

Compacting the Uncompactable!

MESH

Automatically Compacting
Your C++ Application's Memory

Emery Berger
UMass Amherst

with Bobby Powers, David Tench, & Andrew McGregor
University of Massachusetts Amherst

<http://libmesh.org>

[PLDI 2019]



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN

Music by MITCH LEIGH



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



Reconquer all of Spain!



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

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Reconquer
all of Spain!



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Reconquer Allocate
all of Spain!



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(malloc)



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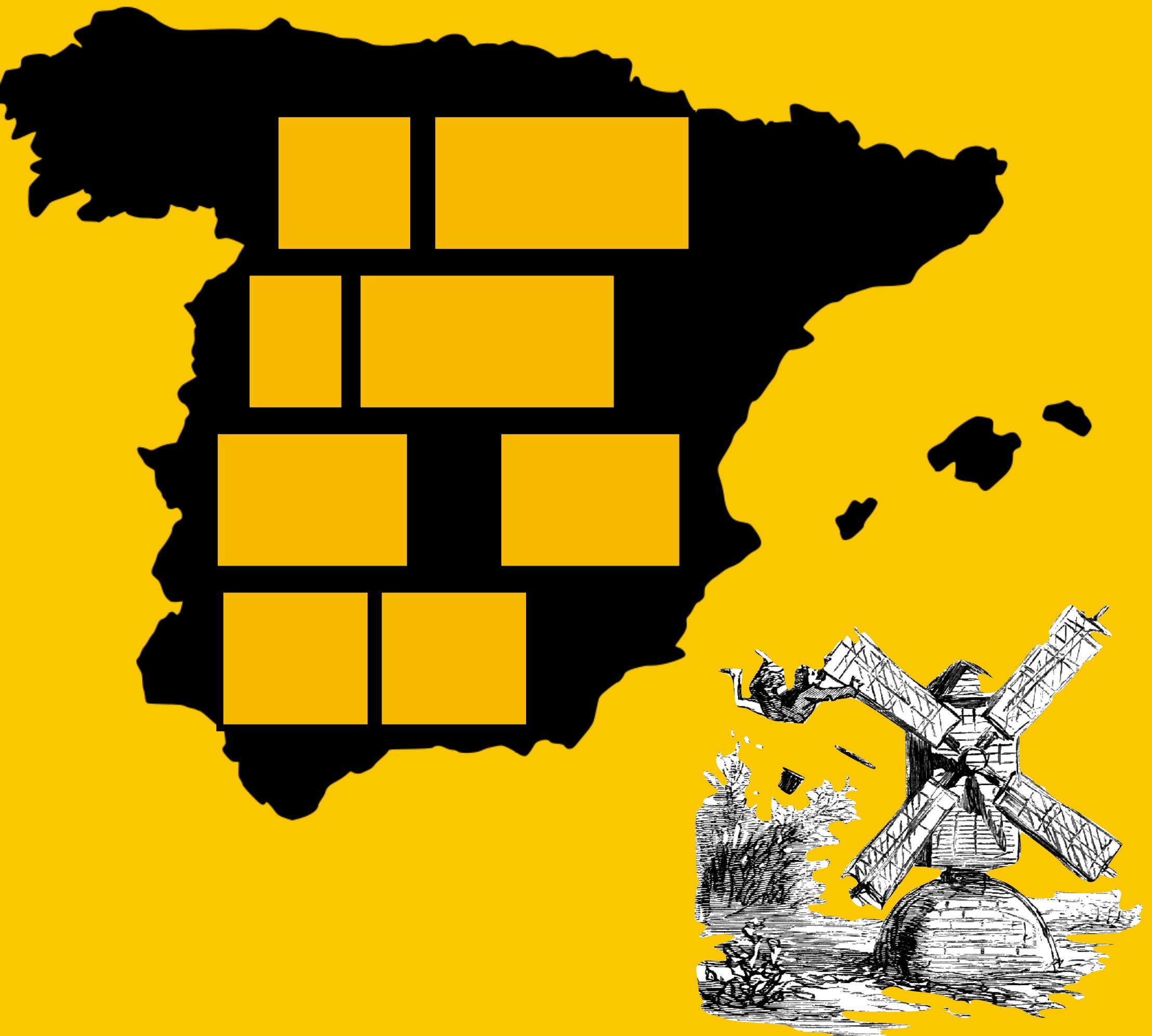


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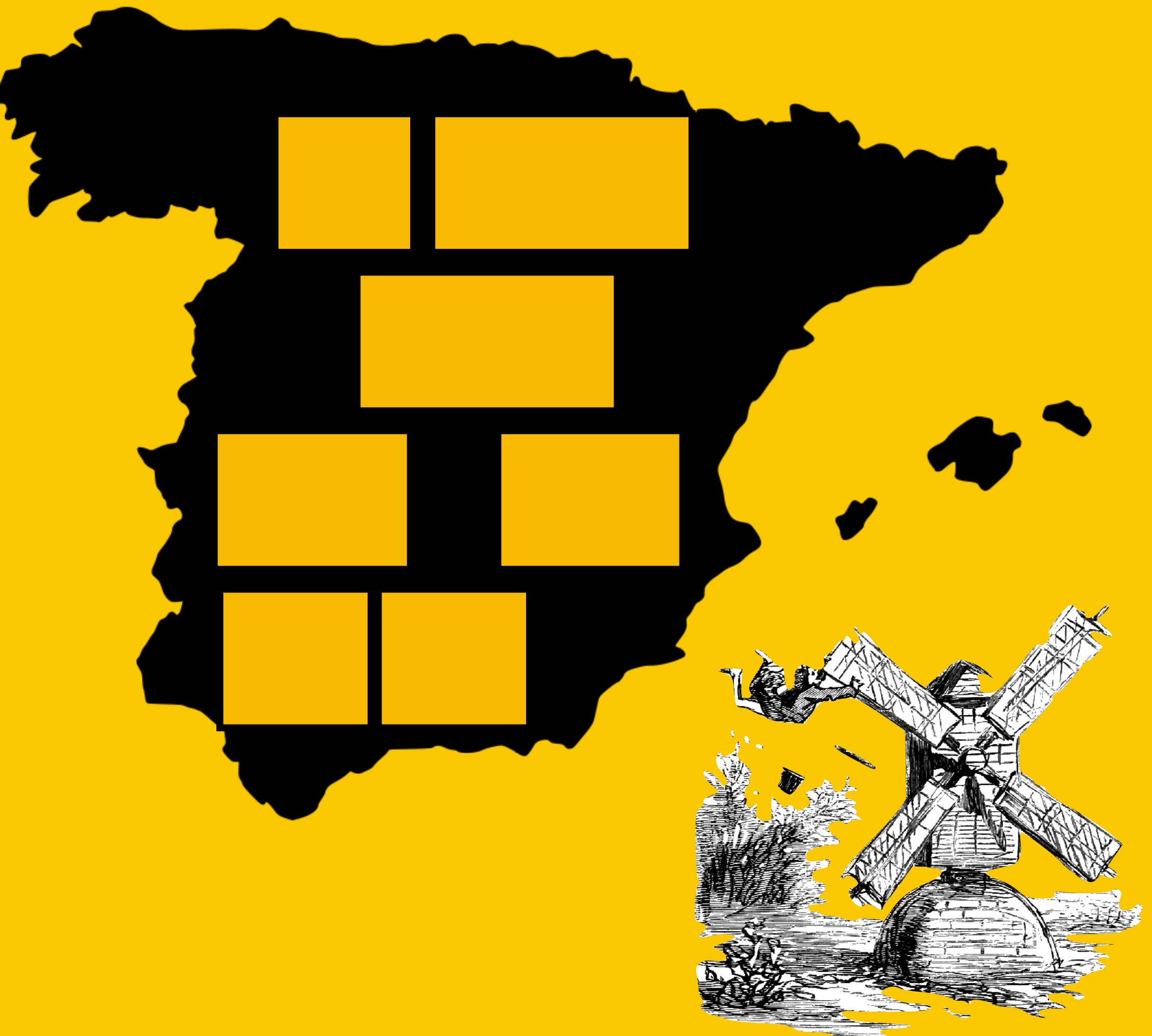


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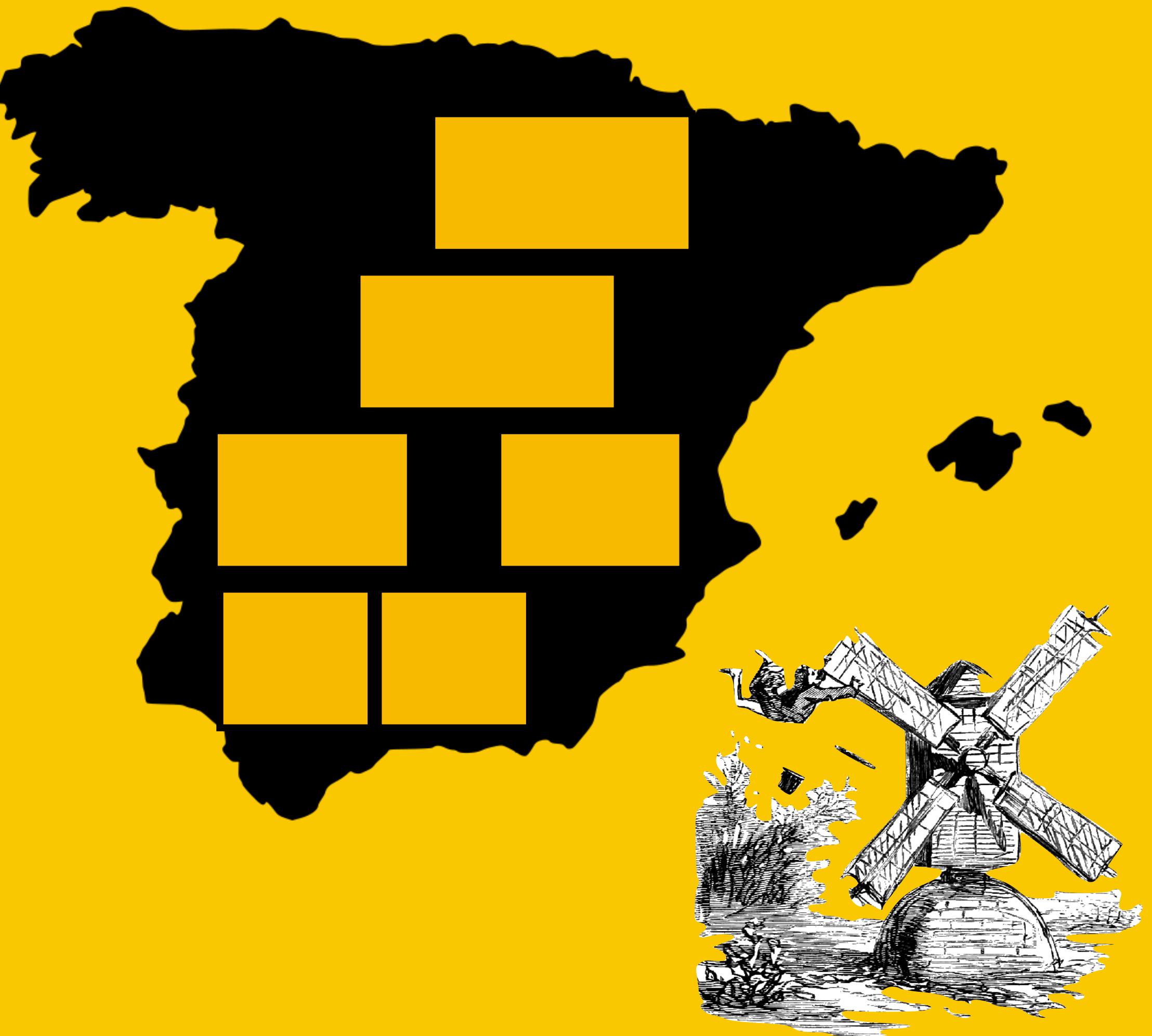
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?

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!



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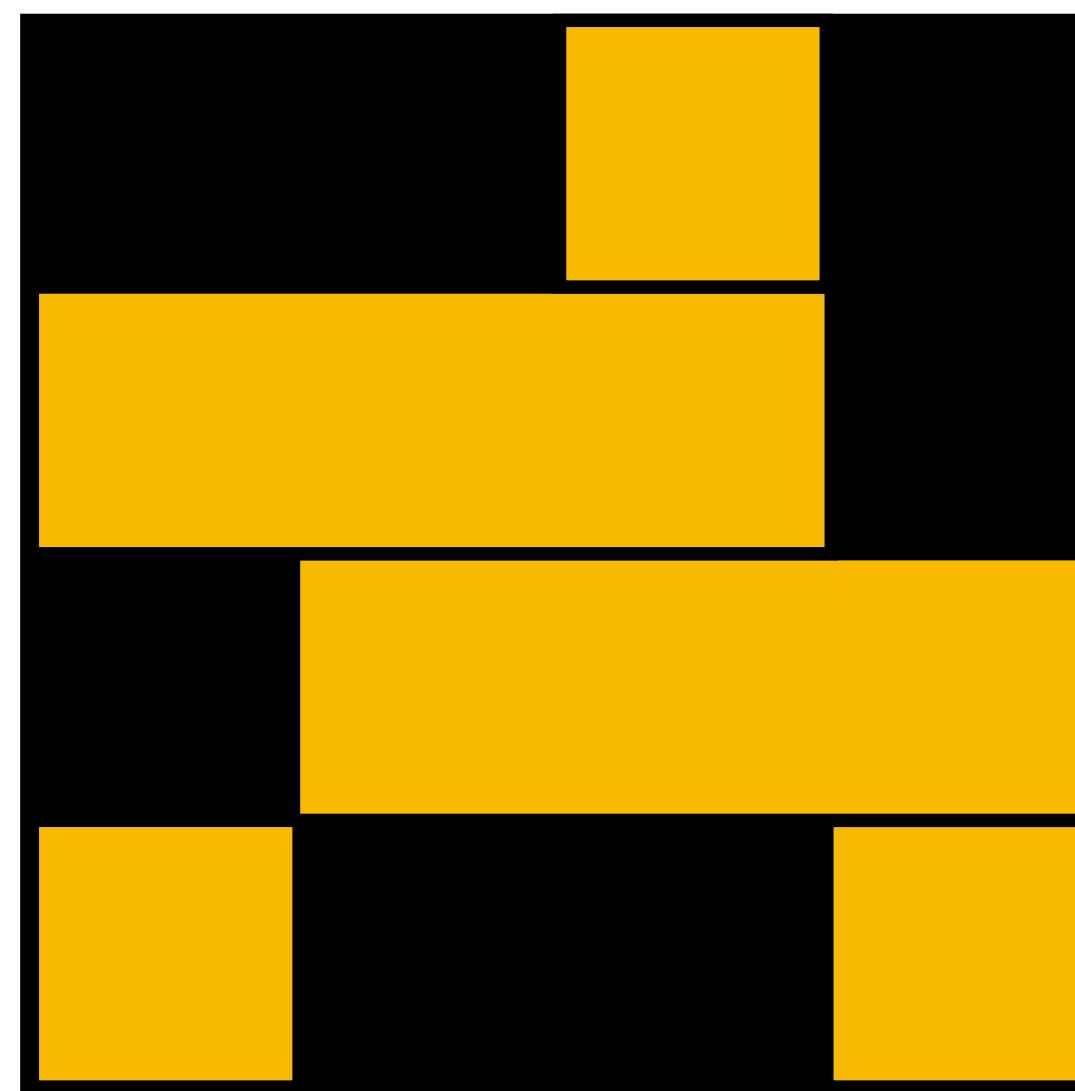
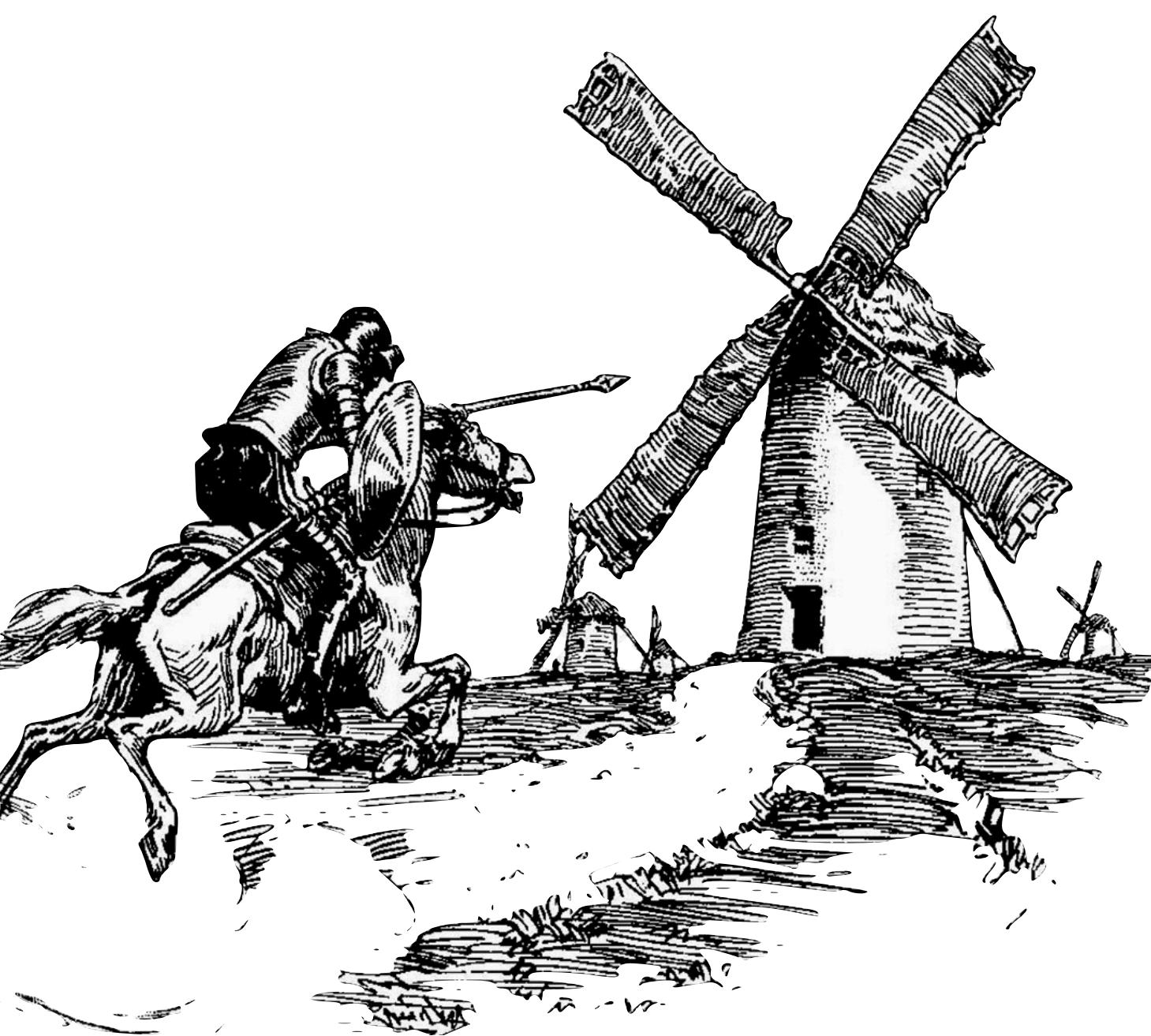


i!





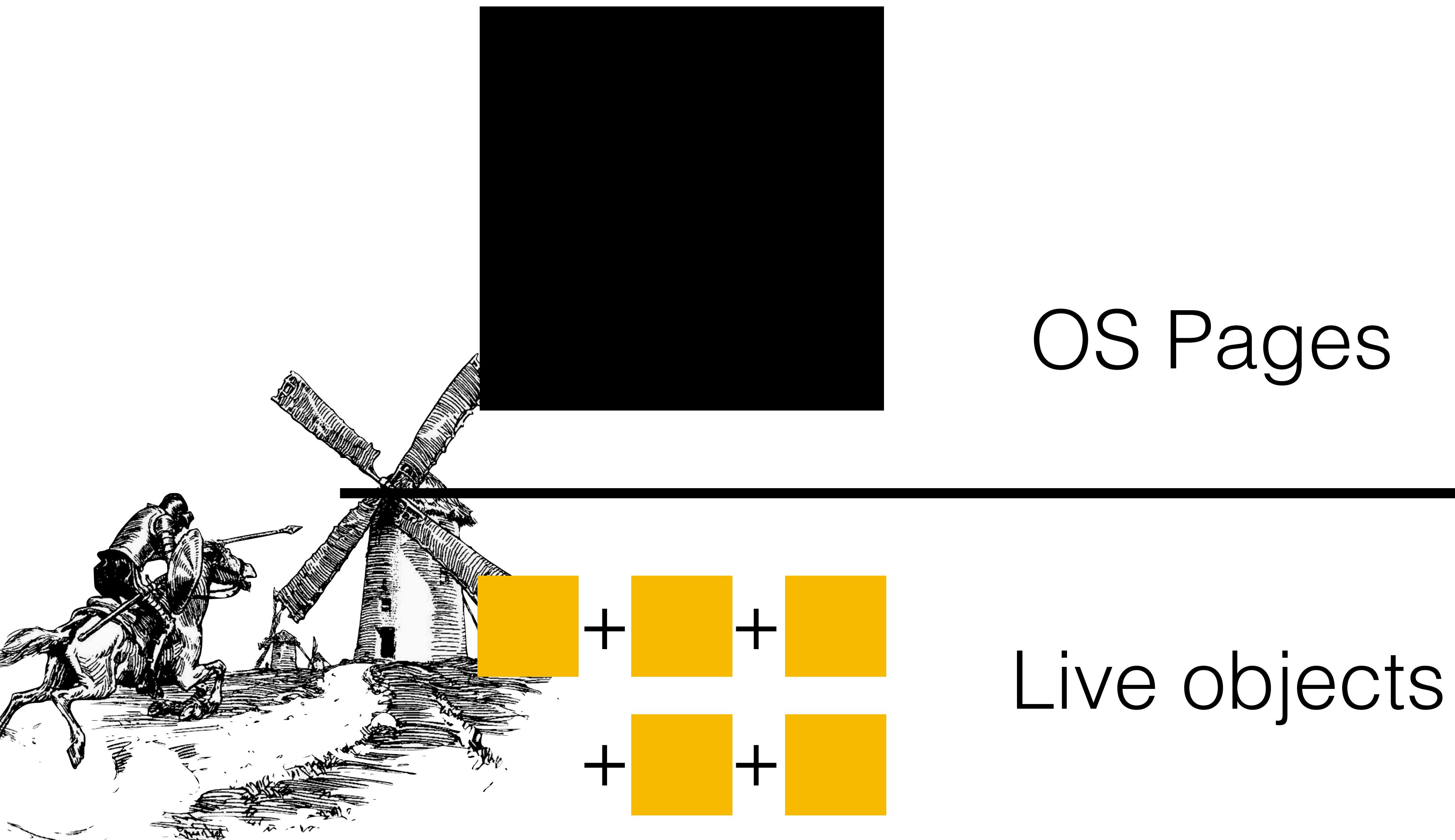
¡Fragmentación!



