





Status Report: Broken Dependency Orderings in the Linux Kernel

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The "Fear" of Broken Dependencies

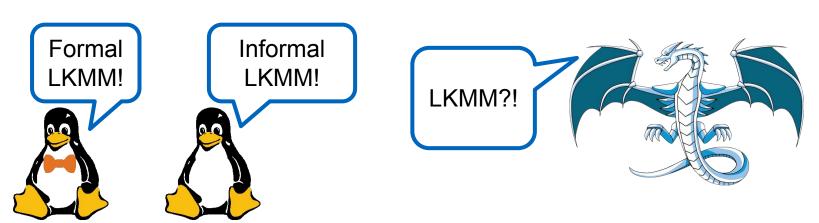
- Linux kernel uses non-standard C.
- Linux-kernel Memory Consistency Model (LKMM) differs from C11 memory model.
- Compilers unaware of LKMM ⇒ potential for miscompilations (?)





The "Fear" of Broken Dependencies

- Linux kernel uses non-standard C.
- Linux-kernel Memory Consistency Model (LKMM) differs from C11 memory model.
- Compilers unaware of LKMM ⇒ potential for miscompilations (?)







Previously on ... Broken Dependency Orderings in the Linux Kernel



"[...] but dammit, I want to see an actual real example arguing for why it would be relevant and why the compiler would need our help."

- Linus Torvalds -





LKMM Dependency Checker

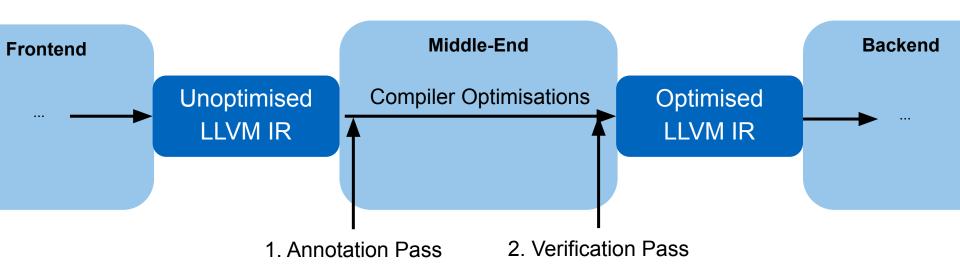
- Annotates and verifies syntactic addr and ctrl dependencies
- Interprocedural analysis up to a given depth
- ClangBuiltLinux







Our Approach







Challenge: Unambiguous and Implementable Dependency Definitions for Analysis of Real Kernel Code





Part 1:

Address Dependencies (and How to Break Them)





But What Is an Address Dependency?

"A read event and another memory access event are linked by an address dependency if the value obtained by the read affects the location accessed by the other event."

- tools/memory-model/Documentation/explanation.txt -





Here's an Address Dependency

```
x = READ_ONCE(foo);
bar = &x[42];
y = READ_ONCE(*bar);
```





```
Is This an Address Dependency? int arr[1];
```

```
x = READ_ONCE(foo);
y = READ_ONCE(arr[x]);
```

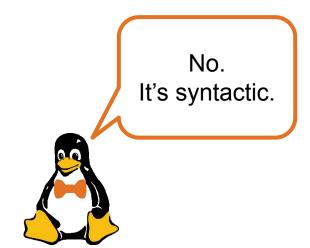




Is This an Address Dependency?

```
int arr[1];
```

```
x = READ_ONCE(foo);
y = READ_ONCE(arr[x]);
```



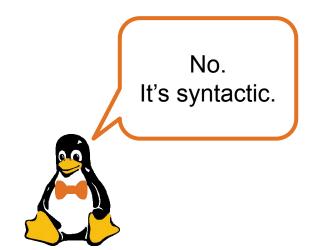




Is This an Address Dependency?

```
int arr[1];
```

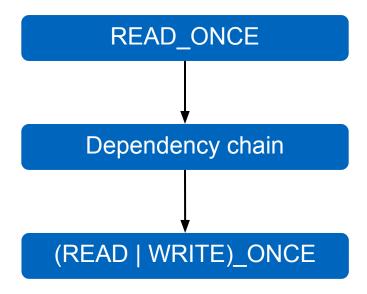
```
x = READ_ONCE(foo);
y = READ_ONCE(arr[0]);
```







