

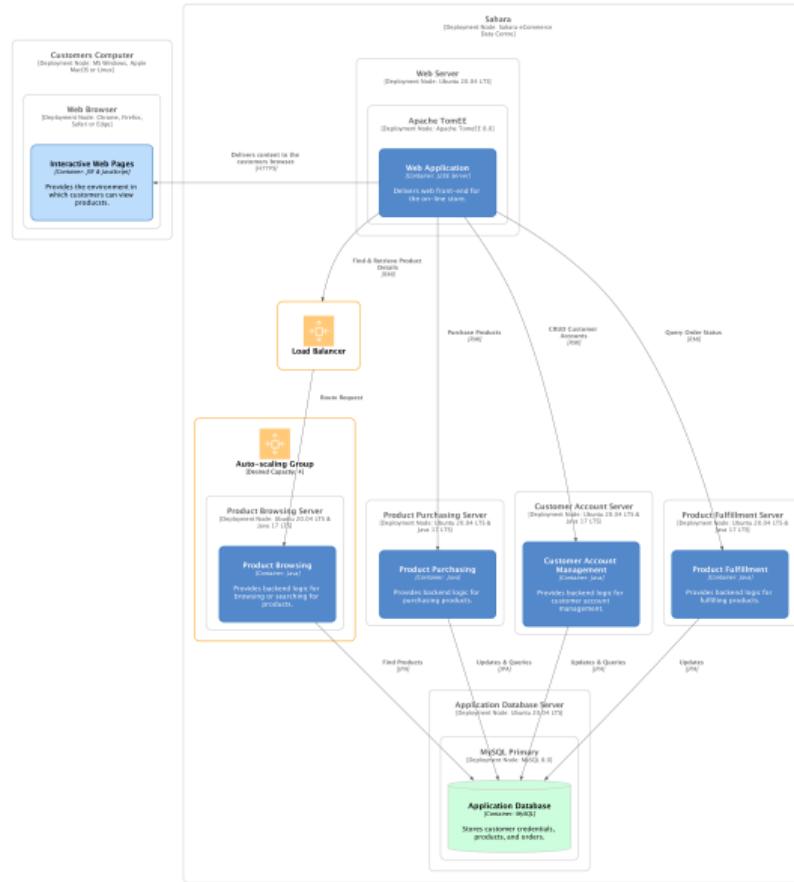
Distributed Computing II

Software Architecture

Brae Webb

April 3, 2023

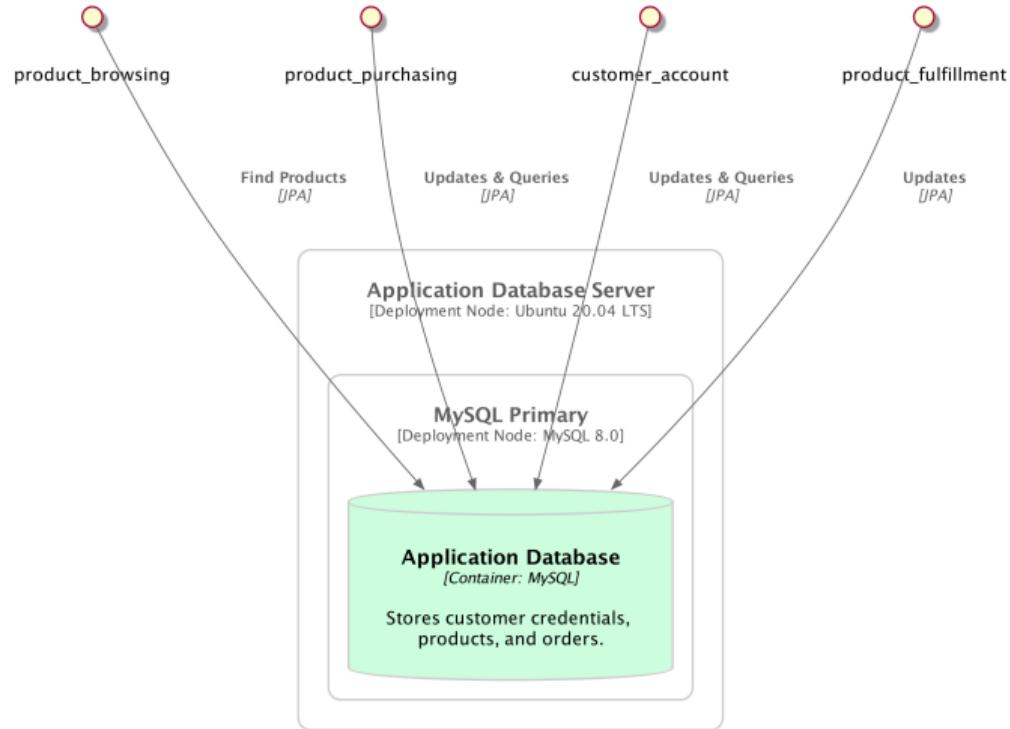
Previously in CSSE6400...



