

### ***C. 1. Recurrence Tree Method – Example 3***

***Solve the following recurrence equation:***

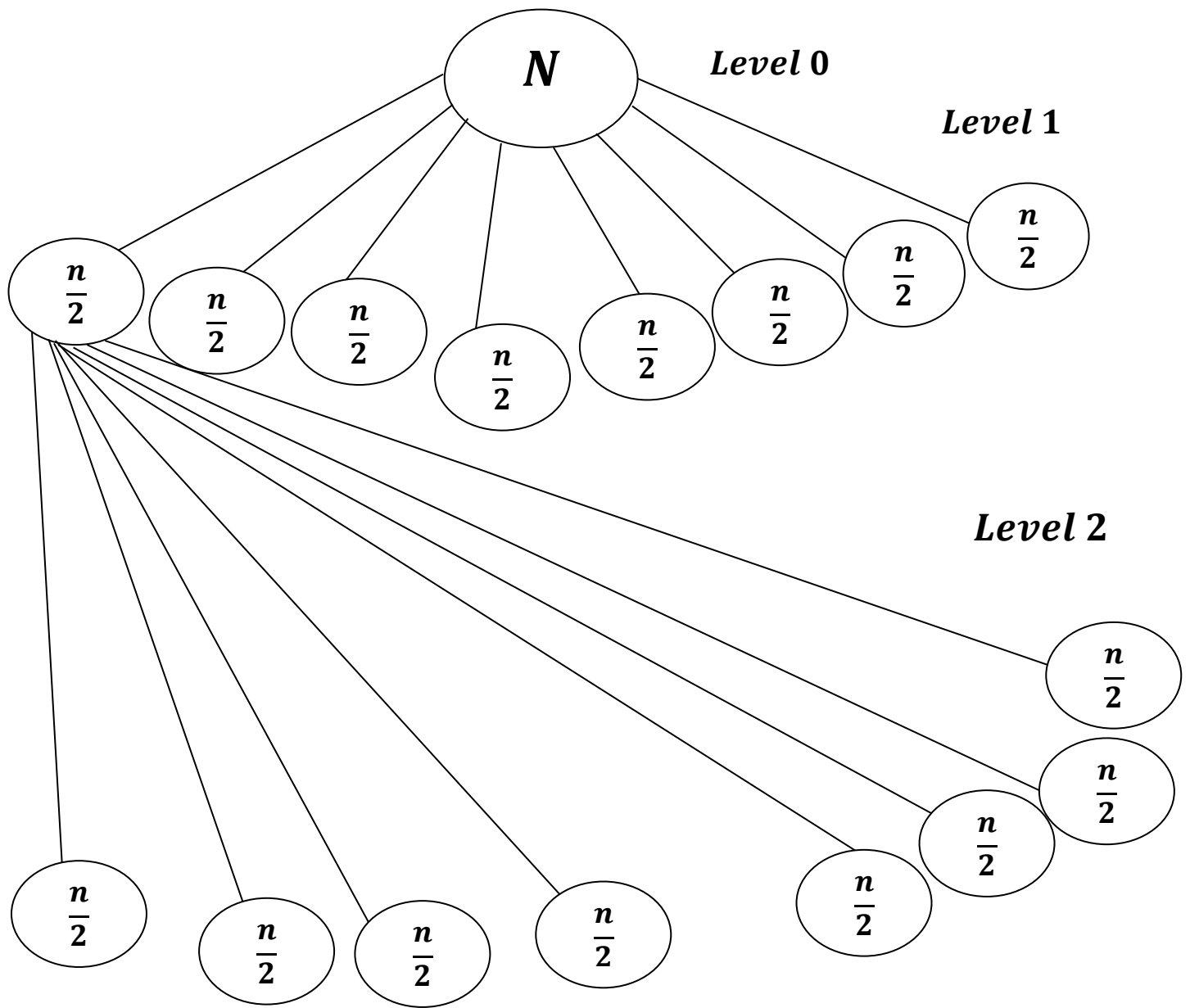
$$t_n = \begin{cases} 1 & \text{for } n = 1 \\ 8T\left(\frac{n}{2}\right) & \text{for } n > 1 \end{cases}$$

***Solution:***

***Initially the input is n. It is sub – divided into eight sub – problems. In the next level, each of these eight sub – problems is again divided into eight sub problems( so a total of  $8 \times 8 = 64$  sub problems).***