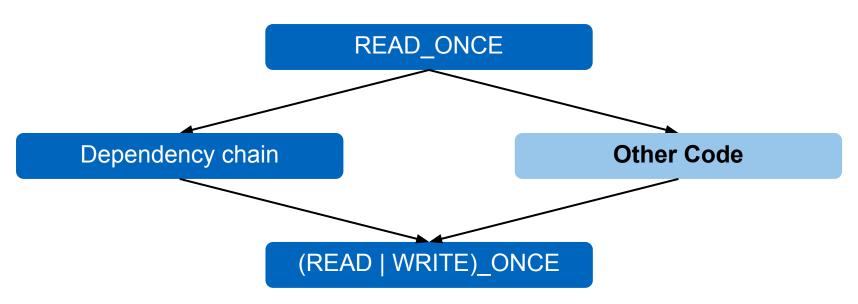
Breaking Address Dependencies







Breaking Address Dependencies - Making the Dep Chain Conditional







Breaking Address Dependencies - Making the Dep Chain Conditional

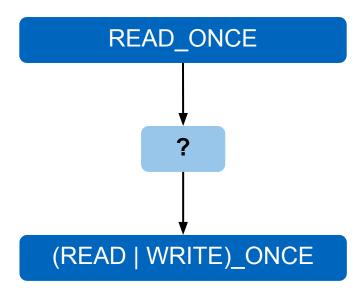
```
x = READ_ONCE(foo);
                         x = READ_ONCE(foo);
                          if (x == baz) /* Oh no! */
bar = &x[42];
y = READ_ONCE(*bar);
                            bar = \&baz[42];
                          else
                            bar = &x[42];
                          y = READ_ONCE(*bar);
```

Source: <u>foo/bar/baz.c::42 - 44</u> (?)





Breaking Address Dependencies - No Dep Chain







Broken Dependency Chain (mm/ksm.c::2032)

```
stable_node = page_stable_node(page); /* READ_ONCE(...) */
if (stable_node)
   /*More code*/
   stable_node->head = &migrate_nodes;
   list_add(&stable_node->list, stable_node->head); /* WRITE_ONCE() */
```

Source: <u>mm/ksm.c::2032 - 2046</u>





```
stable_node = page_stable_node(page); /* READ_ONCE(...) */
if (stable_node)
    /*More code*/
   stable_node->head = &migrate_nodes;
    list_add(&stable_node->list, stable_node->head); /* WRITE_ONCE() */
stable_node = page_stable_node(page); /* READ_ONCE(...) */
if (stable_node)
    /*More code*/
    /* Oh no! */
    list_add(&stable_node->list, &migrate_nodes); /* WRITE_ONCE() */
```





Transformation to Conditional (fs/nfs/delegation.c::617 - 622)

```
if (place_holder)
    delegation = rcu_dereference([...]); /* READ_ONCE() */
if (!delegation || delegation != place_holder_deleg)
    delegation = list_entry_rcu([...]);
for([...], delegation = list_entry_rcu(delegation, [...])) /* READ_ONCE() */
```

Source: fs/nfs/delegation.c 617 - 622



```
if(place_holder == NULL) {
    delegation = list_entry_rcu([...]);
else {
    cmp = rcu_dereference(place_holder->delegation); /* READ_ONCE(...) */
    if(cmp != NULL) {
         if(cmp == place_holder_deleg) /* Oh no! */
             delegation = place_holder_deleg; /* Oh no! (cont.) */
        else
             delegation = list_entry_rcu([...]);
    } else {
        delegation = list_entry_rcu([...]);
for([...], delegation = list_entry_rcu(delegation, [...])) /* READ_ONCE() */
```





Challenge: Unambiguous and Implementable Dependency Definitions for Analysis of Real Kernel Code





Part 2:

Control Dependencies (and How to Break Them)





But What Is a Control Dependency?

