

Syntax			
x, y, z	Variable	$L ::=$	Type label
l	Value label	L_c	class label
m	Method label	L_a	abstract type label
$v ::=$	Value	$S, T, U, V, W ::=$	Type
x	variable	$p.L$	type selection
$t ::=$	Term	$T \{z \Rightarrow \overline{D}\}$	refinement
v	value	$T \wedge T$	intersection type
$\text{val } x = \text{new } c; t$	new instance	$T \vee T$	union type
$t.l$	field selection	\top	top type
$t.m(t)$	method invocation	\perp	bottom type
$p ::=$	Path	$S_c, T_c ::=$	Concrete type
x	variable	$p.L_c \mid T_c \{z \Rightarrow \overline{D}\} \mid T_c \wedge T_c \mid \top$	
$p.l$	selection	$D ::=$	Declaration
$c ::= T_c \{\overline{d}\}$	Constructor	$L : S..U$	type declaration
$d ::=$	Initialization	$l : T$	value declaration
$l = v$	field initialization	$m : S \rightarrow U$	method declaration
$m(x) = t$	method initialization		
$s ::= \overline{x} \mapsto \overline{e}$	Store	$\Gamma ::= \overline{x} : \overline{T}$	Environment

Reduction		$t \mid s \rightarrow t' \mid s'$
$\frac{y \mapsto T_c \{l = v' \mid \overline{m(x)} = t\} \in s}{y.m_i(v) \mid s \rightarrow [v/x_i]t_i \mid s} \quad (\text{MSEL})$	$\text{val } x = \text{new } c; t \mid s \rightarrow t \mid s, x \mapsto c \quad (\text{NEW})$	
$\frac{y \mapsto T_c \{l = v \mid \overline{m(x)} = t\} \in s}{y.l_i \mid s \rightarrow v_i \mid s} \quad (\text{SEL})$	$\frac{t \mid s \rightarrow t' \mid s'}{e[t] \mid s \rightarrow e[t'] \mid s'} \quad (\text{CONTEXT})$	
where evaluation context		$e ::= [] \mid e.m(t) \mid v.m(e) \mid e.l$
Type Assignment		$\Gamma \vdash t : T$
$\frac{x : T \in \Gamma}{\Gamma \vdash x : T} \quad (\text{VAR})$	$\frac{\Gamma \vdash t \ni l : T'}{\Gamma \vdash t.l : T'} \quad (\text{SEL})$	
$\frac{\Gamma \vdash t \ni m : S \rightarrow T}{\Gamma \vdash t.m(t') : T} \quad (\text{MSEL})$	$\frac{\Gamma \vdash T_c \text{ wfe}, T_c \prec_y L : S..U, \overline{D}}{\Gamma, y : T_c \vdash S <: \overline{U}, \overline{d} : \overline{D}, t' : T'} \quad (\text{NEW})$	
Declaration Assignment		$\Gamma \vdash d : D$
$\frac{\Gamma \vdash v : V', V' <: V}{\Gamma \vdash (l = v) : (l : V)} \quad (\text{VDECL})$	$\frac{\Gamma \vdash S \text{ wfe}}{\Gamma, x : S \vdash t : T', T' <: T} \quad (\text{MDECL})$	
Figure 1. The DOT Calculus : Syntax, Reduction, Type / Declaration Assignment		

