

Triangle-Free Graphs

Ramsey Numbers

David Stalfa & Adam Plumer

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Every size n graph has

- an x clique, or
- y independent vertices

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Every size n graph has

- an x clique, or
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An (x,y) graph = a graph with

- no x -clique
- no y independent vertices

$R(x,y) \leq n$ iff there is no (x,y) graph of size n

$(3,k)$ = set of triangle free graphs

$$R(3,4) \leq 9$$

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Proof by contradiction

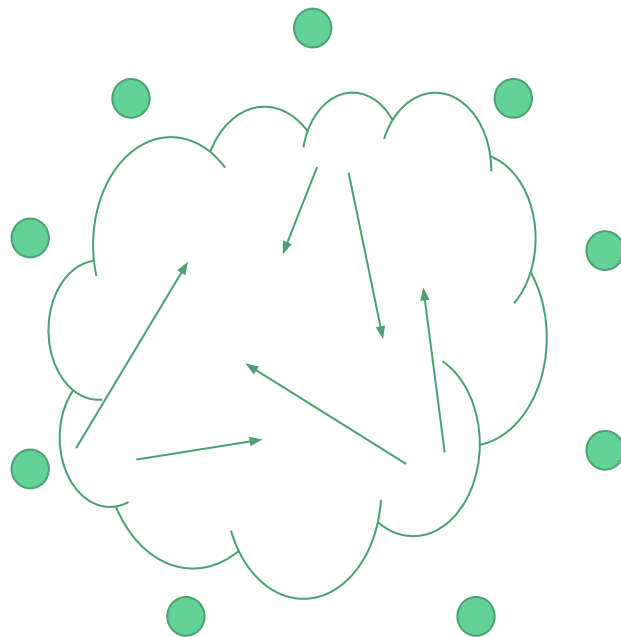
Suppose there exists a (3,4) graph G of size 9

- G has a certain structure
- removing vertices from G would result in a graph from which it is impossible to reconstruct G

Therefore, there is no (3,4) graph of size 9

$$R(3,4) \leq 9$$

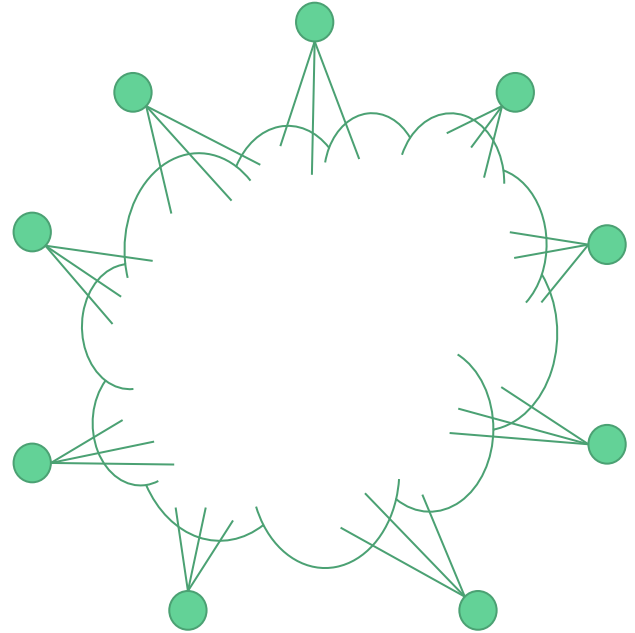
Suppose there is some (3,4) graph G of size 9



$$R(3,4) \leq 9$$

G is regular degree 3

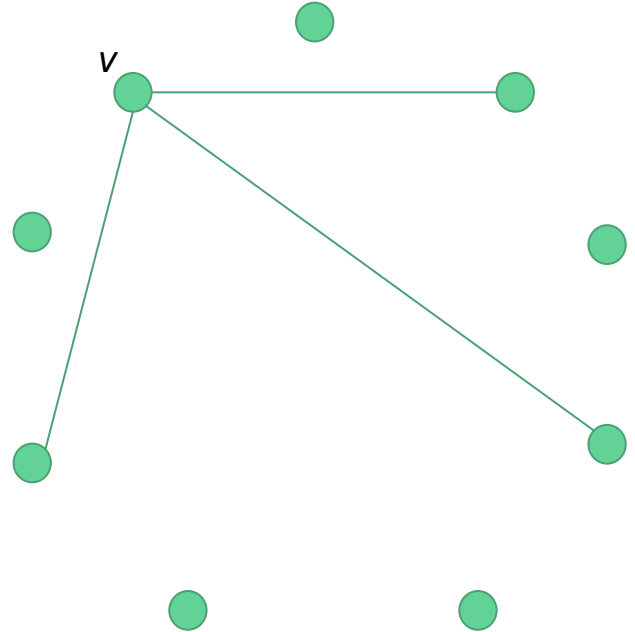
Max degree of G = min degree of G = 3



$$R(3,4) \leq 9$$

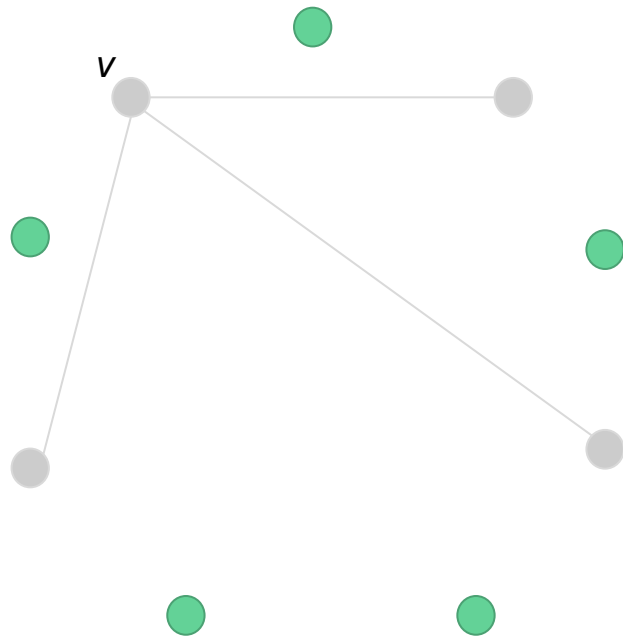
G is regular degree 3

Remove one vertex and its neighborhood



$$R(3,4) \leq 9$$

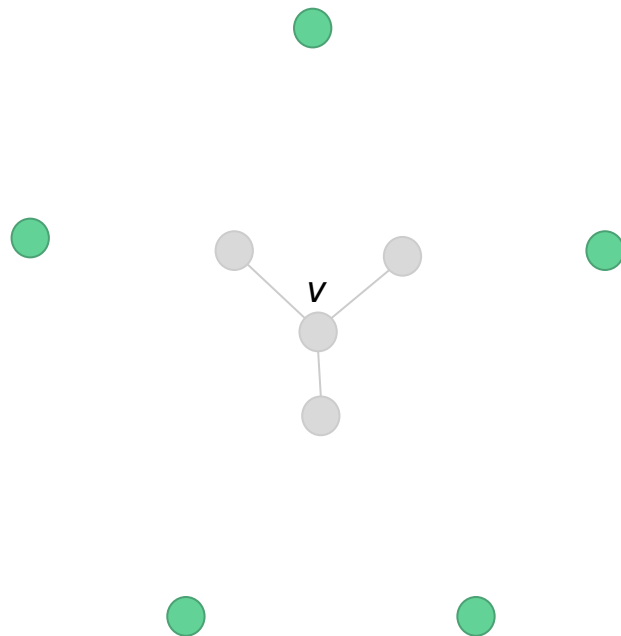
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Leaves a size 5 graph



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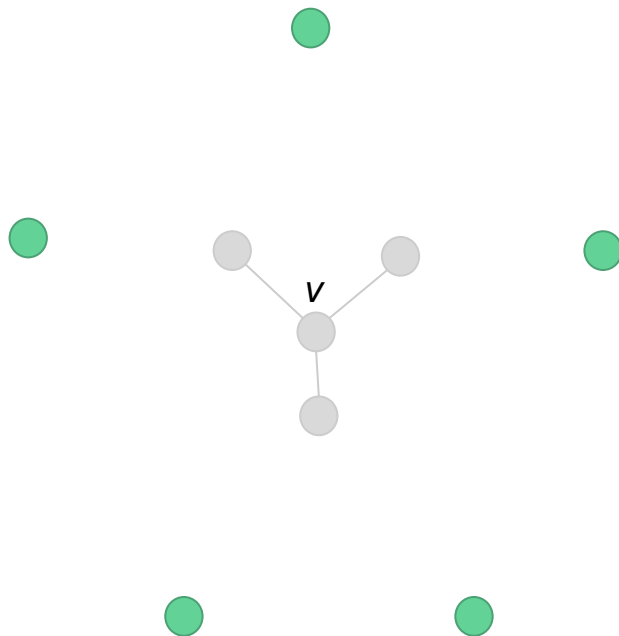
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Leaves a size 5 graph