***Here 8T represents 8 subdivision and  $\frac{n}{2}$  represents  $\frac{n}{2}$  increase of problem size.***

***This process is continued till a pattern is obtained.***



*Recurrence tree at levels 1 and 2 ( and only one problem division is shown)*

<i>Level</i>	<i>No. of problems</i>	<i>Problem Size</i>	<i>Work done = No. of problems × problem size</i>
<b>0</b>	<b>1</b>	<b><math>n</math></b>	<b><math>1 \times n = n</math></b>
<b>1</b>	<b>8</b>	<b><math>\frac{n}{8}</math></b>	<b><math>8 \times \frac{n}{8} = n</math></b>
<b>2</b>	<b><math>8^2</math></b>	<b><math>\frac{n}{8^2}</math></b>	<b><math>8^2 \times \frac{n}{8^2} = n</math></b>
<b>.</b>	<b>.</b>	<b>.</b>	<b>.</b>
<b>.</b>	<b>.</b>	<b>.</b>	<b>.</b>
<b><math>k</math></b>	<b><math>8^k</math></b>	<b><math>\frac{n}{8^k}</math></b>	<b><math>8^k \times \frac{n}{8^k} = n</math></b>
<b><math>\log_2 n</math></b>	<b><math>8^{\log_2 n}</math></b>	<b>1</b>	<b><math>8^{\log_2 n} \times 1 = n^3</math></b>

*The problem size reduces to 1 as  $T(n) = 1$*

*for  $n = 1$  i.e.  $T(1) = 1$  or  $t_1 = 1$ .*

*It can be observed that at every level a problem is divided into eight subproblems or nodes is increasing in the following pattern: 1, 8, 64, ... ( $8^0, 8^1, 8^2, \dots, 8^i$ ).*

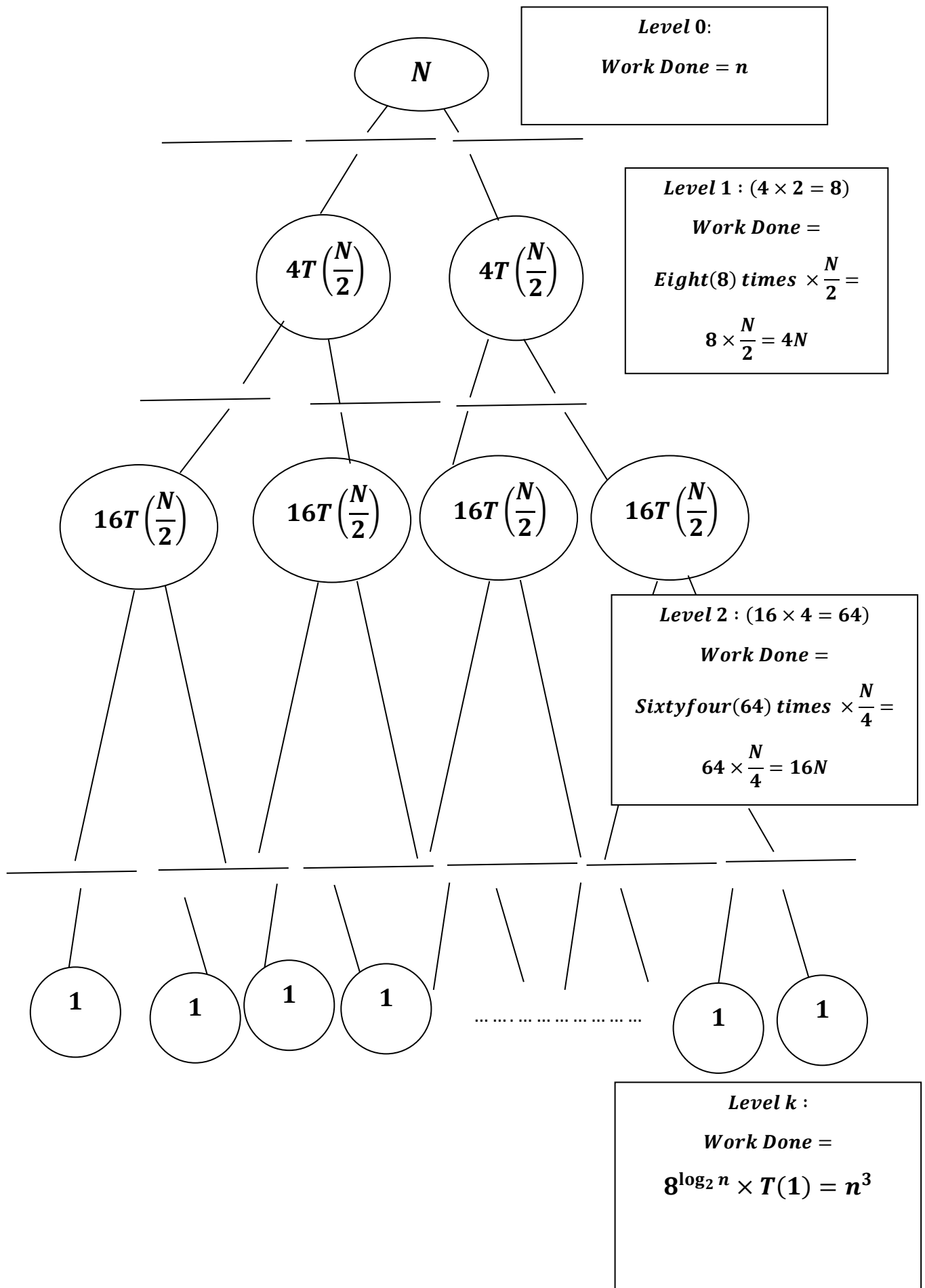
*The problem size is decreasing in a geometric series as follows:  $\left(n, \frac{n}{8}, \frac{n}{8^2}, \dots, \frac{n}{8^k}, \dots, 1\right)$ . Based on the table, the amount of work done at the  $\log_2 n$  can be calculated as follows:*

$$\begin{aligned}
 & 8^{\log_2 n} \times T(1) \\
 &= (2^3)^{\log_2 n} \times 1 \\
 &= (2^{\log_2 n})^3 \\
 &= (n)^3 [a^{\log_a b} = b] \\
 &= n^3
 \end{aligned}$$

