

Q1) Consider the following situation where  $n=2$

<u>Members</u>	<u>ordered list</u>
$m_1$	$\{ l_1, l_2 \}$
$m_2$	$\{ l_1, l_2 \}$
$l_1$	$\{ m_2, m_1 \}$
$l_2$	$\{ m_2, m_1 \}$

The algo given will output the following valid matching

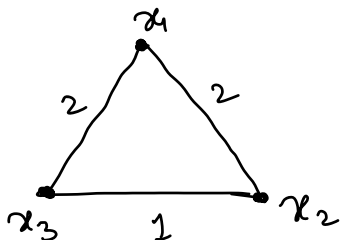
$$\{ (m_1, l_1), (m_2, l_2) \}$$

$$\text{Pref}_{l_1}(m_2) > \text{Pref}_{l_1}(m_1)$$

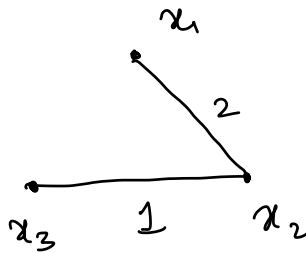
$$\text{Pref}_{m_2}(l_1) > \text{Pref}_{m_2}(l_2)$$

Clearly, the matching is not stable. The given algo does not guarantee a stable matching.

Q2) Consider the following graph

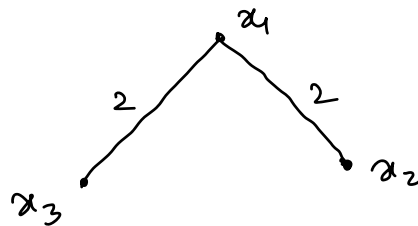


MST rooted at  $x_1$  is the following



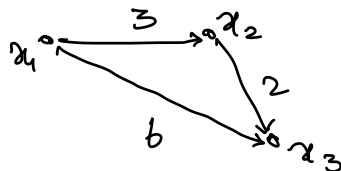
However, the path  $x_1 \rightsquigarrow x_3$  in this MST is not the shortest path from  $x_1$  to  $x_3$  in the overall graph.

Q3) Consider the same graph as above. Running Dijkstra's algo starting from  $x_1$ , we get



Clearly, this is not the MST.

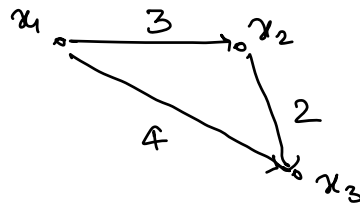
Q4) a) Consider the following graph



Setting the additive factor to  $> 1$ , we see that the shortest path between  $x_1$  &  $x_3$  changes.

b) Scaling the edge weights by a constant multiplicative factor does not change the shortest path.

c) Consider the following graph



Clearly, squaring the edge weights changes the shortest path between  $x_1$  &  $x_3$ .

Note: Any operation that preserves the ordering of the edge weights will not alter the MST.

$$w_i > w_j \Rightarrow w_i + k > w_j + k$$

$$k \cdot w_i > k \cdot w_j$$

$$w_i^2 > w_j^2$$