- no 3-clique

- no 3 independent

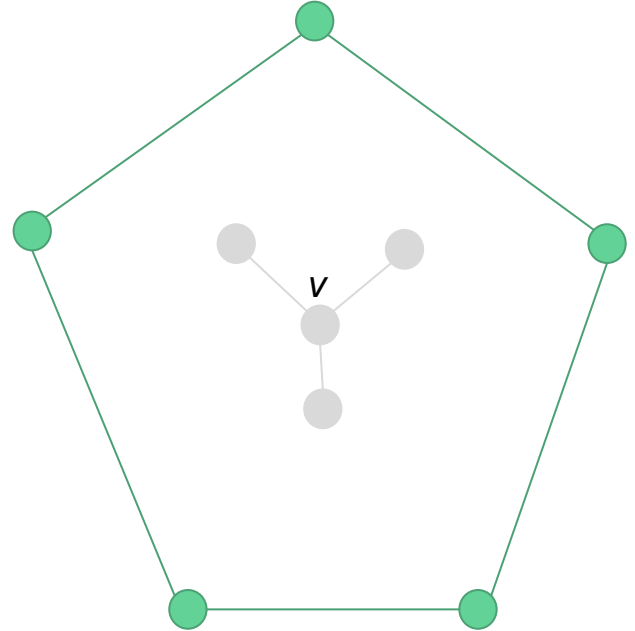
(all are independent of v in G)



$$R(3,4) \leq 9$$

G is regular degree 3

There is only one $(3,3)$ graph of size 5

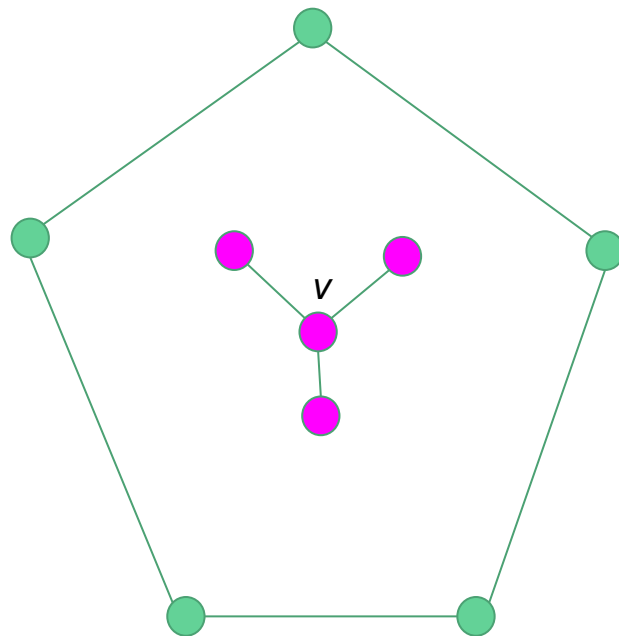


$$R(3,4) \leq 9$$

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There is only one $(3,3)$ graph of size 5

Add v and its neighbors back

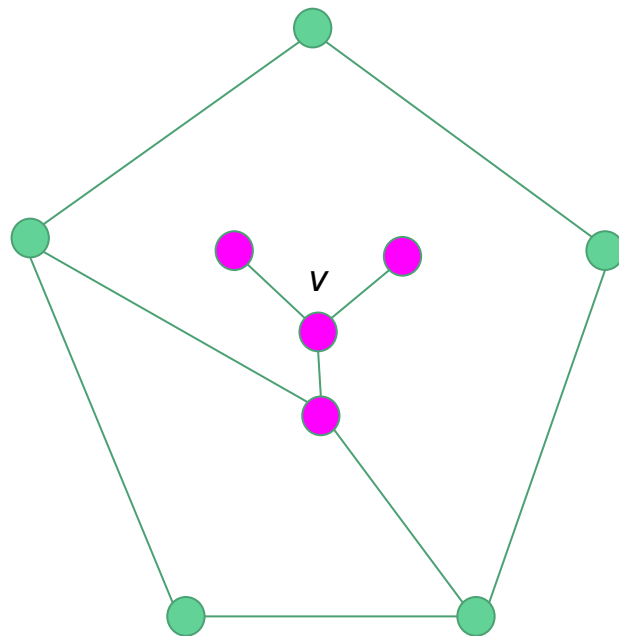


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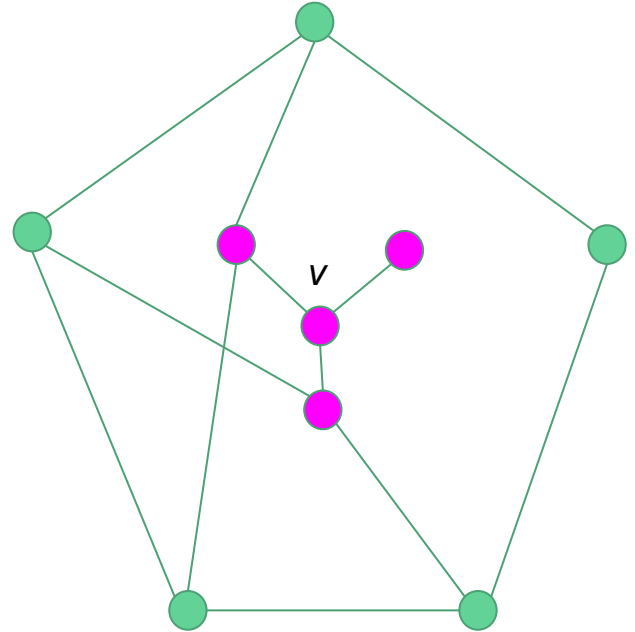


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G is regular degree 3

There is only one $(3,3)$ graph of size 5

Add v and its neighbors back



$$R(3,4) \leq 9$$

Suppose there exists a (3,4) graph G of size 9

- G must be regular degree 3
- removing a neighborhood from G results in a (3,3) graph of size 5
- **G cannot be constructed from a (3,3) graph of size 5**

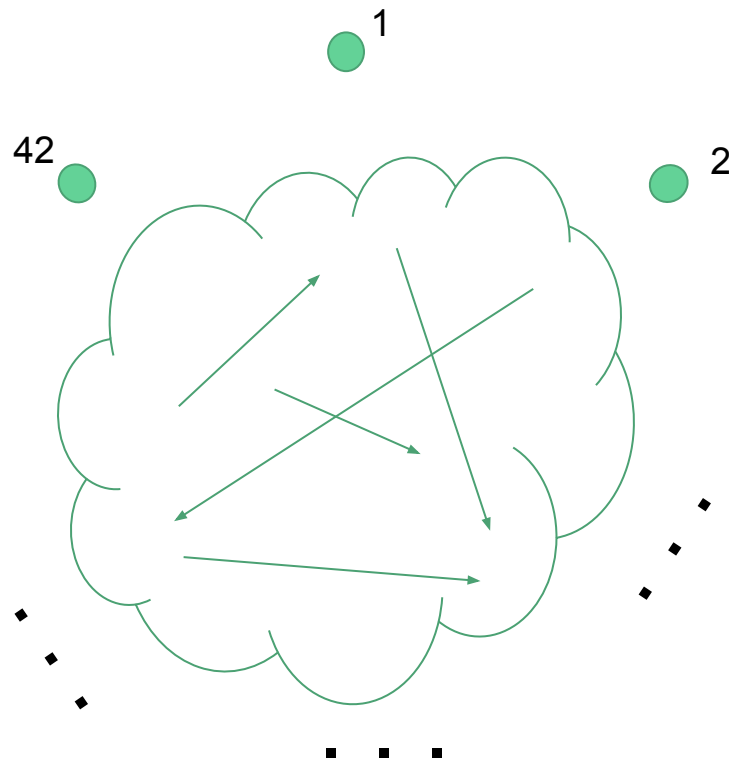
So, there is no (3,4) graph of size 9

Therefore, $R(3,4) \leq 9$

$$R(3,10) \leq 42$$

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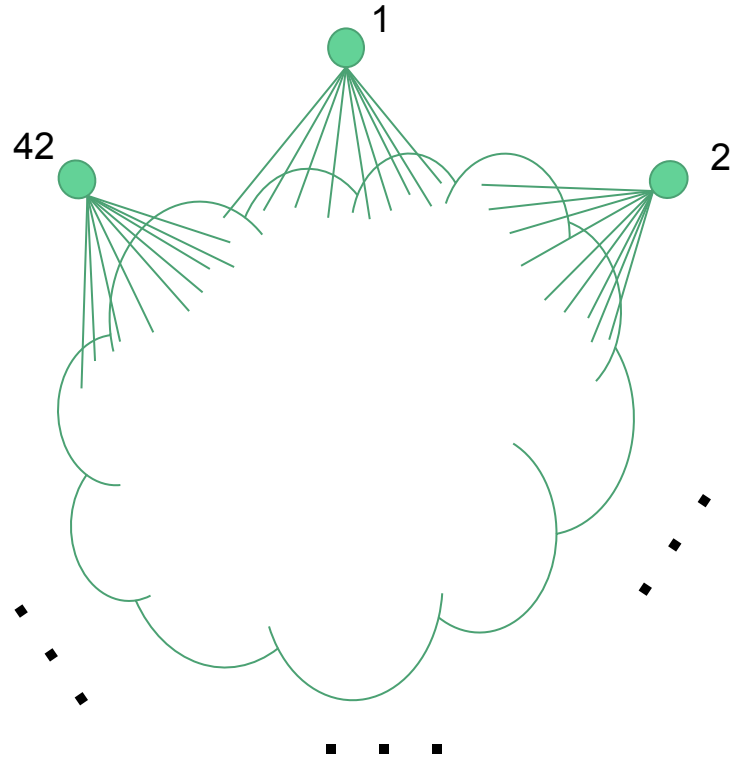
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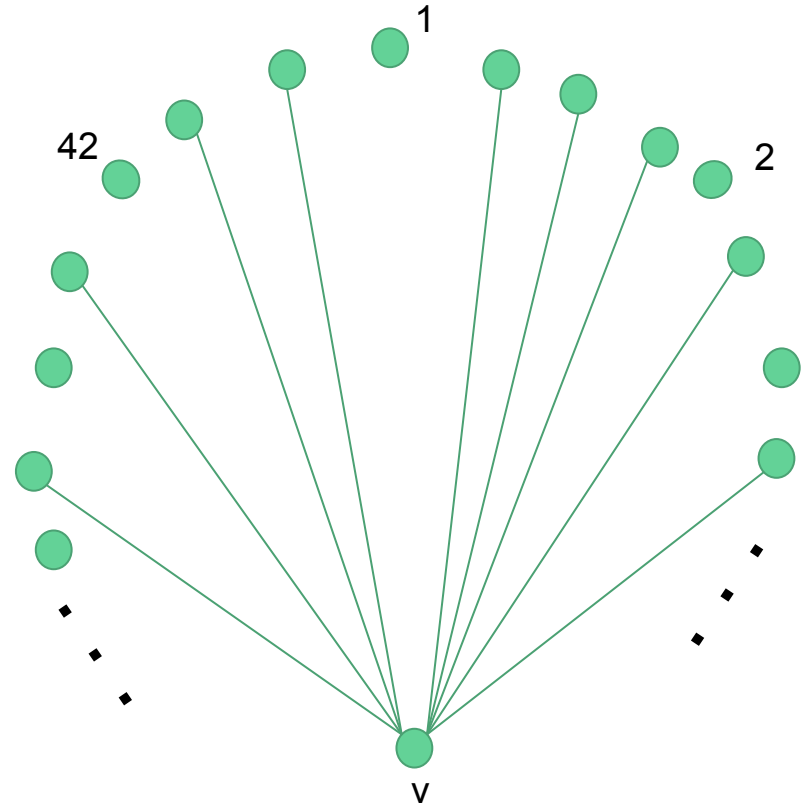
- G must be regular degree 9



