

CS 1510: Algorithm Design  
Homework 3 - Greedy Problems 7  
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## Problem 7

Our greedy algorithm goes through each page  $j$  in the fast memory and locates its first reappearance  $a_j$  in the slow memory, then evicts the page which reappears the furthest away in the slow memory.

We will prove that the algorithm  $G$  solves the problem.

Suppose towards a contradiction that there is an input  $I$  on which the algorithm  $G$  produces unacceptable output. Let  $Opt(I)$  be an optimal output which agrees with  $G(I)$  for the maximum number of steps of all such optimal outputs.

Since  $G$  produces unacceptable output and  $Opt$  produces the correct solution, they must disagree on at least one interval. Let  $k$  be the first such interval where they disagree, so that  $G_k \neq O_k$ .

Construct  $Opt'(I) = Opt(I) - O_k + G_k$ , by swapping the positions of  $G_k$  and  $O_k$  in  $Opt$ . Then clearly  $Opt'(I)$  agrees with  $G$  for one additional term than  $Opt$ . Additionally, since  $G$  evicts the page from the fast memory which reappears furthest in time in the slow memory, then the next appearance of the page  $G_k$  accesses must reappear after  $Opt_k$ . Because of this, in  $Opt'$  the page  $O_k$  represents would need to be reloaded before the page that  $G_k$  represents, thus replacing  $O_k$  with  $G_k$  in  $Opt'$  will not force an additional eviction of the page that  $G_k$  represents.

Also note that since we're only adjusting in memory (fast and slow) the pages represented by  $O_k$  and  $G_k$ , the remaining pages will be unaffected by this swap from  $Opt$  to  $Opt'$ . Since no additional evictions are necessary,  $Opt'$  will have at most the same number of evictions as  $Opt$ . So if  $Opt$  was optimal the  $Opt'$  is optimal as well.

Thus we have shown that there exists  $Opt'(I)$  which is an optimal output which agrees with  $G(I)$  for one step further than  $Opt(I)$ , and so we have a contradiction with the fact that  $Opt(I)$  agrees with  $G(I)$  for the maximum number of steps. So our original assumption that there was an input on which  $G$  produced unacceptable output was incorrect, and we are done.