*Alternatively,*

$$\begin{aligned}
 & 8^{\log_2 n} \times T(1) \\
 &= (2^3)^{\log_2 n} \times 1 \\
 &= (2^{\log_2 n})^3 \\
 &= (n^{\log_2 2})^3 [a^{\log_a b} = b^{\log_a a} = b] \\
 &= n^3
 \end{aligned}$$

*Now to calculate amount of work done at other level lets divide the above again:*



*Re – writing the above table as:*

<i>Level</i>	<i>No. of problems</i>	<i>Problem Size</i>	<i>Work done = No. of problems × problem size</i>
<b>0</b>	<b>1</b>	<b><math>n</math></b>	<b><math>1 \times n = n</math></b>
<b>1</b>	<b>8</b>	<b><math>\frac{n}{2}</math></b>	<b><math>8 \times \frac{n}{2} = 4n</math></b>
<b>2</b>	<b><math>8^2</math></b>	<b><math>\frac{n}{4}</math></b>	<b><math>8^2 \times \frac{n}{4} = 16n</math></b>
<b>.</b>	<b>.</b>	<b>.</b>	<b>.</b>
<b>.</b>	<b>.</b>	<b>.</b>	<b>.</b>
<b><math>k</math></b>	<b><math>8^k</math></b>	<b><math>\frac{n}{2^k}</math></b>	<b><math>8^k \times \frac{n}{2^k} = n \times 2^{2k}</math></b>
<b><math>\log_2 n</math></b>	<b><math>8^{\log_2 n}</math></b>	<b>1</b>	<b><math>8^{\log_2 n} \times 1 = n^3</math></b>

*Hence at level 0 , work done is  $N$ . At the next level , the cost is 8 times  $\left(\frac{N}{2}\right) = 4N$  ; at level 2 , the cost is 64 times  $\left(\frac{N}{4}\right) = 16N$  and so on.*

*The work done is increasing in the following pattern:  
1, 4, 16, ..... Therefore the total cost is the work done  
at the last level and work done at all other level (0, 1, 2, 3, ...,  
( $\log_2 n - 1$ ).*

*Thus, the total cost of the tree can be estimated as  
follows:*

$$\Rightarrow \sum_{i=0}^{\log_2 n - 1} 4^i n + 8^{\log_2 n} \times T(1)$$

*Where,  $8^{\log_2 n} \times T(1)$  is the last level for  $\log_2 n$ .*

*and for  $\log_2 n - 1$  i.e. till  $\log_2 n - 1$  we have :*  $\sum_{i=0}^{\log_2 n - 1} 4^i n$

*Hence we got :*  $\sum_{i=0}^{\log_2 n - 1} 4^i n + 8^{\log_2 n} \times T(1)$

*And we know :  $8^{\log_2 n} \times T(1) = n^3$ , hence:*

$$\Rightarrow \sum_{i=0}^{\log_2 n - 1} 4^i n + n^3$$

$$\Rightarrow n \times \sum_{i=0}^{\log_2 n - 1} 4^i + n^3$$

*And we know the geometric series :*

$$\sum_{k=1}^n k = 1 + x + x^2 + \dots + x^n = \frac{x^{n+1} - 1}{x - 1}$$

*And hence for  $\sum_{i=0}^{\log_2 n - 1} 4^i$ , we get*

$$\Rightarrow n \times \left( \frac{4^{(\log_2 n - 1) + 1} - 1}{4 - 1} \right) + n^3$$

$$\Rightarrow n \times \left( \frac{4^{\log_2 n} - 1}{4 - 1} \right) + n^3$$

$$\Rightarrow n \times \left( \frac{(n)^2 - 1}{4 - 1} \right) + n^3 \left[ \because 4^{\log_2 n} = (2^{\log_2 n})^2 = n^2 \right]$$

$$\Rightarrow n \times \left( \frac{n^2 - 1}{3} \right) + n^3$$



$$\Rightarrow \frac{n^3 - n}{3} + n^3$$

$$\Rightarrow \frac{n^3 - n + 3n^3}{3}$$

$$\Rightarrow \frac{4n^3 - n}{3}$$

*Hence total cost of the tree as :  $\Theta\left(\frac{4n^3 - n}{3}\right)$*

$$\Rightarrow \Theta\left(\frac{4n^3}{3} - \frac{n}{3}\right)$$

$$\Rightarrow \Theta\left(\frac{4n^3}{3}\right)$$

$$\Rightarrow \frac{4}{3} \times \Theta(n^3)$$

$$\Rightarrow \Theta(n^3)$$

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