

# Number Theory Coursework

The result proved in this coursework has also been proved using the same method in the Lean theorem prover. This is available at <https://github.com/ChrisHughes24/numbertheory>

## 1 Preliminaries

For this proof one related theorem is required. For any  $n \in \mathbb{N}_{\geq 0}$ , let  $v_p(n)$  be the maximum integer  $k$  such that  $p^k$  divides  $n$ . Then for any integers  $n > 0$ , and  $k > 0$ ,  $v_p\binom{n}{k} \geq v_p(n) - v_p(k)$ .

A theorem by Legendre says that for any  $n$ ,  $v_p(n!) = \sum_{i=0}^{\infty} \left\lfloor \frac{n}{p^i} \right\rfloor$ .

To compute a bound on  $v_p\binom{n}{k}$ , observe that

$$\begin{aligned} v_p\binom{n}{k} &= v_p(n!) - v_p(k!) - v_p((n-k)!) \\ &= \sum_{i=0}^{\infty} \left\lfloor \frac{n}{p^i} \right\rfloor - \left\lfloor \frac{k}{p^i} \right\rfloor - \left\lfloor \frac{n-k}{p^i} \right\rfloor \end{aligned} \tag{1}$$

Observe that if  $p^i$  divides  $n$  but not  $k$  then the corresponding term in the sum is 1, so the sum is at least as big as the number of  $i$  that divide  $n$  but not  $k$ , which is  $v_p(n) - v_p(k)$

## 2 Main proof

**Lemma 1** *If  $m$  is even then  $x + 1$  is a power of 2*

Consider congruences mod  $p$  when  $m$  is even, and  $p$  divides  $x^m + 1$ . Then  $p$  also divides  $x + 1$

$$x^m + 1 \equiv (-1)^m + 1 \equiv 2 \pmod{p} \tag{2}$$

So  $2 = 0 \pmod{p}$  and therefore  $p = 2$  and  $x + 1$  is a power of 2.

**Lemma 2** *If  $m$  is even then  $x^m + 1 \equiv 2 \pmod{4}$*

This is because  $x$  must be odd, so  $x^m$  is an odd square, so  $x^m$  is  $1 \pmod{4}$ .

**Lemma 3**  *$m$  is odd*

Assume  $m$  is even.

By Lemma 1,  $x^m + 1$  is a power of 2. But by Lemma 2  $x^m + 1$  is 2 mod 4. 2 is the only power of 2 congruent to 2 mod 4. So  $x^m + 1 = 2$ , but this contradicts  $x > 1$  and  $m > 1$ . So  $m$  is odd.

**Lemma 4**  $x + 1$  divides  $x^m + 1$

$-1$  is a root of the polynomial  $X^m + 1$  since  $m$  is odd. So by the factor theorem  $x + 1$  divides  $x^m + 1$ .

The idea of the rest of the proof is to prove that  $\frac{x^m+1}{x+1}$  divides  $m$ , and is therefore less than or equal to  $m$ , which puts bounds on  $x$  and  $m$ . The divisibility is proven by proving every prime power that divides  $\frac{x^m+1}{x+1}$  also divides  $m$ .

**Lemma 5** Let  $p$  be a prime and suppose  $p^r$  divides  $m$ . Let  $t$  be a positive integer. Let  $i$  be an integer at least 2. The  $p^{r+t+1}$  divides  $\binom{m}{i}p^{ti}$

If  $p = 2$  then since  $m$  is odd,  $r = 0$ . So  $p^{r+t+1} = p^{t+1}$  which divides  $p^{ti}$ .

Otherwise, we use the fact that  $v_p\left(\binom{m}{i}\right) \geq v_p(m) - v_p(i)$ .

First prove that  $p^{ti-t} > i$ .  $p^{ti-t} = p^t p^{t(i-2)} \geq 3p^{t(i-2)} \geq p^{t(i-2)} + 2 > t(i-2) + 2 \geq i$ , So  $p^{ti-t}$  does not divide  $i$ .

Therefore  $v_p\left(\binom{m}{i}\right) > v_p(m) - (ti - t)$ . So  $v_p\left(\binom{m}{i}p^{ti}\right) \geq v_p(m) + t$ .

**Lemma 6** If  $p$  is prime and  $p^t$  divides  $x + 1$ , but  $p^{t+1}$  does not divide  $x + 1$  and  $p^{s+t}$  divides  $x^m + 1$ , then  $p^s$  divides  $m$ .

If  $t = 0$  then  $s$  must also be equal to zero, since any prime dividing  $x^m + 1$  also divides  $x + 1$ .

First prove that  $p^r$  divides  $m$ , when  $r \leq s$  by induction on  $r$ . We can assume  $t$  is positive.

The case  $r = 0$  is trivial.

Now assume  $p^r$  divides  $m$  and deduce  $p^{r+1}$  divides  $m$ , provided  $r < s$ .

Write  $x = kp^t - 1$  where  $p$  does not divide  $k$ .

Then  $x^m + 1 = (kp^t - 1)^m + 1 = mkp^t - \sum_{i=2}^m \binom{m}{i}(-k)^i p^{it}$ .

Using Lemma 5,  $p^{r+t+1}$  divides the sum. It also divides  $x^m + 1$ , so  $p^{r+t+1}$  must divide  $mkp^t$ , so  $p^{r+1}$  divides  $m$ .

**Lemma 7**  $\frac{x^m+1}{x+1} \leq m$

By Lemma 6, every prime power that divides  $\frac{x^m+1}{x+1}$  also divides  $m$ , So  $\frac{x^m+1}{x+1}$  divides  $m$  and is therefore less than or equal to  $m$ .

**Lemma 8**  $m=3$

Suppose for a contradiction that  $m \geq 4$ . Write  $m = m' + 4$ , where  $m' \in \mathbb{N}_0$ . Similarly write  $x = x' + 2$ .

Now show  $(m' + 4)(x' + 2) + m' + 4 < (x' + 2)^{m'+4} + 1$ .

$$\begin{aligned}
 & (m' + 4)(x' + 2) + x'(m' + 4) \\
 & = (3m' + 12) + x'(m' + 4) \\
 & < (3(x' + 2)^{m'} + 12(x' + 2)^{m'}) + x'((x' + 2)^{m'} + 4(x' + 2)^{m'}) \\
 & \leq (x' + 2)^{m'+4} + 1
 \end{aligned} \tag{3}$$

But by Lemma 7,  $(m' + 4)(x' + 2) + m' + 4 \geq (x' + 2)^{m'+4} + 1$ , so  $m$  cannot be greater than 3.  $m$  is odd and  $m > 1$  so  $m = 3$

**Lemma 9**  $x=2$

Suppose for a contradiction that  $x \geq 3$ . Write  $x = x' = 3$  Then

$$\begin{aligned}
 & mx + m \\
 & = 3x' + 12 \\
 & < x'^3 + 9x'^2 + 27x' + 27 \\
 & = x^m + 1
 \end{aligned} \tag{4}$$

So  $x < 3$ , but  $x > 1$ , so  $x = 2$ .

**Lemma 10** *The solutions are  $x = 2$ ,  $m = 3$  and  $n \geq 2$*

Given  $x = 2$  and  $m = 3$  simple calculations verify that this is a solution if and only if  $n \geq 2$ .