Well-formed types		$\Gamma \vdash T \text{ wf}$
$\frac{\Gamma \vdash T \text{ wf}}{\Gamma, z : T \{z \Rightarrow \overline{D}\} \vdash \overline{D} \text{ wf}} \quad (\text{RFN-WF})$	$\Gamma \vdash \top \text{ wf} \quad (\top\text{-WF})$	
$\frac{\Gamma \vdash T \{z \Rightarrow \overline{D}\} \text{ wf}}{\Gamma \vdash T \{z \Rightarrow \overline{D}\} \text{ wf}} \quad (\text{REFL-WF})$	$\Gamma \vdash \perp \text{ wf} \quad (\perp\text{-WF})$	
$\frac{\Gamma \vdash p \ni L : S..U, S \text{ wf}, U \text{ wf}}{\Gamma \vdash p.L \text{ wf}} \quad (\text{TSEL-WF}_1)$	$\frac{\Gamma \vdash p \ni L : \perp..U}{\Gamma \vdash p.L \text{ wf}} \quad (\text{TSEL-WF}_2)$	
$\frac{\Gamma \vdash T \text{ wf}, T' \text{ wf}}{\Gamma \vdash T \wedge T' \text{ wf}} \quad (\wedge\text{-WF})$	$\frac{\Gamma \vdash T \text{ wf}, T' \text{ wf}}{\Gamma \vdash T \vee T' \text{ wf}} \quad (\vee\text{-WF})$	
Well-formed declarations		$\Gamma \vdash D \text{ wf}$
$\frac{\Gamma \vdash S \text{ wf}, U \text{ wf}}{\Gamma \vdash L : S..U \text{ wf}} \quad (\text{TDECL-WF})$	$\frac{\Gamma \vdash S \text{ wf}, U \text{ wf}}{\Gamma \vdash m : S \rightarrow U \text{ wf}} \quad (\text{MDECL-WF})$	
$\frac{\Gamma \vdash T \text{ wf}}{\Gamma \vdash l : T \text{ wf}} \quad (\text{VDECL-WF})$		
Well-formed and expanding types		$\Gamma \vdash T \text{ wfe}$
$\frac{\Gamma \vdash T \text{ wf}, T \prec_z \overline{D}}{\Gamma \vdash T \text{ wfe}} \quad (\text{WFE})$		

Figure 5. The DOT Calculus : Well-Formedness

$\text{dom}(\overline{D} \wedge \overline{D}') = \text{dom}(\overline{D}) \cup \text{dom}(\overline{D}')$	$\text{dom}(\overline{D} \vee \overline{D}') = \text{dom}(\overline{D}) \cap \text{dom}(\overline{D}')$
$(D \wedge D')(L) = L : (S \vee S')..(U \wedge U')$	if $(L : S..U) \in \overline{D}$ and $(L : S'..U') \in \overline{D}'$
$(D \wedge D')(L) = D(L)$	if $L \notin \text{dom}(\overline{D}')$
$(D \wedge D')(L) = D'(L)$	if $L \notin \text{dom}(\overline{D})$
$(D \wedge D')(m) = m : (S \vee S') \rightarrow (U \wedge U')$	if $(m : S \rightarrow U) \in \overline{D}$ and $(m : S' \rightarrow U') \in \overline{D}'$
$(D \wedge D')(m) = D(m)$	if $m \notin \text{dom}(\overline{D}')$
$(D \wedge D')(m) = D'(m)$	if $m \notin \text{dom}(\overline{D})$
$(D \wedge D')(l) = l : T \wedge T'$	if $(l : T) \in \overline{D}$ and $(l : T') \in \overline{D}'$
$(D \wedge D')(l) = D(l)$	if $l \notin \text{dom}(\overline{D}')$
$(D \wedge D')(l) = D'(l)$	if $l \notin \text{dom}(\overline{D})$
$(D \vee D')(L) = L : (S \wedge S')..(U \vee U')$	if $(L : S..U) \in \overline{D}$ and $(L : S'..U') \in \overline{D}'$
$(D \vee D')(m) = m : (S \wedge S') \rightarrow (U \vee U')$	if $(m : S \rightarrow U) \in \overline{D}$ and $(m : S' \rightarrow U') \in \overline{D}'$
$(D \vee D')(l) = l : T \vee T'$	if $(l : T) \in \overline{D}$ and $(l : T') \in \overline{D}'$

Sets of declarations form a lattice with the given meet \wedge and join \vee , the empty set of declarations as the top element, and the bottom element \overline{D}_\perp . Here \overline{D}_\perp is the set of declarations that contains for every term label l the declaration $l : \perp$, for every type label L the declaration $L : \top.. \perp$ and for every method label m the declaration $m : \top \rightarrow \perp$.