Question

What is the *problem*?

Database



Disclaimer

This is *not* a database course.



Advanced Database Systems (INFS3200)

Course level

Undergraduate

Faculty

Engineering, Architecture & Information Technology

School

Info Tech & Elec Engineering

Units

2

Duration

One Semester

Class contact

2 Lecture hours, 1 Tutorial hour, 1 Practical or Laboratory hour

Incompatible

INFS7907

Prerequisite

INFS2200

Assessment methods

Current course offerings

| Course offerings | Location | Mode | Course Profile |
|------------------|----------|----------|--------------------------------|
| Semester 1, 2022 | St Lucia | Internal | COURSE PROFILE |
| Semester 1, 2022 | External | External | COURSE PROFILE |
| Semester 2, 2022 | External | External | PROFILE UNAVAILABLE |
| Semester 2, 2022 | St Lucia | Internal | PROFILE UNAVAILABLE |

Please Note: Course profiles marked as not available may still be in development.

Course description

Distributed database design, query and transaction processing, data integration, data warehousing, data cleansing, management of spatial data, and data from large scale distributed devices.

Archived offerings

| Course offerings | Location | Mode | Course Profile |
|------------------|----------|-------------------|--------------------------------|
| Semester 1, 2021 | St Lucia | Flexible Delivery | COURSE PROFILE |
| Semester 1, 2021 | External | External | COURSE PROFILE |
| Semester 2, 2021 | External | External | COURSE PROFILE |
| Semester 2, 2021 | St Lucia | Internal | COURSE PROFILE |
| Semester 1, 2020 | St Lucia | Internal | COURSE PROFILE |

Question

How do we fix database scaling issues?

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Answer

- Replication

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- Replication
- Partitioning

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- Replication
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- Independent databases

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Question

What is *replication*?

Definition 1. Replication

Data copied across multiple different machines.



| product_id | name | stock | price |
|------------|---------------------------------|-------|---------|
| 1234 | Nicholas Cage Reversible Pillow | 10 | \$10.00 |
| 4321 | Lifelike Elephant Inflatable | 5 | \$50.00 |



| product_id | name | stock | price |
|------------|---------------------------------|-------|---------|
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Definition 2. Replica

Database node which stores a copy of the data.

Question

What are the advantages of *replication*?

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Answer

- Scale out our database to cope with *load*.

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- Provide *fault tolerance* from a single database instance failure.

Question

What are the advantages of *replication*?

Answer

- Scale out our database to cope with *load*.
- Provide *fault tolerance* from a single database instance failure.
- Locate databases *closer to end-users*.

Question

How do we replicate our data?

First approach

Leader-follower Replication



Leader-based Replication

On write Writes sent to leader, change is propagated via change stream.

Leader-based Replication

On write Writes sent to leader, change is propagated via change stream.

On read Any replica can be queried.



