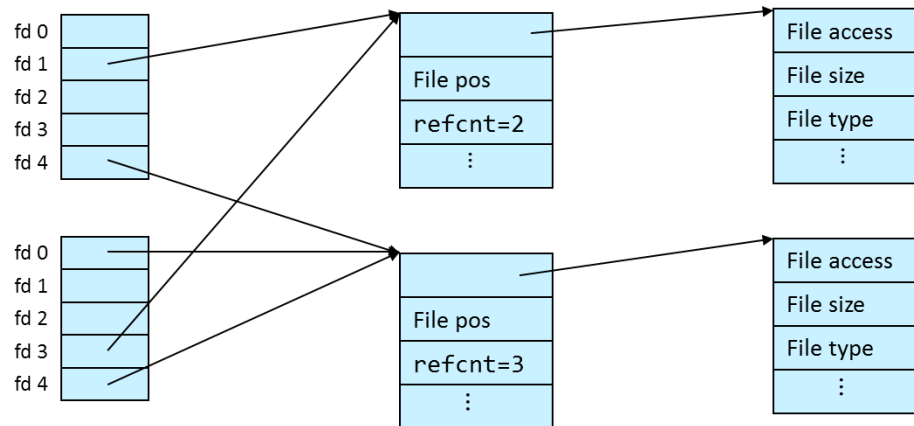


# Input/Output

## Files and Directories



# Class Outline

- **File Metadata**
- **Directories**
- **Kernel File Management**
- **Summary**

```
struct stat {
    dev_t      st_dev;
    ino_t      st_ino;
    mode_t     st_mode;
    nlink_t    st_nlink;
    uid_t      st_uid;
    gid_t      st_gid;
    dev_t      st_rdev;
    off_t      st_size;
    blksize_t  st_blksize;
    blkcnt_t   st_blocks;
    struct timespec st_atim;
    struct timespec st_mtim;
    struct timespec st_ctim;
};
```

# File Metadata

# File Metadata

- Metadata is data about data, in this case information about a file
  - filename
  - file type
  - file size
  - file creation, modification, and access time
  - file access permissions
  - ...
- Per-file metadata maintained by kernel
  - not all metadata is stored in the same place
    - filename: stored in the directory that contains the files
    - everything else is stored in the inode of the file
      - inode: internal storage unit used by Unix file systems

# File Metadata

- File metadata can be accessed with the `*stat*` family of Unix I/O API calls

```
struct stat {
    dev_t          st_dev;          // device
    ino_t          st_ino;         // inode
    mode_t         st_mode;        // protection and file type
    nlink_t        st_nlink;       // number of hard links
    uid_t          st_uid;         // user ID of owner
    gid_t          st_gid;         // group ID of owner
    dev_t          st_rdev;        // device id (if special file)
    off_t          st_size;        // total size, in bytes
    unsigned long  st_blksize;     // preferred blocksize for filesystem
    // I/O
    unsigned long  st_blocks;      // number of 512b blocks allocated
    // high-precision timestamps (precision: nanoseconds; since Linux kernel
    // >=2.6)
    struct timespec st_atim;       // time of last access
    struct timespec st_mtim;       // time of last modification
    struct timespec st_ctim;       // time of last status change
    // low-precision timestamps (precision: seconds; always available)
    time_t         st_atime;       // time of last access
    time_t         st_mtime;       // time of last modification
    time_t         st_ctime;       // time of last status change
};
```

# File Metadata

## ■ struct stat

- manual: stat(2), inode(7)
- st\_mode encodes the type and access permissions of a file
  - ▶ use S\_IS\*\*\*(stat.st\_mode) macros for convenience  
if (S\_ISREG(stat.st\_mode)) // regular file
  - ▶ use S\_IR/W/X\*\*\* masks to check the access permissions  
if (stat.st\_mode & S\_IRUSR) // owner has read permission
- st\_uid / st\_gid
  - ▶ use getpuid(stat.st\_uid) / getgrgid(sb.st\_gid) to retrieve user/group information
- st\_size vs st\_blocks
  - ▶ file is sparse if  $st\_size / 512 > st\_blocks$

