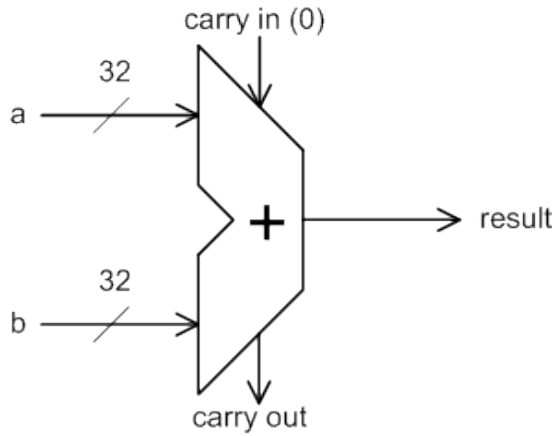
When N = 32

- If we use the two's complement method
  - 0000 0000 0000 0000 0000 0000 0000 0000 = 0
  - 0000 0000 0000 0000 0000 0000 0000 0001 = +1
  - 0000 0000 0000 0000 0000 0000 0000 0010 = +2
  - …
  - 0111 1111 1111 1111 1111 1111 1111 1110 = +2,147,483,646
  - 0111 1111 1111 1111 1111 1111 1111 1111 = +2,147,483,647
  - 1000 0000 0000 0000 0000 0000 0000 0000 = - 2,147,483,648
  - 1000 0000 0000 0000 0000 0000 0000 0001 = - 2,147,483,647
  - 1000 0000 0000 0000 0000 0000 0000 0010 = - 2,147,483,646
  - …
  - 1111 1111 1111 1111 1111 1111 1111 1101 = - 3
  - 1111 1111 1111 1111 1111 1111 1111 1110 = - 2
  - 1111 1111 1111 1111 1111 1111 1111 1111 = - 1

Addition

- We are quite familiar with adding two numbers in decimal
  - What about adding two binary numbers?
- If we use the two's complement method to represent binary numbers, addition can be done in a straightforward way



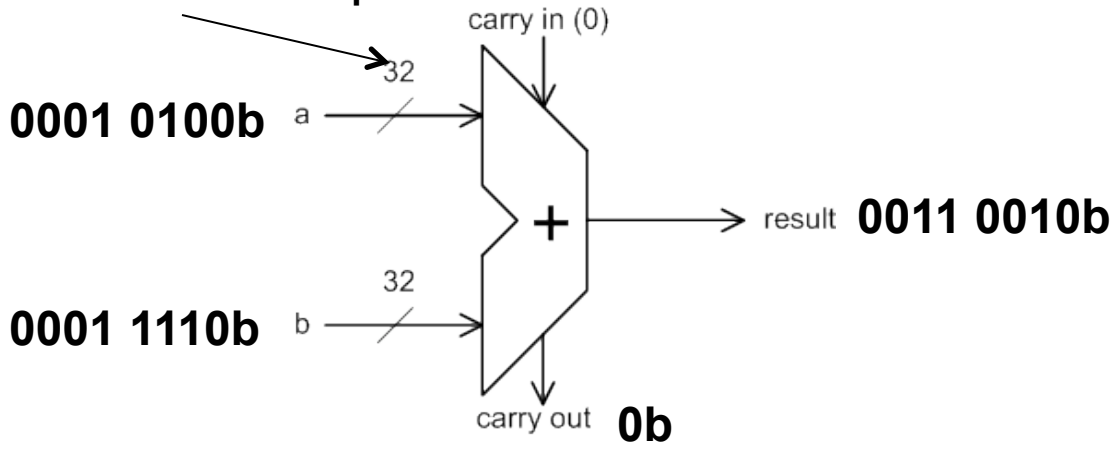
Suppose:
N=8
a=20
b=30

What is result and carry out?

