

# CS/COE0447: Computer Organization and Assembly Language

## Logic Design Introduction (Brief?)

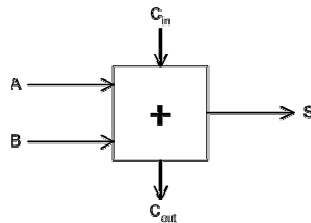
### Appendix C: The Basics of Logic Design

modified by Bruce Childers  
Sangyeun Cho

Dept. of Computer Science  
University of Pittsburgh

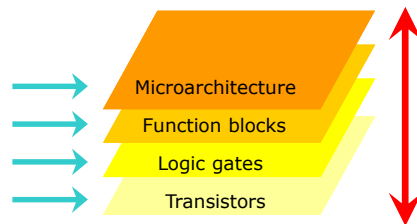
## Logic design?

- Digital hardware is implemented by way of *logic design*
- Digital circuits process and produce two discrete values: 0 and 1
- Example: 1-bit full adder (FA)



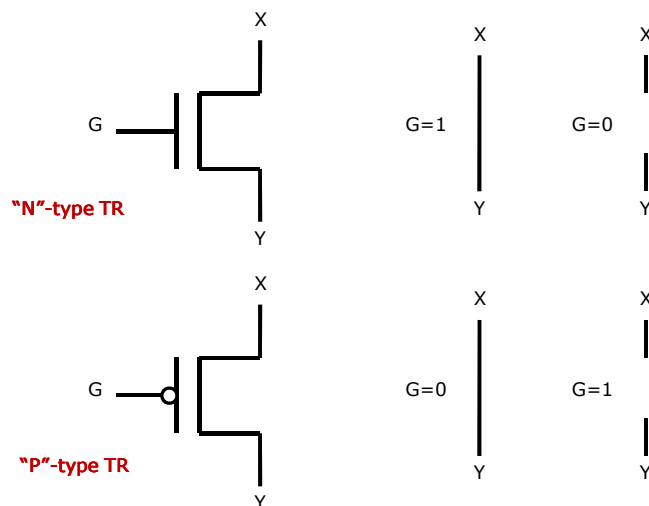
## Layered design approach

- Logic design is done using **logic gates**
- Often we design a desired function using high-level languages and somewhat higher level than logic gates
- Two approaches in design
  - Top down
  - Bottom up