# File Metadata

- struct stat
  - timestamps
    - ▶ st\_?tim: nanosecond precision since Linux 2.6
    - ▶ st\_?time: Unix / Linux <2.6 low-resolution timestamps (precision: seconds)
    - ▶ traditionally, Unix filesystems did not record a file's creation date (birth date)
      - newer file systems may support it, retrieve via statx() system call
  - st\_ctim[e] vs st\_mtim[e]
    - ▶ ctim[e]: timestamp of *last update to file inode* (file meta data)
    - ▶ mtim[e]: timestamp of *last update to file data*
    - ▶ in general, mtim[e] ≤ ctim[e]
  - st\_atim[e]
    - ▶ time stamp of *last access to file*
    - ▶ often disabled for performance reasons (mount filesystem with -o noatime)

# File Metadata

- C standard library functions to access file metadata (header `sys/stat.h`)

Operation	API	Remarks
Retrieve file metadata	<code>int stat(const char *pathname, struct stat *statbuf)</code>	
	<code>int lstat(const char *pathname, struct stat *statbuf)</code>	does not follow symbolic links
	<code>int fstat(int fd, struct stat *statbuf)</code>	
	<code>int fstatat(int dirfd, const char *pathname, struct stat *statbuf, int flags)</code>	useful when stat-ing files in a directory
	<code>int statx(int dirfd, const char *pathname, int flags, unsigned int mask, struct statx *statxbuf)</code>	get extended file status



```
| -a
|   ` -b
|       | -c
|       |   ` -d
|       |   ` -f
| -dir1
|   ` -seventyfive
| -dir2
|   | -three
|   ` -seven
| ` -dir3
|     | -five
|     ` -twentytwo
```

# Directories

# Directories

■ C standard library functions to deal with directories (headers `dirent.h`, `sys/types.h`)

Operation	API	Variants
Open	<code>DIR* opendir(const char *name)</code>	<code>fdopendir</code>
Read entry	<code>struct dirent* readdir(DIR *dirp)</code>	
Close	<code>int closedir(DIR *dirp)</code>	
Retrieve descriptor	<code>int dirfd(DIR *dirp)</code>	
Make directory	<code>int mkdir(const char *pathname, mode_t mode)</code>	<code>mkdirat</code>

# Example: Accessing Directories and File Metadata

```
#include <...>

int main(int argc, char *argv[])
{
    DIR *dir = opendir(argc > 1 ? argv[1] : ".");
    if (dir == NULL) { perror("Cannot open directory"); return EXIT_FAILURE; }

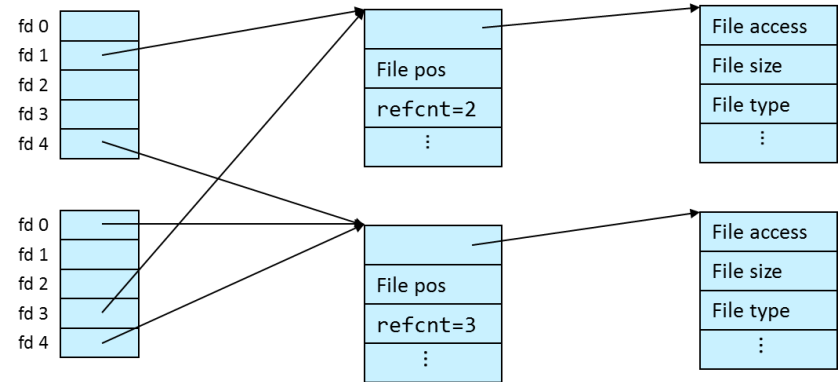
    int dd = dirfd(dir);

    struct dirent *e;
    errno = 0;
    while ((e = readdir(dir)) != NULL) {
        struct stat sb;
        if (fstatat(dd, e->d_name, &sb, 0) < 0) {
            perror("Cannot stat file");
        } else {
            printf("  %-32s %10ld\n", e->d_name, sb.st_size);
        }
        errno = 0;
    }
    if (errno != 0) perror("Cannot enumerate directory");

    closedir(dir);

    return EXIT_SUCCESS;
}
```