Addition

- N=8, a=20, b=30
- Do binary addition to get result and carryout
- Convert A and B to binary? How?
  - a=20=4+16=2^2+2^4 => a is 0001 0100b
  - b=30=16+8+4+2=$2^4+2^3+2^2+2^1$ => b is 0001 1110b

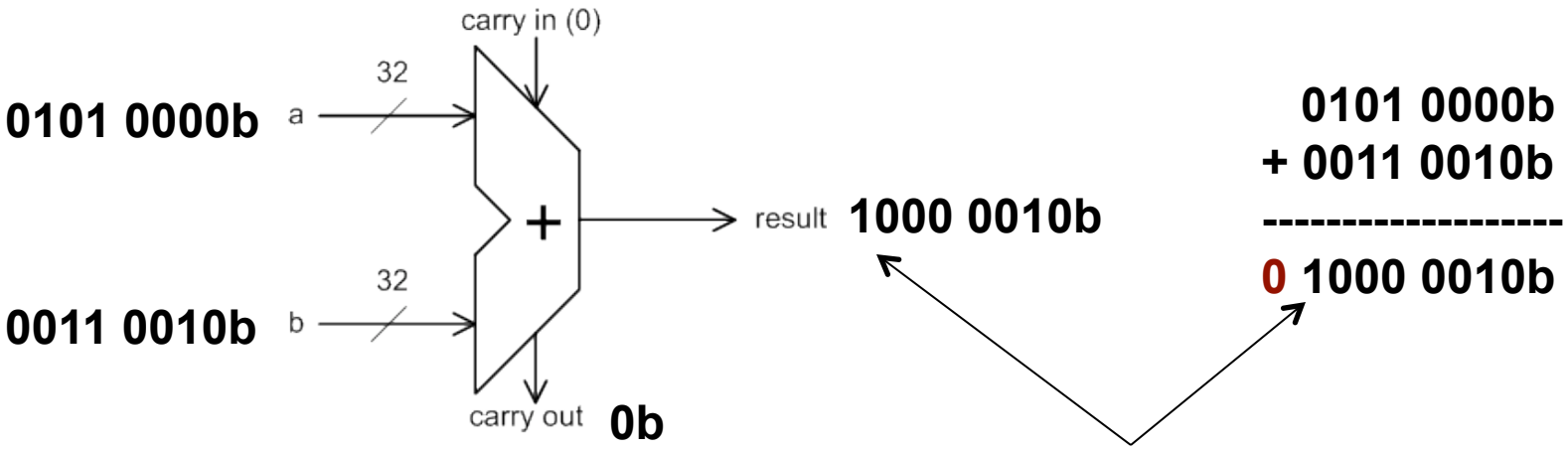
NOTE: N=8 in this example



0001 0100b
+ 0001 1110b
----
0 0011 0010b

Addition

- N=8, a=80, b=50
- Do binary addition to get result and carryout
- Convert A and B to binary? How?
  - A=80=64+16=2^6+2^4 => a is 0101 0000b
  - b=50=32+16+2=2^5+4^3+2^1 => b is 0011 0010b



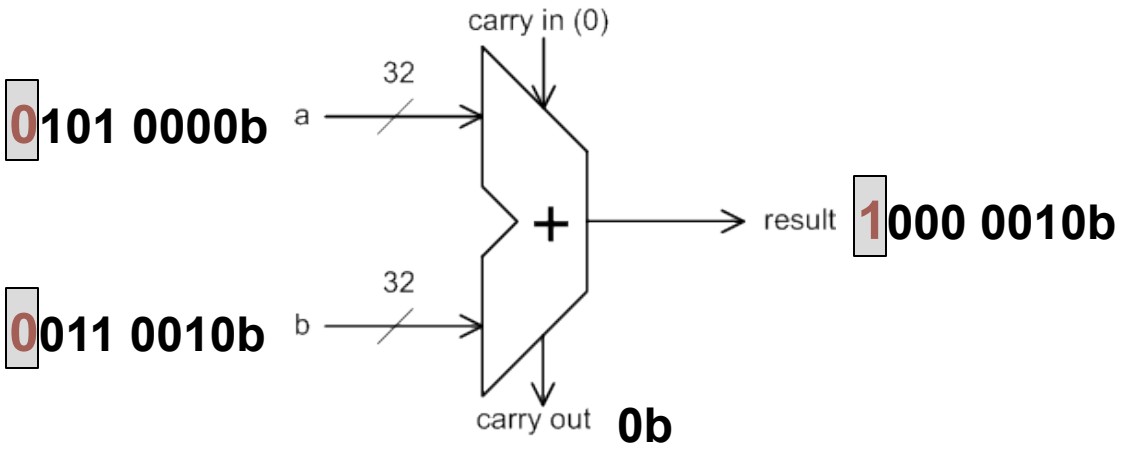
Overflow

- Because we use a limited number of digits to represent a number, the result of an operation may not fit
- No overflow when result remains in expected range
  - We ***add two numbers with different signs***
  - We ***subtract a number from another number with the same sign***
- When can overflow happen?