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Suppose there exists a $(3, 10)$ graph G of size 42

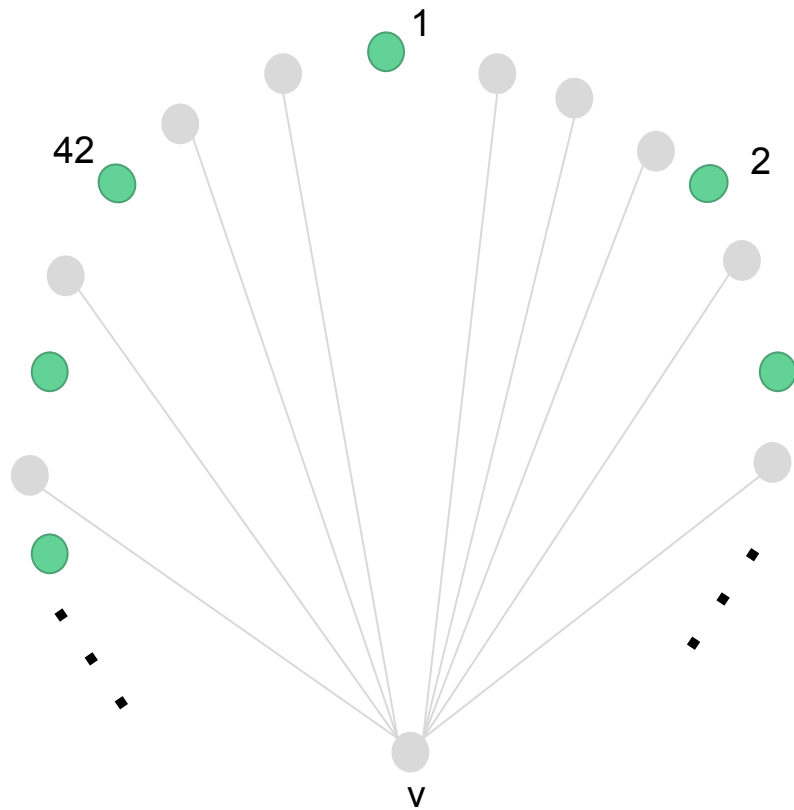
- G must be regular degree 9
- Remove a vertex v and its neighborhood



$$R(3,10) \leq 42$$

Suppose there exists a $(3, 10)$ graph G of size 42

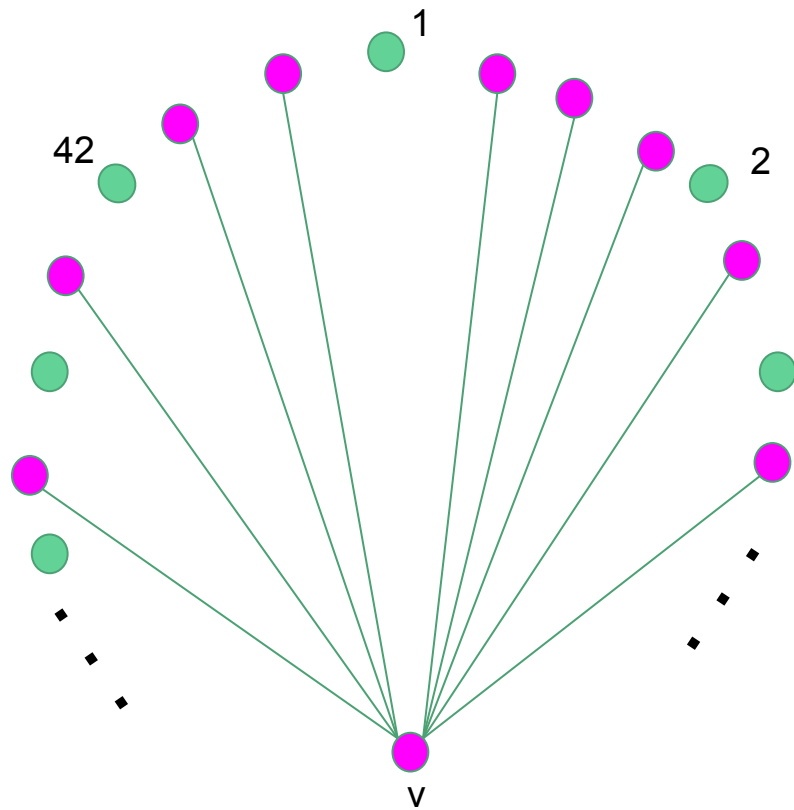
- G must be regular degree 9
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- Results in a $(3, 9)$ graph H of size 32



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Suppose there exists a $(3, 10)$ graph G of size 42

- G must be regular degree 9
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- Results in a $(3, 9)$ graph H of size 32
- G cannot be reconstructed from any such H

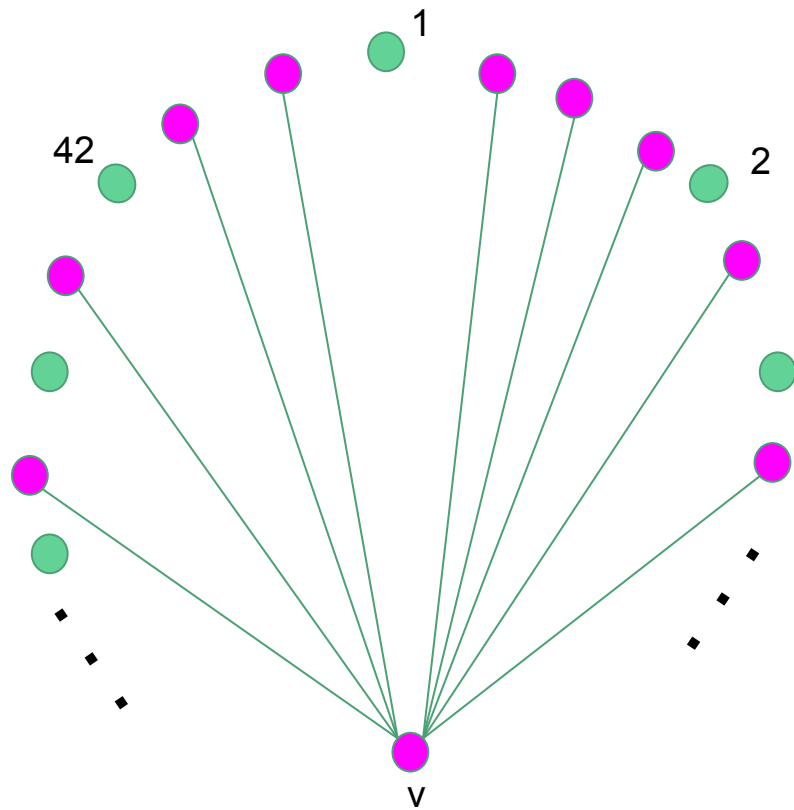


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Therefore, there is no such G



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These steps are computationally hard

Therefore, there is no such G

Maximal triangle free (mtf) method

mtf graph = a graph in which the addition of any edge yields a 3-clique

There is some (x,y) graph of size n iff there is some mtf (x,y) graph of size n

Maximal triangle free (mtf) method

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There is some (x,y) graph of size n iff there is some mtf (x,y) graph of size n

There are often far fewer mtf $(3, k)$ graphs of size n
than there are $(3, k)$ graphs of size n