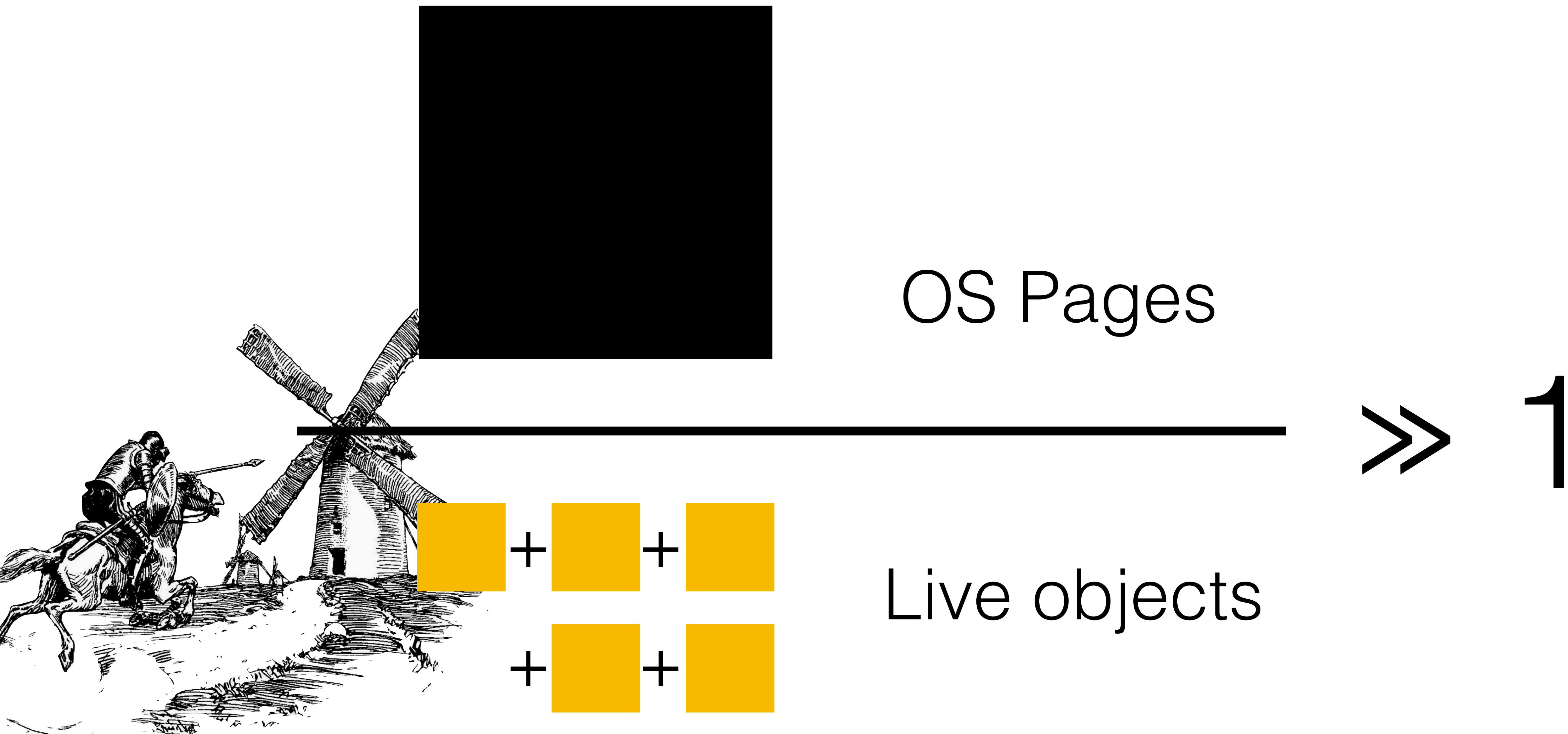
OS Page (“Spain”)

Live objects (malloc'd)

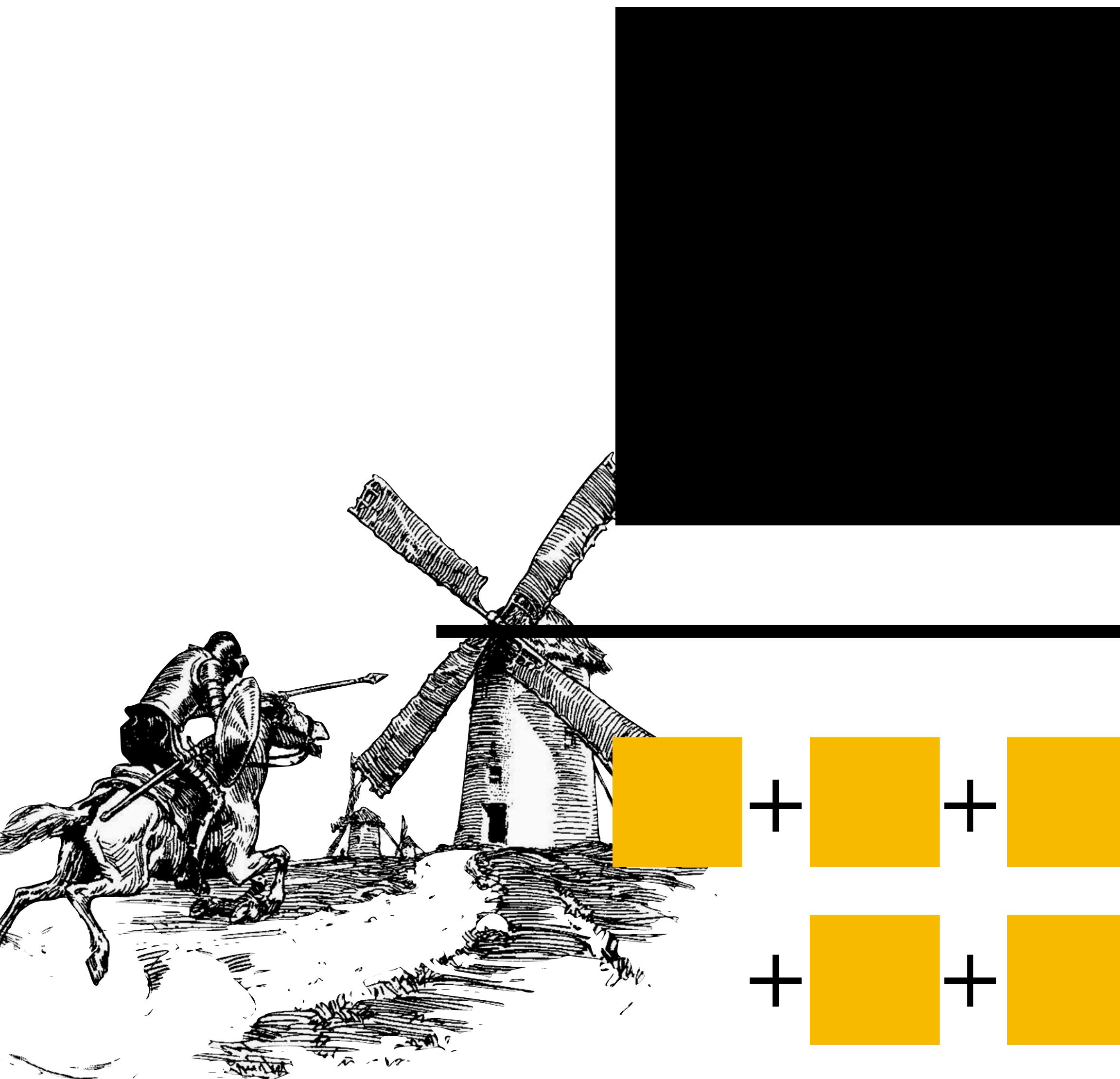
¡Fragmentación!



¡Fragmentación!



¡Fragmentación!

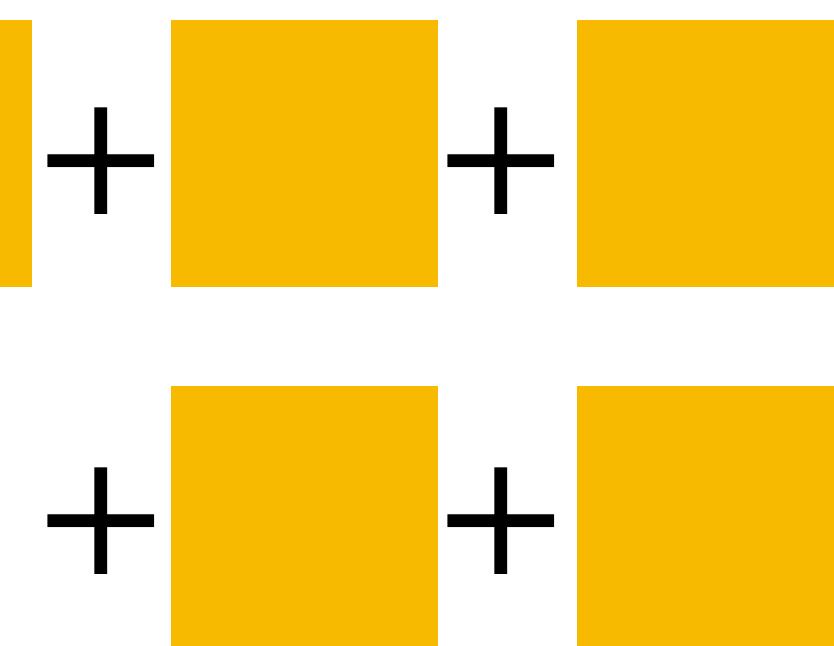


OS Pages

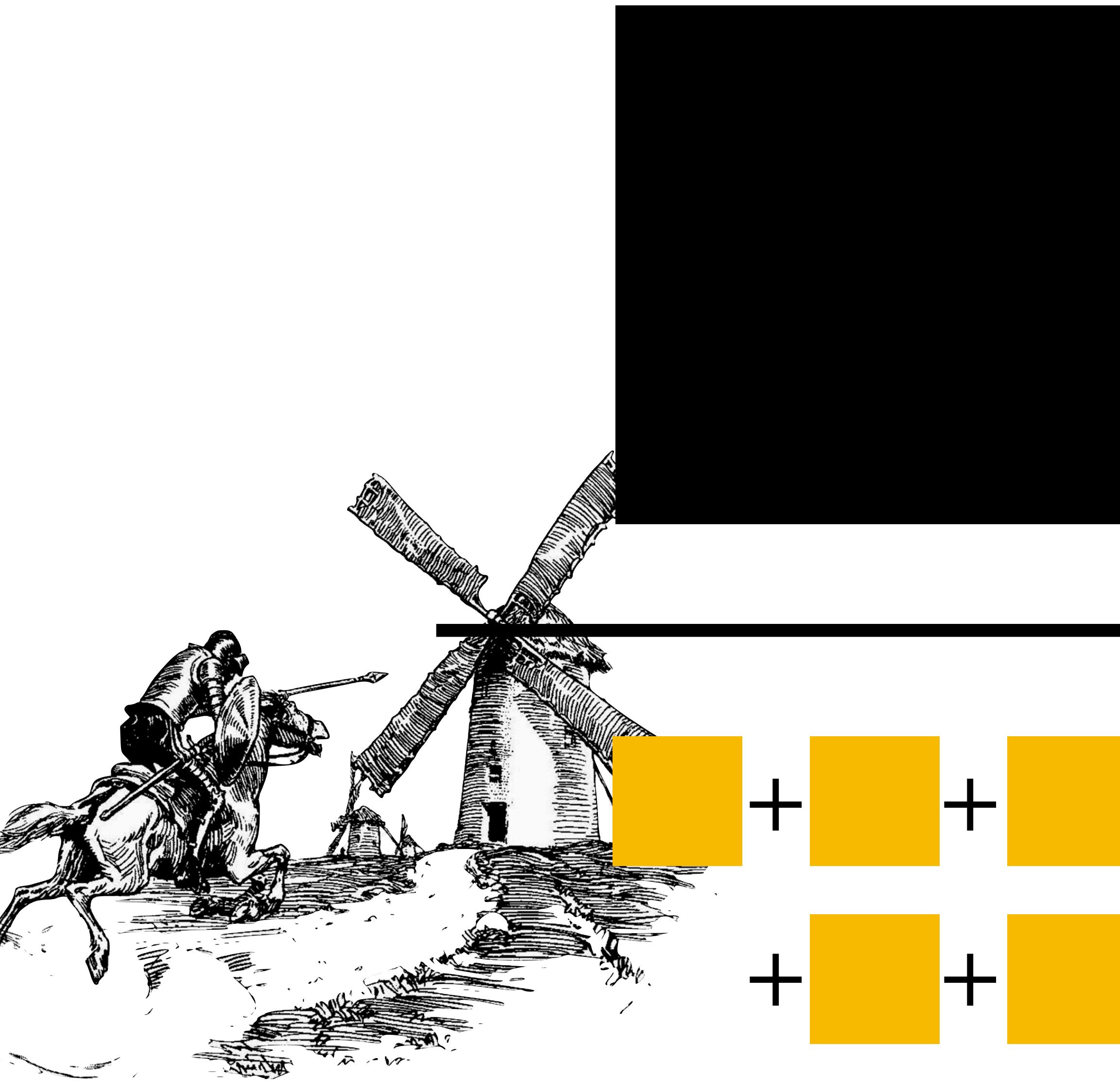
$O(\log \frac{\text{---}}{\text{---}})$ 13x!

[Robson '77]

Live objects



¿Compactación?



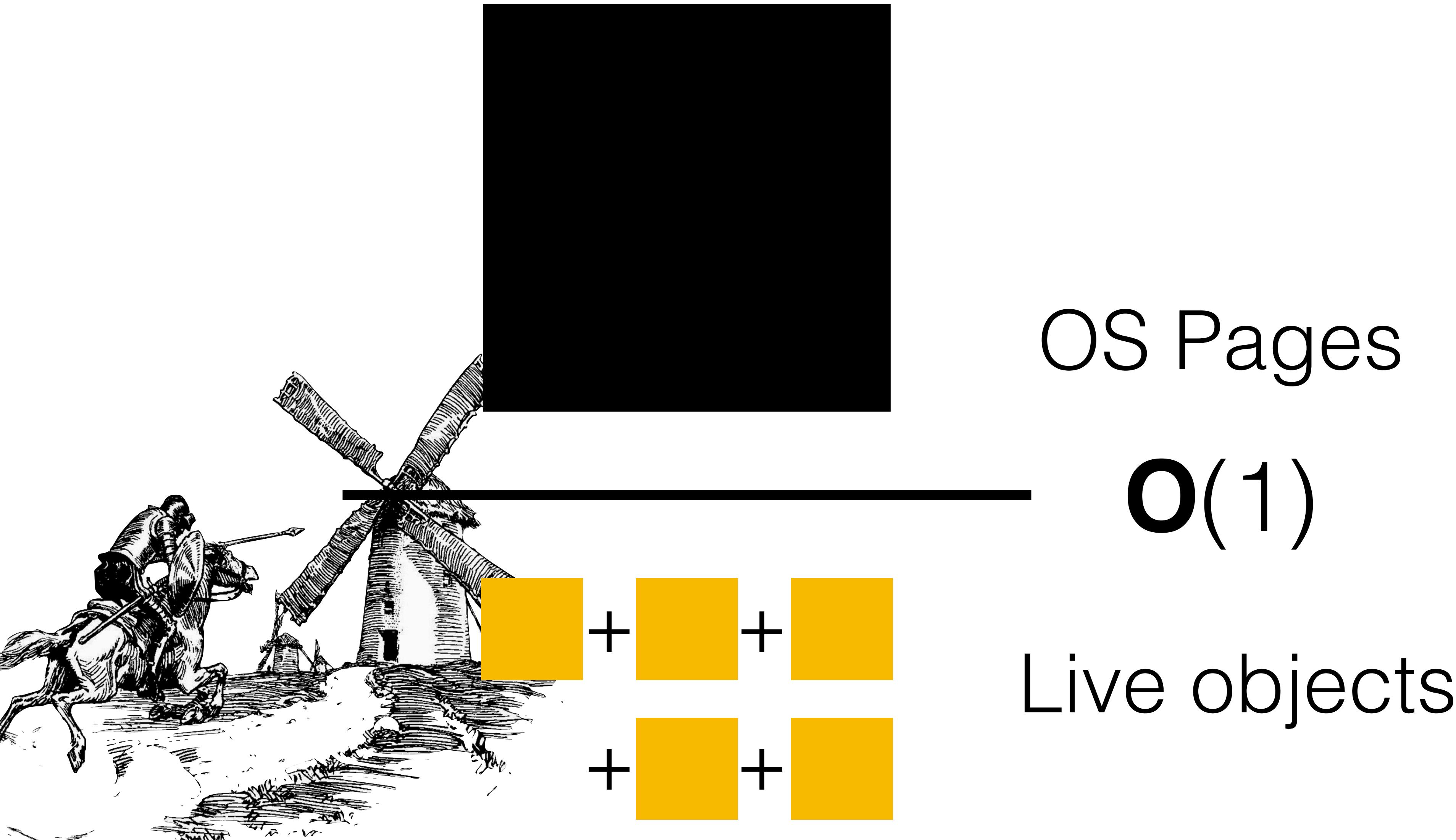
OS Pages

$O(\log \frac{\text{---}}{\text{---}})$ 13x!

[Robson '77]

Live objects

¿Compactación?