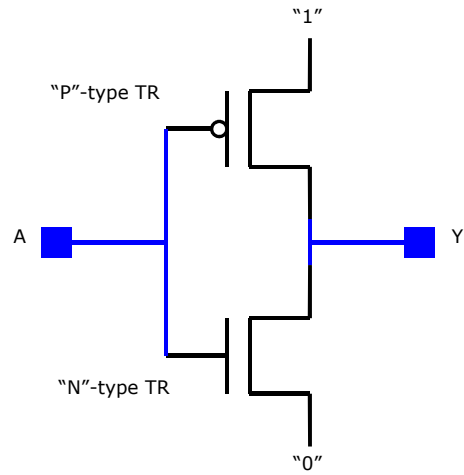


*We'll do logic bottom up*

## Transistor as a switch



## An inverter

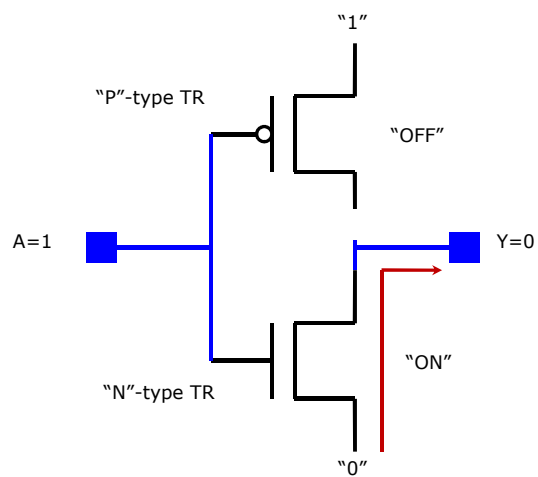


CS/CoE1541: Intro. to Computer Architecture

University of Pittsburgh

5

## When A = 1

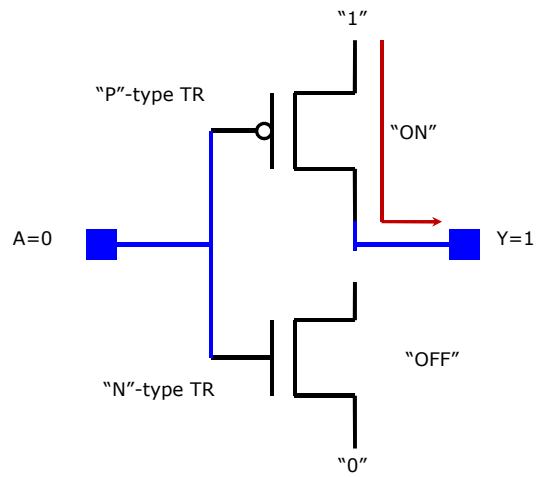


CS/CoE1541: Intro. to Computer Architecture

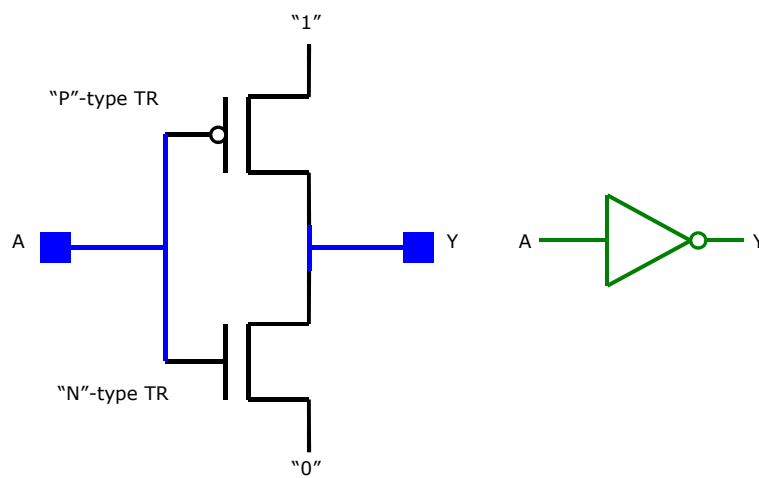
University of Pittsburgh

6

## When $A = 0$

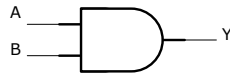


## Abstraction



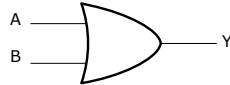
## Logic gates

2-input AND



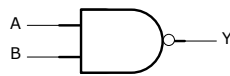
$$Y = A \& B$$

2-input OR



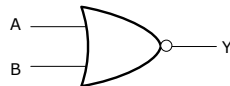
$$Y = A \mid B$$

2-input NAND



$$Y = \sim(A \& B)$$

2-input NOR

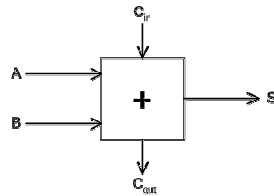


$$Y = \sim(A \mid B)$$

## Describing a function

- $\text{Output}_A = F(\text{Input}_0, \text{Input}_1, \dots, \text{Input}_{N-1})$
- $\text{Output}_B = F'(\text{Input}_0, \text{Input}_1, \dots, \text{Input}_{N-1})$
- $\text{Output}_C = F''(\text{Input}_0, \text{Input}_1, \dots, \text{Input}_{N-1})$
- ...
  
- **Methods**
  - Truth table
  - Sum of products
  - Product of sums

## Truth table



Input			Output	
A	B	C <sub>in</sub>	S	C <sub>out</sub>
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

## Sum of products

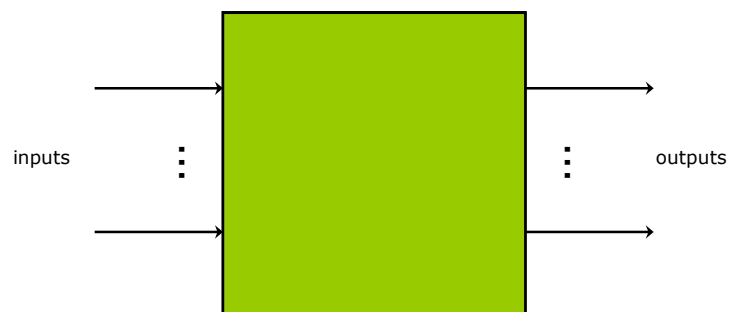
Input			Output	
A	B	C <sub>in</sub>	S	C <sub>out</sub>
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