mtf method Find (3,4) graphs of size 6

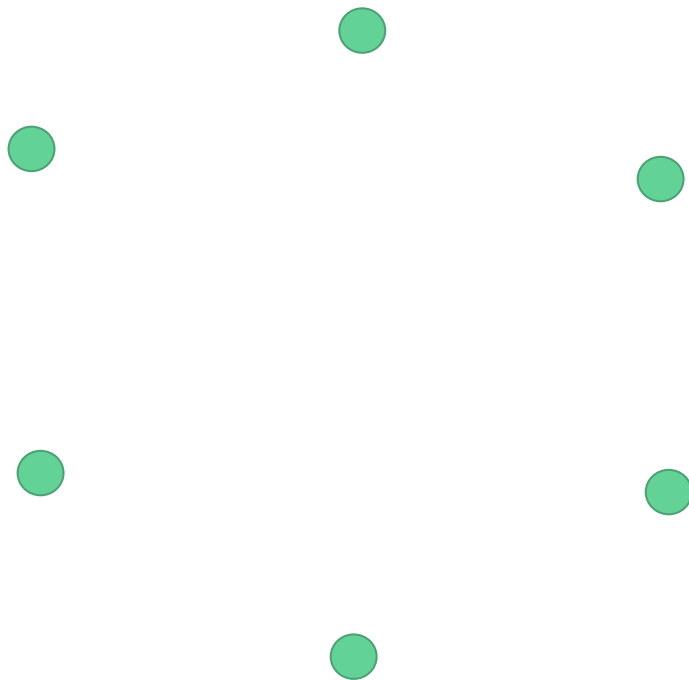
Build up all mtf graphs of size 6

If an mtf graph is a (3,4) graph

 Recursively remove edges and check for new 4-independent sets

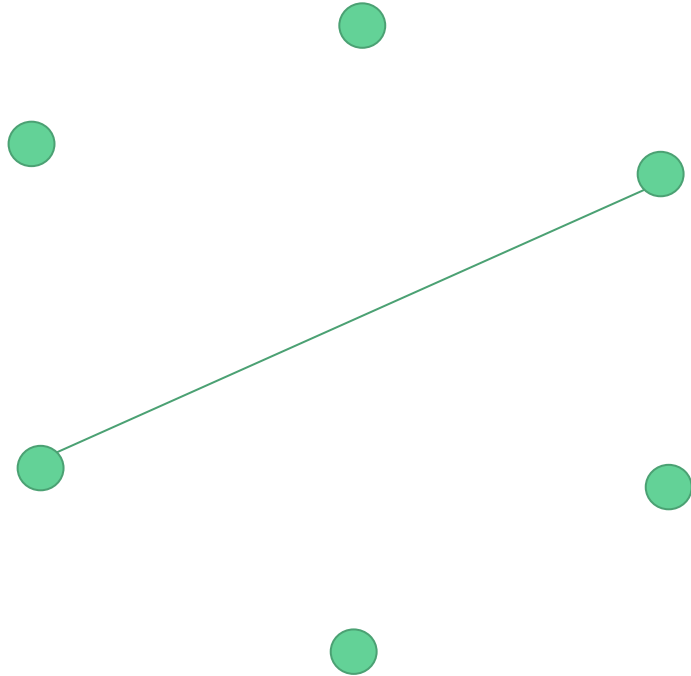
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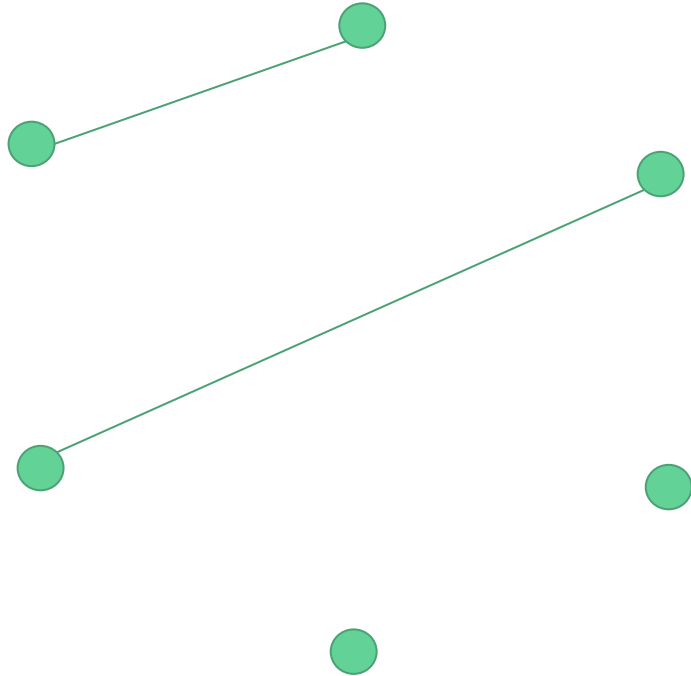
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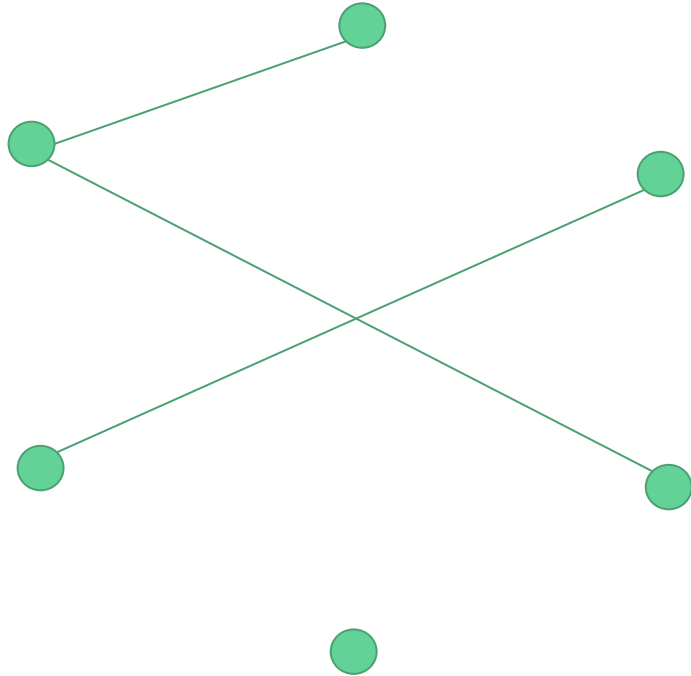
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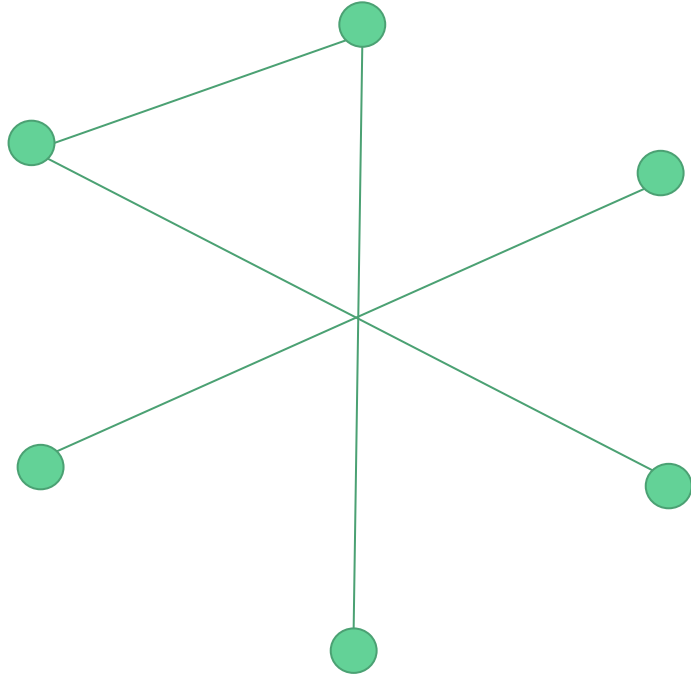
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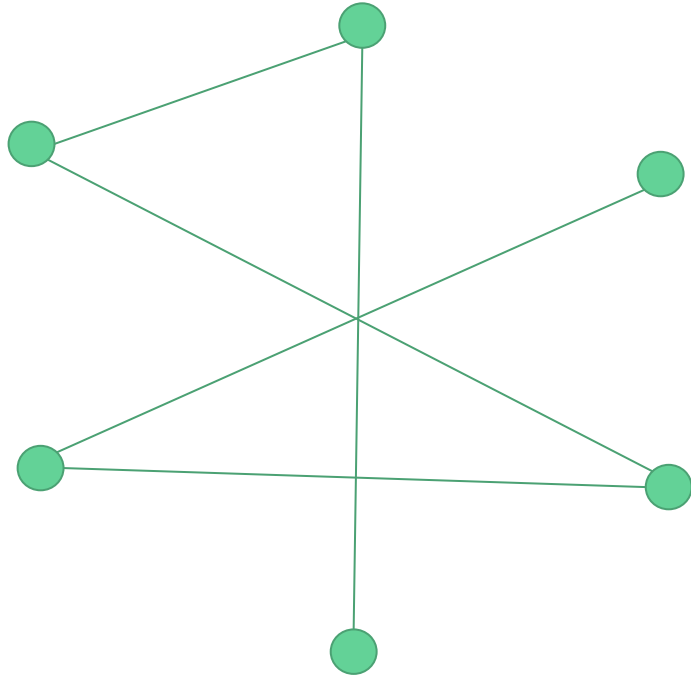
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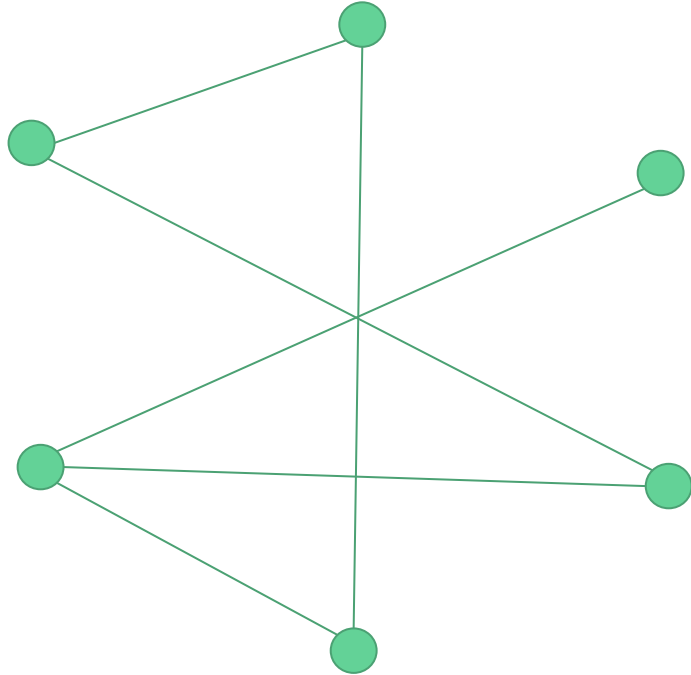
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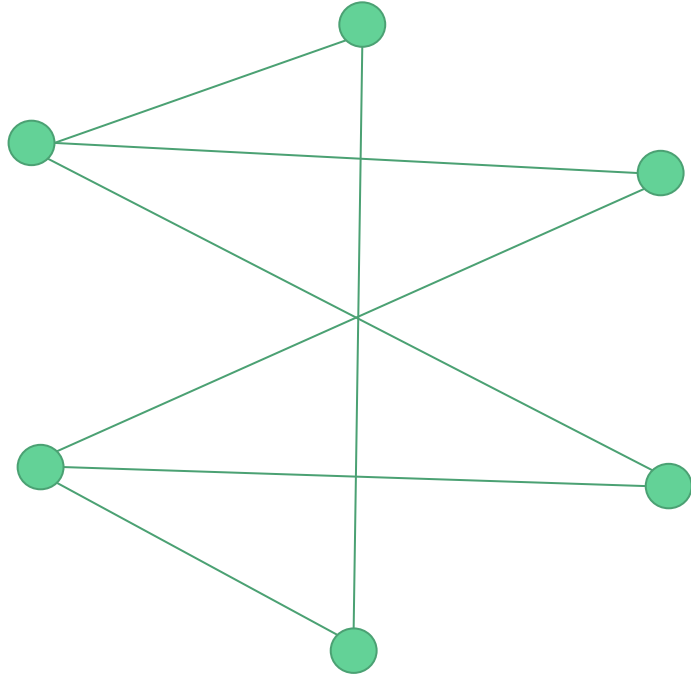
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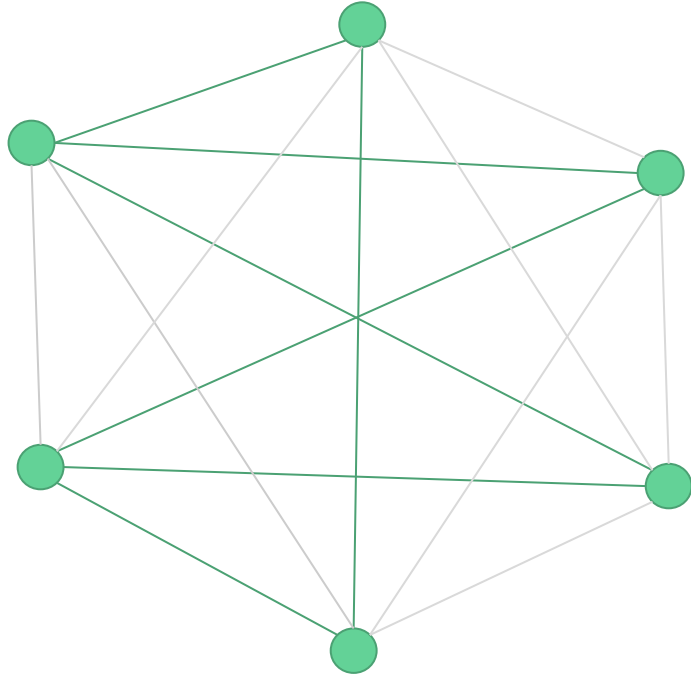
This is an mtf graph