"Finally, a read event X and a write event Y are linked by a control dependency if Y syntactically lies within an arm of an if statement and X affects the evaluation of the if condition via a data or address dependency (or similarly for a switch statement)"

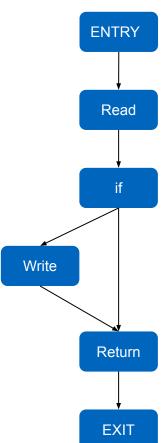
- tools/memory-model/Documentation/explanation.txt (in a few weeks) -





Here's a Control Dependency

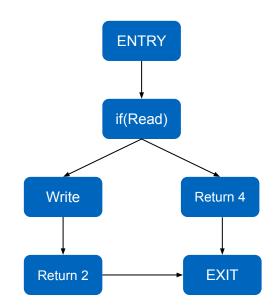
```
x = READ_ONCE(foo);
if(x == 42)
  WRITE_ONCE(bar, 42);
return true;
                       Agreed.
```







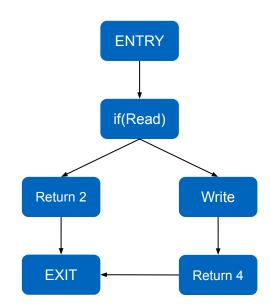
```
if(READ_ONCE(x))
    return 4;
WRITE_ONCE(y, 21);
return 2;
```





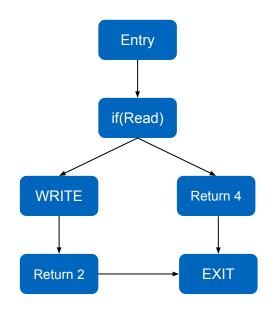


```
if(!READ_ONCE(x)) {
    WRITE_ONCE(y, 42);
    return 2;
}
return 4;
```





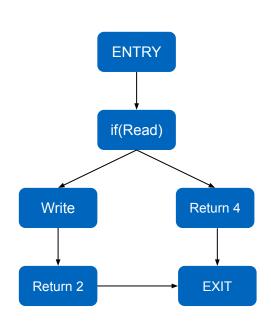








```
if(READ_ONCE(x))
   return 42;
WRITE_ONCE(y, 42);
/* The answer has to be Yes! */
return 0;
                    No.
                                Yes!
```







But What Is a Control Dependency?

"Let G be a control flow graph. Let X and Y be nodes in G.

Y is control dependent on X iff

(1) there exists a directed path P from X to Y with any Z in P

(excluding X and Y) post-dominated by Y and

(2) X is not post-dominated by Y."





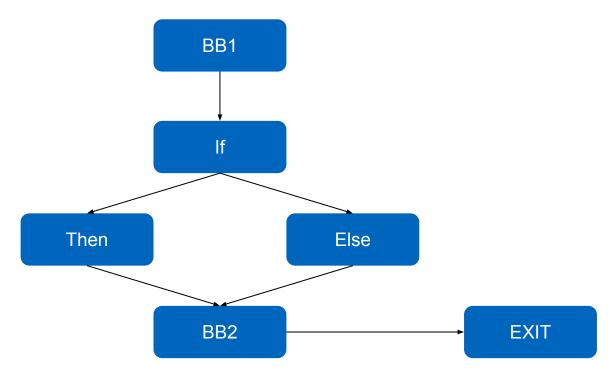
But What Is Post Dominance?

"A node V is post-dominated by a node W in G
if every directed path
from V to STOP (not including V)
contains W."





