# Return Value Optimization

Harder Than It Looks

#### **Outline**

- What is the "return slot"? [3–34]
- Rules of thumb for RVO [35–37]
- How RVO actually works in standard C++ [38–45]
- Test cases and war stories [46–54]
- Writing a new Clang diagnostic [55–72]

Hey look! Slide numbers!

#### What is the "return slot"?

```
int apple()
{
    return 42;
}
int pear()
{
    return 1 + apple();
}
```

Let's start with the simplest possible function.

```
int apple()
{
    return 42;
}
int pear()
{
    return 1 + apple();
}
```

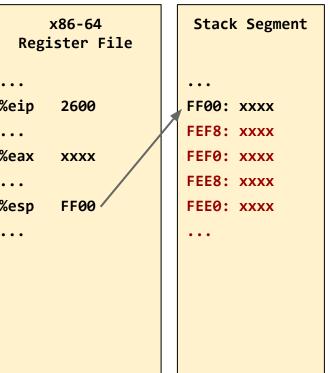
```
clang -00 -fomit-frame-pointer -S

_Z5applev:
   movl $42, %eax
   retq

_Z4pearv:
   callq _Z5applev
   addl $1, %eax
   retq
```

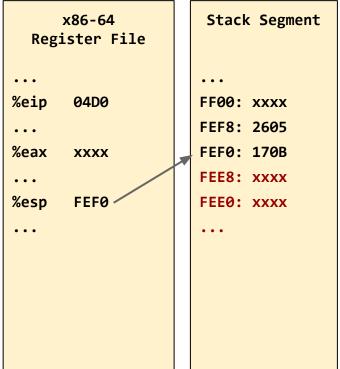
On x86-64, the function's return value usually goes into the %eax register.

```
int apple()
                                               Register File
                        Z5applev:
                       mov1 $42, %eax
                 04d0:
  return 42;
                                             %eip
                                                   2600
                 04d5:
                       reta
                                             %eax
                                                   XXXX
                        Z4pearv:
                                             . . .
int pear()
                       callq Z5applev
                 1706:
                                                   FF00
                                             %esp
                 170b: addl $1, %eax
  return 1 +
                 170e: retq
    apple();
                 2600:
                         callq Z4pearv
   eip = 2600
                 2605:
                         do something else
```



```
x86-64
                                                                    Stack Segment
int apple()
                                                   Register File
                          Z5applev:
                          movl $42, %eax
                   04d0:
  return 42;
                                                 %eip
                                                       1706
                                                                   FF00: xxxx
                   04d5:
                          retq
                                                                   FEF8: 2605
                                                 %eax
                                                                   FEF0: xxxx
                                                       XXXX
                          Z4pearv:
                                                                   FEE8: XXXX
                                                 . . .
                   1706:
                          callq Z5applev
    eip = 1706
                                                       FEF8
                                                 %esp
                                                                   FEE0: xxxx
                   170b: addl $1, %eax
                                                                   . . .
  return 1 +
                   170e: retq
    apple();
                   2600:
                            callq Z4pearv
                            do something else
                   2605:
```

```
int apple()
                       Z5applev:
                        movl $42, %eax
                 04d0:
   eip = 04D0
                                            %eip
                 04d5:
                        retq
                                            %eax
                        Z4pearv:
                                             . . .
int pear()
                 1706:
                       callq Z5applev
                                            %esp
                 170b: addl $1, %eax
  return 1 +
                 170e: retq
    apple();
                 2600:
                         callq Z4pearv
                         do something else
                 2605:
```



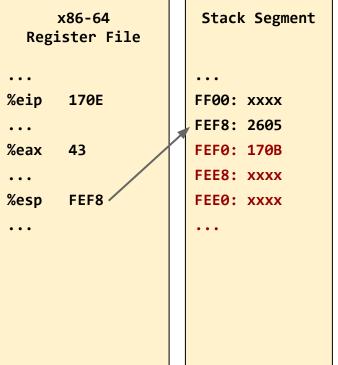
```
int apple()
                                                Register File
                        Z5applev:
                         movl $42, %eax
                  04d0:
                                              %eip
                                                   04D5
                  04d5:
   eip = 04D5
                         retq
                                              %eax
                                                   42
                        Z4pearv:
                                              . . .
int pear()
                  1706:
                        callq Z5applev
                                                    FEF0
                                              %esp
                  170b: addl $1, %eax
  return 1 +
                  170e: retq
    apple();
                  2600:
                          callq Z4pearv
                  2605:
                          do something else
```

x86-64 Stack Segment FF00: xxxx FEF8: 2605 FEF0: 170B FEE8: xxxx FEE0: xxxx . . .

```
int apple()
                                                Register File
                        Z5applev:
                        movl $42, %eax
                  04d0:
  return 42;
                                              %eip
                                                   170B
                  04d5:
                        reta
                                              %eax
                                                   42
                        Z4pearv:
                                              . . .
int pear()
                  1706:
                        callq Z5applev
                                                    FEF8
                                              %esp
   eip = 170B
                  170b: addl $1, %eax
                                                               . . .
  recurn I +
                  170e: retq
    apple();
                  2600:
                          callq Z4pearv
                          do something else
                  2605:
```