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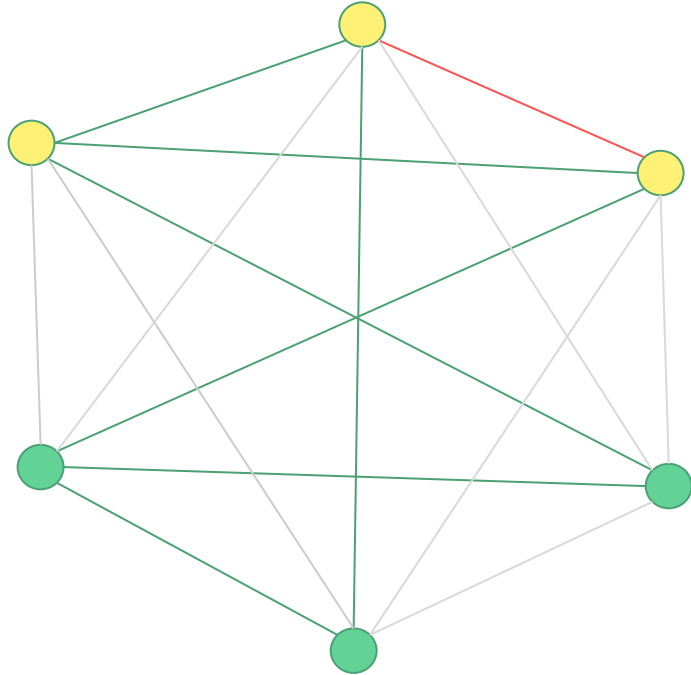
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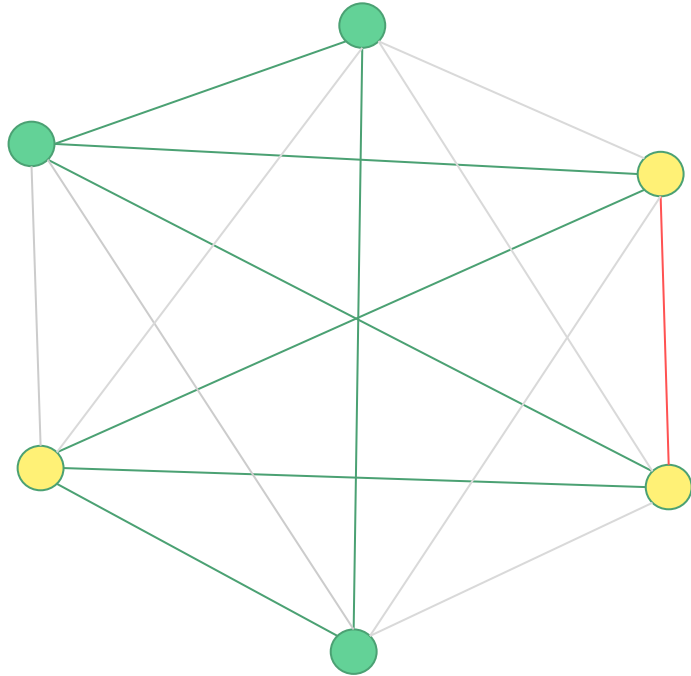
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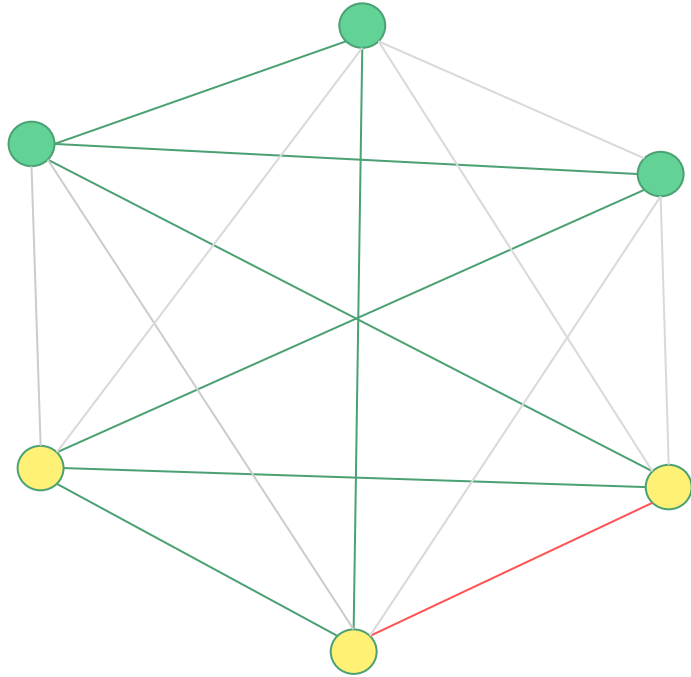
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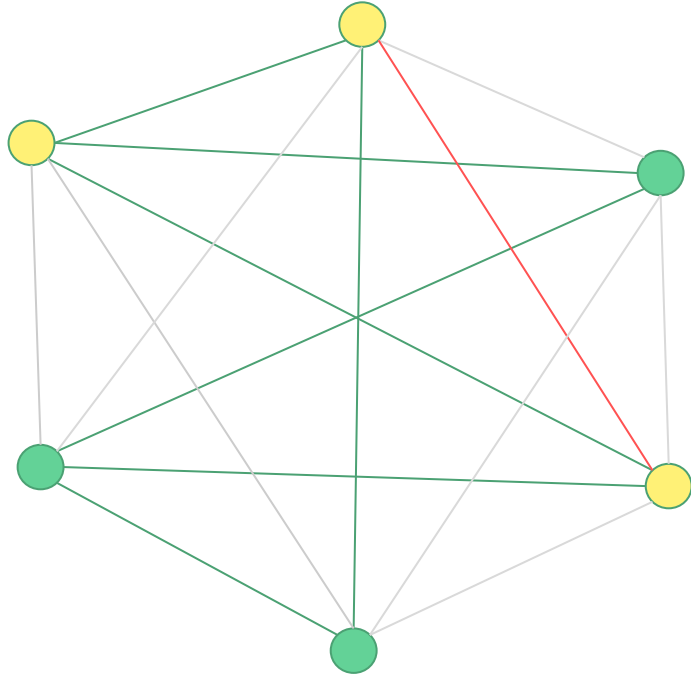
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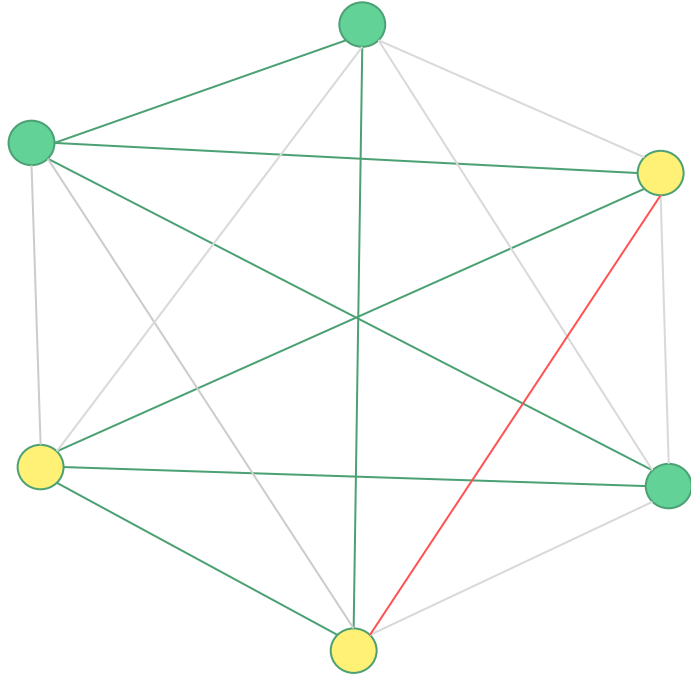
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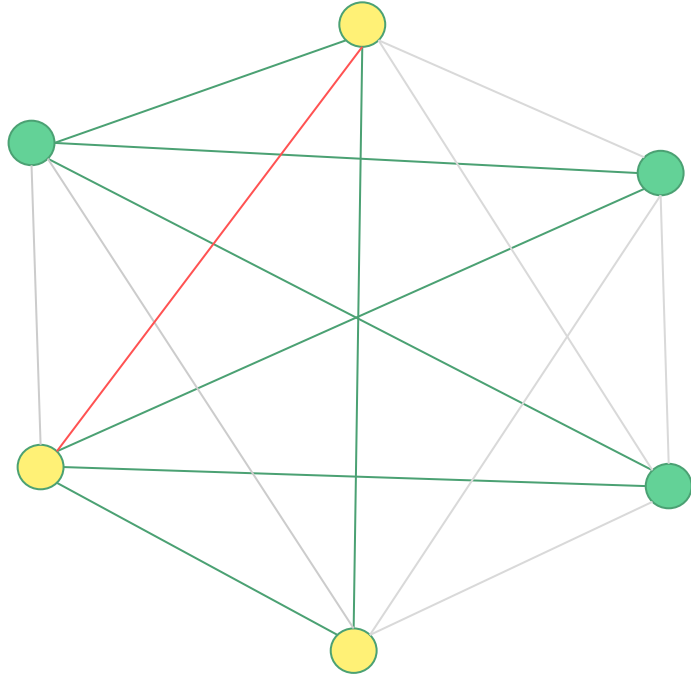
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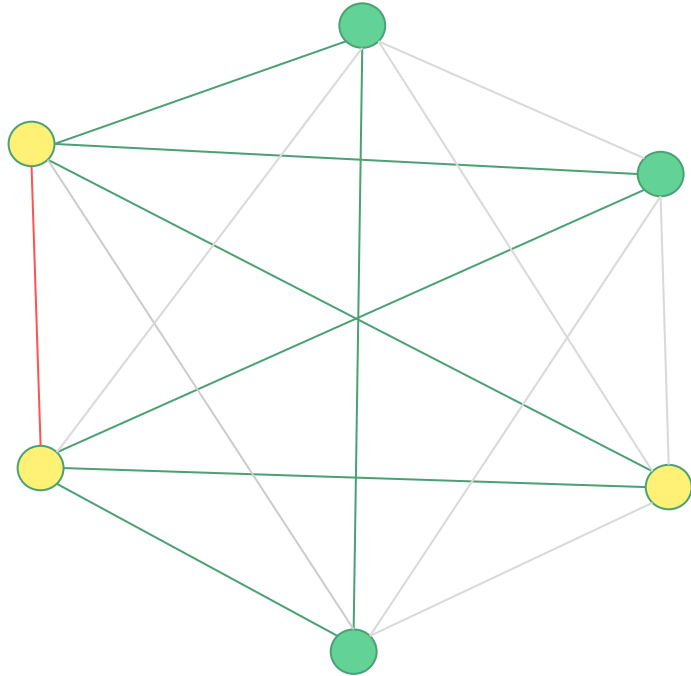