THE
C
PROGRAMMING
LANGUAGE

Brian W. Kernighan • Dennis M. Ritchie

PRENTICE HALL SOFTWARE SERIES



NG
E

M. Ritchie



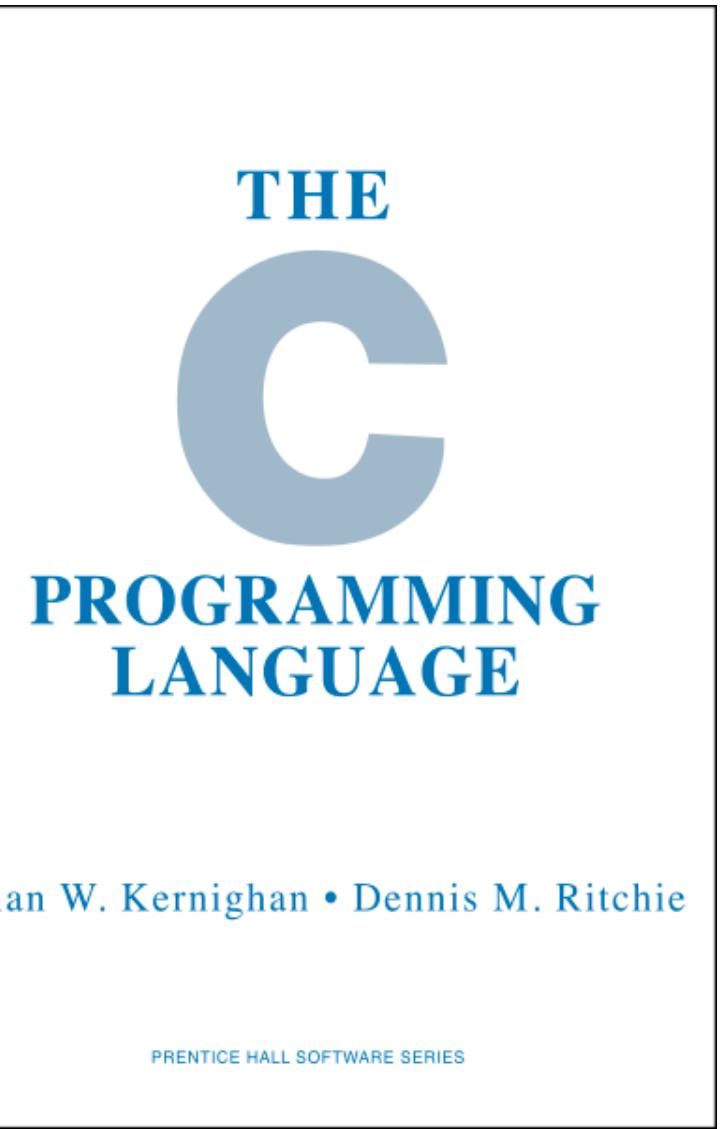
NG
E

M. Ritchie



NG
E

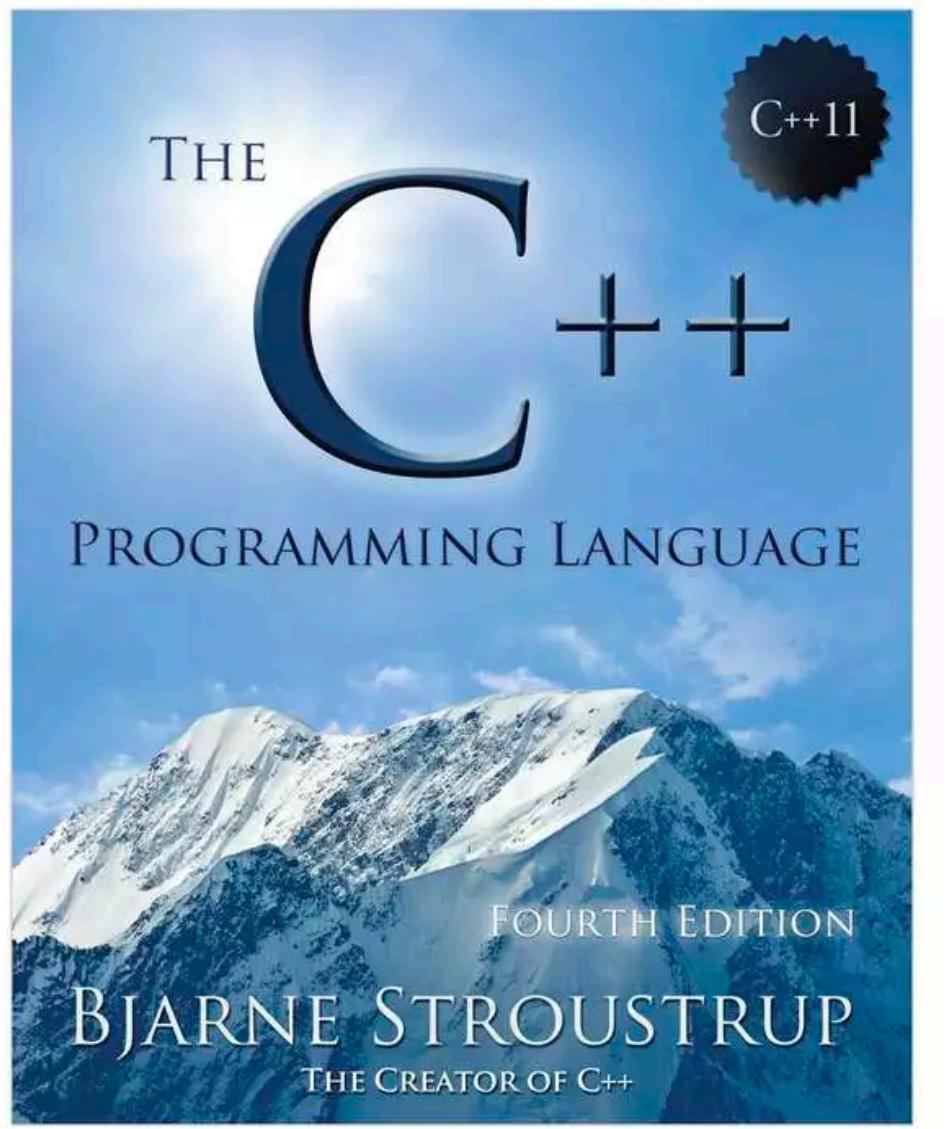
M. Ritchie



THE
C
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LANGUAGE

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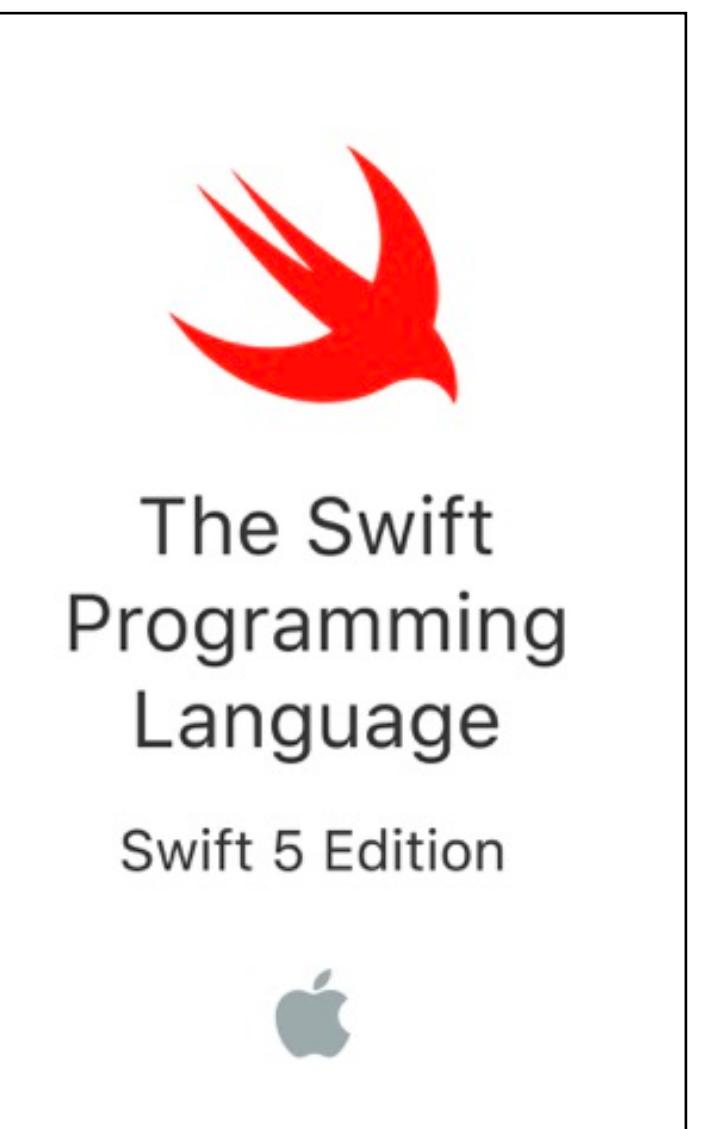
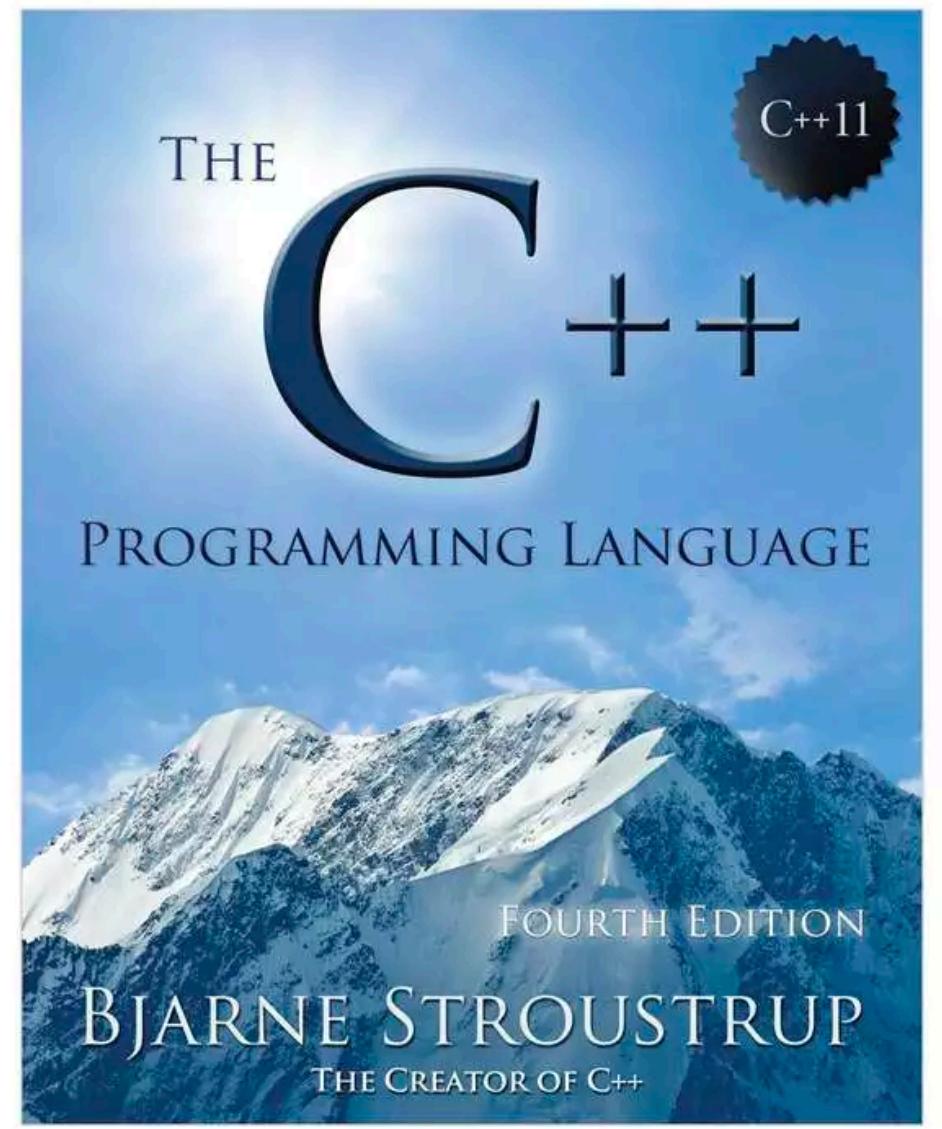
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THE
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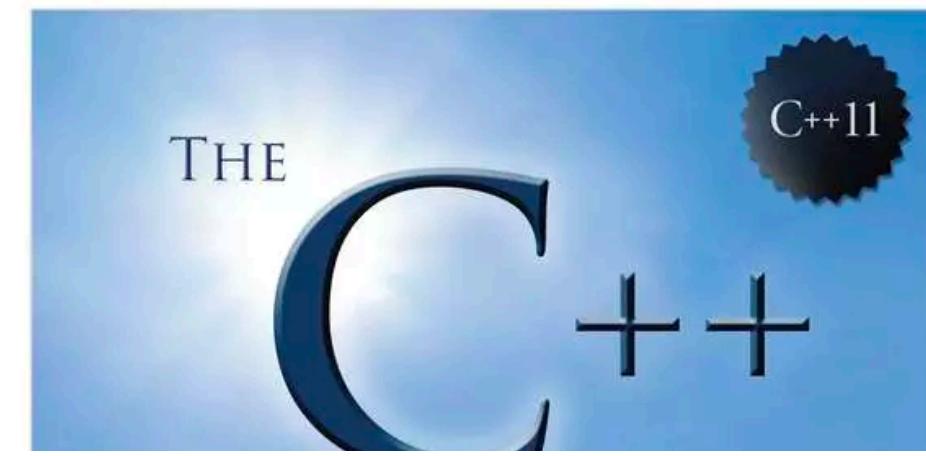
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The Swift
Programming
Language

Swift 5 Edition





[Lattner, 2016]

Why not a tracing GC?

- Native interoperability with unmanaged code
- Deterministic destruction provides:
 - No “finalizer problems” like resurrection, threading, etc.
 - Deterministic performance: can test/debug performance stutters
- Performance:
 - GC use ~3-4x more memory than ARC to achieve good performance
 - Memory usage is very important for mobile and cloud apps
 - Incremental/concurrent GCs slow the mutator like ARC does

Quantifying the Performance of Garbage Collection vs. Explicit Memory Management
Matthew Hertz, Emery D. Berger. OOPSLA'05

Native interoperability with unmanaged code

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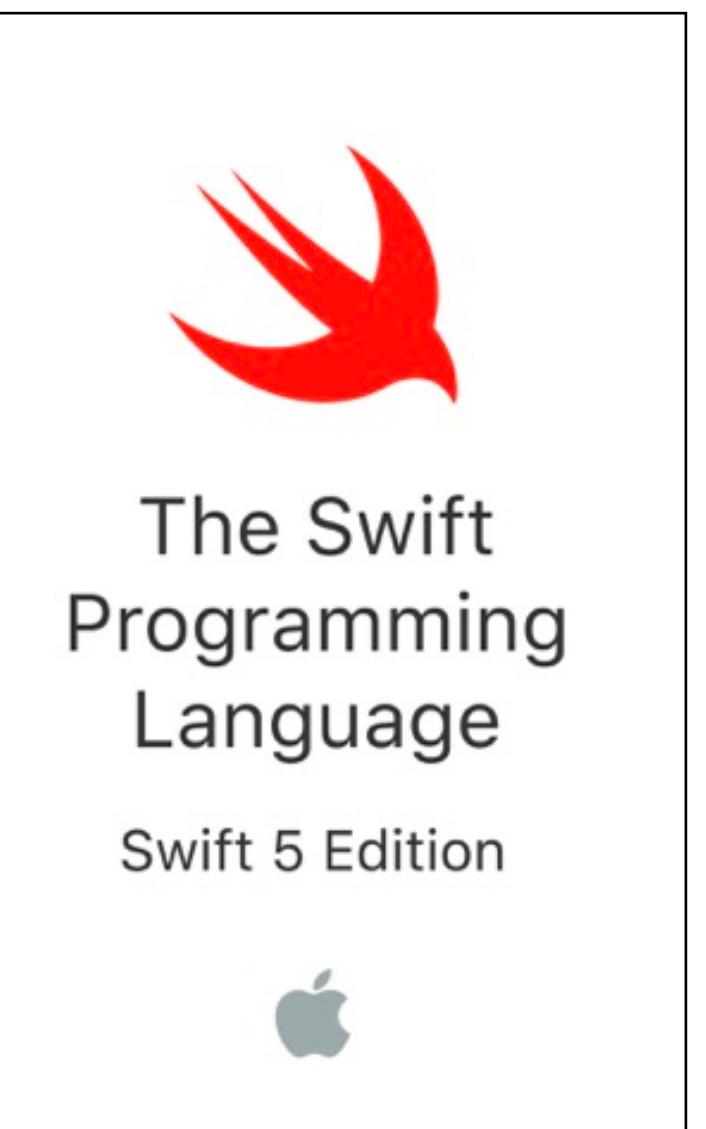
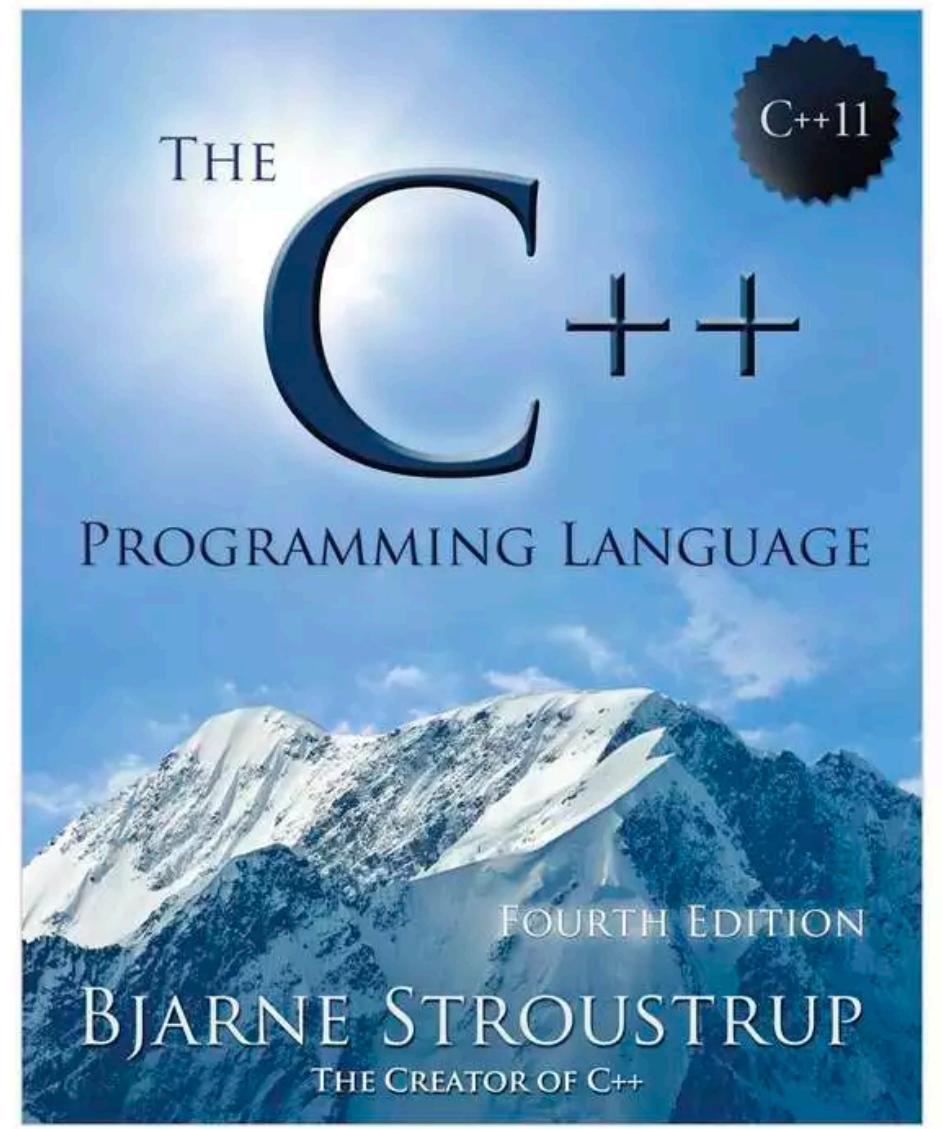
Quantifying the Performance of Garbage Collection vs. Explicit Memory Management

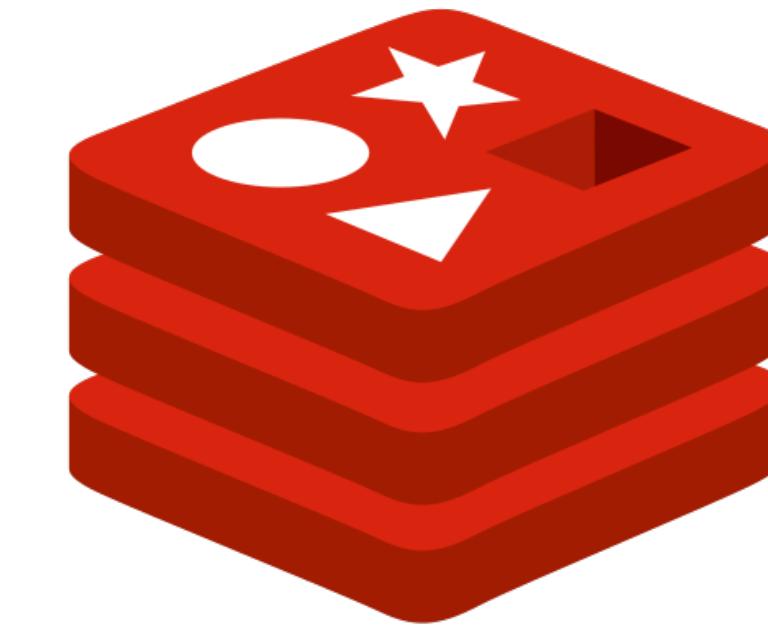
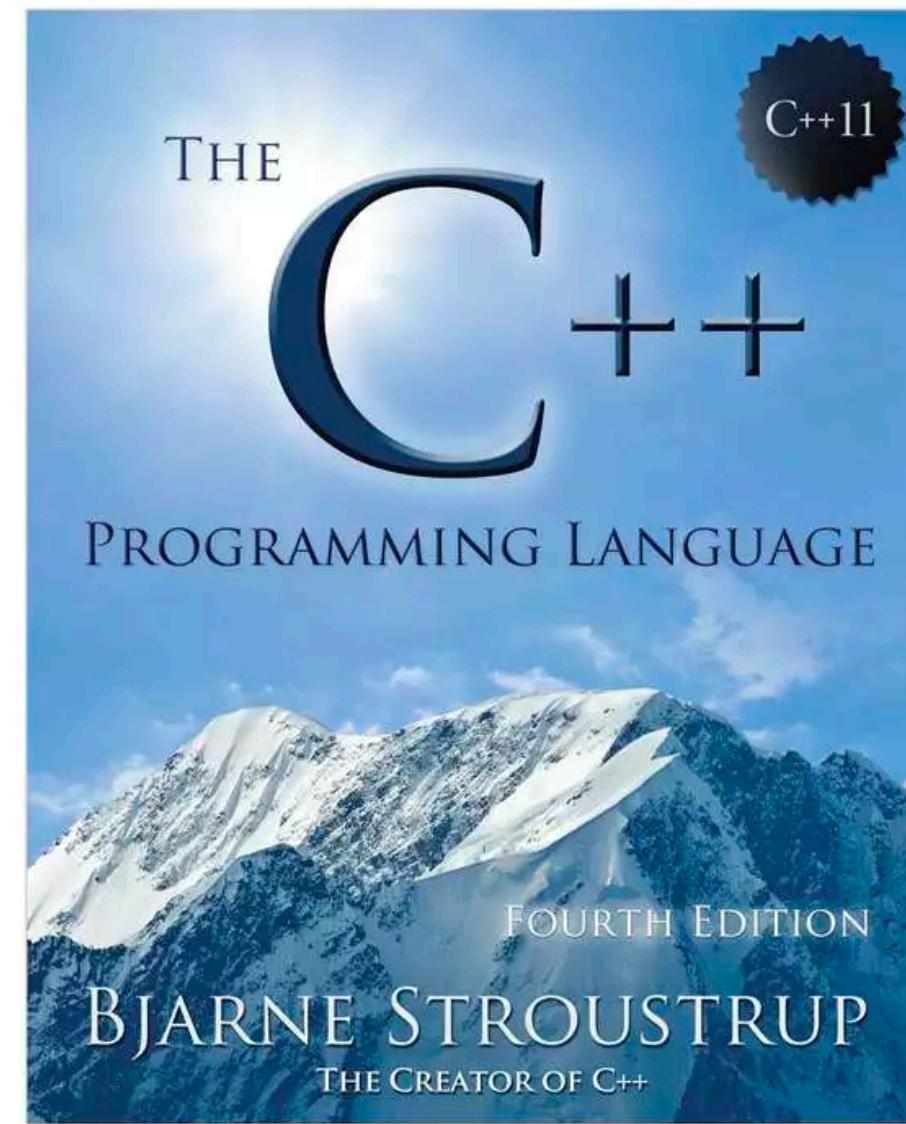
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THE
C
PROGRAMMING
LANGUAGE