x86-64 Stack Segment FF00: xxxx FEF8: 2605 FEF0: 170B FEE8: XXXX FEE0: xxxx

```
int apple()
                       Z5applev:
                       movl $42, %eax
                 04d0:
  return 42;
                                            %eip
                 04d5:
                       retq
                                            %eax
                       Z4pearv:
                                            . . .
int pear()
                 1706:
                       callq Z5applev
                                            %esp
                 170b: addl $1, %eax
   eip = 170E
                 170e: retq
    apple();
                 2600:
                         callq Z4pearv
                         do something else
                 2605:
```



```
int apple()
                      Z5applev:
                      movl $42, %eax
                04d0:
  return 42;
                04d5:
                      retq
                       Z4pearv:
int pear()
                1706:
                      callq Z5applev
                170b: addl $1, %eax
  return 1 +
                170e: retq
    apple();
                2600:
                        callq Z4pearv
                2605:
                        do something else
   eip = 2605
```

```
x86-64
                        Stack Segment
  Register File
%eip
       2605
                       FF00: xxxx
                       FEF8: 2605
       43
%eax
                       FEF0: 170B
                      FEE8: xxxx
. . .
       FF00
%esp
                       FEE0: xxxx
                       . . .
```

Now we're back where we started, except that %eip is 2605 instead of 2600, and %eax is correctly 43!

This is the magic of a *call stack*.

# But what do we do with big objects?

```
struct Fruit {
    int data[5]; // 20 bytes
};
Fruit apple()
    return Fruit{1,2,3,4,5}; // Can't return this in %eax!
... apple() ...
```