Propogating changes

Synchronous vs. Asynchronous

ProductBrowsing

Leader

Follower1

1 UPDATE products SET stock=4

2 UPDATE

3 OK

4 OK

ProductBrowsing

Leader

Follower1

ProductBrowsing

Leader

Follower1

1 UPDATE products SET stock=4

2 UPDATE

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4 OK

ProductBrowsing

Leader

Follower1



Synchronous propagation

- Writes must propagate to *all followers* before being successful.

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- *Any* replica goes down, *all* replicas are un-writable.

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- Writes must propagate to *all followers* before being successful.
- *Any* replica goes down, *all* replicas are un-writable.
- Writes must *wait* for propagation to all replicas.

Asynchronous propagation

- Writes *don't* have to *wait* for propagation.

Asynchronous propagation

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- If the leader goes down before propagating, the *write is lost*.

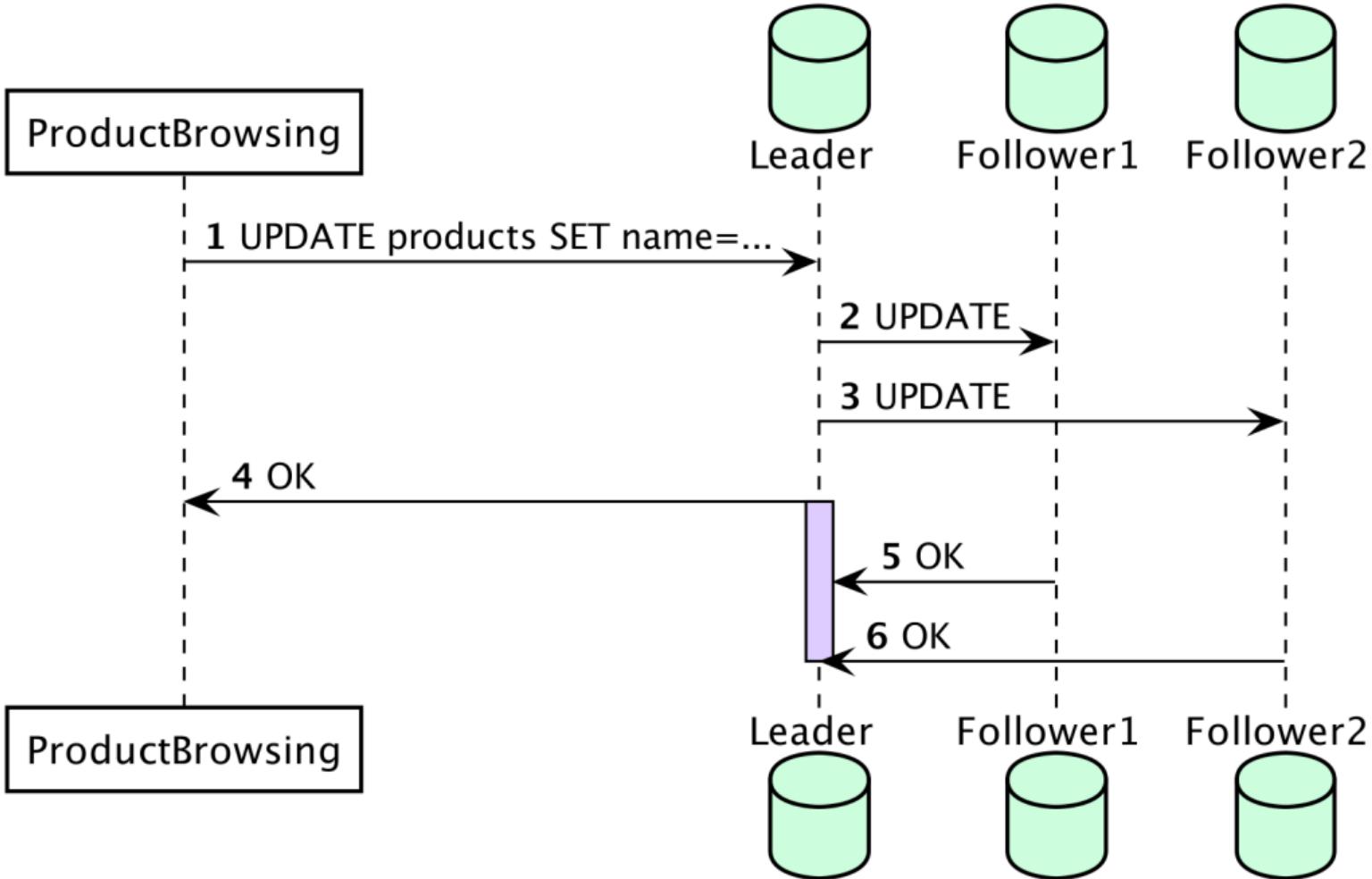
Asynchronous propagation

- Writes *don't* have to *wait* for propagation.
- If the leader goes down before propagating, the *write is lost*.
- Replicas can have out-dated or *stale* data.

Definition 3. Replication Lag

The time taken for replicas to update *stale* data.





Eventually, all replicas must become consistent

The system is *eventually consistent*

Eventual Consistency

Problems?