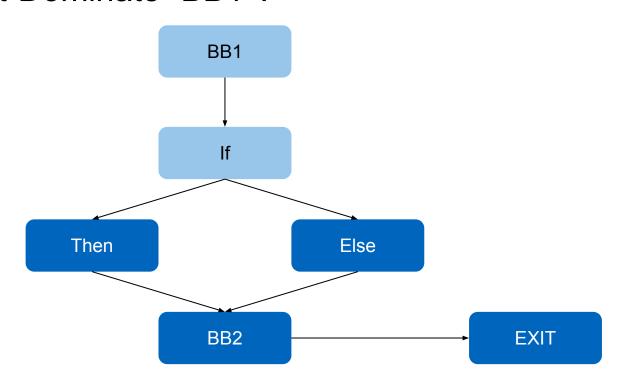
Illustrating Post-Dominance







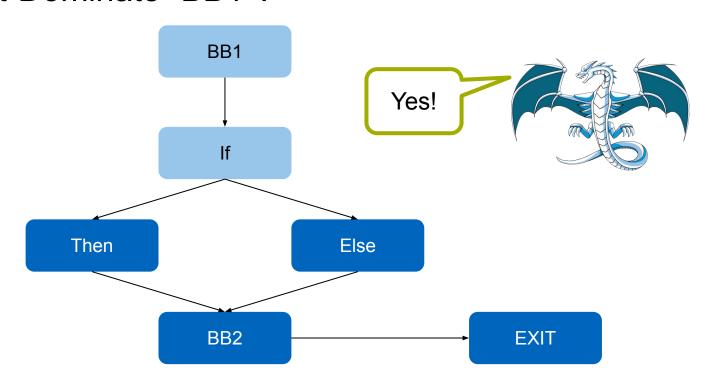
Does "If" Post-Dominate "BB1"?







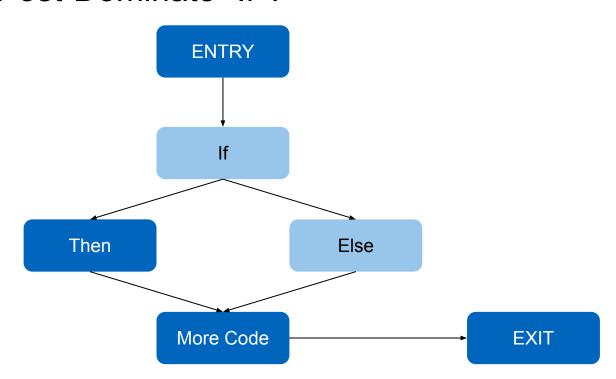
Does "If" Post-Dominate "BB1"?







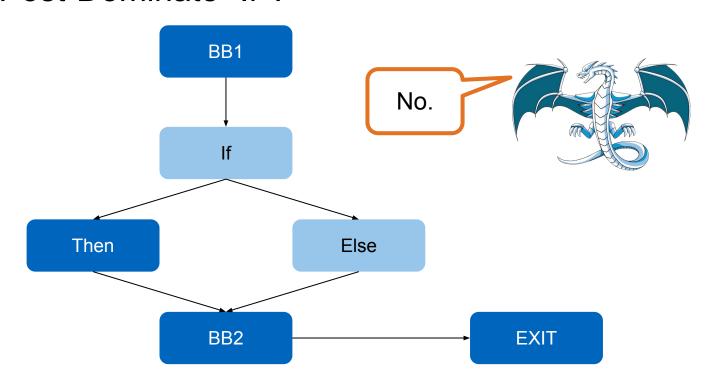
Does "Then" Post-Dominate "If"?







Does "Then" Post-Dominate "If"?



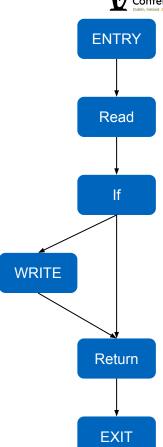




Here's a Control Dependency

```
x = READ_ONCE(foo);
if(x == 42)
    WRITE_ONCE(bar, 24);
return true;
```