a	b	overflow possible?
+	+	yes
+	-	no
-	+	no
-	-	yes

Overflow

- What is special about the cases where overflow happened?
  - ***The input values signs are the same; so, can go outside range***
- Overflow detection
  - Adding two positive numbers yields a negative number
  - Adding two negative numbers yields a positive number



Check signs
a is positive
b is positive
result isn't!

Overflow

- Can detect by inspecting sign bits of inputs and output
- Alternatively, can also detect by watching “carries”

```
 0 1110 000x     (carries from previous bit add)
   0101 0000     a
+  0011 0010     b
----------------
 0 1000 0010
```

Notice the carry into sign bit is different than the final carryout

When carry into sign bit doesn’t equal carryout implies overflow

What happens on overflow?

- The CPU can
  - Generate an exception (what is an exception?)
  - Set a flag in the status register (what is the status register?)
  - Do nothing

- Languages may have different notions about overflow

- Do we have overflows in the case of unsigned, always positive numbers?
  - Example: addu, addiu, subu

MIPS example

- I looked at the MIPS32 instruction set manual
- ADD, ADDI instructions generate an ***exception*** on overflow
- ADDU, ADDIU are ***silent***

```
li     $t0,0x40000000
add    $t1,$t0,$t0
```

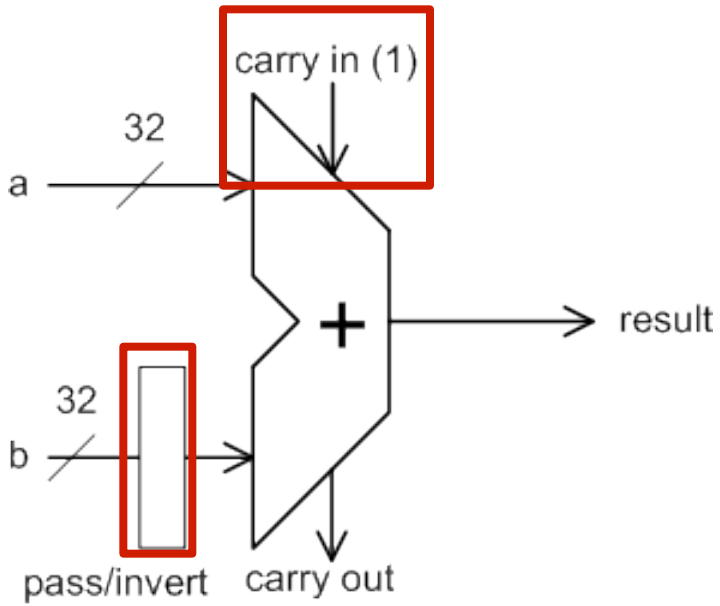
MARS give error

```
li     $t0,0x40000000
addu   $t1,$t0,$t0
```