Brae Webb
@braewebb



Brae Webb

@braewebb

Name: Brae

Cancel

Save



Brae Webb

@braewebb

Name: Brae

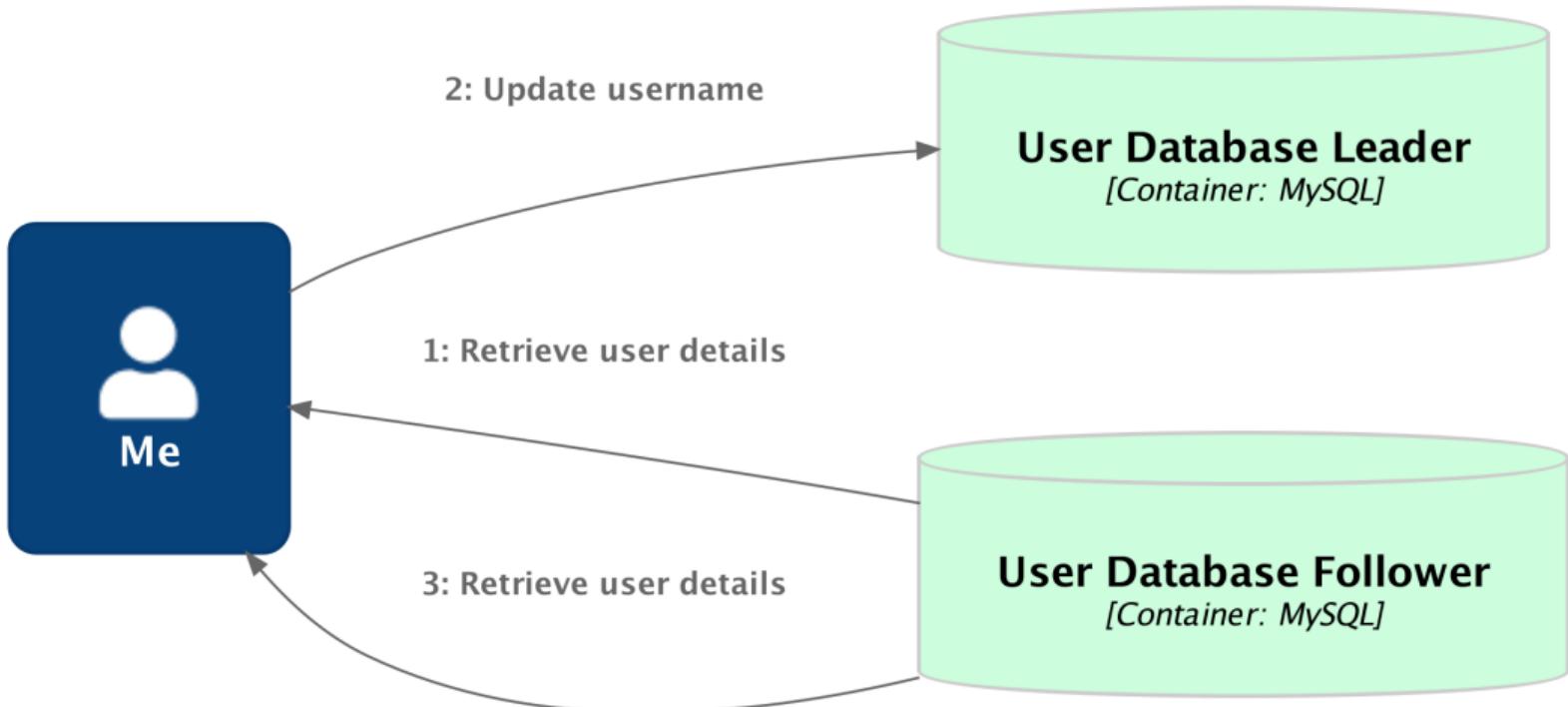
Cancel

Save



Brae Webb

@braewebb



Definition 4. Read-your-writes Consistency

Users always see the updates that *they have made*.



Brae Webb
@braewebb

My fist post



Brae Webb

@braewebb

My fist post



Brae Webb

@braewebb

My first post



Brae Webb

@braewebb

My fist post



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Definition 5. Monotonic Reads

Once a user reads an updated value, they don't later see the old value.

Summary

- Leader-follower databases allow *reads to scale* more effectively.
- Asynchronous propagation weakens consistency to *eventually consistent*.
- Leader-follower databases still have a *leader write bottle-neck*.

Second approach

Multi-leader Replication