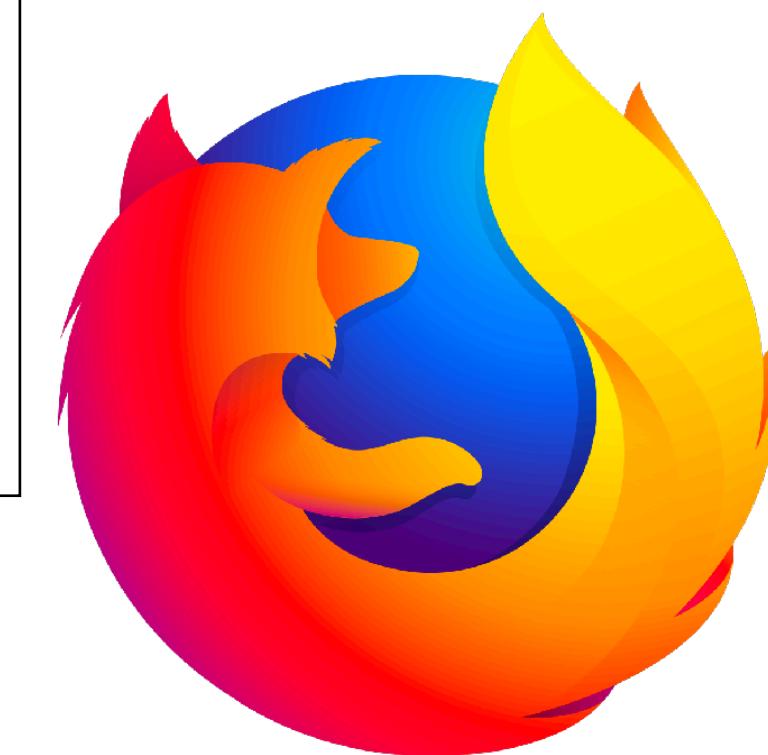
Brian W. Kernighan • Dennis M. Ritchie

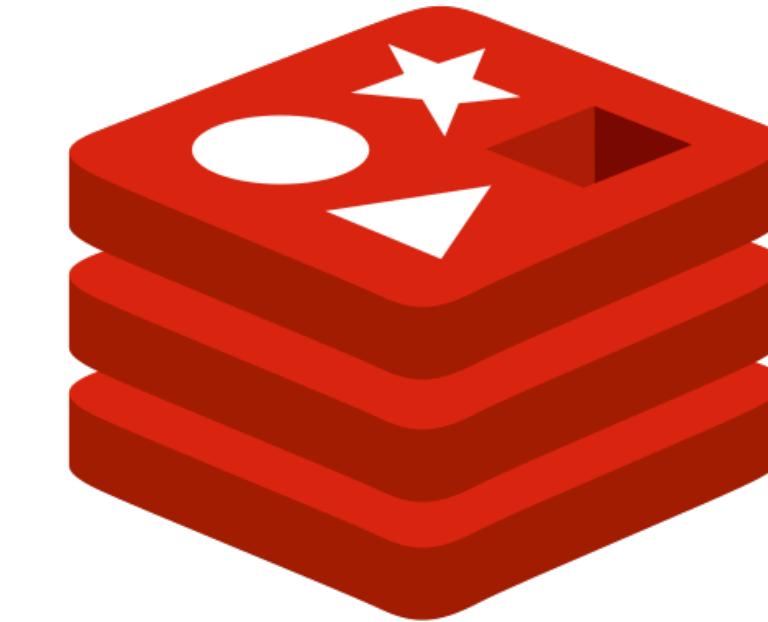
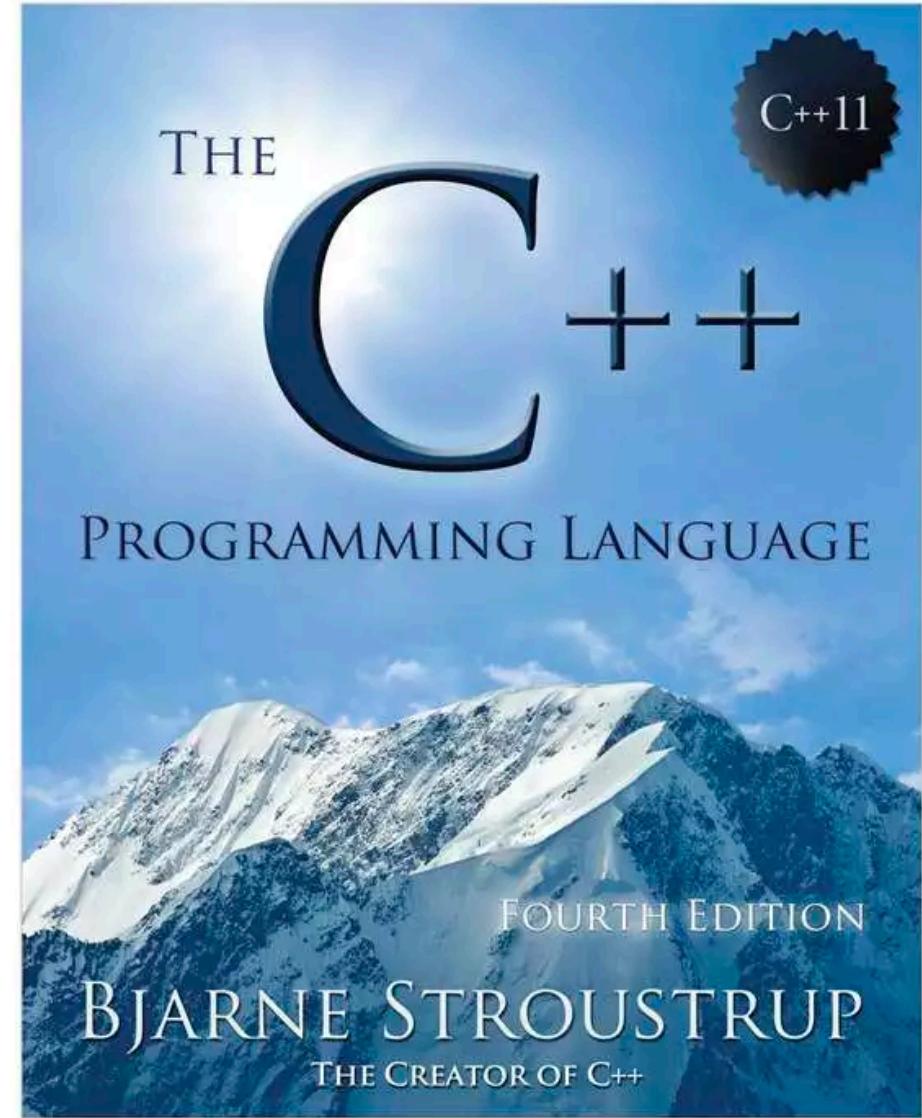
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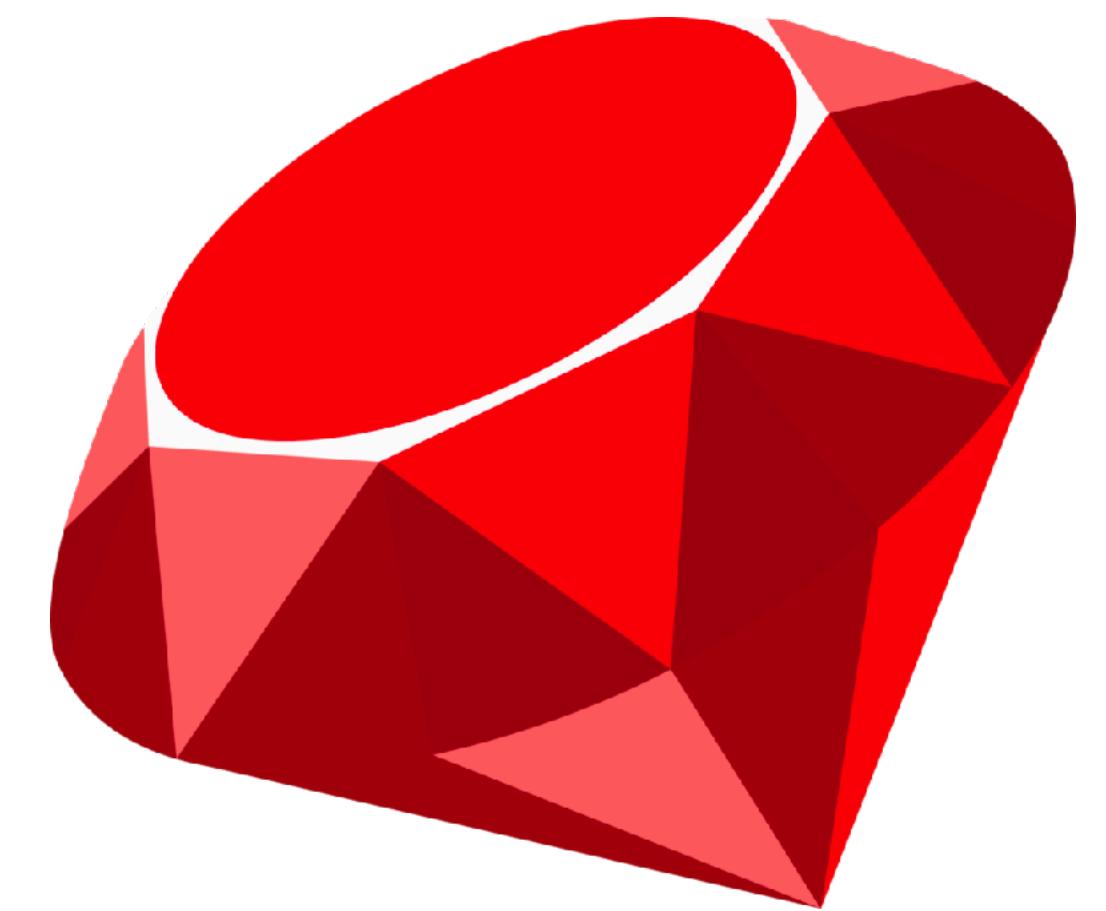


redis





redis

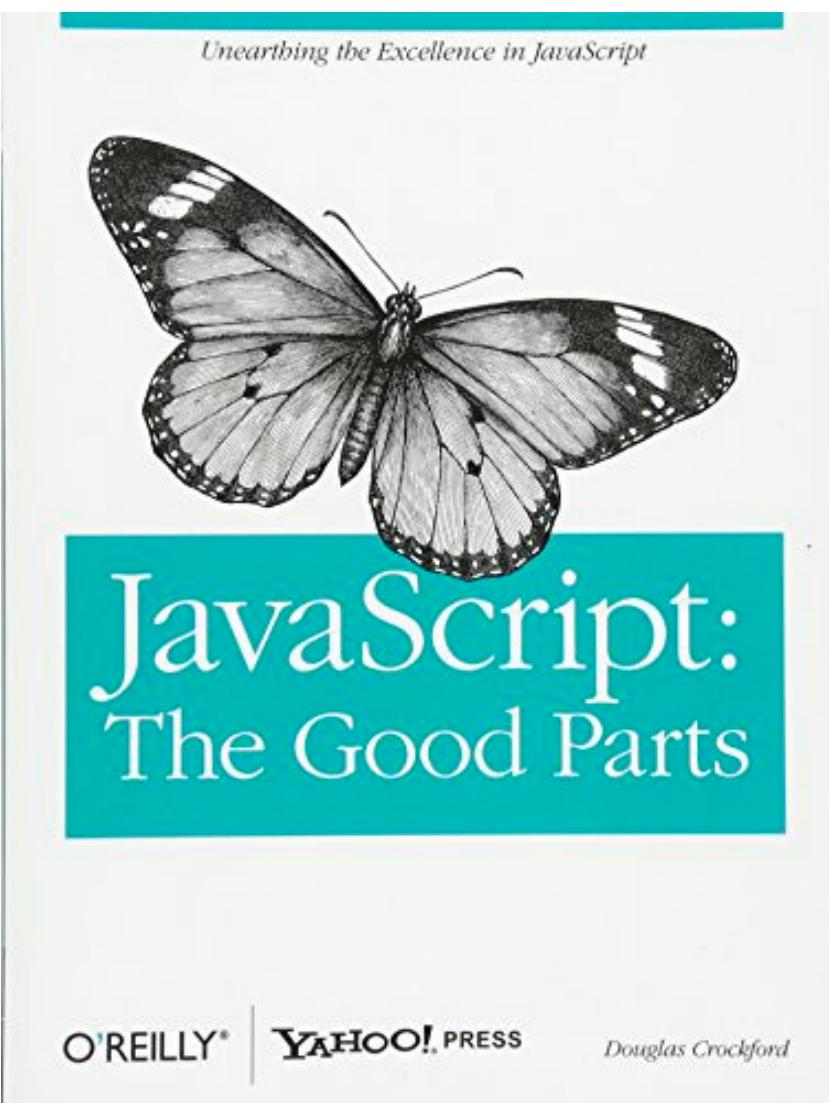
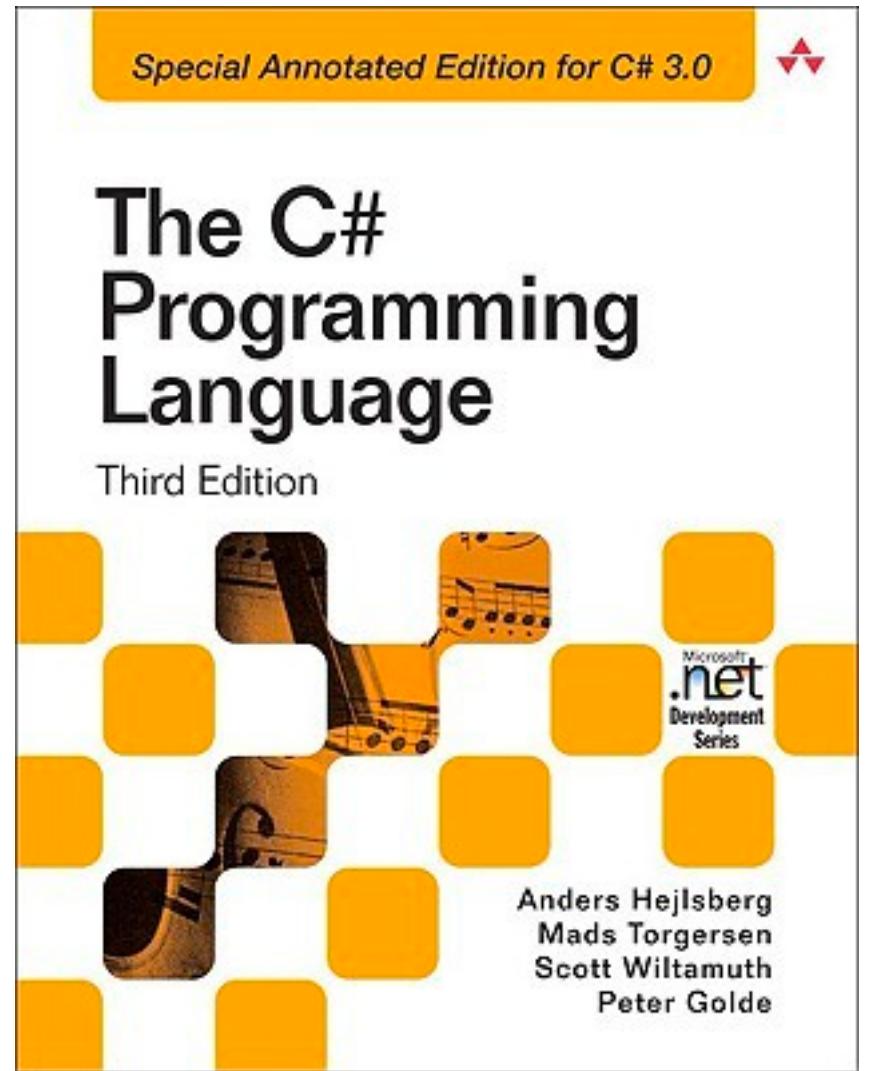
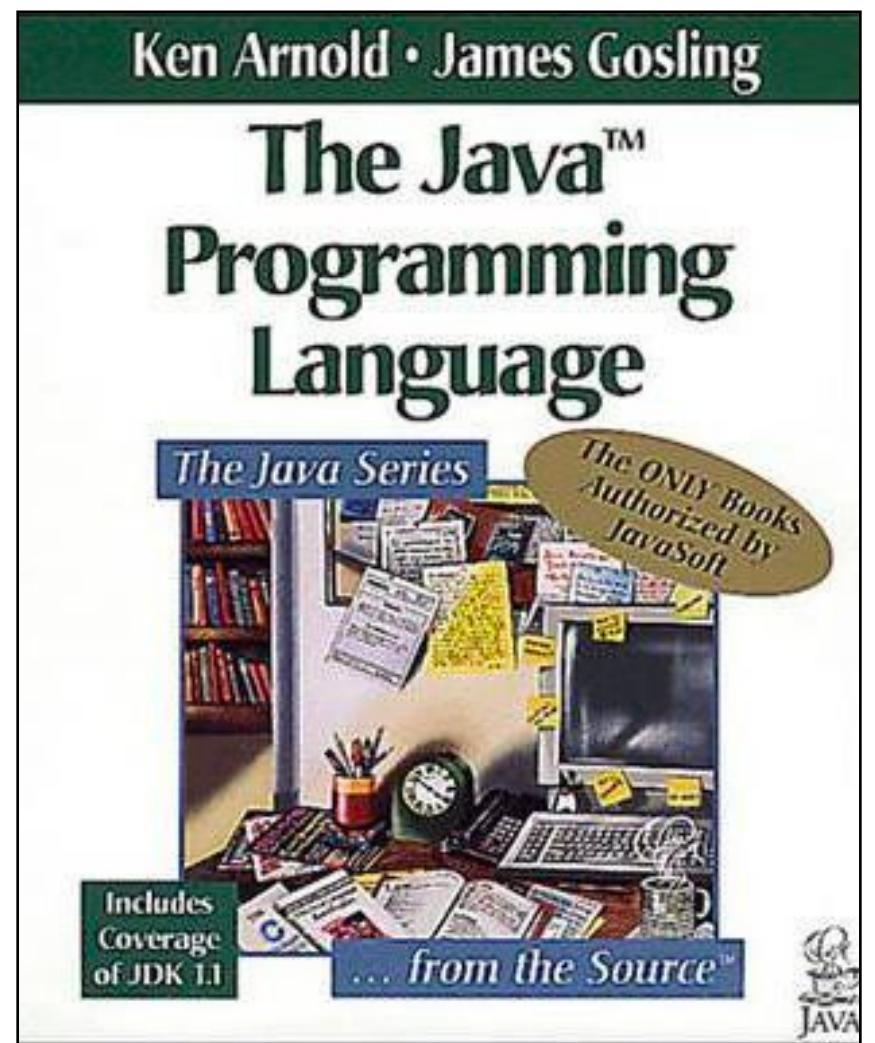


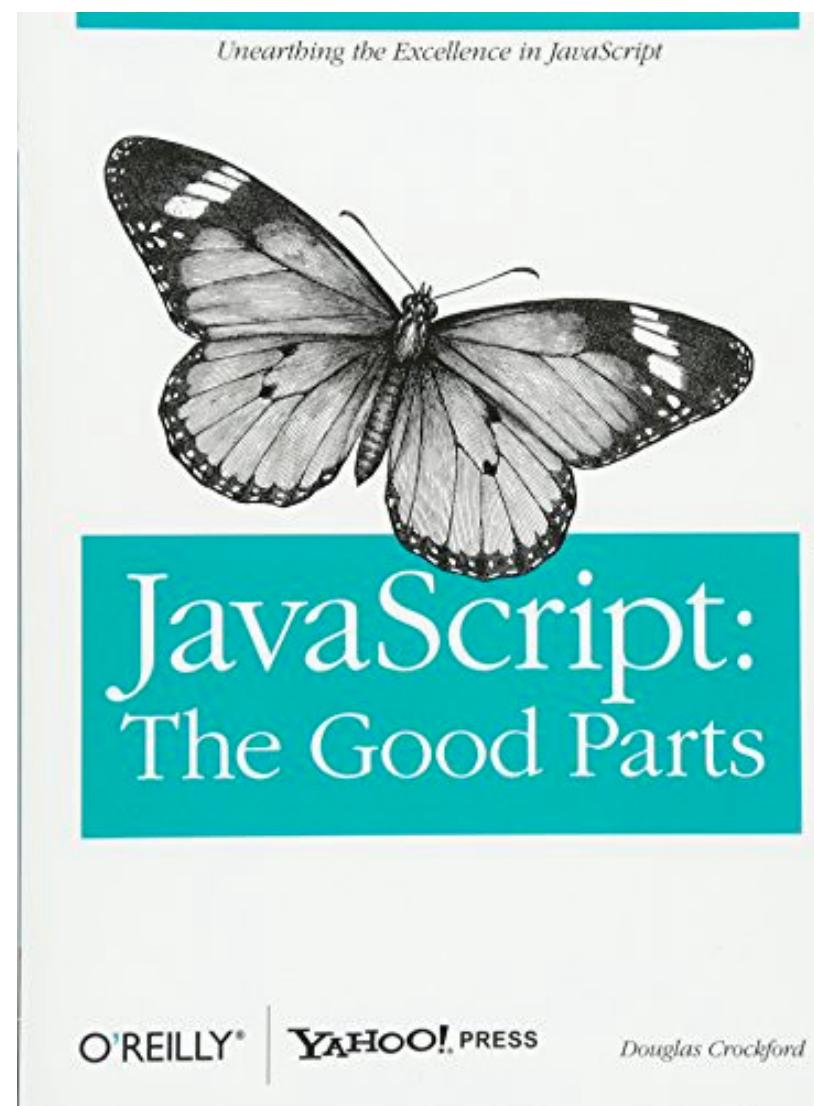
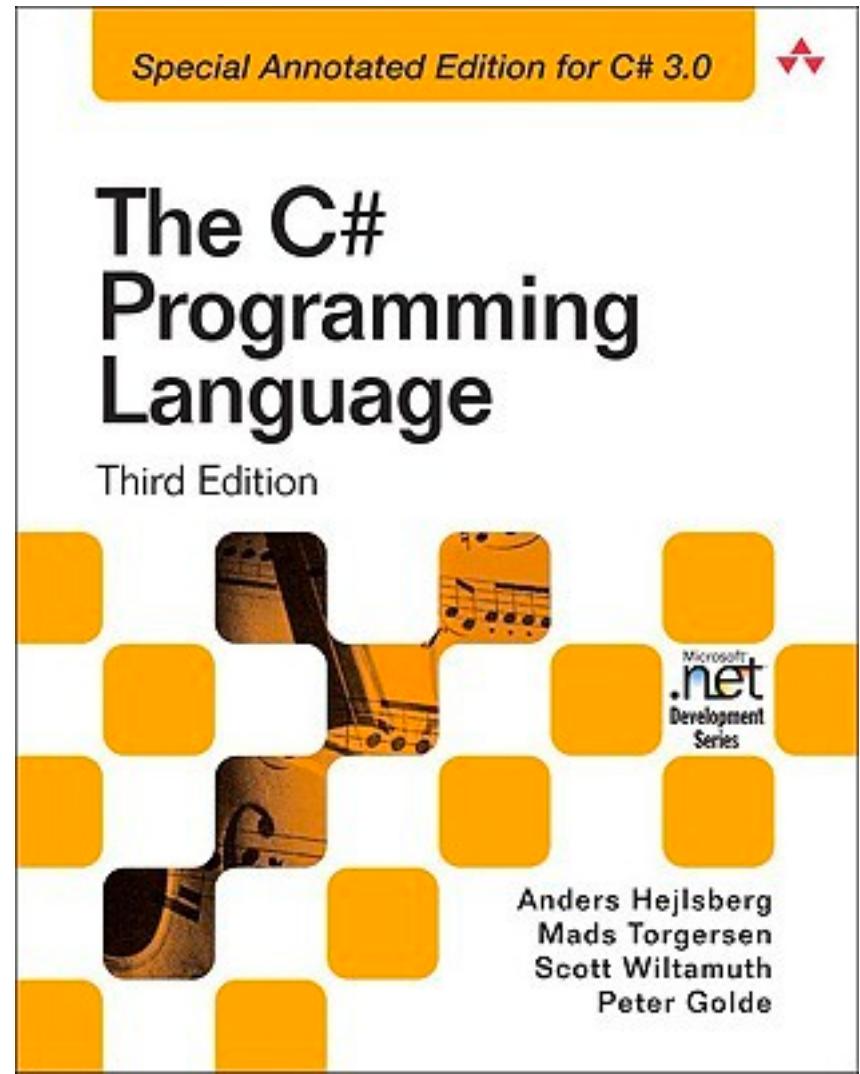
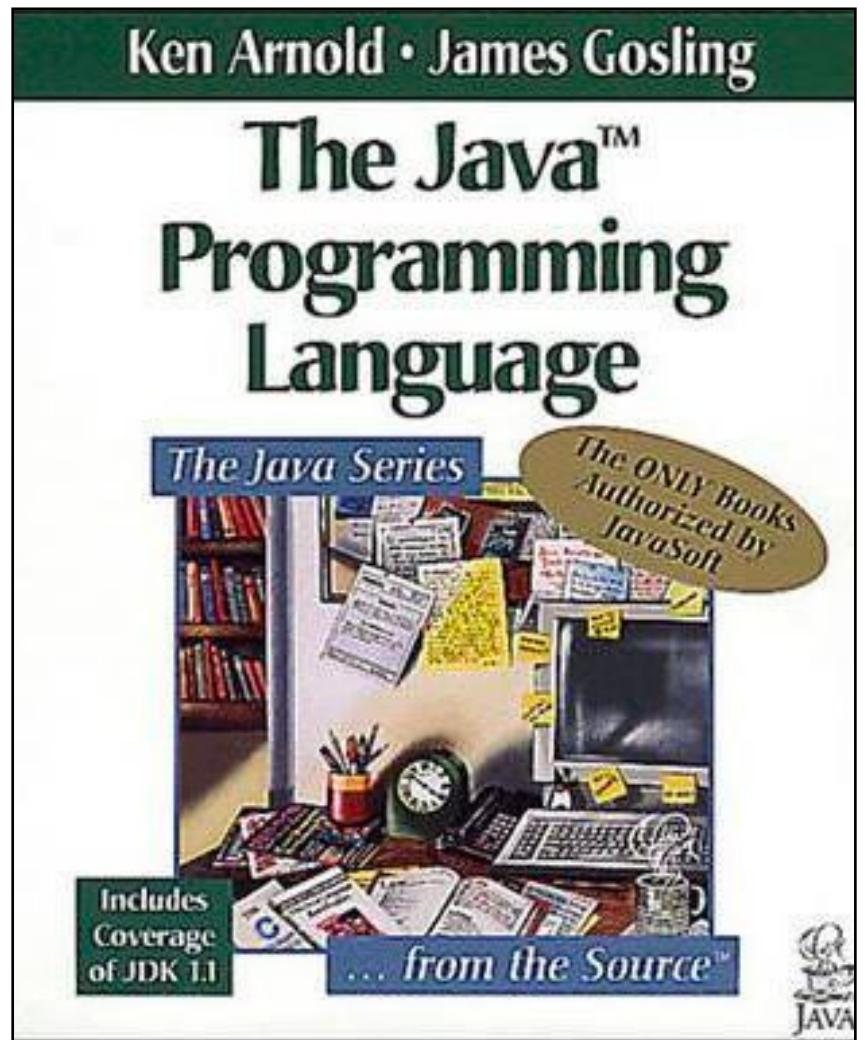
MySQL™