Figure 2. The DOT Calculus : Declaration Lattice

Membership		$\Gamma \vdash t \ni D$
$\frac{\Gamma \vdash p : T, T \prec_z \overline{D}}{\Gamma \vdash p \ni [p/z]\overline{D}_i} \quad (\text{PATH-}\ni)$	$\frac{z \notin \text{fn}(D_i) \quad \Gamma \vdash t : T, T \prec_z \overline{D}}{\Gamma \vdash t \ni D_i} \quad (\text{TERM-}\ni)$	
Expansion		$\Gamma \vdash T \prec_z \overline{D}$
$\frac{\Gamma \vdash T \prec_z \overline{D}'}{\Gamma \vdash T \{z \Rightarrow \overline{D}\} \prec_z \overline{D}' \wedge \overline{D}} \quad (\text{RFN-}\prec)$	$\frac{\Gamma \vdash p \ni L : S..U, U \prec_z \overline{D}}{\Gamma \vdash p.L \prec_z \overline{D}} \quad (\text{TSEL-}\prec)$	
$\frac{\Gamma \vdash T_1 \prec_z \overline{D}_1, T_2 \prec_z \overline{D}_2}{\Gamma \vdash T_1 \wedge T_2 \prec_z \overline{D}_1 \wedge \overline{D}_2} \quad (\wedge\text{-}\prec)$	$\frac{\Gamma \vdash T_1 \prec_z \overline{D}_1, T_2 \prec_z \overline{D}_2}{\Gamma \vdash T_1 \vee T_2 \prec_z \overline{D}_1 \vee \overline{D}_2} \quad (\vee\text{-}\prec)$	
$\Gamma \vdash \top \prec_z \{\}$	$\Gamma \vdash \perp \prec_z \overline{D}_\perp$	$(\perp\text{-}\prec)$
Figure 3. The DOT Calculus : Membership and Expansion		

Subtyping		$\Gamma \vdash S <: T$
$\frac{\Gamma \vdash T \text{ wfe}}{\Gamma \vdash T <: \overline{T}} \quad (\text{REFL})$	$\frac{\Gamma \vdash T \text{ wfe}}{\Gamma \vdash \perp <: \overline{T}} \quad (\perp\text{-}\prec)$	
$\frac{\Gamma \vdash T \{z \Rightarrow \overline{D}\} \text{ wfe}, S <: T, T \prec_z \overline{D}'}{\Gamma \vdash S <: T \{z \Rightarrow \overline{D}\}} \quad (\prec\text{-RFN})$	$\frac{\Gamma \vdash T \{z \Rightarrow \overline{D}\} \text{ wfe}, T <: T'}{\Gamma \vdash T \{z \Rightarrow \overline{D}\} <: T'} \quad (\text{RFN-}\prec)$	
$\frac{\Gamma \vdash p \ni L : S..U, S <: U, S' <: S}{\Gamma \vdash S' <: p.L} \quad (\prec\text{-TSEL})$	$\frac{\Gamma \vdash p \ni L : S..U, S <: U, U <: U'}{\Gamma \vdash p.L <: U'} \quad (\text{TSEL-}\prec)$	
$\frac{\Gamma \vdash T <: T_1, T <: T_2}{\Gamma \vdash T <: T_1 \wedge T_2} \quad (\prec\text{-}\wedge)$	$\frac{\Gamma \vdash T_1 <: T, T_2 <: T}{\Gamma \vdash T_1 \vee T_2 <: T} \quad (\vee\text{-}\prec)$	
$\frac{\Gamma \vdash T_2 \text{ wfe}, T <: T_1}{\Gamma \vdash T <: T_1 \vee T_2} \quad (\prec\text{-}\vee_1)$	$\frac{\Gamma \vdash T_1 \text{ wfe}, T <: T_2}{\Gamma \vdash T <: T_1 \wedge T_2} \quad (\wedge_1\text{-}\prec)$	
$\frac{\Gamma \vdash T_1 \text{ wfe}, T <: T_2}{\Gamma \vdash T <: T_1 \vee T_2} \quad (\prec\text{-}\vee_2)$	$\frac{\Gamma \vdash T_1 \wedge T_2 <: T}{\Gamma \vdash T_1 \text{ wfe}, T_2 <: T} \quad (\wedge_2\text{-}\prec)$	
$\frac{\Gamma \vdash T \text{ wfe}}{\Gamma \vdash T <: \overline{T}} \quad (\prec\text{-}\top)$		
Declaration subsumption		$\Gamma \vdash D <: D'$
$\frac{\Gamma \vdash S' <: S, T <: T'}{\Gamma \vdash (L : S..T) <: (L : S'..T')} \quad (\text{TDECL-}\prec)$	$\frac{\Gamma \vdash S' <: S, T <: T'}{\Gamma \vdash (m : S \rightarrow T) <: (m : S' \rightarrow T')} \quad (\text{MDECL-}\prec)$	
$\frac{\Gamma \vdash T <: T'}{\Gamma \vdash (l : T) <: (l : T')} \quad (\text{VDECL-}\prec)$		

Figure 4. The DOT Calculus : Subtyping and Declaration Subsumption