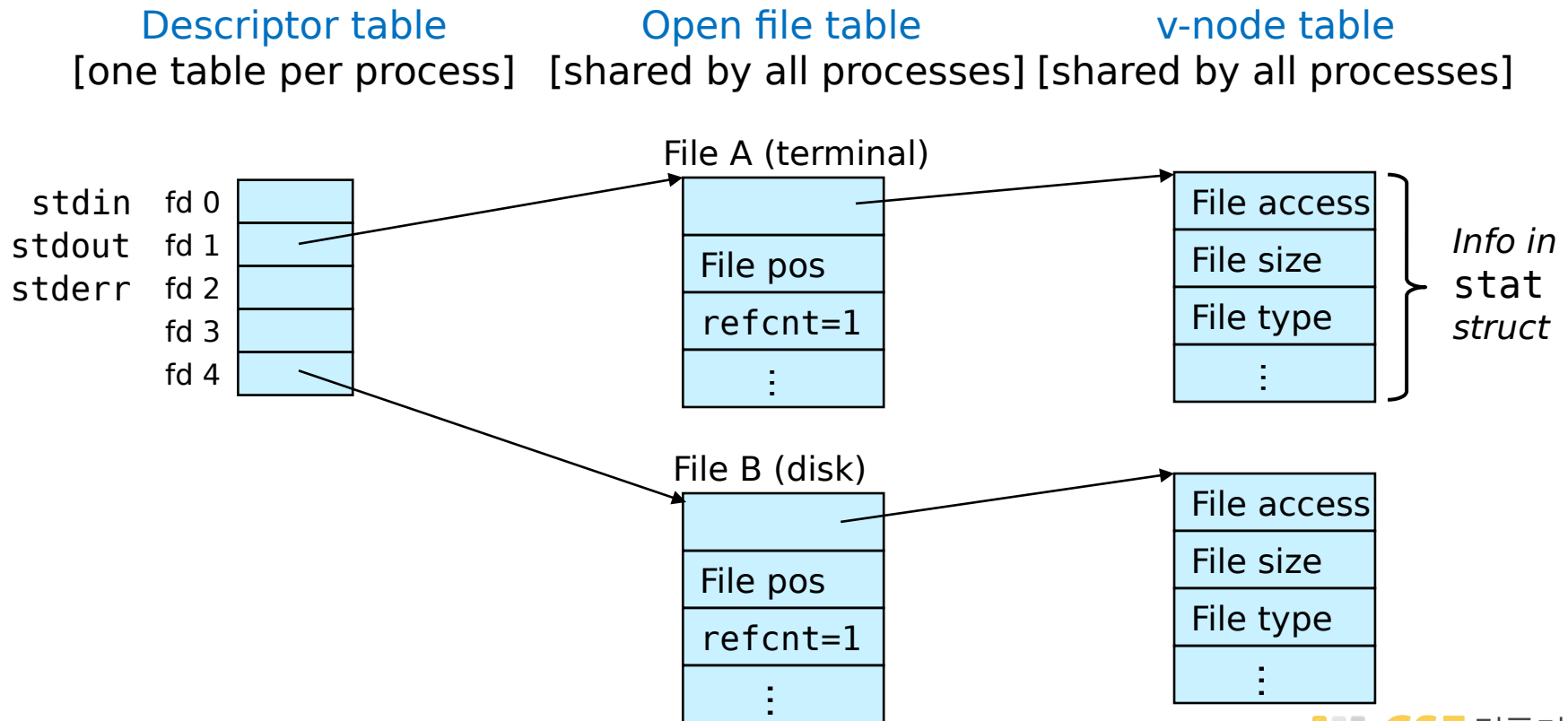
```
$ ./statter /
    opt
4096
    sys
0
    proc
0
    devnull
64
    usr
4096
    lib64
12288
    home
4096
    .
4096
    data
4096
    var
4096
```



# Kernel File Management

# How the Unix Kernel Represents Open Files

- Two descriptors referencing two distinct open disk files. Descriptor 1 (stdout) points to terminal, and descriptor 4 points to open disk file



# File Sharing

- Two distinct descriptors sharing the same disk file through two distinct open file table entries
  - E.g., Calling open twice with the same filename argument