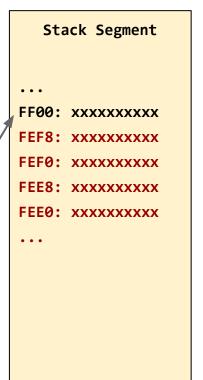
#### But what do we do with big objects?

```
struct Fruit {
    int data[5]; // 20 bytes
};
Fruit apple()
    return Fruit{1,2,3,4,5};
... apple() ...
```

```
clang -00 -fomit-frame-pointer -S
Z5applev:
  movq %rdi, %rax
  movl $1, (%rdi)
                            For big objects,
  mov1 $2, 4(%rdi)
                           the compiler adds
  mov1 $3, 8(%rdi)
                           a hidden function
  mov1 $4, 12(%rdi)
                             parameter: a
  mov1 $5, 16(%rdi)
                              "return slot
  reta
                           address," passed
                               in %rdi.
  subq $24, %rsp
  movq %rsp, %rdi
  callq _Z5applev
  # now %rsp[0..20] hold our answer
```

```
Z5applev:
Fruit apple()
                     04d0:
                           movq %rdi, %rax
                     04d3: movl $1, (%rdi)
  return Fruit{
                     04d9: movl $2, 4(%rdi)
                     04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                     04e7: movl $4, 12(%rdi)
  };
                           mov1 $5, 16(%rdi)
                     04ee:
                     04f5:
                           reta
                     2600:
                           subq $24, %rsp
                     2604:
                           movq %rsp, %rdi
                           callq Z5applev
                     2607:
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
 Register File
%eip
       2600
%rdi
       XXXX
. . .
       FF00
%esp
```



```
Z5applev:
Fruit apple()
                           movq %rdi, %rax
                     04d0:
                     04d3: movl $1, (%rdi)
  return Fruit{
                     04d9: movl $2, 4(%rdi)
                     04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                     04e7: movl $4, 12(%rdi)
  };
                           movl $5, 16(%rdi)
                     04ee:
                     04f5:
                           reta
                     2600:
                            subq $24, %rsp
     eip = 2604
                     2604:
                           movq %rsp, %rdi
                           callq Z5applev
                     2607:
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
                       Stack Segment
 Register File
%eip
       2604
                    FF00: XXXXXXXXXX
                    FEF8: XXXXXXXXXX
%rdi
                    FEF0: XXXXXXXXXX
       XXXX
                    FEE8: XXXXXXXXXX
. . .
%esp
       FEE8
                    FEE0: XXXXXXXXXX
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                           movq %rdi, %rax
                     04d0:
                     04d3: movl $1, (%rdi)
  return Fruit{
                     04d9: movl $2, 4(%rdi)
                     04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                     04e7: movl $4, 12(%rdi)
  };
                           movl $5, 16(%rdi)
                     04ee:
                     04f5:
                           reta
                     2600:
                           subq $24, %rsp
  apple()
                     2604:
                           movq %rsp, %rdi
                     2607: callq Z5applev
... | eip = 2607
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
                        Stack Segment
 Register File
%eip
       2607
                     FF00: XXXXXXXXXX
                     FEF8: XXXXXXXXXX
%rdi
       FEE8
                     FEF0: XXXXXXXXXX
                     FEE8: XXXXXXXXXX
. . .
%esp
       FEE8 -
                     FEE0: XXXXXXXXXX
. . .
                     . . .
```

```
Z5applev:
                         movq %rdi, %rax
                  04d0:
   eip = 04D0
                  04d3: movl $1, (%rdi)
return Fruit{
                  04d9: movl $2, 4(%rdi)
                  04e0: movl $3, 8(%rdi)
  1,2,3,4,5
                  04e7: movl $4, 12(%rdi)
};
                         movl $5, 16(%rdi)
                  04ee:
                  04f5:
                         reta
                  2600:
                         subq $24, %rsp
apple()
                  2604: movq %rsp, %rdi
                  2607: callq Z5applev
                         # examine %rsp[0..20]
                  260c:
                     . . .
```

```
x86-64
                       Stack Segment
 Register File
%eip
       04D0
                    FF00: XXXXXXXXXX
                    FEF8: XXXXXXXXXX
%rdi
       FEE8
                    FEF0: XXXXXXXXXX
                    FEE8: XXXXXXXXXX
. . .
%esp
       FEE0
                    FEE0: 000000260C
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                    04d0:
                           movq %rdi, %rax
     eip = 04D3
                    04d3: movl $1, (%rdi)
  re-carminate
                    04d9: movl $2, 4(%rdi)
                    04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                    04e7: movl $4, 12(%rdi)
  };
                           movl $5, 16(%rdi)
                    04ee:
                    04f5:
                           reta
                    2600:
                           subq $24, %rsp
  apple()
                    2604:
                           movq %rsp, %rdi
                    2607: callq Z5applev
                           # examine %rsp[0..20]
                    260c:
                        . . .
```

```
x86-64
                       Stack Segment
 Register File
%eip
       04D3
                    FF00: XXXXXXXXXX
                    FEF8: XXXXXXXXXX
%rdi
       FEE8
                    FEF0: XXXXXXXXXX
                    FEE8: XXXXXXXXXX
. . .
%esp
       FEE0
                    FEE0: 000000260C
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                     04d0:
                           movq %rdi, %rax
                     04d3: movl $1, (%rdi)
  receip = 04D9
                     04d9: movl $2, 4(%rdi)
                     04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                     04e7: movl $4, 12(%rdi)
  };
                            mov1 $5, 16(%rdi)
                     04ee:
                     04f5:
                            reta
                     2600:
                            subq $24, %rsp
  apple()
                     2604:
                           movq %rsp, %rdi
                           callq Z5applev
                     2607:
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
                        Stack Segment
 Register File
%eip
       04D9
                    FF00: XXXXXXXXXX
                    FEF8: XXXXXXXXXX
%rdi
       FEE8
                     FEF0: XXXXXXXXXX
                    FEE8: xx00000001
. . .
%esp
       FEE0
                    FEE0: 000000260C
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                           movq %rdi, %rax
                     04d0:
                     04d3: movl $1, (%rdi)
  return Fruit{
                     04d9: movl $2, 4(%rdi)
                     04e0: movl $3, 8(%rdi)
     eip = 04E0
                     04e7: movl $4, 12(%rdi)
  };
                           mov1 $5, 16(%rdi)
                     04ee:
                     04f5:
                           reta
                     2600:
                           subq $24, %rsp
  apple()
                     2604: movq %rsp, %rdi
                     2607: callq Z5applev
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
                       Stack Segment
 Register File
%eip
       04E0
                    FF00: XXXXXXXXXX
                    FEF8: XXXXXXXXXX
%rdi
       FEE8
                    FEF0: XXXXXXXXXX
                    FEE8: 0200000001
. . .
%esp
       FEE0
                    FEE0: 000000260C
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                           movq %rdi, %rax
                     04d0:
                     04d3: movl $1, (%rdi)
  return Fruit{
                     04d9: movl $2, 4(%rdi)
                     04e0: movl $3, 8(%rdi)
  }; eip = 04E7
                     04e7: movl $4, 12(%rdi)
                           mov1 $5, 16(%rdi)
                     04ee:
                     04f5:
                           reta
                     2600:
                           subq $24, %rsp
  apple()
                     2604:
                           movq %rsp, %rdi
                           callq Z5applev
                     2607:
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
                        Stack Segment
 Register File
%eip
       04E7
                    FF00: XXXXXXXXXX
                    FEF8: XXXXXXXXXX
%rdi
       FEE8
                     FEF0: xx00000003
                    FEE8: 0200000001
. . .
%esp
       FEE0
                    FEE0: 000000260C
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                           movq %rdi, %rax
                     04d0:
                     04d3: movl $1, (%rdi)
  return Fruit{
                    04d9: movl $2, 4(%rdi)
                    04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                    04e7: movl $4, 12(%rdi)
                           mov1 $5, 16(%rdi)
                    04ee:
                     04f5:
                           reta
                     2600:
                           subq $24, %rsp
  apple()
                     2604:
                           movq %rsp, %rdi
                           callq Z5applev
                     2607:
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
                        Stack Segment
 Register File
%eip
       04EE
                    FF00: XXXXXXXXXX
                    FEF8: XXXXXXXXXX
%rdi
       FEE8
                     FEF0: 0400000003
                    FEE8: 0200000001
. . .
%esp
       FEE0
                    FEE0: 000000260C
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                     04d0:
                           movq %rdi, %rax
                     04d3: movl $1, (%rdi)
  return Fruit{
                     04d9: movl $2, 4(%rdi)
                     04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                     04e7: movl $4, 12(%rdi)
  };
                           movl $5, 16(%rdi)
                     04ee:
     eip = 04F5
                     04f5:
                           reta
                     2600:
                           subq $24, %rsp
  apple()
                     2604:
                           movq %rsp, %rdi
                           callq Z5applev
                     2607:
                           # examine %rsp[0..20]
                     260c:
                        . . .
```

```
x86-64
                        Stack Segment
 Register File
%eip
       04F5
                    FF00: XXXXXXXXXX
                    FEF8: xx00000005
%rdi
       FEE8
                     FEF0: 0400000003
                    FEE8: 0200000001
. . .
%esp
       FEE0
                    FEE0: 000000260C
. . .
                     . . .
```

```
Z5applev:
Fruit apple()
                           movq %rdi, %rax
                     04d0:
                     04d3: movl $1, (%rdi)
  return Fruit{
                    04d9: movl $2, 4(%rdi)
                    04e0: movl $3, 8(%rdi)
     1,2,3,4,5
                    04e7: movl $4, 12(%rdi)
  };
                           mov1 $5, 16(%rdi)
                    04ee:
                    04f5:
                           reta
                     2600:
                           subq $24, %rsp
  apple()
                     2604:
                           movq %rsp, %rdi
                           callq Z5applev
                     2607:
                           # examine %rsp[0..20]
                     260c:
     eip = 260C
```