Why multi-leader?

- If you have multiple leaders, you can write to any, allowing *writes to scale*.

Why multi-leader?

- If you have multiple leaders, you can write to any, allowing *writes to scale*.
- A leader going down doesn't prevent writes, giving *better fault-tolerance*.

Question

What might go wrong?

Question

What might go wrong?

Answer

Write conflicts



Where possible

Avoid write conflicts



Where impossible

Convergence

Convergence Strategies

- Assign each *write* a unique ID.

Convergence Strategies

- Assign each *write* a unique ID.
- Assign each *leader replica* a unique ID.

Convergence Strategies

- Assign each *write* a unique ID.
- Assign each *leader replica* a unique ID.
- Custom resolution logic.



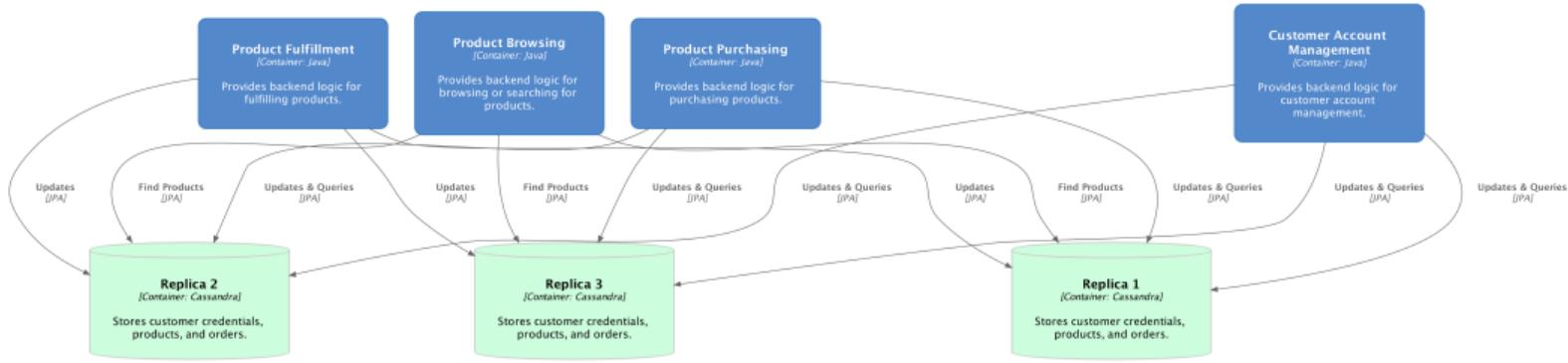
Resolving Conflicts

On Write When a conflict is first noticed, take proactive resolution action.