- $S = A'B'C_{in} + A'BC_{in}' + AB'C_{in}' + ABC_{in}$
- $C_{out} = A'BC_{in} + AB'C_{in} + ABC_{in}' + ABC_{in}$

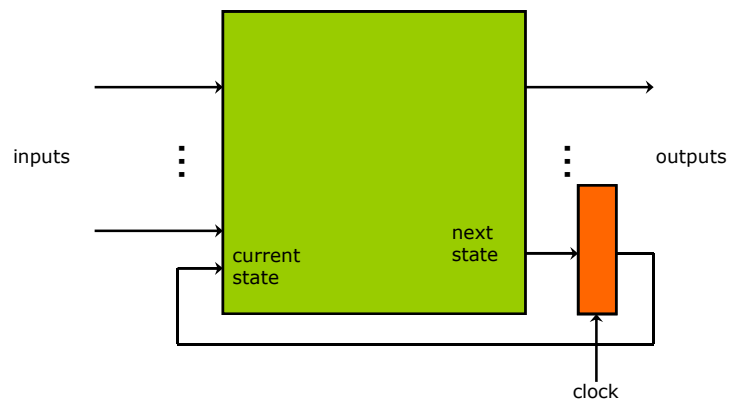
## Combinational vs. sequential logic

- Combinational logic = **function**
  - A function whose outputs are dependent only on the current inputs
  - As soon as inputs are known, outputs can be determined
- Sequential logic = **combinational logic + memory**
  - Some memory elements (i.e., "state")
  - Outputs are dependent on the current state and the current inputs
  - Next state is dependent on the current state and the current inputs

## Combinational logic



## Sequential logic



## Combinational logic

- Any combinational logic can be implemented using sum of products or product of sums
- Input-output relationship can be defined in the truth table format
- From the truth table, each output function can be derived
- Boolean expressions can be further manipulated (e.g., to reduce cost) using various Boolean algebraic rules



## Boolean algebra

- Boole, George (1815~1864): mathematician and philosopher; inventor of Boolean Algebra, the basis of all computer arithmetic
- Binary values:  $\{0,1\}$
- Two binary operations: AND ( $\times/\cdot$ ), OR ( $+$ )
- One unary operation: NOT ( $\sim$ )

## Boolean algebra

- Binary operations: AND ( $\times/\cdot$ ), OR ( $+$ )
  - Idempotent
    - $a \cdot a = a + a = a$
  - Commutative
    - $a \cdot b = b \cdot a$
    - $a + b = b + a$
  - Associative
    - $a \cdot (b \cdot c) = (a \cdot b) \cdot c$
    - $a + (b + c) = (a + b) + c$
  - Distributive
    - $a \cdot (b + c) = a \cdot b + a \cdot c$
    - $a + (b \cdot c) = (a + b) \cdot (a + c)$

## Boolean algebra

- De Morgan's laws

- $\sim(a \cdot b) = \sim a + \sim b$
- $\sim(a + b) = \sim a \cdot \sim b$

It is not true I ate the sandwich and the soup.

same as:

I didn't eat the sandwich or I didn't eat the soup.

- More...

- $a + (a \cdot b) = a$
- $a \cdot (a + b) = a$
- $\sim \sim a = a$
- $a + \sim a = 1$
- $a \cdot (\sim a) = 0$

It is not true that I went to the store or the library.

same as:

I didn't go to the store and I didn't go to the library.

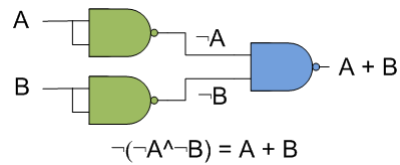
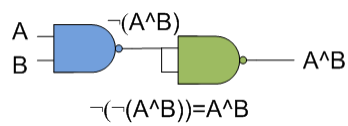
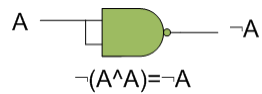
## Expressive power

- With AND/OR/NOT, we can express any function in Boolean algebra

- Sum (+) of products ( $\cdot$ )

- What if we have NAND/NOR/NOT?
- What if we have NAND only?
- What if we have NOR only?