Now we're back where we started, except %eip is 260C instead of 2607, and (FEE8 .. FEFC) have been correctly initialized to {1,2,3,4,5}.

Notice how we told apple where to write its output; it trusted us.

The *caller* owns the return slot.

# Moving into the return slot

```
struct Fruit {
    int data[5];
    Fruit(Fruit&&);
};
Fruit apples and oranges(bool condition)
    Fruit x = ...;
    Fruit y = ...;
    return std::move(condition ? x : y);
```

#### Moving into the return slot

```
struct Fruit {
    int data[5];
    Fruit(Fruit&&);
};
Fruit apples and oranges(bool condition)
    Fruit x = ...;
    Fruit y = ...;
    return std::move(condition ? x : y);
```

```
Stack Segment
 Return Slot
  (20 bytes)
Return Address
   Fruit x
  (20 bytes)
   Fruit y
  (20 bytes)
```

# Copy elision

```
struct Fruit {
    int data[5];
};

Fruit nothing_but_apples()
{
    Fruit x = ...;
    return x;
}
```

#### Copy elision

```
struct Fruit {
    int data[5];
};

Fruit nothing_but_apples()
{
    Fruit x = ...;
    return x;
}
```

```
_Z18nothing_but_applesv:
    subq $24, %rsp
    movq %rdi, %rax
    movl $1, (%rsp)
    movl $2, 4(%rsp)
    movl $3, 8(%rsp)
    movl $4, 12(%rsp)
    movl $5, 16(%rsp)
    movq (%rsp), (%rdi)
    movq 8(%rsp), 8(%rdi)
    movl 16(%rsp), 16(%rdi)
    addq $24, %rsp
    retq
```

This code (slightly altered for presentation) constructs Fruit x at %rsp[0..20] and then copies it into the return slot at %rdi[0..20].

#### Stack Segment

Return Slot (20 bytes)

Return Address

Fruit x (20 bytes)

. . .

#### Copy elision

But we can do better! Here we construct Fruit x directly in the return slot at %rdi[0..20].

#### Stack Segment

Return Slot, also x (20 bytes)

**Return Address** 

• • •

#### Copy elision rules in C++

- In C++03 through C++14, copy elision was *permitted* in many circumstances, which we'll discuss in a moment.
- In C++17, copy elision became *mandatory* in many circumstances.
- But there are cases where copy elision is not always possible...

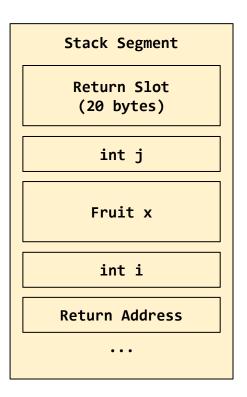
#### When can't we elide?

```
Fruit apples_to_apples(int i, Fruit x, int j)
{
    return x;
}
```

In this (slightly altered) example, the caller passes in Fruit x at one stack address, and the return slot at a different stack address.

We must get the data out of x and into the return slot somehow!

We can't elide the copy because we don't control x's physical location.



#### When can't we elide?

```
static Fruit x;
Fruit apples_to_apples()
{
    return x;
}
```

Reductio ad absurdum:
We can't elide this copy because we don't control x's physical location.

Data Section

Fruit x (20 bytes)

Stack Segment

Return Slot (20 bytes)

Return Address

. . .

# One more important case...



# Slicing to base class

```
struct Durian : Fruit {
    double smell;
};

Fruit slapchop()
{
    Durian x = ...;
    return x;
}
```



We can't elide this copy because, while we do control x's physical location, x is of the "wrong" type for constructing into the return slot.

#### Stack Segment

Return Slot (20 bytes)

**Return Address** 

Durian x (32 bytes)

. . .