MARS doesn't give error
$t1=0x80000000

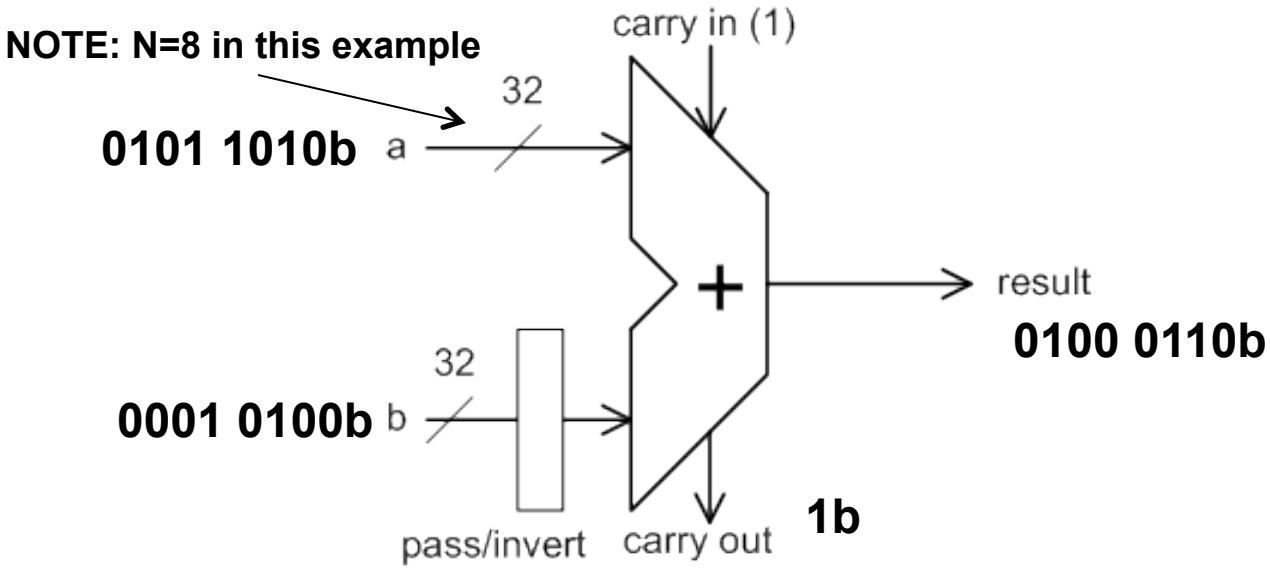
Subtraction

- We know how to add
- We know how to negate a number
- We will use the above two known operations to perform subtraction
- A – B = A + (-B)
- The hardware used for addition can be extended to handle subtraction!



Subtraction

- N=8, a=90, b=20
- Do binary subtraction (A+(-B)) to get result and carryout
- Convert A and B to binary? How?
  - a=90 is 0101 1010b
  - b=20 is 0001 0100b



find –b
invert 0001 0100b
= 1110 1011b
+ 0000 0001b

1110 1100b

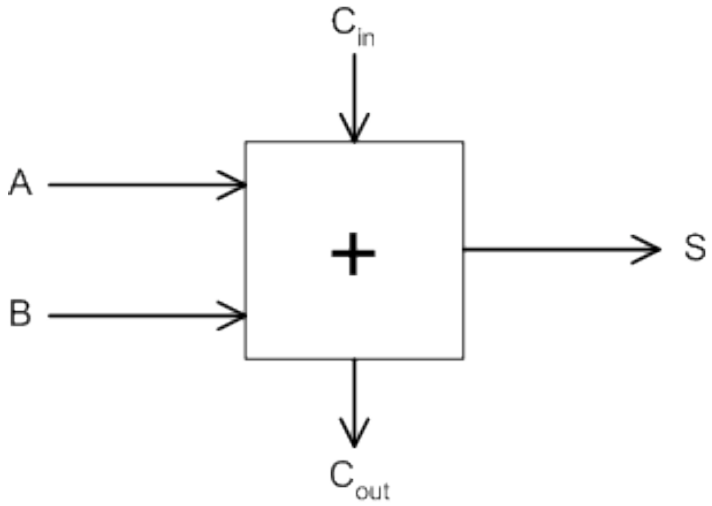
Now, add a
0101 1010b
+ 1110 1100b
---
1 0100 0110b

1-bit adder

- We will look at a single-bit adder
  - Will build on this adder to design a 32-bit adder