Office

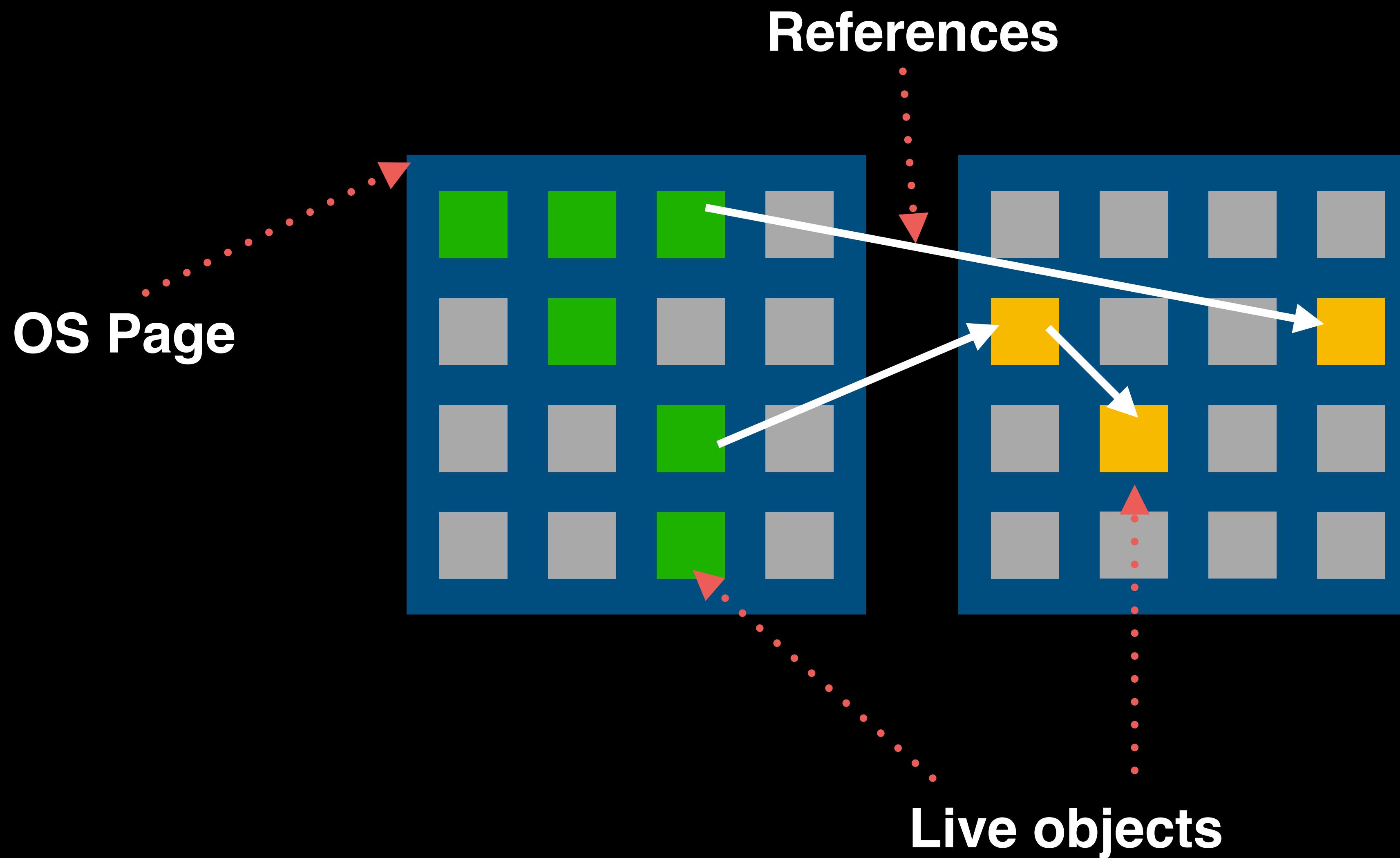


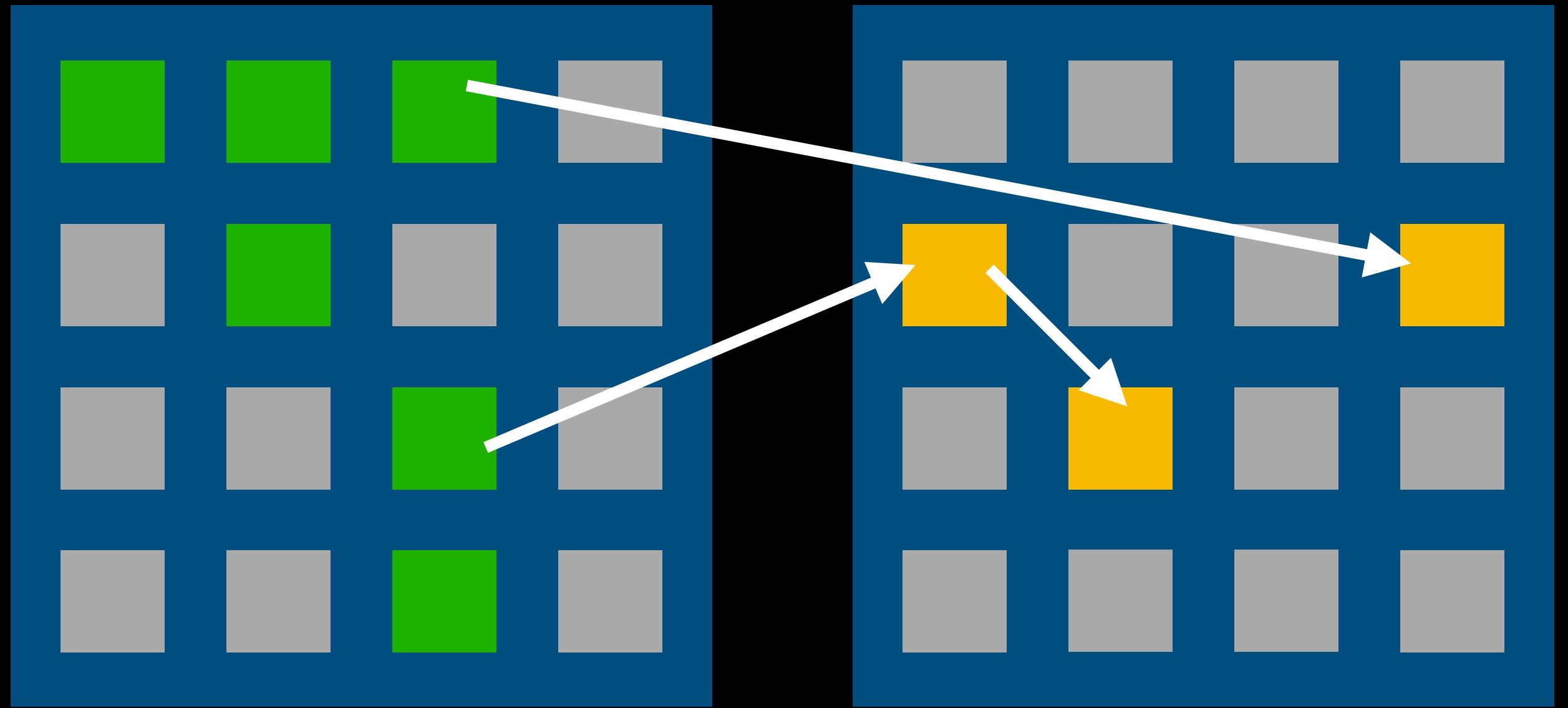


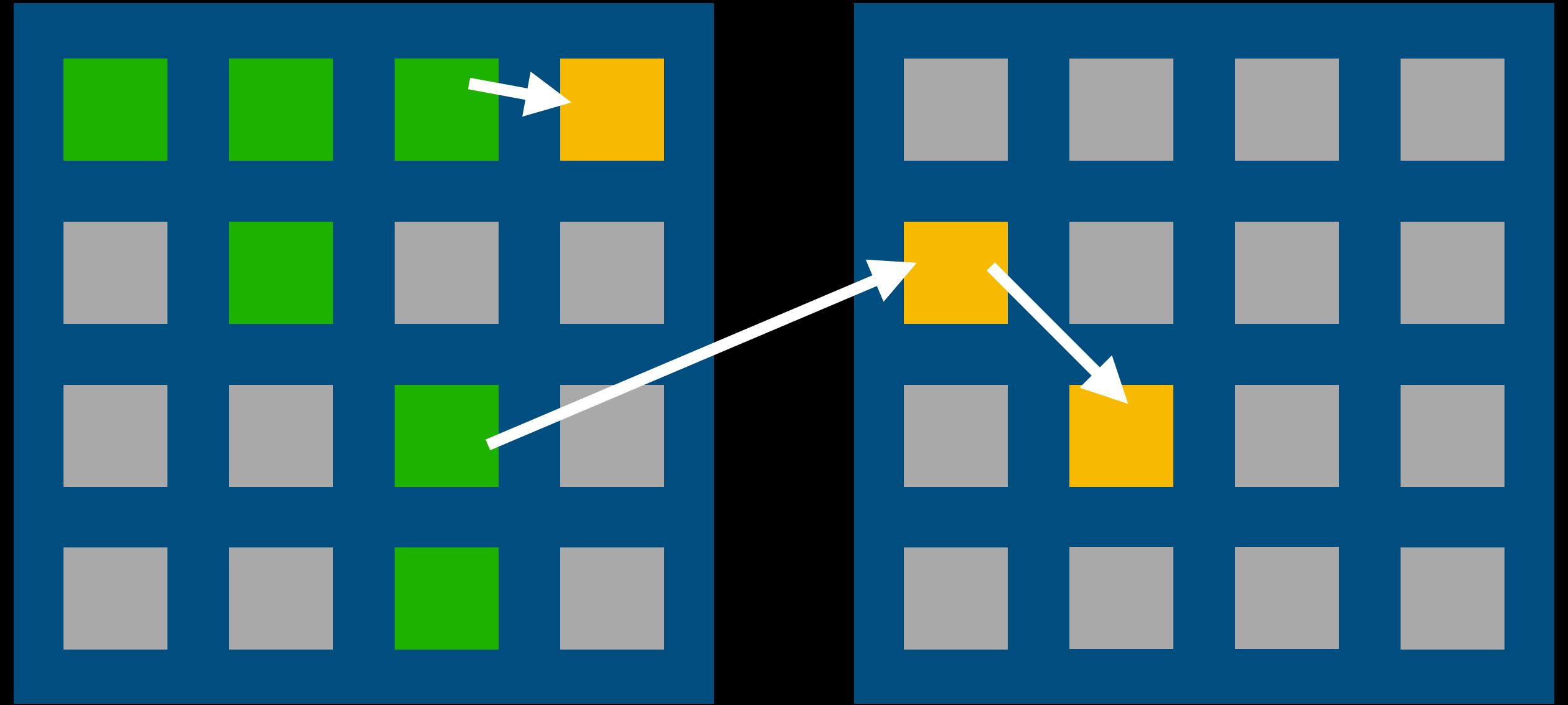


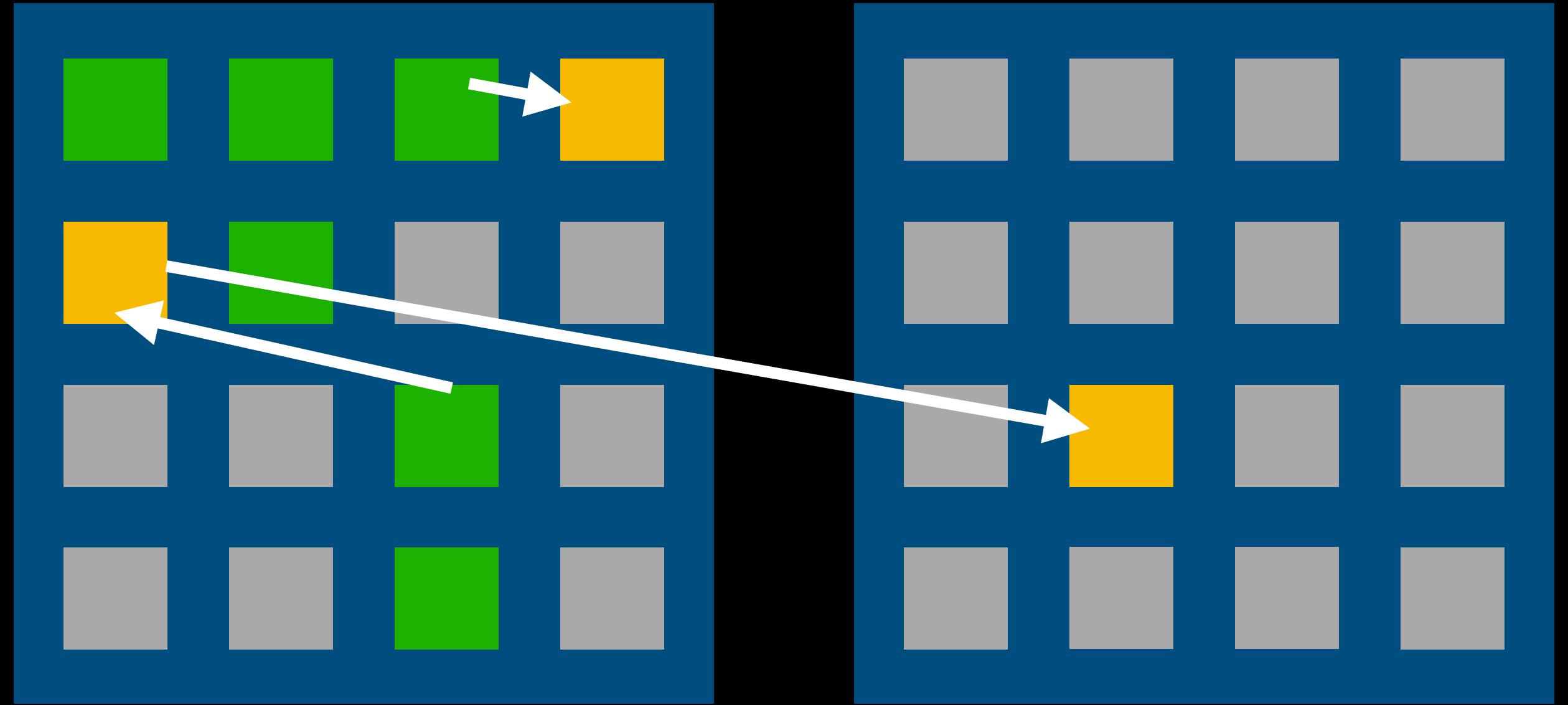
¡Compactación!