#### Rules of thumb for RVO

Here's what I used to tell people when they asked about RVO:

 "Unnamed RVO" (URVO): Returning a temporary (a prvalue) will trigger copy-elision.

```
return Fruit{1,2,3,4,5};
return my_helper_function();
```

• "Named RVO" (NRVO): Returning a local variable "by name" will trigger copy-elision, except in the corner cases we've covered.

```
return x;
```

And even if copy-elision doesn't happen, implicit move happens...

#### Rules of thumb for RVO

Here's what I used to tell people when they asked about RVO:

• "Implicit move": When returning a local variable "by name" doesn't trigger copy-elision, the compiler's overload resolution will still automatically treat the name x as an rvalue!

```
std::string identity(std::string x) {
    return x; // x will be implicitly moved-from, not copied!
}
```

 Because of C++11's "implicit move," writing return std::move(x) is almost always a pessimization — it never helps, and it might hurt by disabling NRVO.



## Story time!

```
extern "C" {
    typedef enum { CEK NONE, CEK INT, CEK STRING, CEK LIST } cexpr kind;
    typedef struct cexpr cexpr;
    cexpr *cexpr new int(int value);
    cexpr *cexpr new string(const char *data);
    cexpr *cexpr new empty list();
    cexpr *cexpr clone(cexpr *self);
    cexpr *cexpr free(cexpr *self);
    cexpr kind cexpr get kind(cexpr *self);
    int cexpr int value(cexpr *self);
    int cexpr list length(cexpr *self);
    void cexpr list append(cexpr *self, cexpr *v);
```

### C++ wrapper around C API

```
class Cexpr {
protected:
   cexpr *impl ;
   Cexpr(cexpr *p) : impl (p) {}
public:
   Cexpr(const Cexpr& rhs) : impl (cexpr clone(rhs.impl )) {}
   Cexpr(Cexpr&& rhs) : impl (rhs.impl ) { rhs.impl = nullptr; }
   ~Cexpr() { cexpr free(impl ); }
};
class CexprInt : public Cexpr {
public:
    CexprInt(int value) : impl_(cexpr_new_int(value)) {}
    int value() const { return cexpr int value(impl ); }
};
```

### C++ wrapper around C API

```
class CexprList : public Cexpr {
public:
   CexprList() : impl (cexpr new empty list()) {}
    int length() const { return cexpr list length(impl ); }
   void append(const Cexpr& v) {
        cexpr *p = cexpr clone(v.impl );
        cexpr list append(impl , p);
   void append(Cexpr&& v) {
        cexpr *p = std::exchange(v.impl , nullptr);
        cexpr list append(impl , p);
```

## C++ functions that return Cexpr

```
class ConfigManager {
   virtual Cexpr as cexpr() const = 0;
};
class ConfigManagerImpl : public ConfigManager {
    Cexpr as cexpr() const override {
        CexprList cfg;
        cfg.append(CexprInt(17));
        cfg.append(CexprInt(42));
        cfg.append(CexprString("hike"));
        return cfg;
```

### C++ functions that return Cexpr

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```

Now QA determines that such-and-such an operation is taking much longer than it should. We profile the code and discover that we're spending an awful lot of cycles in cexpr\_clone().

(Not literally in this config code, but somewhere using exactly the same pattern of construction and return.)

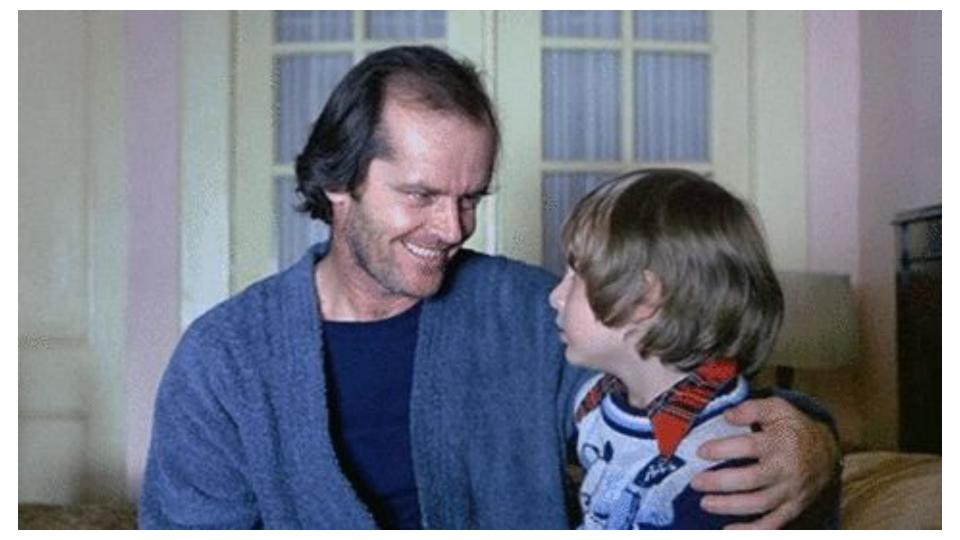
### C++ functions that return Cexpr

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(Not literally in this config code, but somewhere using exactly the same pattern of construction and return.)

"But you told me that return x would treat x as an rvalue! Why is it copying?"



## The actual rules of "implicit move"

Before 2013, the rules were slightly different — but even the C++11 rules were posthumously amended by Defect Report CWG1579.