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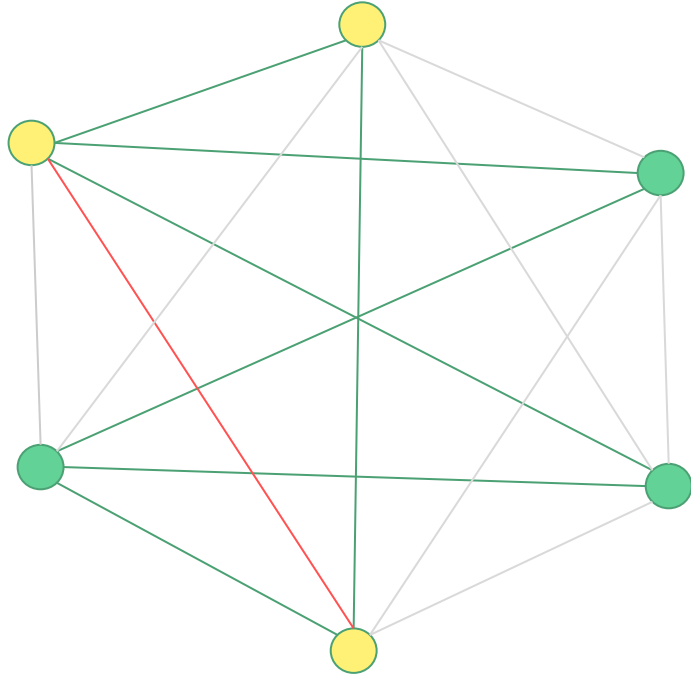
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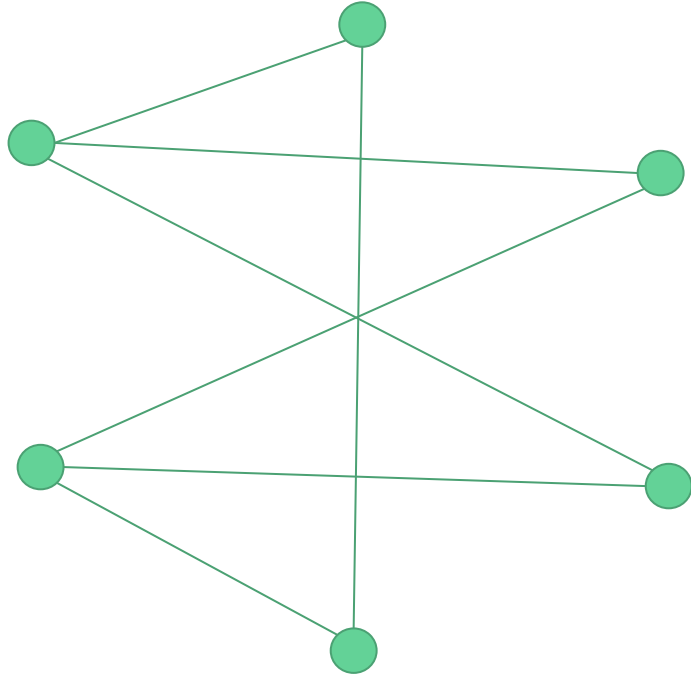
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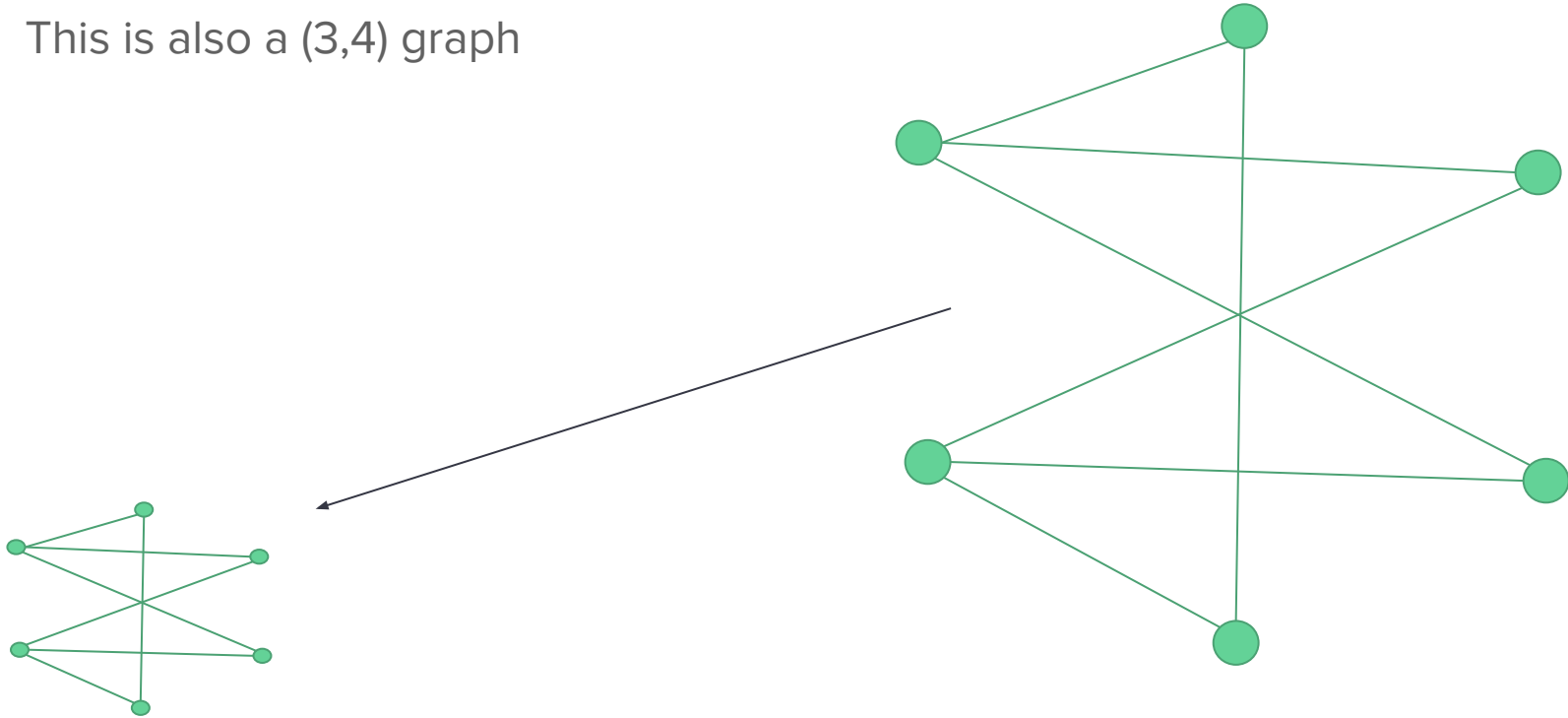
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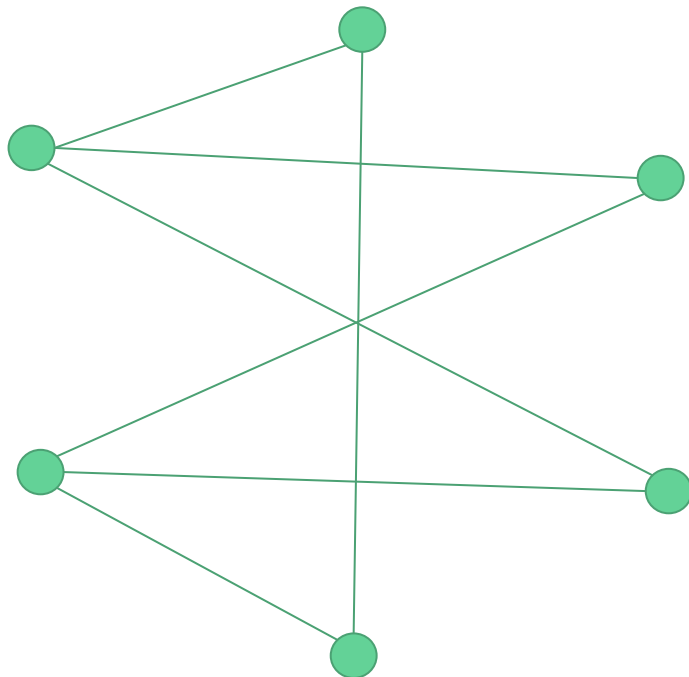
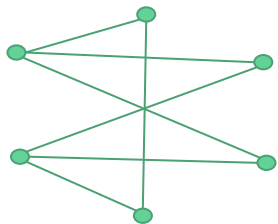
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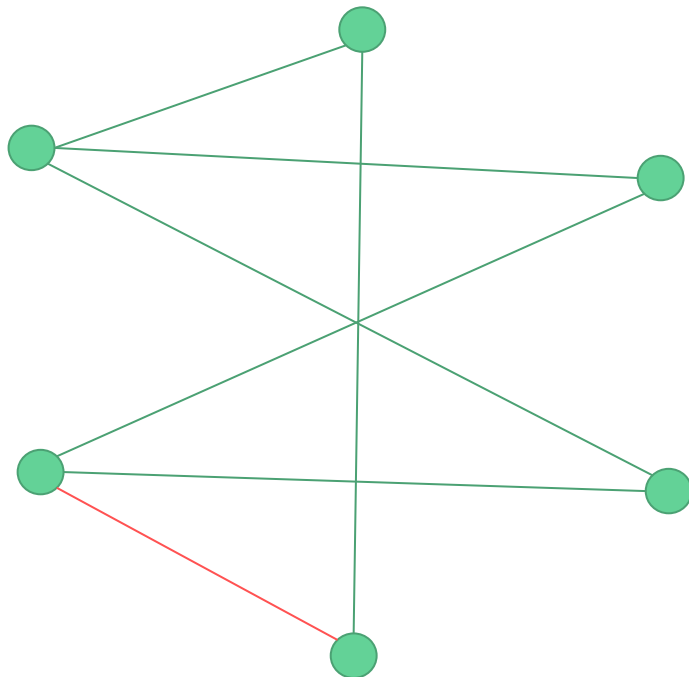
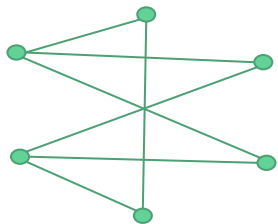
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Recursively remove edges to try to find more (3,4) graphs