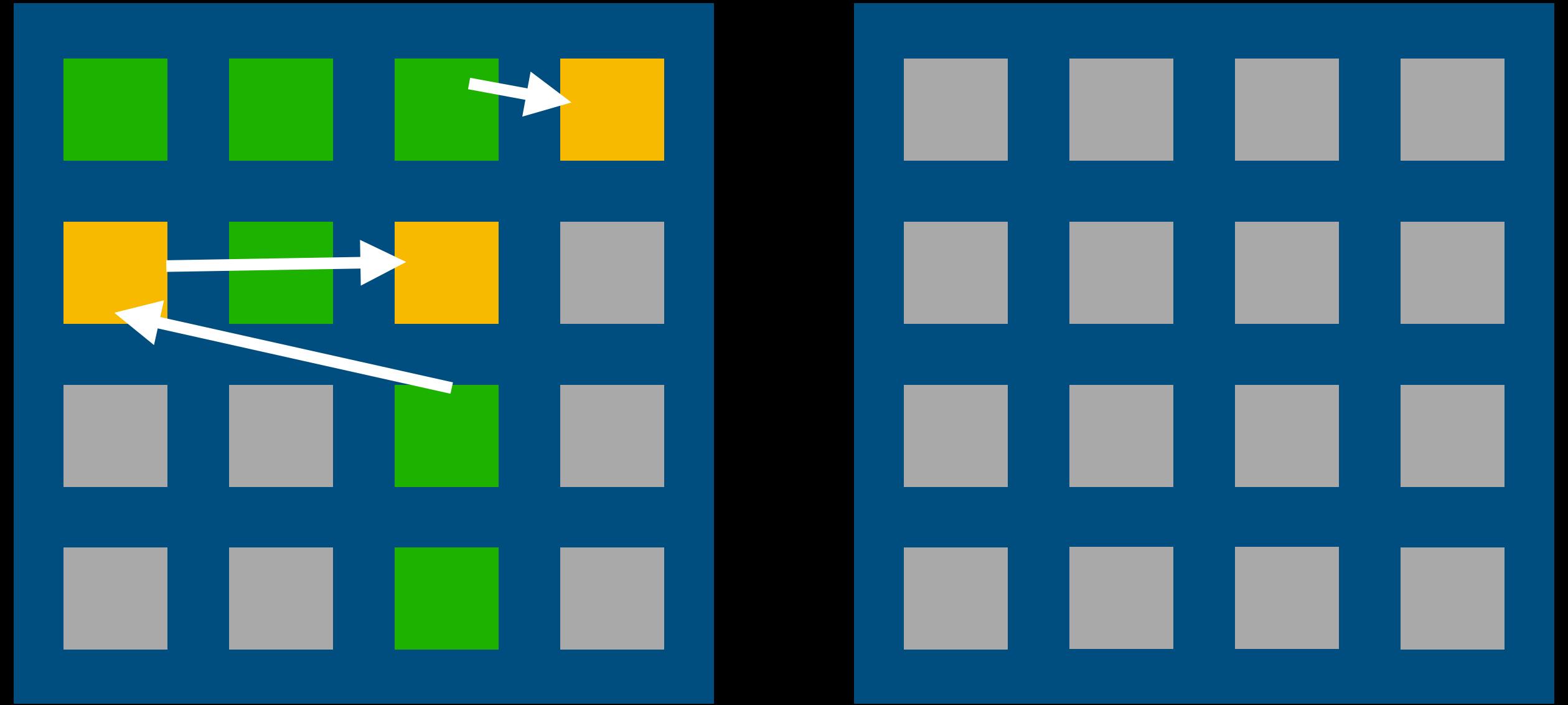
Compaction In Action

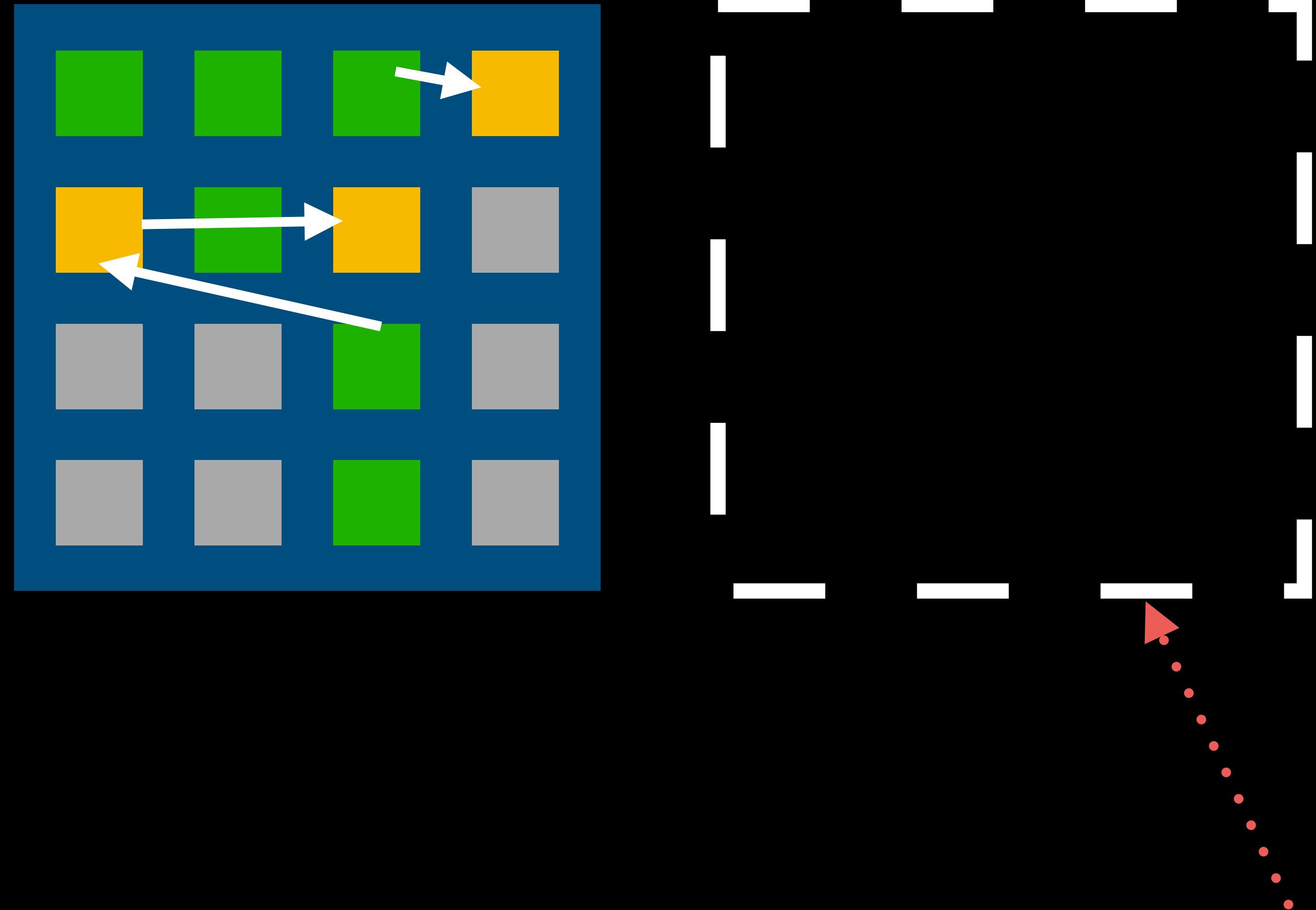












Page returned to the OS

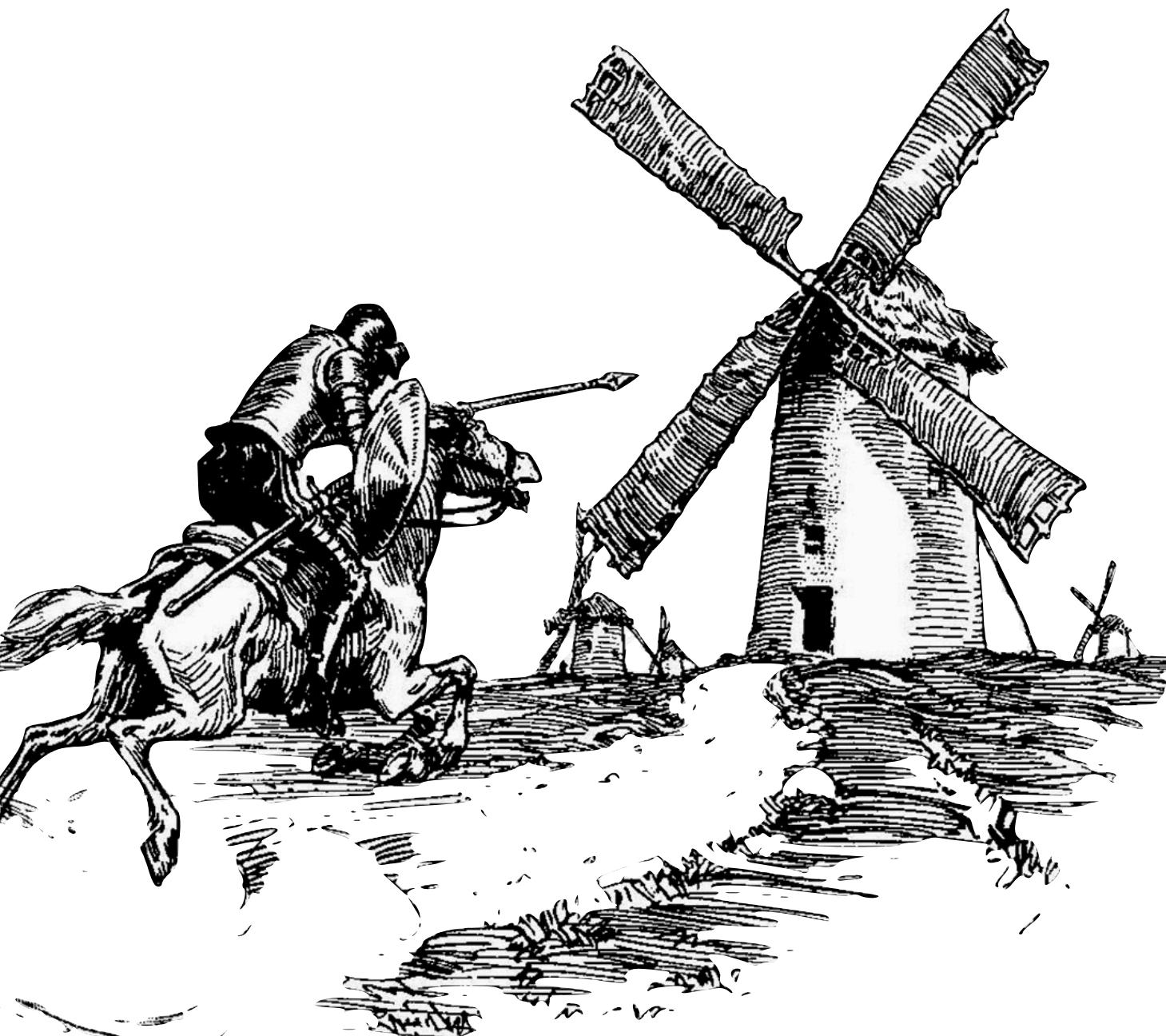
```
$ clang++ -o yolo main.cc
```



```
$ clang++ -o yolo main.cc  
$ strip yolo
```



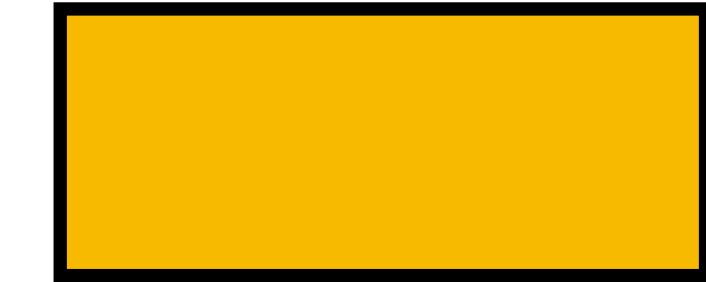
```
$ clang++ -o yolo main.cc  
$ strip yolo  
$ ./yolo
```



**No way to precisely
distinguish pointers
from integers**

```
$ clang++ -o yolo main.cc  
$ strip yolo  
$ ./yolo
```

0xDEADC000



**No way to precisely
distinguish pointers
from integers**

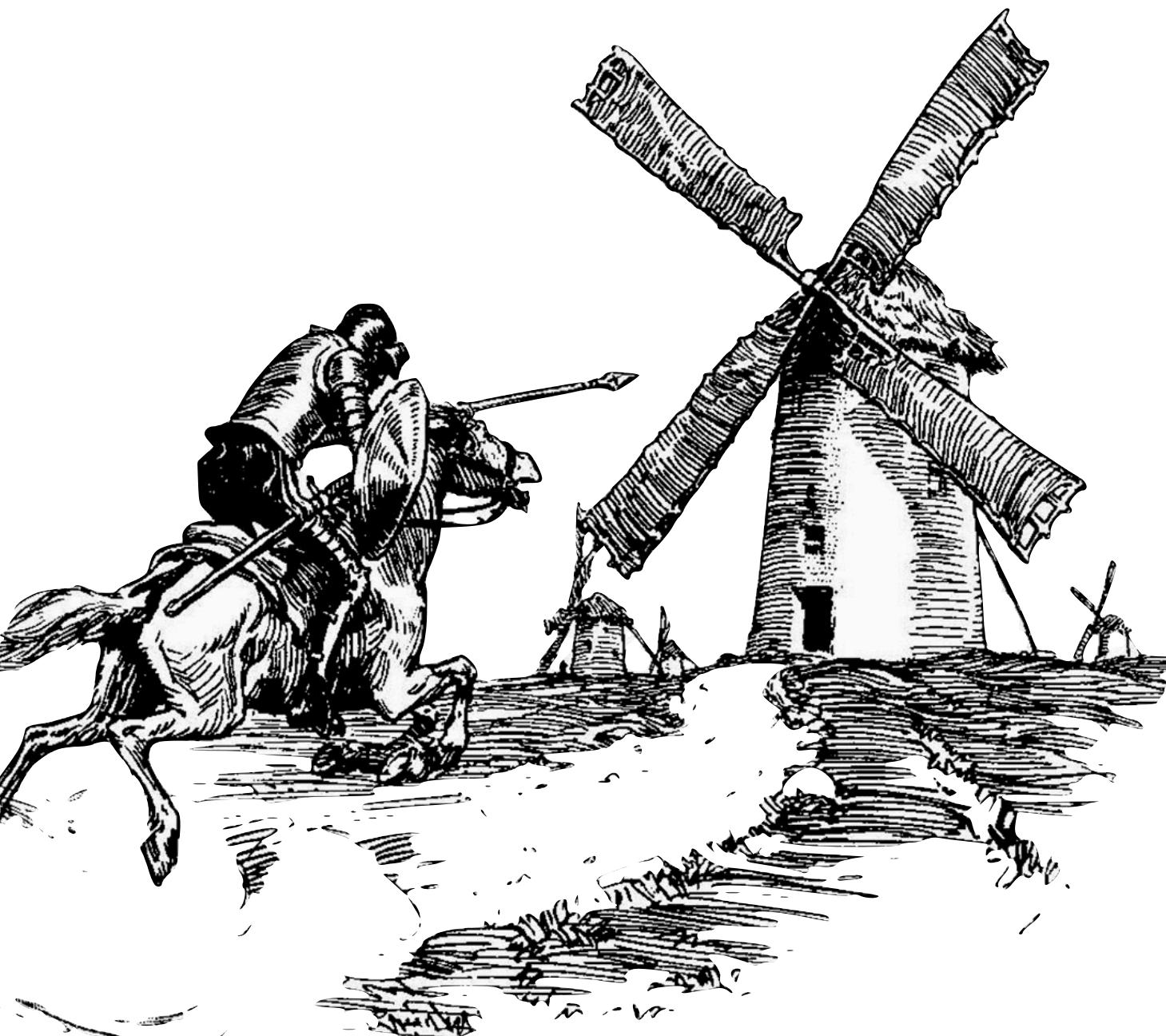
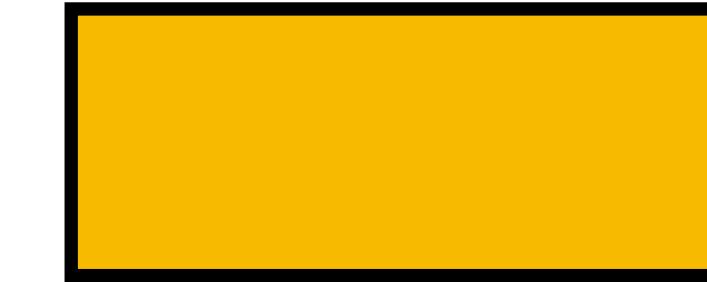


```
$ clang++ -o yolo main.cc
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```
$ strip yolo
```

```
$ ./yolo
```

0xDEADC000



**No way to precisely
distinguish pointers
from integers**

0xBEEFC000



0xDEADC000

```
$ clang++ -o yolo main.cc
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```
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```

```
$ ./yolo
```

0xDEADC000



0xBEEFC000



No way to precisely
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from integers



```
$ clang++ -o yolo main.cc  
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$ ./yolo
```

0xDEADC000



0xDEADC000



0xBEEFC000



No way to precisely
distinguish pointers
from integers



```
union tiny
{
    int * ptr;
    uintptr_t flag;
};

tiny x;

// initialize
x.ptr = new int;

// set flag true
x.flag |= 1;
```

```
union tiny
{
    int * ptr;
    uintptr_t flag;
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tiny x;
```

0xDEADC001



0xDEADC000



```
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```

0xBEEFC000



```
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```

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0xDEADC001



0xDEADC000



0xBEEFC000



```
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// set flag true
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```