[class.copy.elision] /3 If the expression in a return statement is a (possibly parenthesized) id-expression that names an object with storage duration declared in automatic the body parameter-declaration-clause of the innermost enclosing function or lambda-expression [... then ...] overload resolution to select the constructor for the copy is first performed as if the object were designated by an rvalue. If the first overload resolution fails or was not performed, or if the type of the first parameter of the selected constructor is not an rvalue reference to the object's type (possibly cv-qualified), overload resolution is performed again, considering the object as an Ivalue.

## What went wrong here? Slicing.

```
Cexpr as_cexpr() const override {
    CexprList cfg;
    cfg.append(CexprInt(17));
    cfg.append(CexprInt(42));
    cfg.append(CexprString("hike"));
    return cfg;
}
```

The constructor that we want to call here is Cexpr(Cexpr&&).

But cfg's type is CexprList, not Cexpr!

"If the first overload resolution fails or was not performed, or if the type of the first parameter of the selected constructor is not an rvalue reference to *the object's type* (possibly cv-qualified), overload resolution is performed again, considering the object as an Ivalue."

The slicing here is 100% intentional; but it still silently disables the optimization we thought the compiler was giving us.

## War story #2: inplace\_function

```
template<class R, class... Args, size t Capacity, size t Alignment>
                                                                       "Converting constructor": I know
class inplace function<R(Args...), Capacity, Alignment>
                                                                         how to construct myself from
                                                                         the source type, stealing his
   vtable t *vptr ;
    alignas(Alignment) char storage [Capacity];
                                                                              guts if necessary.
public:
    template<size_t SrcCap, size_t SrcAlign>
    inplace_function(inplace_function<R(Args...), SrcCap, SrcAlign> const & rhs)
        requires(Capacity >= SrcCap && Alignment >= SrcAlign)
        vptr = rhs.vptr ;
        vptr ->copy(storage , rhs.storage );
    template<size t SrcCap, size t SrcAlign>
    inplace_function(inplace_function<R(Args...), SrcCap, SrcAlign>&& rhs)
        requires(Capacity >= SrcCap && Alignment >= SrcAlign)
        vptr_ = std::exchange(rhs.vptr_, empty_function_vptr);
        vptr ->move(storage , rhs.storage );
```

#### Maintainer proposed this instead:

```
template<class R, class... Args, size t Capacity, size t Alignment>
class inplace function<R(Args...), Capacity, Alignment>
    inplace_function(auto vptr, auto func, void *rhs_storage) : vptr_(vptr) {
       func(storage , rhs storage);
public:
   template<size t DstCap, size t DstAlign>
    operator inplace_function<R(Args...), DstCap, DstAlign>() const &
        requires(DstCap >= Capacity && DstAlign >= Alignment)
       return {vptr_, vptr_->copy, storage_};
   template<size t DstCap, size t DstAlign>
    operator inplace function<R(Args...), DstCap, DstAlign>() &&
        requires(DstCap >= Capacity && DstAlign >= Alignment)
        auto vptr = std::exchange(vptr_, empty_function vptr);
        return {vptr, vptr->move, storage_};
```

"Conversion operator": I know how to convert myself to the destination type, lending him my own guts if necessary.

This seems to match our mental model a bit better.

### What happens here?

```
using IPF32 = inplace function<void(), 32>;
using IPF64 = inplace function<void(), 64>;
IPF32 get upstream callback() {
    return [msg = std::string(...)](){
        std::cout << msg << std::endl;</pre>
IPF64 ConfigManagerImpl::callback() {
    auto result = get_upstream_callback();
    assert(!result.empty());
    return result;
                                     Suddenly we're spending a lot of time in
                                       std::string's copy constructor!
```

## What went wrong here? Converting.

```
using IPF32 = inplace function<void(), 32>;
using IPF64 = inplace function<void(), 64>;
IPF32 get upstream callback() {
    return [msg = std::string(...)](){
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IPF64 ConfigManagerImpl::callback() {
    auto result = get upstream callback();
    assert(!result.empty());
    return result;
```

Here we convert IPF32 to IPF64.

This works fine before the patch:
 overload resolution finds
 IPF64::IPF64(IPF32&&).

 After the patch, it finds
 IPF32::operator IPF64()&&
 instead

"If the first overload resolution fails or was not performed, or if the type of the first parameter of the **selected constructor** is not an rvalue reference to the object's type (possibly cv-qualified), overload resolution is performed again, considering the object as an Ivalue."

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```
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    return [msg = std::string(...)](){
        std::cout << msg << std::endl;</pre>
IPF64 ConfigManagerImpl::callback() {
    auto result = get upstream callback():
    assert(!result.empty());
                                  For extra fun, GCC
    return result;
```

Here we convert IPF32 to IPF64. This works fine before the patch:

overload resolution finds IPF64::IPF64(IPF32&&).

After the patch, it finds IPF32::operator IPF64()&& instead.

"If the first overload resolution fails or was not performed, or if the type of the first parameter of the **selected**constructor is not an rvalue reference to the object's type

for extra fun, GCC (nossibly cv-qualified), overload and ICC do move (not resolution is performed again, copy) in this case. considering the object as an Ivalue."

### Why are constructors so important?

Why does [class.copy.elision] /3 put such emphasis on finding constructors that take *rvalue references to exact class types?* What's so important about this particular special case that they decided to bake it into the standard? And not even bake it into C++14, but actually retroactively insert the same rule into C++11!?

## Why are constructors so important?

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```
std::unique_ptr<ConfigManager> create() {
    auto p = std::make_unique<ConfigManagerImpl>();
    return p;
}
```

## Why are constructors so important?

Why does [class.copy.elision] /3 put such emphasis on finding constructors that take *rvalue references to exact class types?* What's so important about this particular special case that they decided to bake it into the standard? And not even bake it into C++14, but actually retroactively insert the same rule into C++11!?