Descriptor table

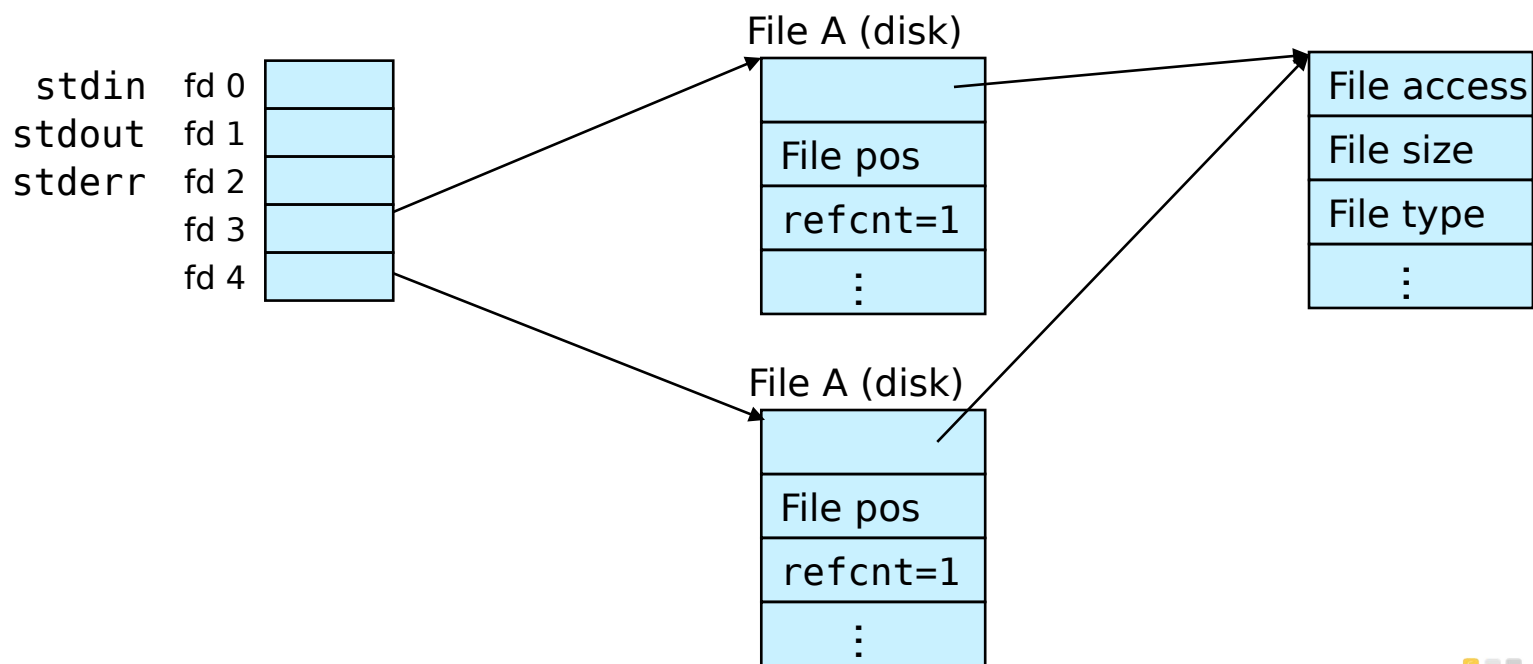
[one table per process]

Open file table

[shared by all processes]

v-node table

[shared by all processes]



# I/O Redirection

- I/O redirection is one of the core concepts of Unix

```
$ ls > output.txt
```

```
$ ls | sort -R
```

```
$ cat < input.txt
```

- How does a shell implement I/O redirection?

# File Descriptor Manipulation

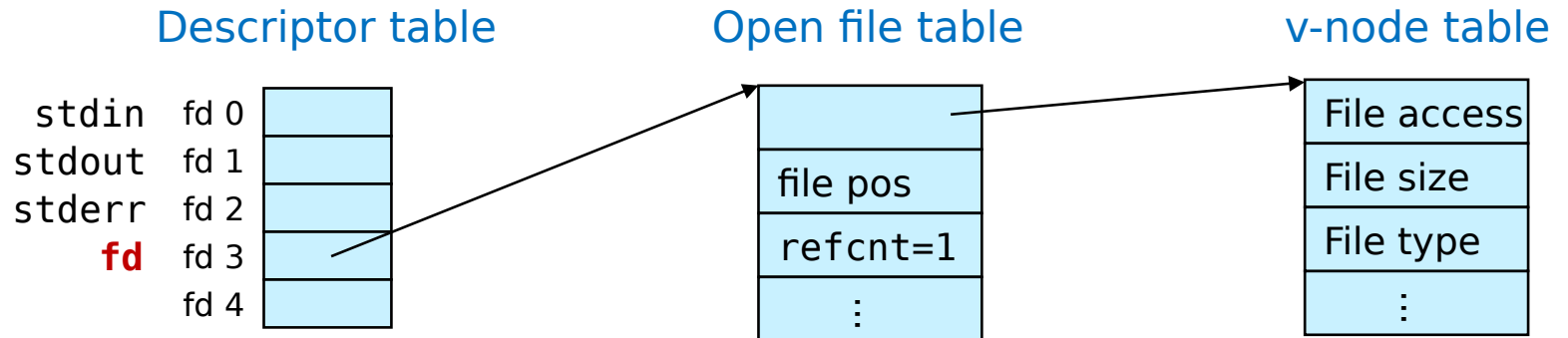
■ C standard library functions allow manipulation of file descriptors (header `unistd.h`)

