow f$
 $D \rightarrow g|h|i.$

5) $S \rightarrow a_1 a_2$
 $a_1 \rightarrow a_3 a_4 | a_5 a_4 | a_5 | a_4$
 $a_2 \rightarrow a_6 a_7$
 $a_5 \rightarrow a|b$
 $a_3 \rightarrow c|d|e$
 $a_6 \rightarrow f$
 $a_4 \rightarrow g|h|i.$