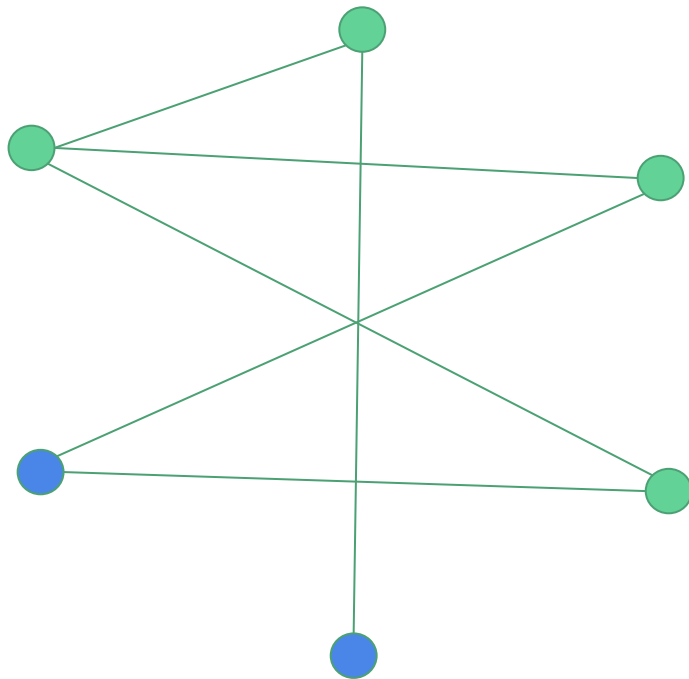
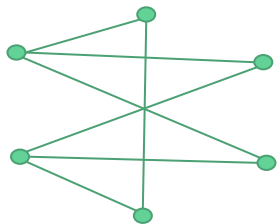
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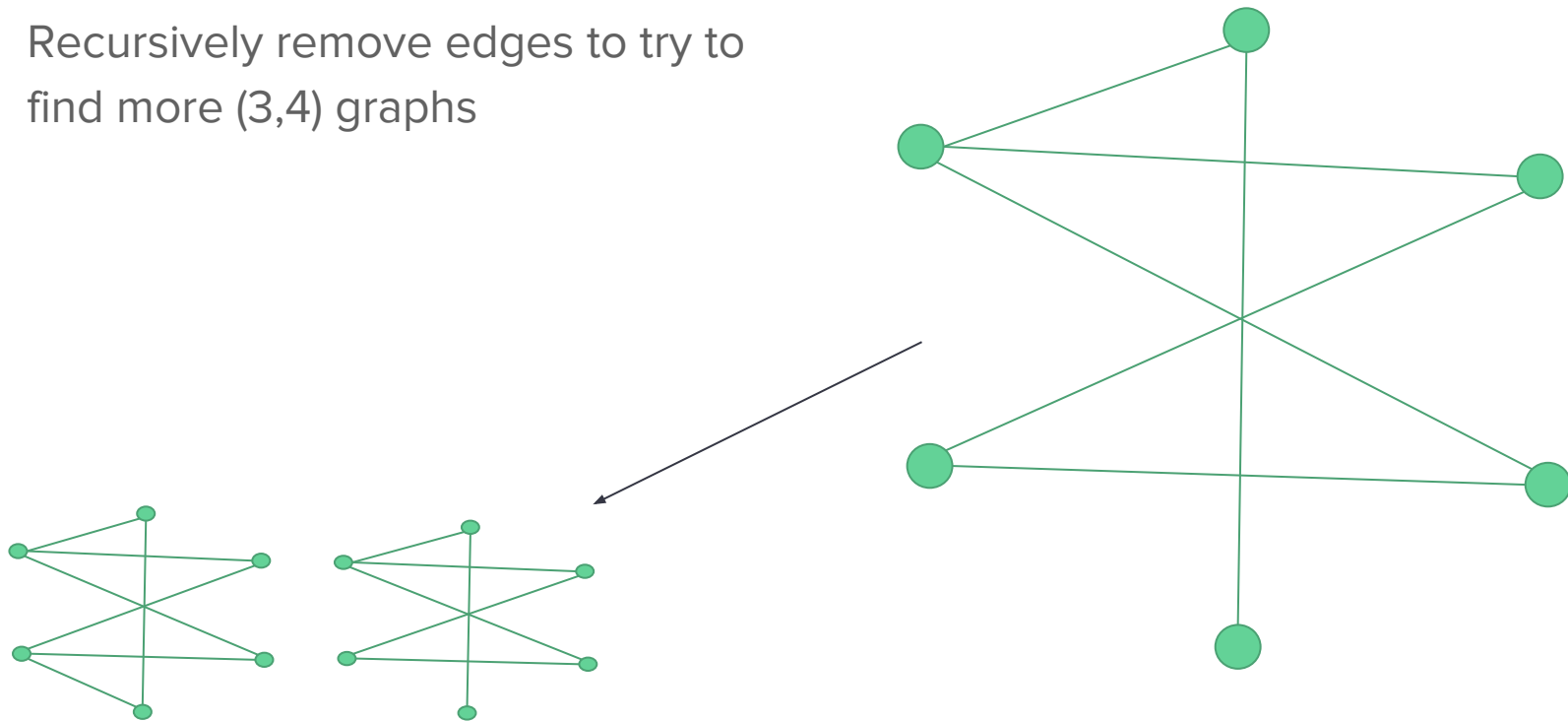
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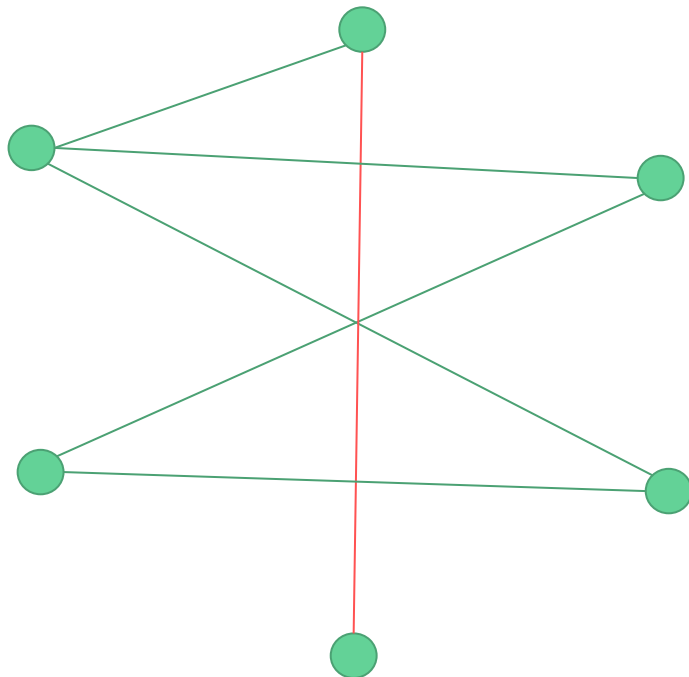
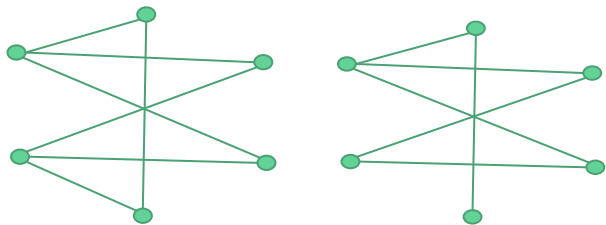
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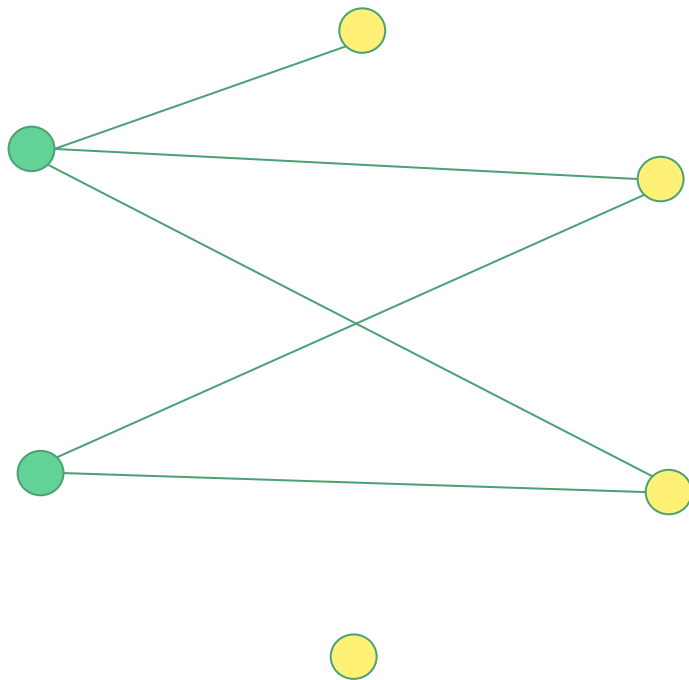
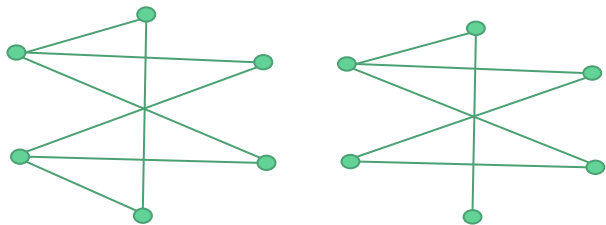
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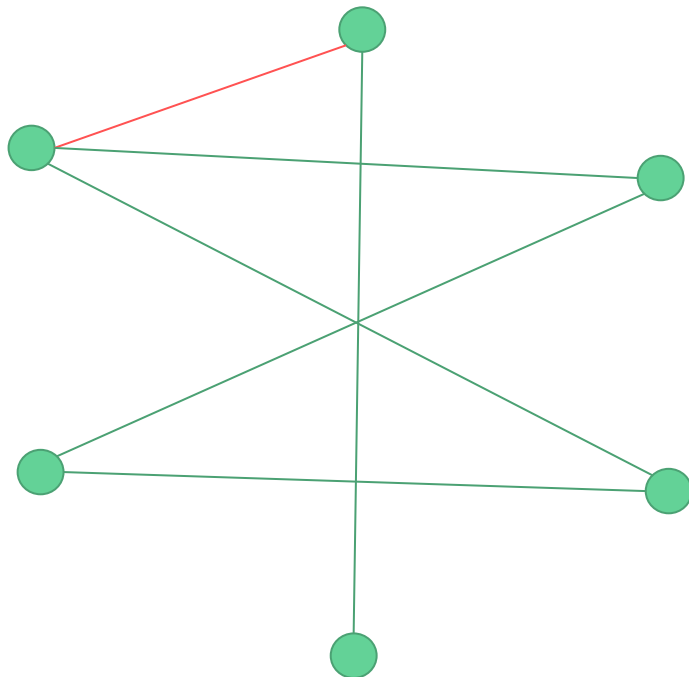
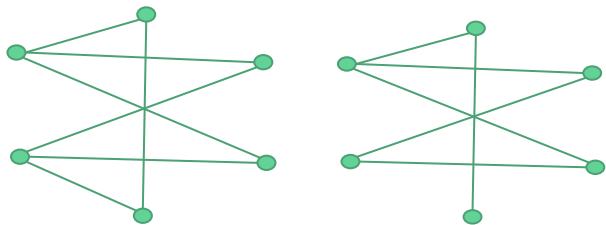
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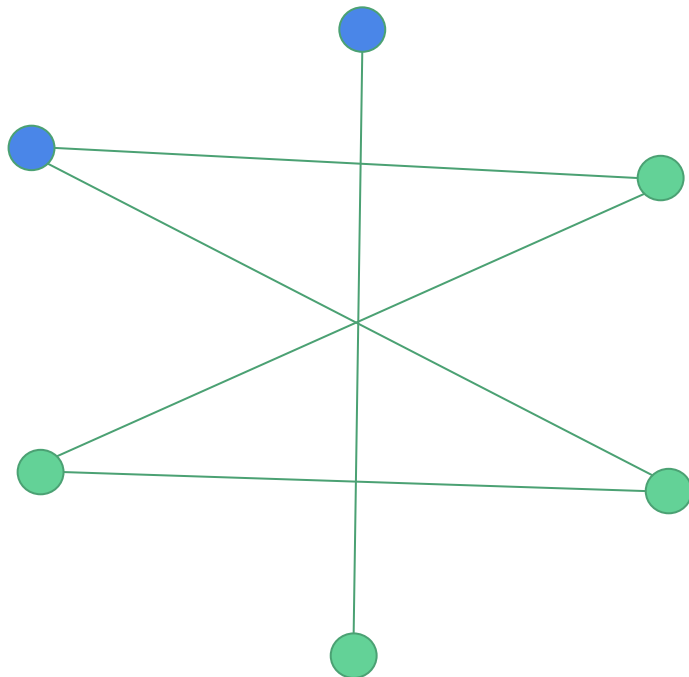
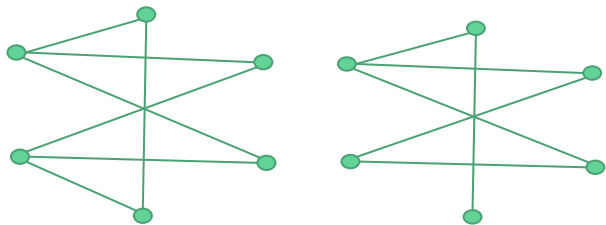