Operation	API	Remarks
Duplicate file descriptor	<code>int dup(int oldfd)</code>	Returned fd points to same entry in open file table as oldfd
Duplicate file descriptor to specific fd	<code>int dup2(int oldfd, int newfd)</code>	Entry in newfd is overwritten with value in oldfd
Retrieve file descriptor of file stream	<code>int fileno(FILE *stream)</code>	

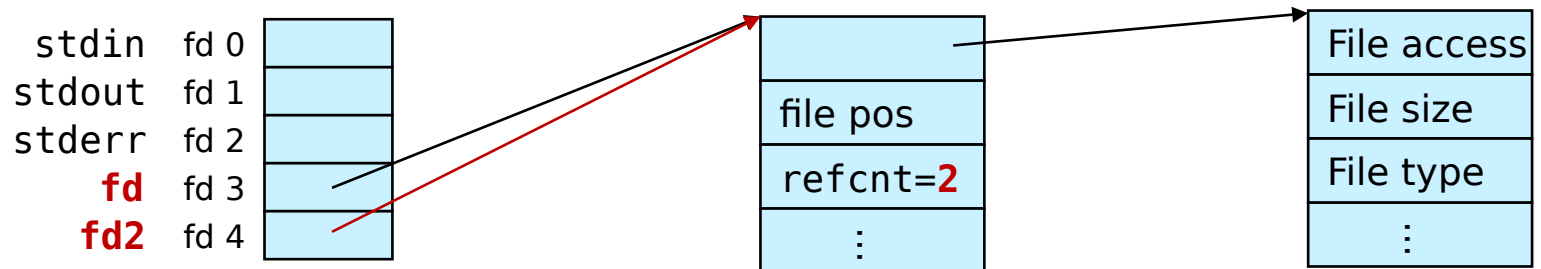


# File Descriptor Manipulation: dup()

■ `fd = open("file.txt", O_RDONLY, 0);`

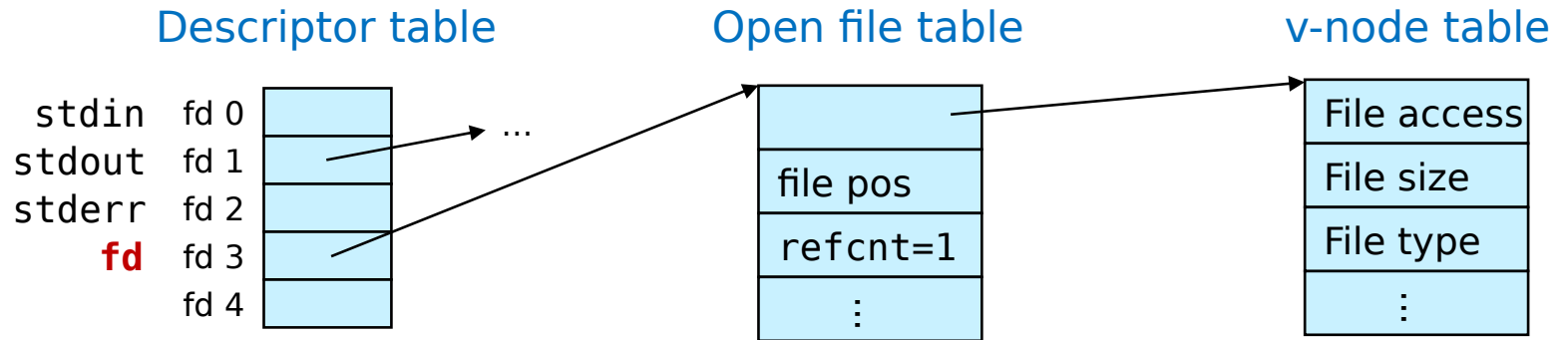


■ `fd2 = dup(fd);`

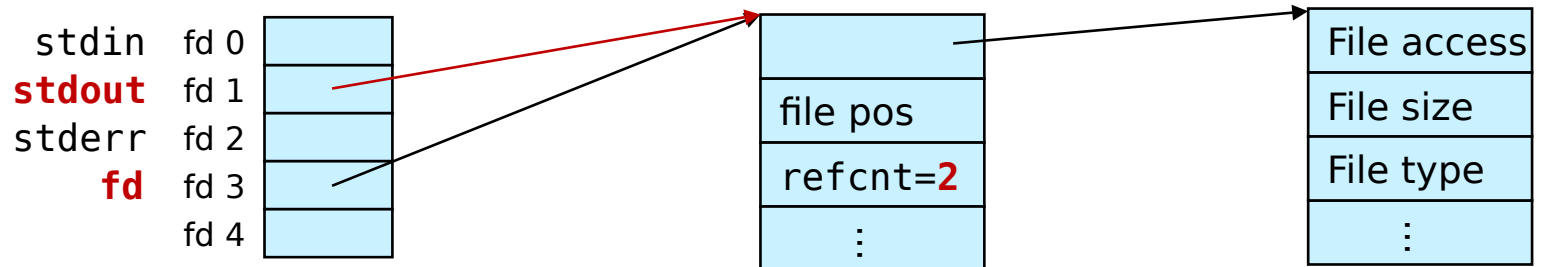


# File Descriptor Manipulation: dup2()

■ `fd = open("output.txt", O_WRONLY, 0);`



■ `dup2(fd, STDOUT_FILENO);`



# Fun with File Descriptors

- What is the output of this program if the input file contains "System Programming"?

```
#include <...>

int main(int argc, char *argv[])
{
    int fd1, fd2, fd3;
    char c1, c2, c3;
    char *fname = argv[1];

    fd1 = open(fname, O_RDONLY, 0);
    fd2 = open(fname, O_RDONLY, 0);
    fd3 = open(fname, O_RDONLY, 0);
    if ((fd1 == -1) || (fd2 == -1) || (fd3 == -1)) {