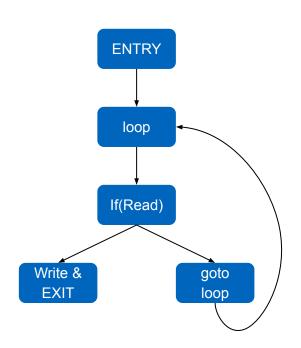








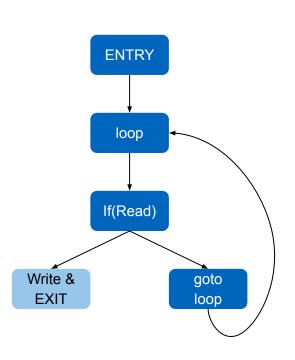
```
loop:
if(READ_ONCE(x)) {
  WRITE_ONCE(y, 42);
   return 0;
goto loop;
```







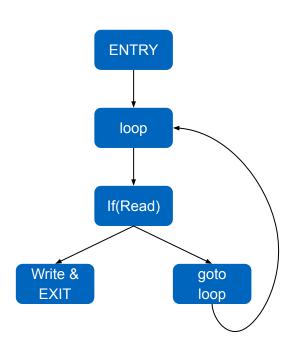
```
loop:
if(READ_ONCE(x)) {
   WRITE_ONCE(y, 42);
   return 0;
                No.
goto loop;
```







```
loop:
                                Yes!
if(READ_ONCE(x))
   WRITE_ONCE(y, 42);
   return 0;
goto loop;
```







```
ENTRY
loop:
                                        Yes!
if(READ_ONCE(x))
                                                        dool
   WRITE_ONCE(y, 42); /* Yes! */
                                                       If(Read)
    return 0;
                                                 Write &
                                                              goto
                                                 EXIT
                                                              dool
                    No.
goto loop;
```

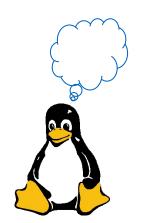




But What Is a Control Dependency?

There is a control dependency from a marked read A to a marked write B if there is a condition C s.t. B is within the scope of C and C depends on A.

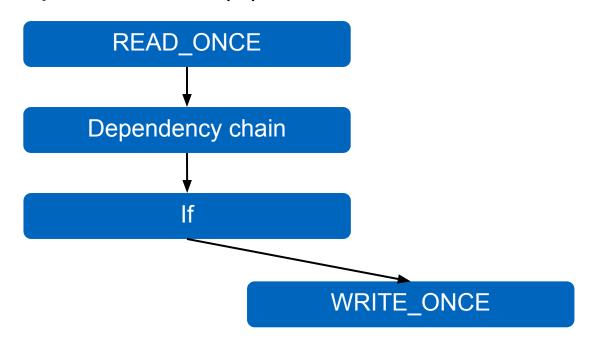
- WIP Control Dependency Definition -







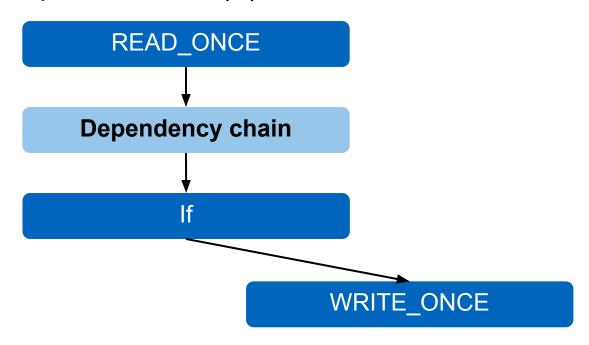
Breaking Control Dependencies (?)







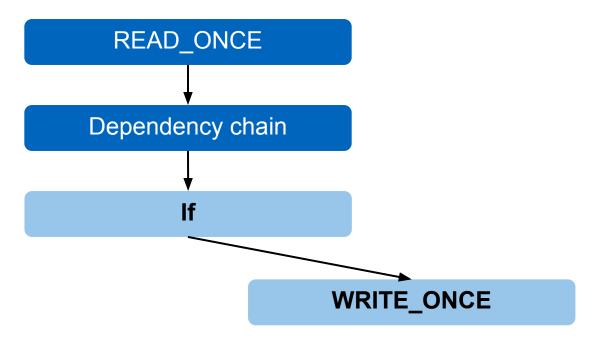
Breaking Control Dependencies (?)