On Read When a conflict is next read, ask for a resolution.

Third Approach

Leaderless Replication



How do they work?

Each read/write is sent to *multiple* replicas.





How are changes propagated?

- Read Repair

How are changes propagated?

- Read Repair
- Anti-entropy Process

Question

How do we know it's consistent?



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Answer

Quorum Reads and Writes

Quorum Consistency

$$w + r > n$$

n total replicas

w amount of replicas to *write* to

r amount of replicas to *read* from



Replica 1

[Container: Cassandra]



Replica 2

[Container: Cassandra]



Replica 3

[Container: Cassandra]

n total replicas

w amount of replicas to *write* to

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Summary

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- *Replication* copies data to multiple replicas.
- *Leader-based* replication is most common and simplest.
- Replication introduces *eventual consistency*.
- *Multi-leader* replication scales writes as well as reads but introduces *write conflicts*.
- *Leaderless* replication is another approach which keeps the problems of multi-leader.

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- *Replication*
- Partitioning
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Definition 6. Partitioning

Split the data of a system onto multiple nodes, these nodes are *partitions*.



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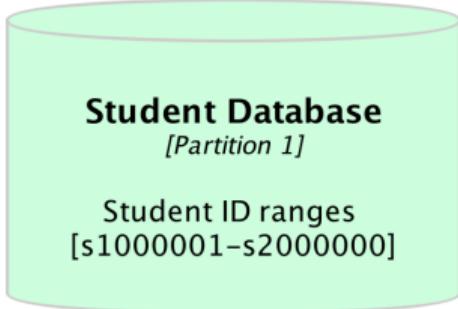
| product_id | name | stock | price |
|------------|---------------------------------|-------|---------|
| 1234 | Nicholas Cage Reversible Pillow | 10 | \$10.00 |

Question

How should we decide which data is stored where?



| student_id | name | ... |
|------------|--------------|-----|
| s0746283 | Bobby Tables | ... |
| ... | ... | ... |



| student_id | name | ... |
|------------|-----------|-----|
| s1637285 | Brae Webb | ... |
| ... | ... | ... |

Question

What is the problem with this?

Question

What is the problem with this?

Answer

Over time some partitions become inactive, while others receive almost all load.

Question

How should we decide which data is stored where?

Question

How should we decide which data is stored where?

Answer

Maximize spread of requests, avoiding *skewing*.

Question

Have we seen this before?

Question

Have we seen this before?

Answer

Hashing?

Question

What is the problem with this?

Question

What is the problem with this?

Answer

Range queries are inefficient, i.e. get all students between s4444444 and s4565656

Question

How do we route queries?

Query-insensitive Load Balancer

Randomly route to any node, responsibility of the node to re-route to the correct node.



Query-sensitive Load Balancer

A load balancer which understands which queries should be forwarded to which node.



Client-aware Queries

Place the responsibility on clients to choose the correct node.



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- A *consistent method* to chose which node is required.
- Partitioning by *primary key* can create *skewing*.
- Partitioning by *hash* makes range queries less efficient.
- Three approaches to *routing requests*.

Disclaimer

We have ignored the hard parts of replication.

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- *Partitioning*
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- *Independent databases*

Distributed state creates a lot of *complexity*

When programmers are faced with complexity

They create abstractions

One key database abstraction

Transactions

Definition 7. Transaction

A group of operations performed as if they were one.

ACID

A tomic

C onsistent

I solated

D urable

The pushback

NoSQL and microservice architectures pushed back against transactions.

For more on transactions



Sept 25-26, 2015
thestrangeloop.com



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 - Leader-based, multi-leader, and leaderless

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- Transactions
 - ACID properties

Summary

- Replications
 - Leader-based, multi-leader, and leaderless
 - Eventual consistency
 - Write conflicts
- Partitioning
 - Consistent method to pick nodes for data
 - Avoiding skewing
- Transactions
 - ACID properties
 - Pushback causing headaches