        fprintf(stderr, "Cannot open input file.\n"); return
EXIT_FAILURE;
    }

    dup2(fd2, fd3);

    read(fd1, &c1, 1);
    read(fd2, &c2, 1);
    read(fd3, &c3, 1);

    printf("c1 = %c, c2 = %c, c3 = %c\n", c1, c2, c3);
    return 0;
}
```

Output: c1 = S, c2 = S, c3 = y

fwfd1.c

# Fun with File Descriptors

- What are the contents of the generated output file?

```
#include <...>

int main(int argc, char *argv[])
{
    int fd1, fd2, fd3;
    char *fname = argv[1];

    if ((fd1 = open(fname, O_CREAT|O_TRUNC|O_RDWR, S_IRUSR|S_IWUSR)) == -1) {
        perror("Cannot open output file"); return EXIT_FAILURE;
    }

    write(fd1, "CSAP", 4);

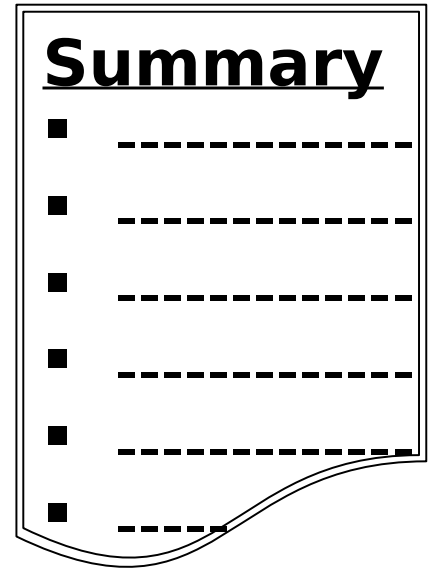
    fd3 = open(fname, O_APPEND|O_WRONLY, 0); File content: CSAPSNU22800
    write(fd3, "M1522", 5);

    fd2 = dup(fd1); // Allocates descriptor
    write(fd2, "SNU", 3);
    write(fd3, "800", 3);

    return 0;
}
```

fwfd2.c

# Summary



# Most Frequently Used Unix I/O System Calls

```
#include <unistd.h>
#include <fcntl.h>
#include <sys/types.h>
#include <sys/stat.h>
```