The C++ Standard Library *loooves* move-enabled converting constructors.

All of these examples use explicitly declared converting constructors. So when they dealt with this issue in the core language, they didn't think about any of the other ways you could get implicit conversion (slicing, conversion operators, or combinations of these).

- Can we get the compiler to warn us when we write return x but return std::move(x) would have been more efficient?
- We can definitely do the opposite. **clang++ -Wmove** will warn about places we wrote return std::move(x) when return x would have been more efficient.

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- We can definitely do the opposite. **clang++ -Wmove** will warn about places we wrote return std::move(x) when return x would have been more efficient.

- As of this writing, Clang does *not* have a warning for the opposite, sneakier case, when we wrote return x expecting a move but got a copy instead.
- So I went and added one!

First I tried adding a check to clang-tidy.

This great blog post got me off to a good start:

http://bbannier.github.io/blog/2015/05/02/Writing-a-basic-clang-static-analysis-check.html

```
Protip: git clone --depth=1 http://llvm.org/git/llvm.git
```

But clang-tidy is basically just for syntactic issues. It can do simple things such as observing that an implicit conversion has resolved to Foo(const Bar&) when Foo(Bar&&) exists, but it cannot "re-evaluate" the entire overload resolution step to see whether Foo(Bar&&) would actually have been a better match, or whether Bar::operator Foo()&& would have been selected, etc.

So I abandoned clang-tidy and went straight for a full-blown clang diagnostic.

Adding the new diagnostic was easy because I could look at existing commits and pull requests that added new diagnostics, and copy them.

Especially D7633 which added -Wpessimizing-move.

- Step 1: Add the machinery for the command-line option.
- Step 2: Add a test case.
- Step 3: Actually add the new code to libSema.

## Step 1: Add command-line option

```
+++ b/include/clang/Basic/DiagnosticGroups.td
@@ -380,7 +380,11 @@
def ExplicitInitializeCall : DiagGroup<"explicit-initialize-call">;
def Packed : DiagGroup<"packed">;
def Padded : DiagGroup<"padded">;
def PessimizingMove : DiagGroup<"pessimizing-move">;
+def ReturnStdMove : DiagGroup<"return-std-move">;
def PointerArith : DiagGroup<"return-std-move">;
def PoundWarning : DiagGroup<"#warnings">;
def PoundPragmaMessage : DiagGroup<"#pragma-messages">,
```

## Step 1: Add command-line option

```
+++ b/include/clang/Basic/DiagnosticSemaKinds.td
@@ -5626,6 +5626,19 @@
 def note remove move : Note<"remove std::move call here">;
+def warn return std move : Warning<</pre>
+ "local variable %0 will be copied despite being "
+ "%select{returned|thrown}1 by name">,
   InGroup<ReturnStdMove>, DefaultIgnore;
+def note add std move : Note<
   "call 'std::move' explicitly to avoid copying">;
 def warn string plus int : Warning<</pre>
   "adding %0 to a string does not append to the string">,
   InGroup<StringPlusInt>;
```

It took me several iterations to tease out the syntax of the diagnostic engine's mini-language.

Notice that the number "1" goes *outside* these curly braces.

Not shown: the even weirder syntax "%diff{ (\$ vs \$)|}1,2"

### Step 2: Add the test case

```
+++ b/test/SemaCXX/warn-return-std-move.cpp
@@ -0,0 +1,327 @@
+// RUN: %clang_cc1 -fsyntax-only -Wreturn-std-move -std=c++11 -verify %s
+// RUN: %clang cc1 -fsyntax-only -Wreturn-std-move -std=c++11 \
    -fdiagnostics-parseable-fixits %s 2>&1 | FileCheck %s
                                      The comment // e1 is just for my own benefit when compiling
+struct Base { stuff };
                                      the test by hand: I number the expected warnings e1, e2, e3...
+struct Derived : public Base {};
                                      and then it's easy to see which one is missing from the output.
+Base test() {
                                       The expected-warning, expected-note, CHECK, and RUN
     Derived d2;
                                      comments are actually meaningful to Clang's testing harness.
     return d2; // e1
     // expected-warning@-1{{will be copied despite being returned by name}}
     // expected-note@-2{{to avoid copying}}
     // CHECK: fix-it:"{{.*}}":{[[@LINE-3]]:12-[[@LINE-3]]:14}:"std::move(d2)"
                                       expected-note@-2{{text}} means "expect to see a note
                                         emitted 2 lines up from here, containing at least text."
+and so on
```

#### Step 3: Add the code

First, I found the code that implemented copy elision. This was easy because Clang has really good Doxygen-style comments. Here is the actual function header cut-and-pasted from lib/Sema/SemaStmt.cpp:

#### Step 3: Add the code

Here is what the code looked like before my patch (modulo some software engineering):

```
ExprResult Sema::PerformMoveOrCopyInitialization(const InitializedEntity &Entity,
   const VarDec1 *NRVOCandidate, QualType ResultType, Expr *Value, bool AllowNRVO) {
 ExprResult Res = ExprError();
 if (AllowNRVO) {
   if (!NRVOCandidate)
     NRVOCandidate = getCopyElisionCandidate(ResultType, Value, CES Default);
   if (NRVOCandidate)
     AttemptMoveInitialization(*this, Entity, NRVOCandidate, ResultType, Value, false, Res);
 // Either we didn't meet the criteria for treating an lvalue as an rvalue,
 // above, or overload resolution failed. Either way, we need to try
 // (again) now with the return value expression as written.
 if (Res.isInvalid())
   Res = PerformCopyInitialization(Entity, SourceLocation(), Value);
 return Res:
```

## Step 3: Add the code

Res = PerformCopyInitialization(Entity, SourceLocation(), Value);

And after my patch:

if (Res.isInvalid())

return Res;

```
that it exists, we can warn the user.
if (AllowNRVO) {
  if (!NRVOCandidate)
    NRVOCandidate = getCopyElisionCandidate(ResultType, Value, CES Default);
  if (NRVOCandidate)
    AttemptMoveInitialization(*this, Entity, NRVOCandidate, ResultType, Value, false, Res);
  if (Res.isInvalid()) {
    auto *FakeCandidate = getCopyElisionCandidate(QualType(), Value, CES_AsIfByStdMove);
    if (FakeCandidate) {
      ExprResult FakeRes = ExprError();
      AttemptMoveInitialization(*this, Entity, FakeCandidate, ResultType, Value, true, FakeRes);
      if (!FakeRes.isInvalid()) {
        bool IsThrow = (Entity.getKind() == InitializedEntity::EK Exception);
        Diag(Value->getExprLoc(), diag::warn_return_std_move)
            << Value->getSourceRange()
            << FakeCandidate->getDeclName() << IsThrow;
                                                                                 Recall the format of our diagnostic: "local
                                                                              variable %0 will be copied despite being
                                                                                 %select{returned|thrown}1 by name".
                                                                                         Here are our %0 and %1.
// Either we didn't meet the criteria for treating an lvalue as an rvalue, ...
```

If no acceptable constructor was found, we go on and look for an *unacceptable* 

constructor or conversion operator, via super lenient rules.

We won't *use* FakeRes, but if we find

#### Step 3b: Add more code

Remember that we were going to have "fixits." How do I generate a fixit note?

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Done! Easy!

#### Step 3c: Add more code

We don't want to warn if the "copy constructor" that was found is trivial, because that's just a memcpy — it's already as optimal as it can be. Also, we don't want to look stupid by accidentally suggesting std::move() around an int.

#### Step 3c: Add more code

We don't want to warn if the "copy constructor" that was found is trivial, because that's just a memcpy — it's already as optimal as it can be. Also, we don't want to look stupid by accidentally suggesting std::move() around an int.

```
const VarDecl *FakeCandidate =
    getCopyElisionCandidate(QualType(), Value, CES AsIfByStdMove);
if (FakeCandidate) {
 const QualType QT = FakeCandidate->getType();
 if (QT.getNonReferenceType().getUnqualifiedType()
            .isTriviallyCopyableType(Context)) {
   // Don't suggest 'std::move' around a trivially copyable variable.
  } else {
    ExprResult FakeRes = ExprError();
   AttemptMoveInitialization(*this, Entity, FakeCandidate,
                              ResultType, Value, true, FakeRes);
    if (!FakeRes.isInvalid()) {
     // emit the diagnostic
```

#### Step 3d: Add more code

Oh, and don't suggest std::move in this case either...

```
String test(String *p) {
     String &x = *p;
     return x; // std::move(x) would be so very wrong here!
if (FakeCandidate) {
 const QualType QT = FakeCandidate->getType();
 if (QT.getNonReferenceType().getUnqualifiedType()
            .isTriviallyCopyableType(Context)) {
   // Don't suggest 'std::move' around a trivially copyable variable.
  } else {
   ExprResult FakeRes = ExprError();
   AttemptMoveInitialization(*this, Entity, FakeCandidate,
                             ResultType, Value, true, FakeRes);
```

#### Step 3d: Add more code

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```
String test(String *p) {
     String &x = *p;
     return x; // std::move(x) would be so very wrong here!
if (FakeCandidate) {
 const QualType QT = FakeCandidate->getType();
 if (QT->isLValueReferenceType()) {
   // Don't suggest 'std::move' around an lvalue reference.
  } else if (QT.getNonReferenceType().getUnqualifiedType()
            .isTriviallyCopyableType(Context)) {
   // Don't suggest 'std::move' around a trivially copyable variable.
  } else {
   ExprResult FakeRes = ExprError();
   AttemptMoveInitialization(...
```

# Step 4: Ship it!



I'm working on this part.

https://reviews.llvm.org/D43322

The nice wording of the warning message is mostly thanks to Richard Smith.