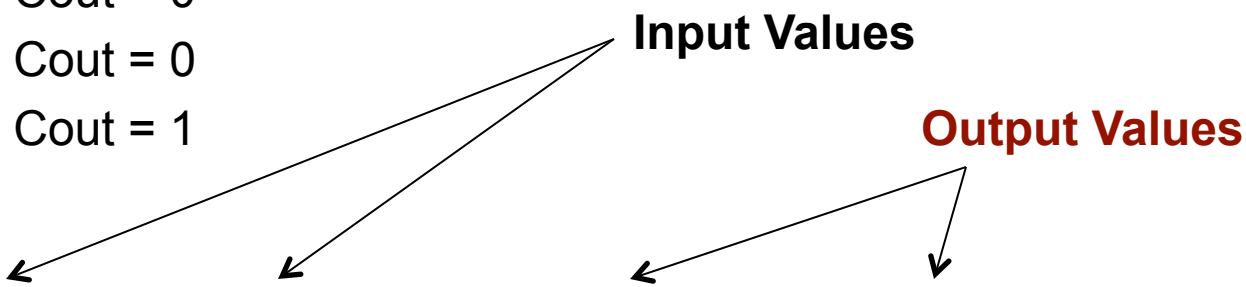
- 3 inputs
  - A: 1st input
  - B: 2nd input
  - C_{in}: carry input

- 2 outputs
  - S: sum
  - C_{out}: carry out



1-bit adder

- What are the binary addition rules?
  - 0 + 0 = 0, Cout = 0
  - 0 + 1 = 1, Cout = 0
  - 1 + 0 = 1, Cout = 0
  - 1 + 1 = 0, Cout = 1



A	B	S	Cout
0	0	**0**	**0**
0	1	**1**	**0**
1	0	**1**	**0**
1	1	**0**	**1**

1-bit adder

- What about Cin?

A	B	Cin	S	Cout
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

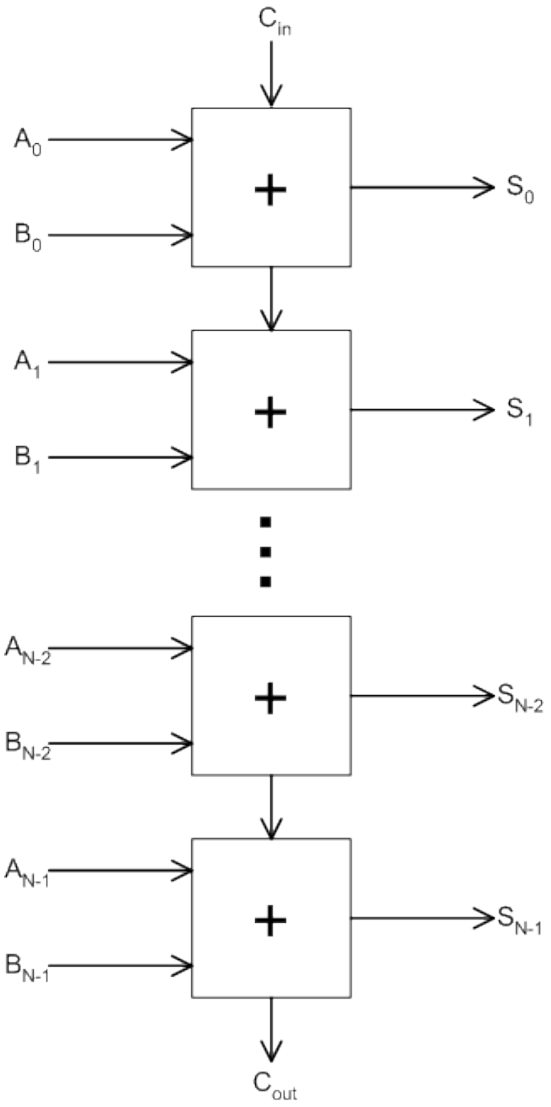
1-bit adder

- What about Cin?

A	B	Cin	S	Cout
0	0	0	**0**	**0**
0	0	1	**1**	**0**
0	1	0	**1**	**0**
0	1	1	**0**	**1**
1	0	0	**1**	**0**
1	0	1	**0**	**1**
1	1	0	**0**	**1**
1	1	1	**1**	**1**

N-bit adder

- An N-bit adder can be constructed with N single-bit adders
  - A carry out generated in a stage is propagated to the next ("ripple-carry adder"
- 3 inputs
  - A: N-bit, 1st input
  - B: N-bit, 2nd input
  - C_{in}: carry input
- 2 outputs
  - S: N-bit sum
  - C_{out}: carry out



N-bit ripple-carry adder