**\$ man -S 2 <func>**

```
int      open(const char *pathname, int flags[, mode_t mode]);
int      creat(const char *pathname, mode_t mode);
```

```
ssize_t  read(int fd, void *buf, size_t count);
ssize_t  write(int fd, const void *buf, size_t count);
```

```
off_t    lseek(int fd, off_t offset, int whence);
int      stat(const char *path, struct stat *buf);
```

```
int      close(int fd);
```

# Most Frequently Used Standard I/O System Calls

```
#include <stdio.h>
```

```
$ man -S 3 <func>
```

```
FILE*    fopen(const char *pathname, const char *mode);
```

```
size_t   fread(void *ptr, size_t size, size_t nmemb, FILE *stream);
```

```
size_t   fwrite(const void *ptr, size_t size, size_t nmemb, FILE  
*stream);
```

```
int       fflush(FILE *stream);
```

```
int       feof(FILE *stream);
```

```
int       ferror(FILE *stream);
```

```
off_t     fseek(FILE *stream, long offset, int whence);
```

```
int       f[get/set]pos(FILE *stream, fpos_t *pos);
```

```
int       fclose(FILE *fp);
```

# Most Frequently Used API to Manage Directories

```
#include <sys/types.h>
#include <dirent.h>
```

**\$ man -S 3 <func>**

```
DIR*      opendir(const char *name);
```

```
struct dirent *readdir(DIR *dirp);
```

```
int       dirfd(DIR *dirp);
```

```
int       closedir(DIR *dirp);
```



# Pros and Cons of Unix I/O

## ■ Pros

- Unix I/O (or system call) is the most general and the lowest overhead form of I/O (e.g., no buffering)
  - ▶ All other I/O packages are implemented using Unix I/O functions.
- Unix I/O provides functions for accessing file metadata.
- Unix I/O functions are async-signal-safe and can be used safely in signal handlers.

## ■ Cons

- Dealing with short counts is tricky and error prone.
- Efficient reading of text lines requires some form of buffering, also tricky and error prone.
- Both of these issues are addressed by the standard I/O packages.

# Pros and Cons of Standard I/O

## ■ Pros

- Buffering increases efficiency by decreasing the number of read and write system calls
- Short counts are handled automatically

## ■ Cons

- Provides no function for accessing file metadata
- Standard I/O functions are not async-signal-safe, and not appropriate for signal handlers.
- Standard I/O is not appropriate for input and output on network sockets
  - ▶ There are poorly documented restrictions on streams that interact badly with restrictions on sockets (CS:APP3e, Sec 10.9)

# Choosing I/O Functions

- General rule: use the highest-level I/O functions you can
  - Many C programmers are able to do all of their work using the standard I/O functions
- When to use standard I/O
  - When working with disk or terminal files
- When to use raw Unix I/O
  - Inside signal handlers, because Unix I/O is async-signal-safe.
  - In rare cases when you need absolute highest performance.

# Aside: Working with Binary Files

## ■ Binary File Examples

- Object code, Images (JPEG, GIF)

## ■ Functions you shouldn't use on binary files

- Line-oriented I/O such as `fgets`, `scanf`, `printf`
  - ▶ Different systems interpret `0x0A` (`'\n'`) (newline) differently:
    - Linux and Mac OS X: `LF(0x0a)` [`'\n'`]
    - HTTP servers & Windows: `CR+LF(0x0d 0x0a)` [`'\r\n'`]
- String functions
  - ▶ `strlen`, `strcpy`
  - ▶ Interprets byte value 0 (end of string) as special

# For Further Information

## ■ The Unix bible:

- W. Richard Stevens & Stephen A. Rago, Advanced Programming in the Unix Environment, 3rd Edition, Addison-Wesley, 2013, ISBN 978-0321637734
  - ▶ Updated from Stevens's 1993 classic text.