Breaking Control Dependencies (?)

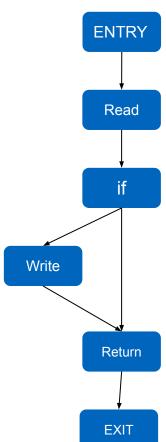






```
#define MAX 1
```

```
x = READ_ONCE(*foo);
if (x % MAX == 0)
    WRITE_ONCE(*bar, 1);
Return true;
```







```
#define MAX 1

x = READ_ONCE(foo);
/* No branch?! */
WRITE_ONCE(*bar, 1);
Return true;
```

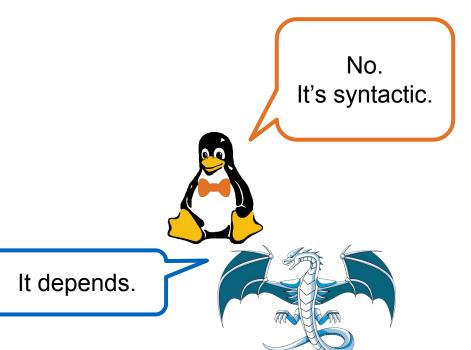






#define MAX 1

```
x = READ_ONCE(*foo);
if (x % MAX == 0)
WRITE_ONCE(*bar, 1);
```





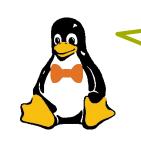


```
x = READ_ONCE(*bar);
if (x == 42)
    WRITE_ONCE(*baz, 1);
else
    WRITE_ONCE(*baz, 2);
WRITE_ONCE(*y, 42);
```





```
x = READ_ONCE(*bar);
if (x == 42)
   /* arm64 compiler can make this a conditional select */
   WRITE_ONCE(*baz, 1);
else
   WRITE_ONCE(*baz, 2);
WRITE_ONCE(*y, 42);
```



Plural - Control Dependencies! There is more than one.

It depends.







Where Does this Leave Us?

"To What Extent Are Compilers Undermining the Linux Kernel Memory Model?"

- LKMM addr dependency checker has found broken dependencies
- LKMM addr dependency checker LLVM RFC hopefully very soon
- WIP LKMM ctrl dependency checker



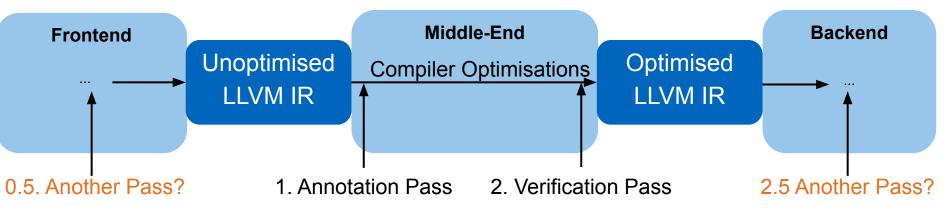
Broken dependencies are real! The LKMM Dependency Checker found broken dependencies!





Limitations

- Syntactic dependency !=> semantic dependency
- Computational feasibility (interprocedural analysis, variadic functions ...)
- Our analysis avoids backedges (right now)
- Frontend and backend optimisations not accounted for







The Hunt for Broken Dependencies - Future Plans

"To What Extent Are Compilers Undermining the Linux Kernel Memory Model?"

- AArch64 defconfigs for various Linux kernel flavours (incl. LTO)
- Random testing with random, but relevant, Linux kernel configs
- LLVM RFCs
- Find ctrl to addr dependency transformations
- Extend our analysis to the backend
- Get the ball rolling on short-term to long-term fixes with sufficient evidence

Get in touch!



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The Hunt for Broken Dependencies - Discussion

"To What Extent Are Compilers Undermining the Linux Kernel Memory Model?"

- Annotating dependency chains?
- -memory-model=lkmm option for compilers?
- Combating false positives:
 - "safe regions" via pragmas or function attributes?
 - Exclusion lists?
- Suggesting fixes to the user? Automatically insert barriers?



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