## Using NAND only



## Using NOR only (your turn)

- Can you do it?
- NOR is  $\neg(A + B)$

### NOT

$$\begin{aligned} &= \neg(A + A) \\ &= \neg A \wedge \neg A \\ &= \neg A \end{aligned}$$

### AND

$$\begin{aligned} &= \neg(\neg(A + A) + \neg(B + B)) \\ &= \neg(\neg A \wedge \neg A + \neg B \wedge \neg B) \\ &= \neg(\neg A + \neg B) \\ &= \neg(\neg A) \wedge \neg(\neg B) \\ &= A \wedge B \end{aligned}$$

### OR

$$\begin{aligned} &= \neg(\neg(A + B) + \neg(A + B)) \\ &= (A + B) \wedge (A + B) \\ &= A + B \end{aligned}$$

## Using NOR only (your turn)

- Can you do it?
- NOR is  $\neg(A + B)$ 
  - I.e., We need to write NOT, AND, and OR in terms of NOR

NOT

$$= \neg A$$

$$= \neg A \wedge \neg A$$

$$= \neg(A + A)$$

AND

$$= A \wedge B$$

$$= \neg(\neg A) \wedge \neg(\neg B)$$

$$= \neg(\neg A + \neg B)$$

$$= \neg(\neg A \wedge \neg A + \neg B \wedge \neg B)$$

$$= \neg(\neg(A + A) + \neg(B + B))$$

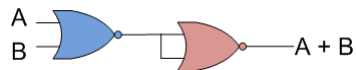
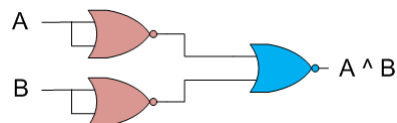
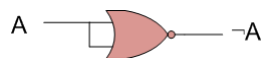
OR

$$= A + B$$

$$= (A + B) \wedge (A + B)$$

$$= \neg(\neg(A + B) + \neg(A + B))$$

## Using NOR only (your turn)

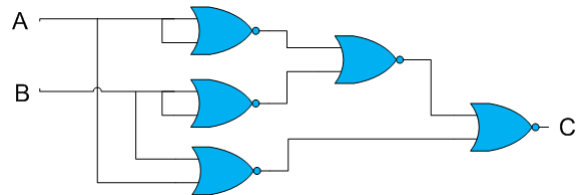
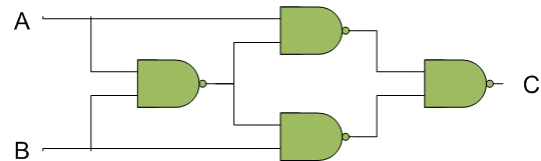


## Now, it's really your turn....

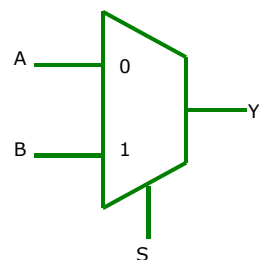
- How about XOR?

A	B	C
0	0	0
0	1	1
1	0	1
1	1	0

$$C = A'B + AB'$$



## Multiplexor (aka MUX)

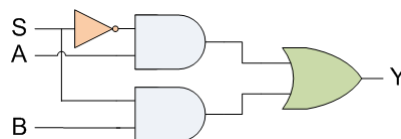


$$Y = (S) ? B : A;$$

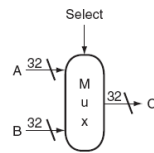
when  $S =$   
0: output A  
1: output B

S	A	B	Y
0	0	x	0
0	1	x	1
1	x	0	0
1	x	1	1

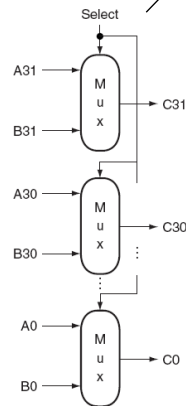
$$Y = S'A + SB$$



# A 32-bit MUX



a. A 32-bit wide 2-to-1 multiplexor



b. The 32-bit wide multiplexor is actually an array of 32 1-bit multiplexors

Use 32 1-bit muxes  
Each mux selects 1 bit  
S is connected to each mux