

# C++ Cheatsheet

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## 1 Initial Template

Uncomment line 5-9 if external library is needed.

```
1 #include <bits/stdc++.h>
2 using namespace std;
3 #define FASTio ios::sync_with_stdio(false);cin.tie(
  NULL);
4 #define DECI fixed<<setprecision(5)
5 // #include <ext/pb_ds/assoc_container.hpp>
6 // #include <ext/pb_ds/tree_policy.hpp>
7 // using namespace __gnu_pbds;
8 // typedef tree<int,null_type,less<int>,rb_tree_tag,
  tree_order_statistics_node_update> indexed_set;
9 // typedef tree<int,null_type,less_equal<int>,
  rb_tree_tag,tree_order_statistics_node_update>
  indexed_multiset;
10 typedef long long ll;
11 typedef unsigned long long ull;
12 typedef long double ld;
13 typedef vector<int> vi;
14 typedef vector<vector<int>> vvi;
15 typedef pair<int,int> pii;
16 typedef priority_queue<int> pqi;
17 typedef deque<int> di;
18 #define pb(k) push_back(k)
19 #define mp(a,b) make_pair(a,b)
20 #define B begin();
21 #define E end();
22 #define nl cout<<"\n"
23 #define DB(x) {static int testInt=1000;if((testInt--
  >0)cout<<"(LINE " <<__LINE__ <<": VALUE " <<x<<)"<t"
  ;}
24 #define LB {static int testIntx=0;if(testIntx<1000)
  cout<<"(LINE " <<__LINE__ <<"," <<testIntx+1<<)"<t";
  else break;testIntx++;}
25 #define TA(arr) {int* lLe=(int*)&arr+1;for(int* xTe=
  arr;xTe!=lLe;xTe++) cout<<*xTe<<" ";nl;}
26 #define nax 1000000007
27 /*****/
28
29 int main() {
30     FASTio
31     int t; cin >> t; while(t--> {
32         LB
33     }
34     return 0;
35 }
```

## 2 STL Library

### 2.1 Containers

vector

deque

list

forward\_list

map

unordered\_map

multimap

unordered\_multimap

set

unordered\_set

multiset

unordered\_multiset

stack

queue

priority\_queue

pair

tuple

tree

### 2.2 Algorithms

sort

reverse

max\_element

min\_element

accumulate

count

find

binary\_search

lower\_bound

upper\_bound

next\_permutation

prev\_permutation

partition

stable\_partition

rotate

min

max

swap

\_gcd

\_builtin\_popcount

## 3 Algorithms

### 3.1 Fibonacci numbers

if  $F_n$  is the  $n$ 'th Fibonacci number, where  $F_0 = 0$  and  $F_1 = 1$ , then

$$F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$$

for any  $n, k \in \mathbb{N}$ .

### 3.2 Geometric Transformation of points

Point  $(x, y, z)$  can be transformed by matrix multiplication

$$\begin{bmatrix} x & y & z & 1 \end{bmatrix} \times \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \\ a_{41} & a_{42} & a_{43} & a_{44} \end{bmatrix} = \begin{bmatrix} x' & y' & z' & 1 \end{bmatrix}$$

Where  $(x', y', z')$  is our answer. If we call the  $4 \times 4$  matrix as  $X$ , then for shifting  $x$  by  $a$  co-ordinate,  $y$  by  $b$  and  $z$  by  $c$  co-ordinate,

$$X = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ a & b & c & 1 \end{bmatrix}$$

Instead of shifting, for scaling

$$X = \begin{bmatrix} a & 0 & 0 & 0 \\ 0 & b & 0 & 0 \\ 0 & 0 & c & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

And finally, for rotating  $\theta$  degrees around the  $x$  axis following the right-hand rule (counter-clockwise direction)

$$X = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \cos \theta & -\sin \theta & 0 \\ 0 & \sin \theta & \cos \theta & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

For 2D rotation of  $(x, y)$  by  $\theta$  degree counterclockwise,

$$\begin{bmatrix} x & y \end{bmatrix} \times \begin{bmatrix} \cos \theta & \sin \theta \\ -\sin \theta & \cos \theta \end{bmatrix} = \begin{bmatrix} x' & y' \end{bmatrix}$$

Where  $(x', y')$  is our answer.

### 3.3 Extended Euclidean Algorithm

Returns the gcd of  $a$  and  $b$  with  $ax + by = \text{gcd}(a, b)$ .

```
1 int gcd(int a, int b, int& x, int& y) {
2     if (b==0) {
3         x=1;
4         y=0;
5         return a;
6     }
7     int u, v;
8     int d=gcd(b, a%b, u, v);
9     x=v;
10    y=u-v*(a/b);
11    return d;
12 }
```

### 3.4 Binary Search

Returns the index of  $x$  in array  $a$ .

```
1 int bin_search(int a[], int n, int x) {
2     int l=0, r=n-1;
3     while(l<=r) {
4         int k=(l+r)/2;
5         if(a[k]==x) {
6             return k;
7         }
8         if(a[k]>x) r=k-1;
9         else l=k+1;
10    }
11    return -1;
12 }
```