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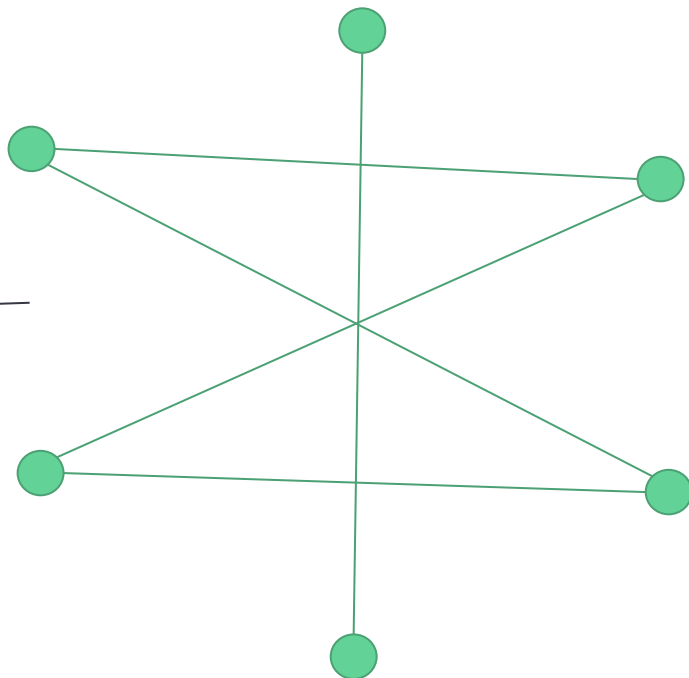
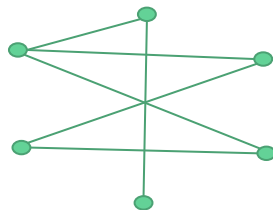
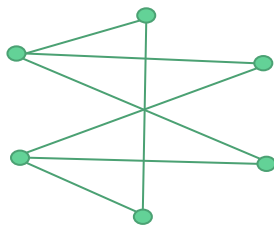
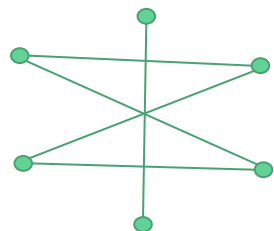
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Fun Facts

Naive upper bound gets pretty close $\rightarrow R(2, 10) + R(3, 9) = 10 + 36 = 46$

Lower bound for $R(3, 10)$ is 40, found in 1987

Since 1987, found 4×10^7 different $r(3, 10, 39)$ graphs

Tight general lower bound found in 1995 ($t^2/\log t$)

Methods for finding lower bound efficiently:

- Artificial bee colony algorithm

- Simulated annealing

New Computational Upper Bounds for Ramsey Numbers $R(3,k)$

Jan Goedgebeur & Stanislaw P. Radziszowski