### 3.5 Processing All Subset

Processes subset of array  $a$ . Initially  $k = 0$  and  $s$  empty.

```
1 void all_subset(int a[], int n, int k, vector<int> s) {
2     if(k==n) {
3         // process subset
4     }
5     else {
6         all_subset(a, n, k+1, s);
7         s.push_back(a[k]);
8         all_subset(a, n, k+1, s);
9         s.pop_back();
10    }
11 }
```

### 3.6 Processing All Permutation

Processes all permutation of array  $a$ , all element should be distinct.  $bm$  is an boolean array with length  $n$ , initially all element is *false*.

```
1 void all_permutation(int a[], int n, vector<int> p, int
2     bm[]) {
3     if(p.size()==n) {
4         // process permutation
5     }
6     else {
7         for(int i=0; i<n; i++) {
8             if(bm[i]) continue;
9             bm[i] = true;
10            p.push_back(a[i]);
11            all_permutation(a, n, p, bm);
12            bm[i] = false;
13            p.pop_back();
14        }
15 }
```

### 3.7 Miller Rabin Primality Test

Checks whether  $n$  prime or not,  $n$  must be in *int* range. Time complexity  $O(\log n)$

```
1 typedef long long ll;
2
3 int binpower(int base, int e, int mod) {
4     int result=1;
5     base%=mod;
6     while(e) {
7         if(e&1)
8             result=(ll)result*base%mod;
9         base=(ll)base*base%mod;
10        e>>=1;
11    }
12    return result;
13 }
14 bool check_composite(int n, int a, int d, int s) {
15     int x=binpower(a, d, n);
16     if(x==1 || x==n-1) return false;
17     for(int r=1; r<s; r++) {
18         x=(ll)x*x%n;
19         if(x ==n-1) return false;
20     }
21     return true;
22 }
```

```

22 };
23 bool MillerRabin(int n) {
24     if (n<2) return false;
25     int r=0;
26     int d=n-1;
27     while((d&1)==0){
28         d>>=1;
29         r++;
30     }
31     for(int a:{2,3,5,7}){
32         if(n==a) return true;
33         if(check_composite(n,a,d,r)) return false;
34     }
35     return true;
36 }

```

### 3.8 DFS algorithm

Runs DFS on a graph with adjacency matrix  $e$ , initially all elements of  $bm$  false.

```

1 void dfs(vector<vector<int>> &e, vector<bool> &bm, int
  v) {
2     bm[v] = true;
3     //process vertex v
4     for (int i: e[v]) {
5         if (!bm[i]) DFS(e,bm,i);
6     }
7 }

```

### 3.9 BFS algorithm

Runs BFS on a graph with adjacency matrix  $e$ , starting point  $s$ , vertex number  $n$ .

```

1 void bfs(vector<vector<int>> e,int n, int s)
2 {
3     vector<bool> bm(n);
4     for(int i=0;i<n;i++) bm[i]=false;
5     queue<int> q;
6     visited[s] = true;
7     q.push(s);
8     while(!q.empty())
9     {
10         s=q.top();
11         //process vertex s
12         q.pop();
13         for (int i: e[s])
14         {
15             if (!bm[i]) {
16                 bm[i] = true;
17                 q.push(i);
18             }
19         }
20     }
21 }

```

### 3.10 modular inverse

Finds modular multiplicative inverse from 1 to  $n$  inclusive mod  $m$ .

```

1 vector<int> mod_inv(int n, int m) {
2     vector<int> inv(n+1);
3     inv[0]=-1; inv[1]=1;
4     for(int i=2; i<=n; i++) inv[i]=m-((m/i)*inv[m%i])%
  m;
5     return inv;
6 }

```

## 4 Useful Results

### 4.1 matrix

The element  $val$  of this struct contains the value of the elements of the matrix, where  $val[i][j]$  represents the value in  $i$ 'th row and  $j$ 'th column.

```

1 struct matrix {
2     vector<vector<int>> val;
3     matrix(int n) {
4         vector<int> temp(n,0);
5         for(int i=0;i<n;i++) val.push_back(temp);
6     }
7     matrix operator+(matrix x) {
8         matrix t_matrix(val.size());
9         for(int i=0;i<val.size();i++) for(int j=0;j<val.
10         size();j++) {
11             t_matrix.val[i][j]=val[i][j]+x.val[i][j];
12         }
13         return t_matrix;
14     }
15     matrix operator-(matrix x) {
16         matrix t_matrix(val.size());
17         for(int i=0;i<val.size();i++) for(int j=0;j<val.
18         size();j++) {
19             t_matrix.val[i][j]=val[i][j]-x.val[i][j];
20         }
21         return t_matrix;
22     }
23     matrix operator*(matrix x) {
24         matrix t_matrix(val.size());
25         for(int i=0;i<val.size();i++) for(int j=0;j<val.
26         size();j++) {
27             int temp=0;
28             for(int k=0;k<val.size();k++) temp+=val[i][k]*(x
29             .val[k][j]);
30             t_matrix.val[i][j]=temp;
31         }
32         return t_matrix;
33     }
34 }

```

### 4.2 Finding directed path with fixed length

Create the adjacency matrix and raise it's power to  $k$ , cell  $(u, v)$  will give the number of distinct path with length  $k$  connecting vertex  $u$  and  $v$  (direction from  $u$  to  $v$ ).