MESH

*Compaction without
Relocation for C/C++*

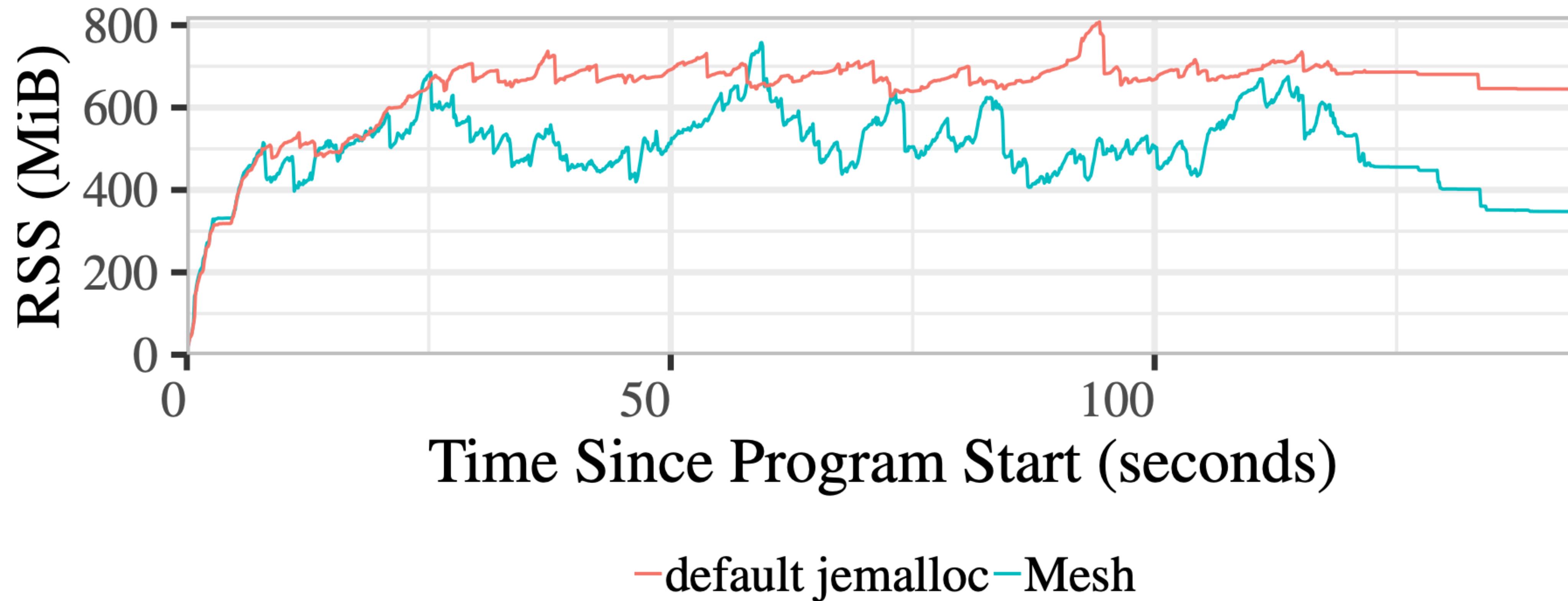
No code changes

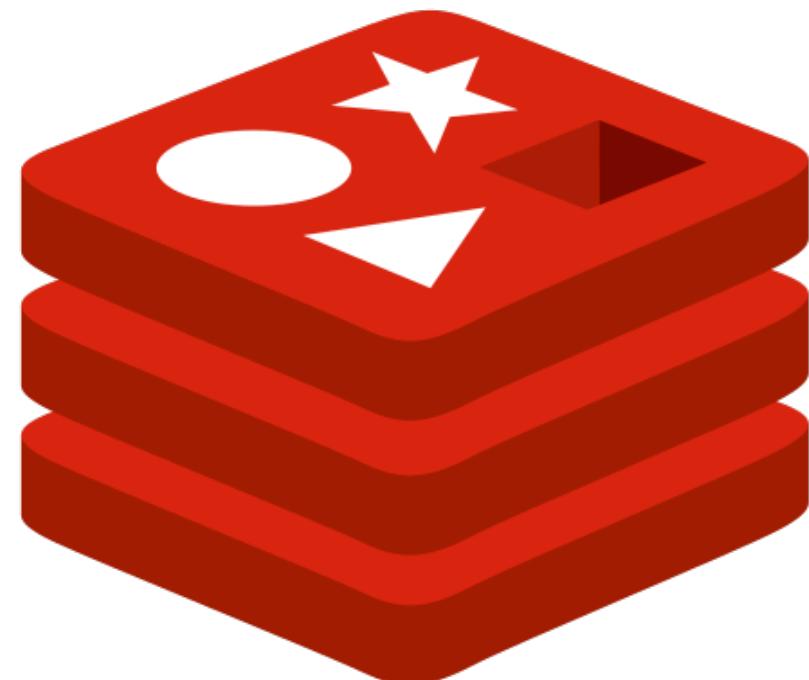
No recompilation

LD PRELOAD and go

17% heap size reduction

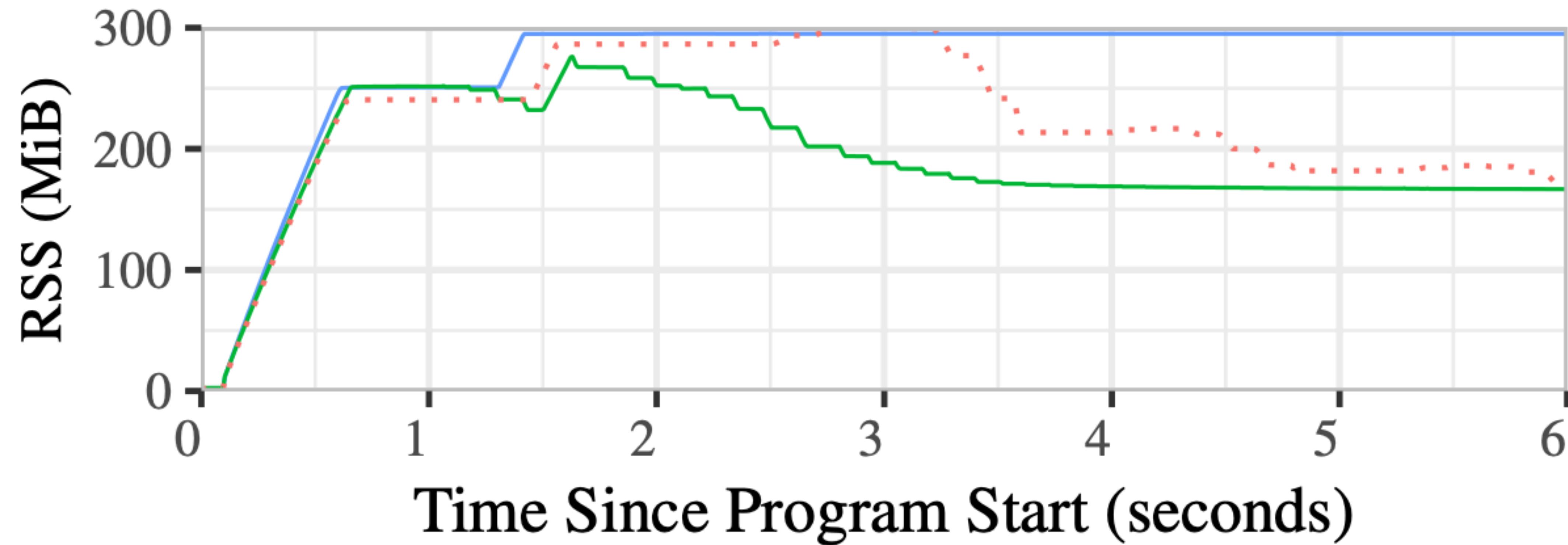
< 1% performance overhead

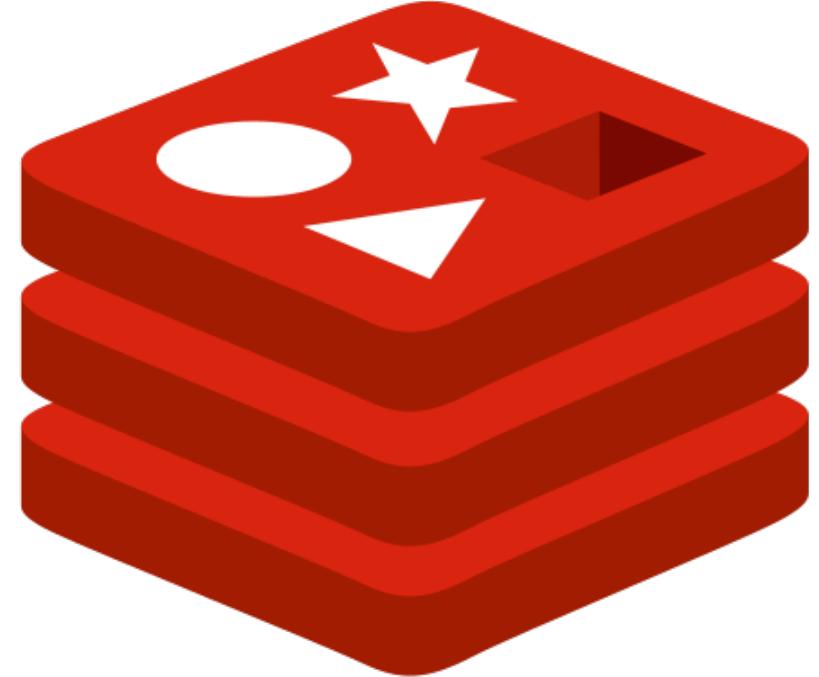




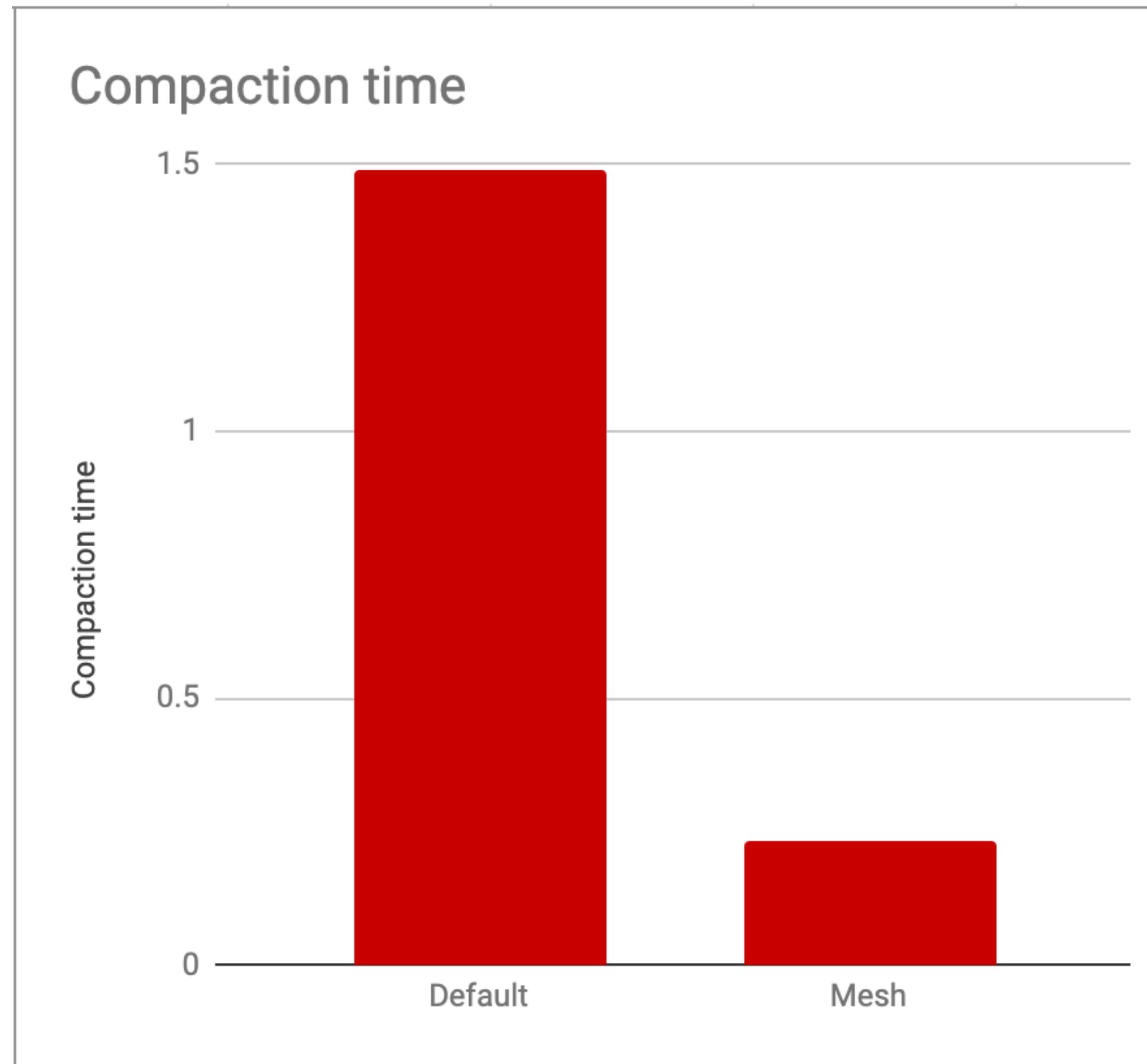
redis

- jemalloc + activatedefrag
- Mesh
- no compaction



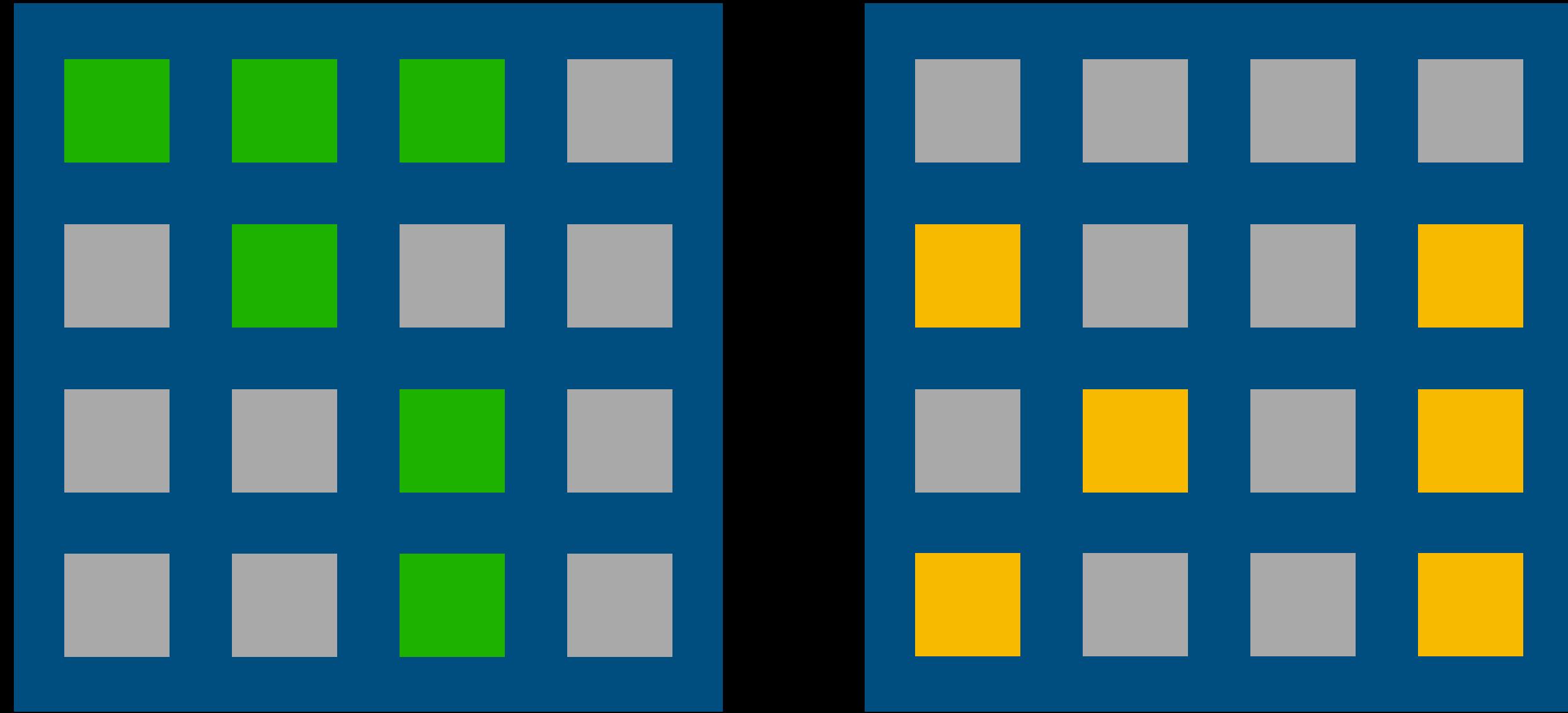


redis



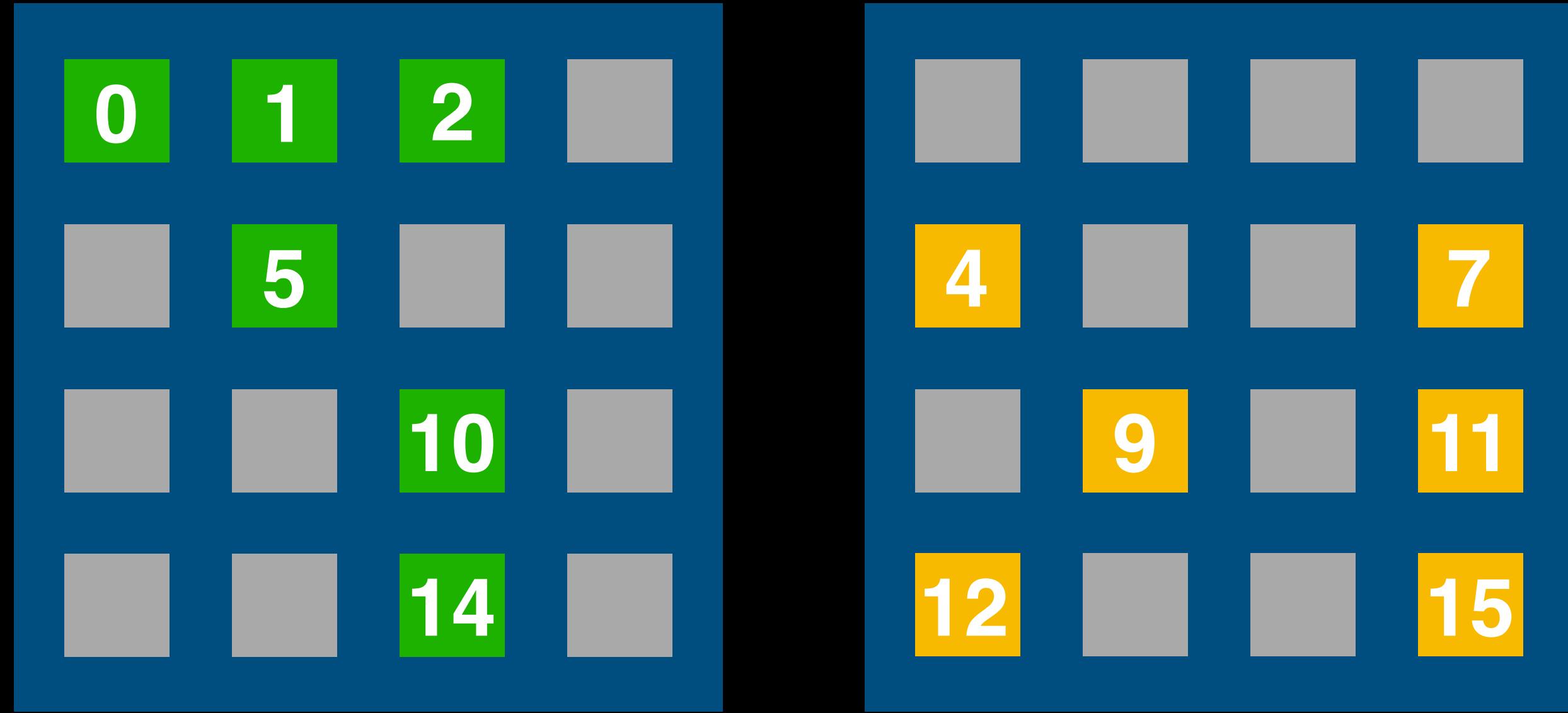
5x reduction in
time spent
compacting

Meshing: compaction without
(virtual!) relocation



Pages are **Meshable** when they:

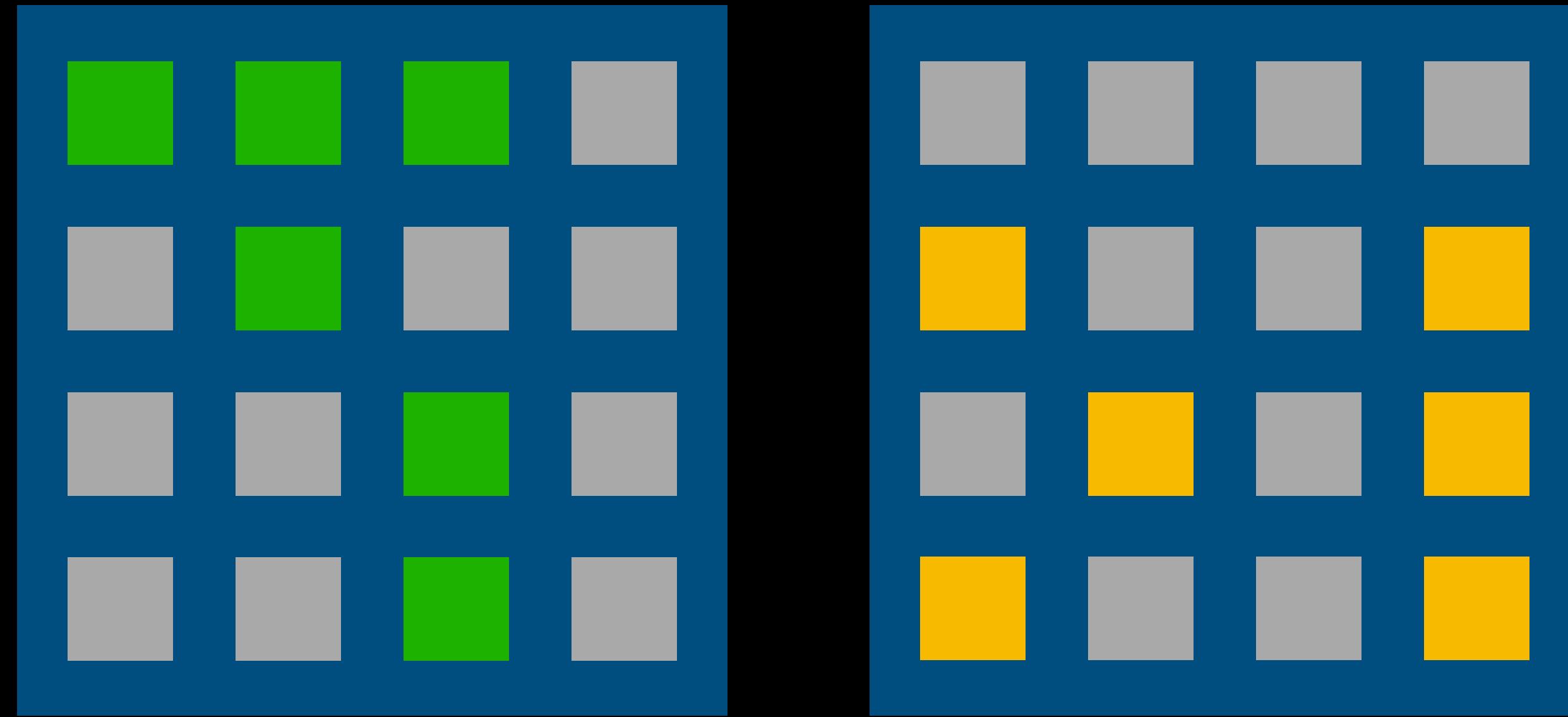
- Hold objects of the same size class



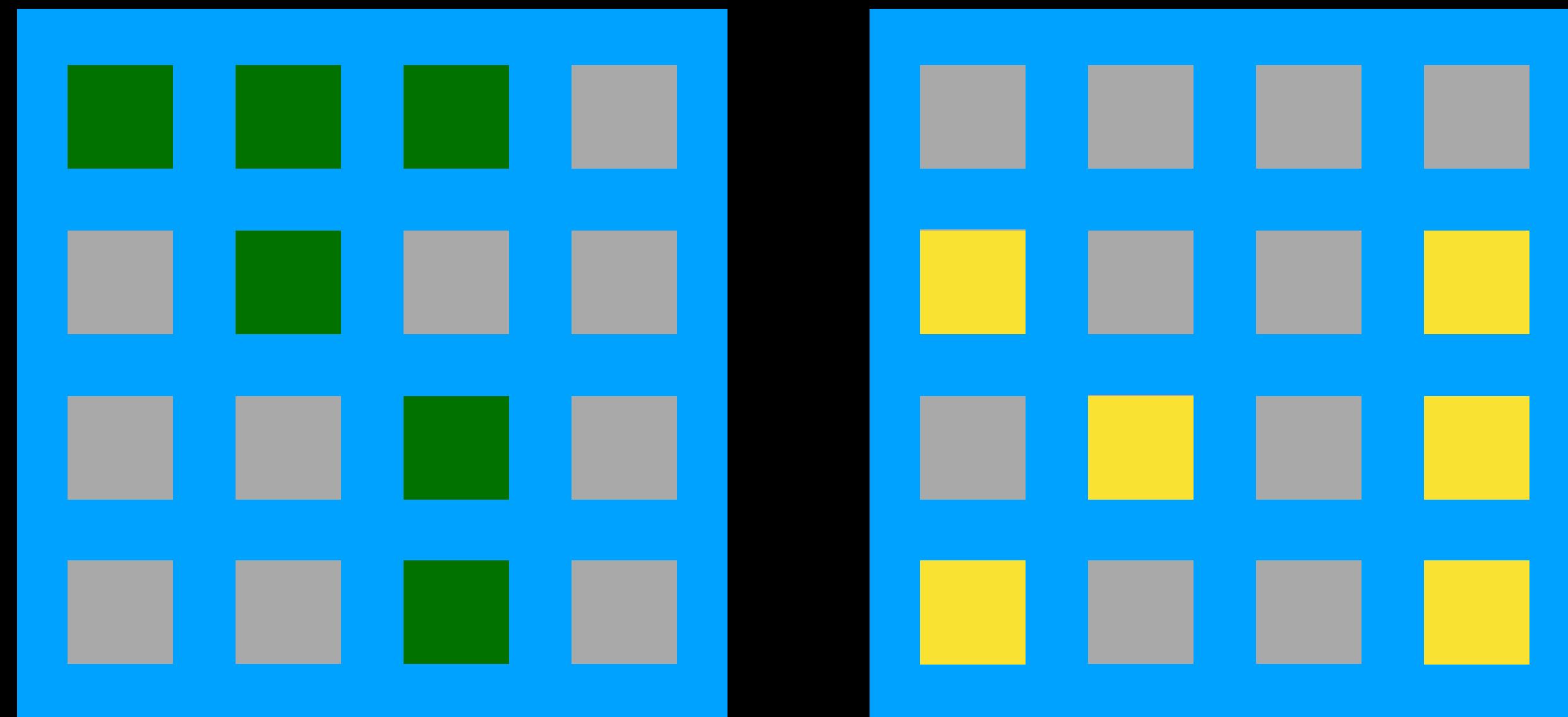
Pages are **Meshable** when they:

- Hold objects of the same size class
- Have non-overlapping object offsets

Mes***hing***

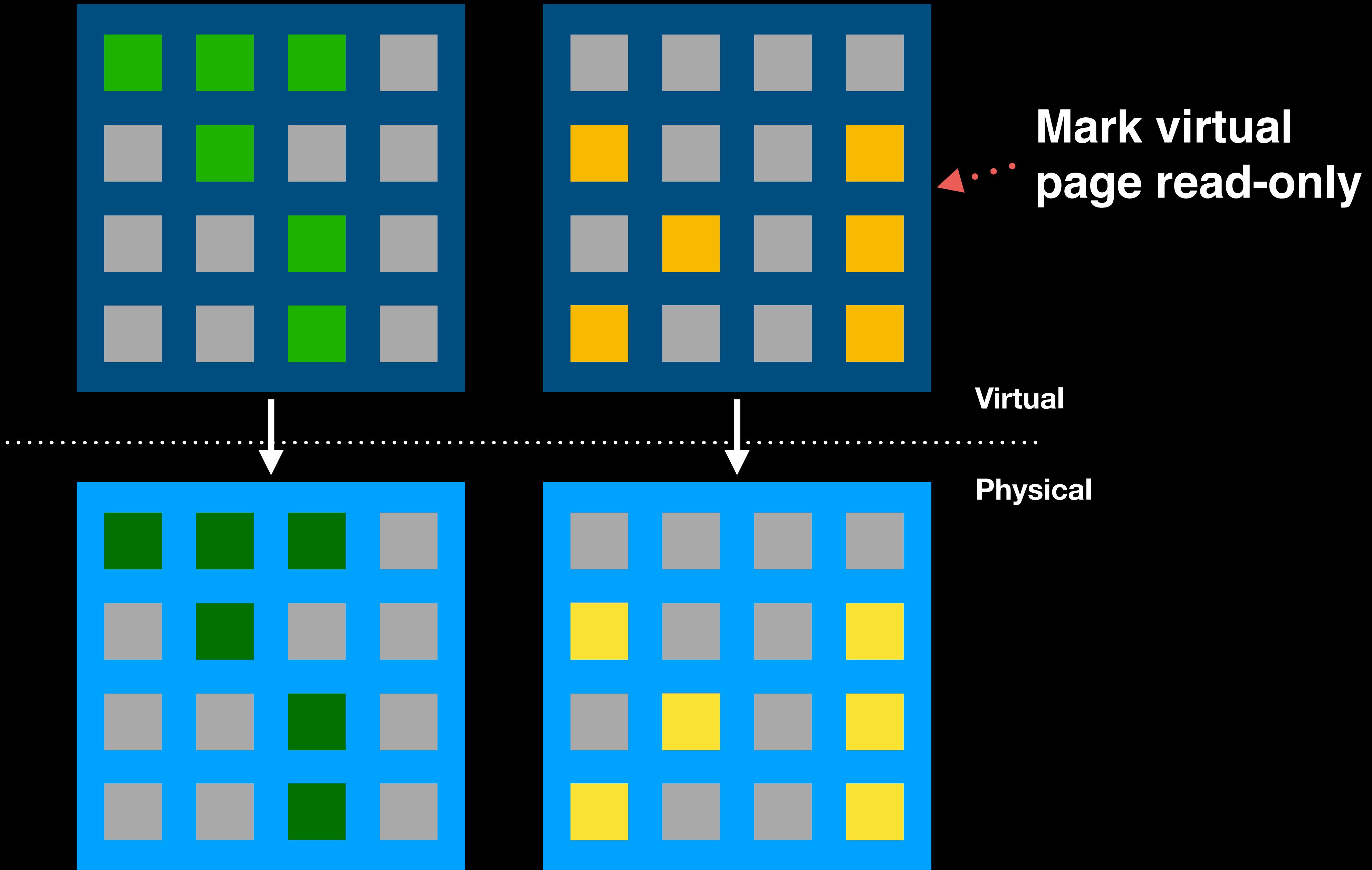


Virtual

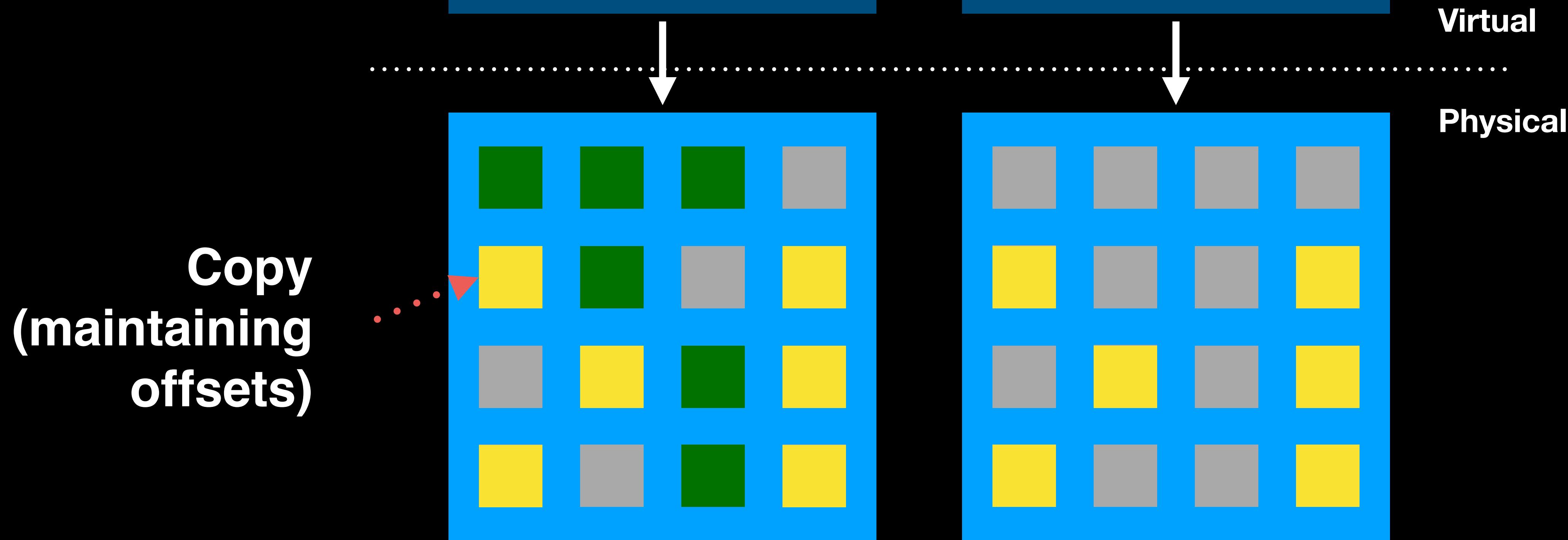


Physical

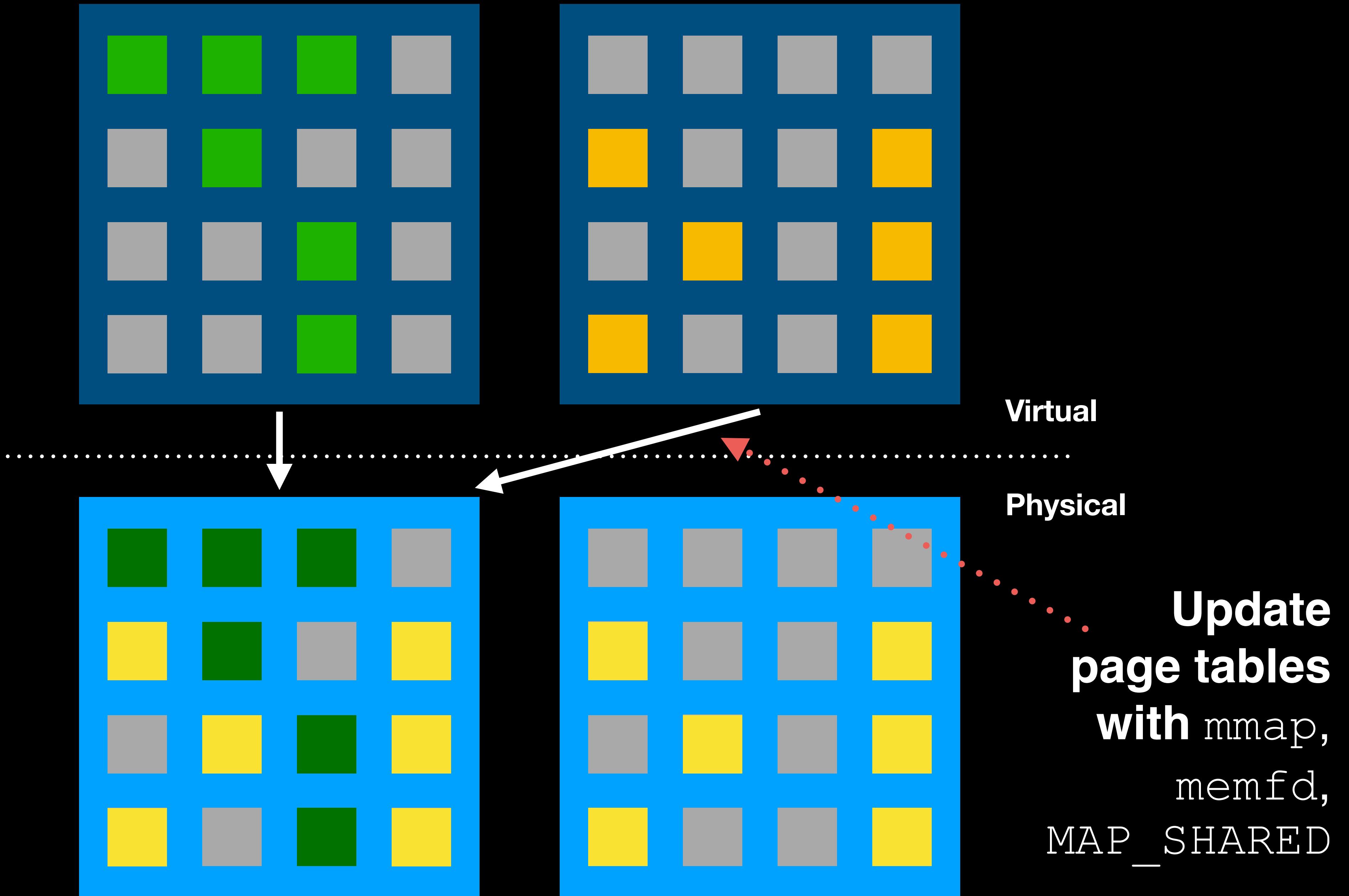
Meshing



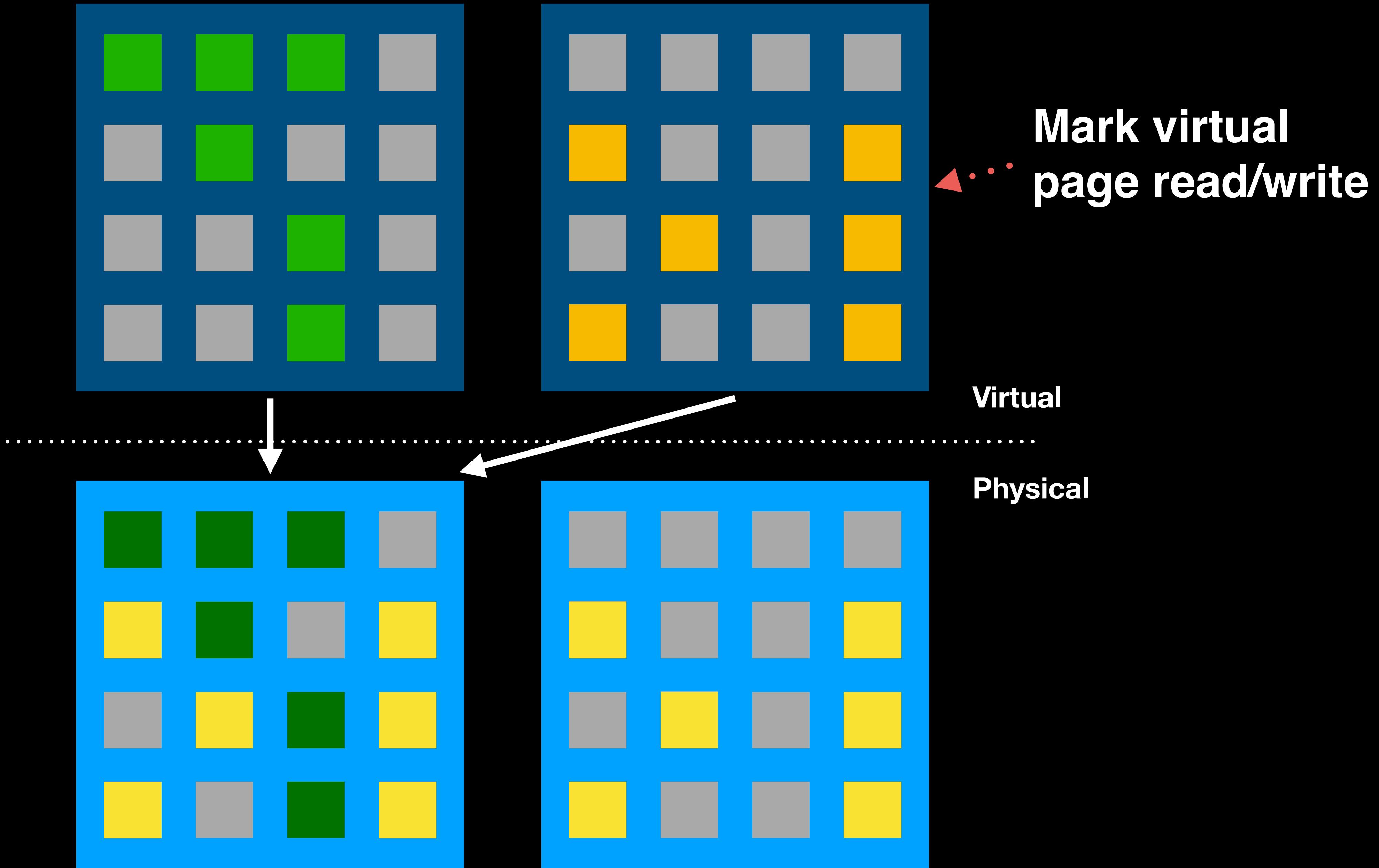
Meshing



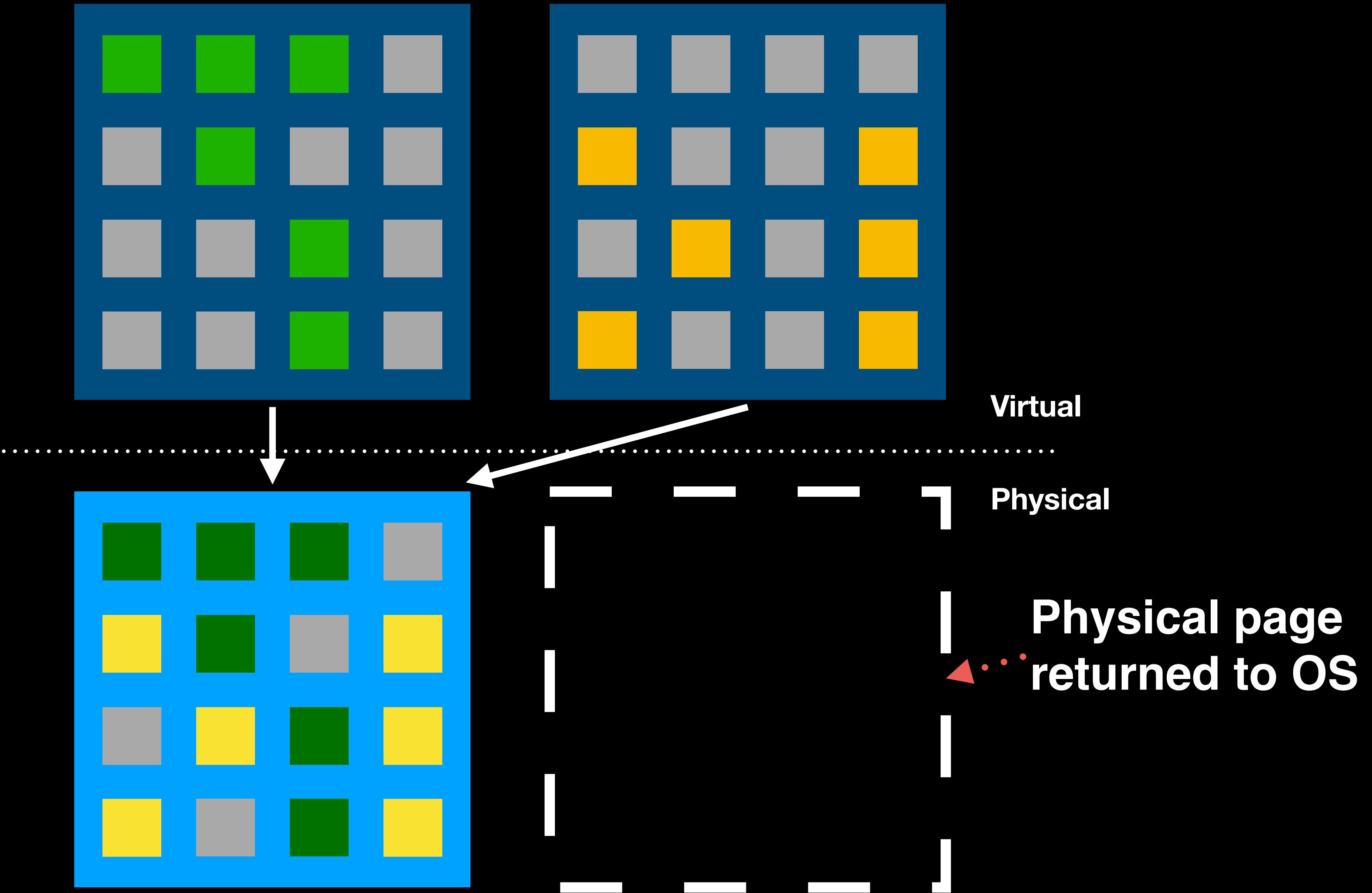
Meshing



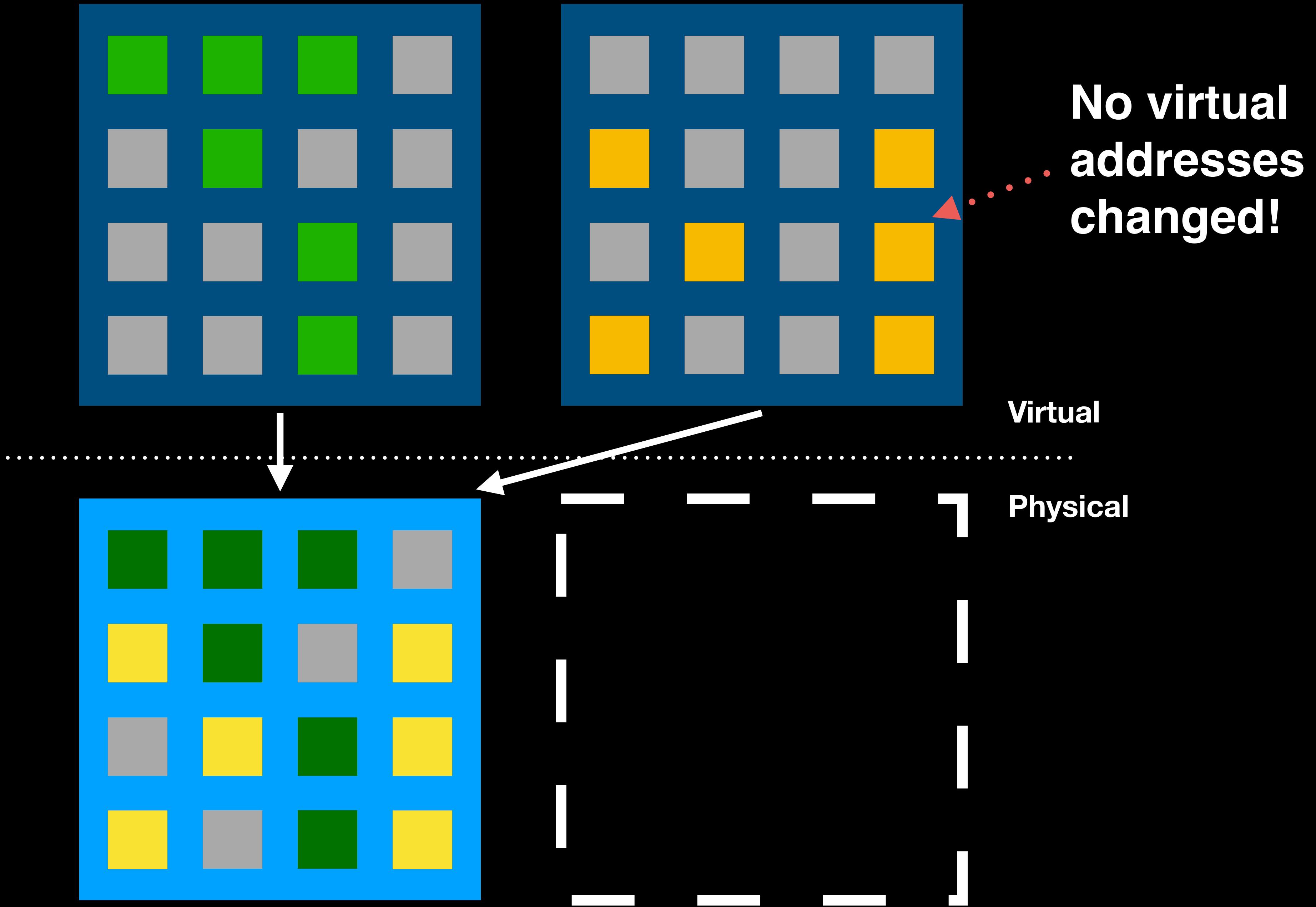
Meshing



*Mes*hing



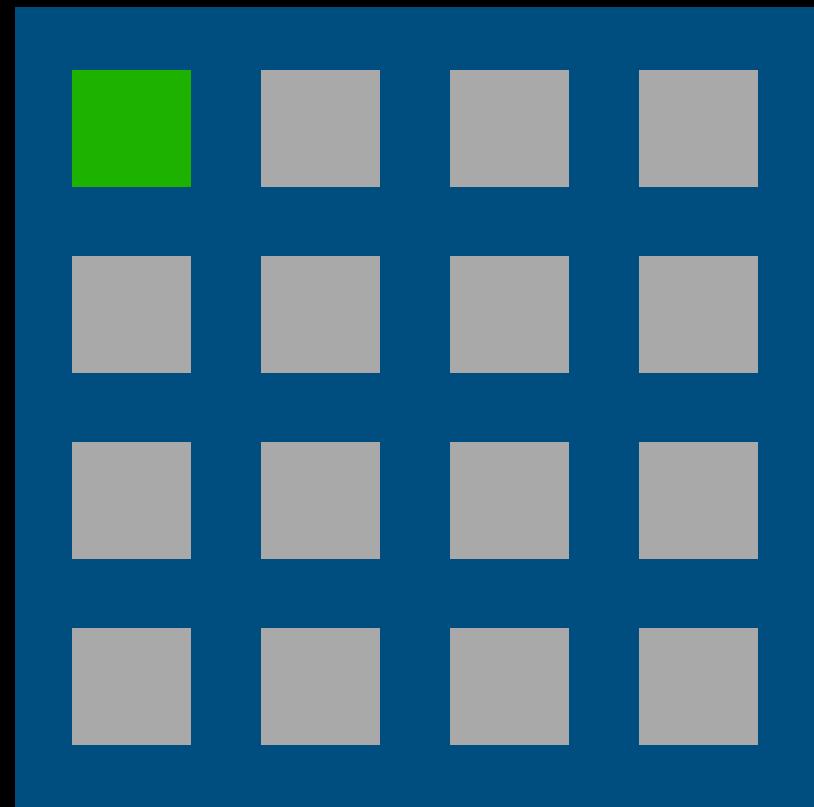
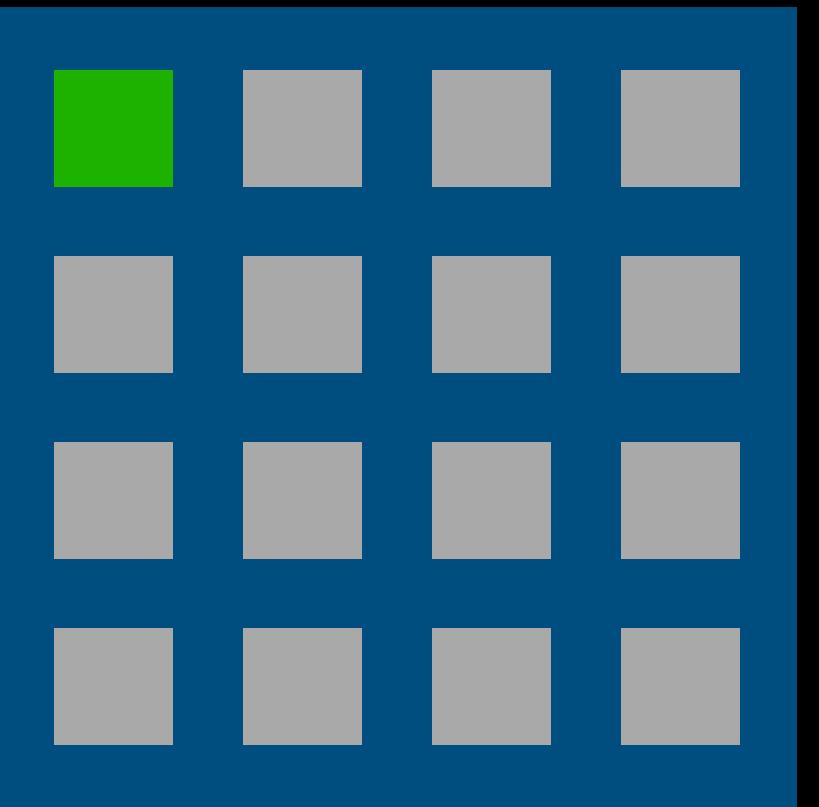
Meshing



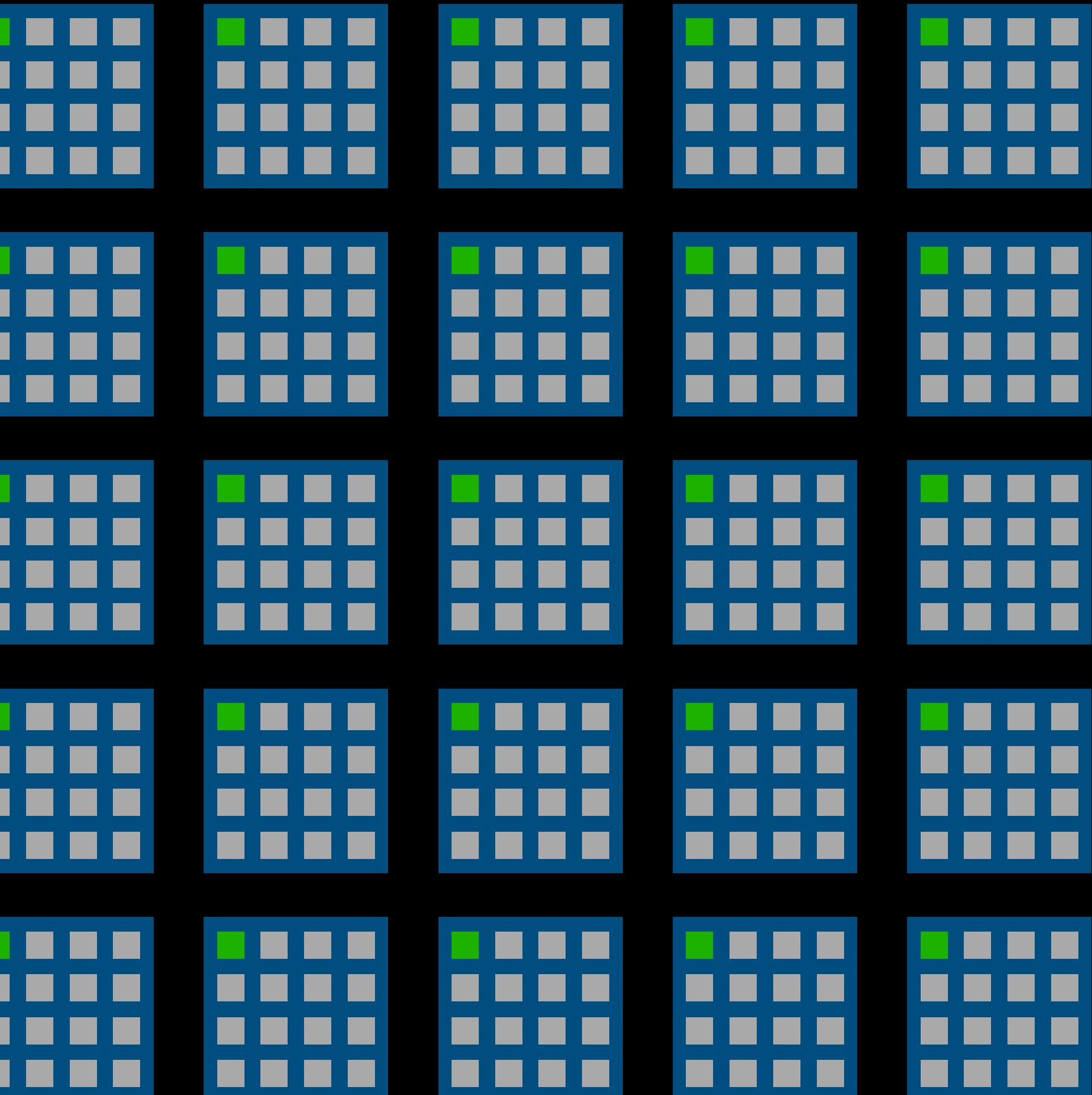
Worst Case:

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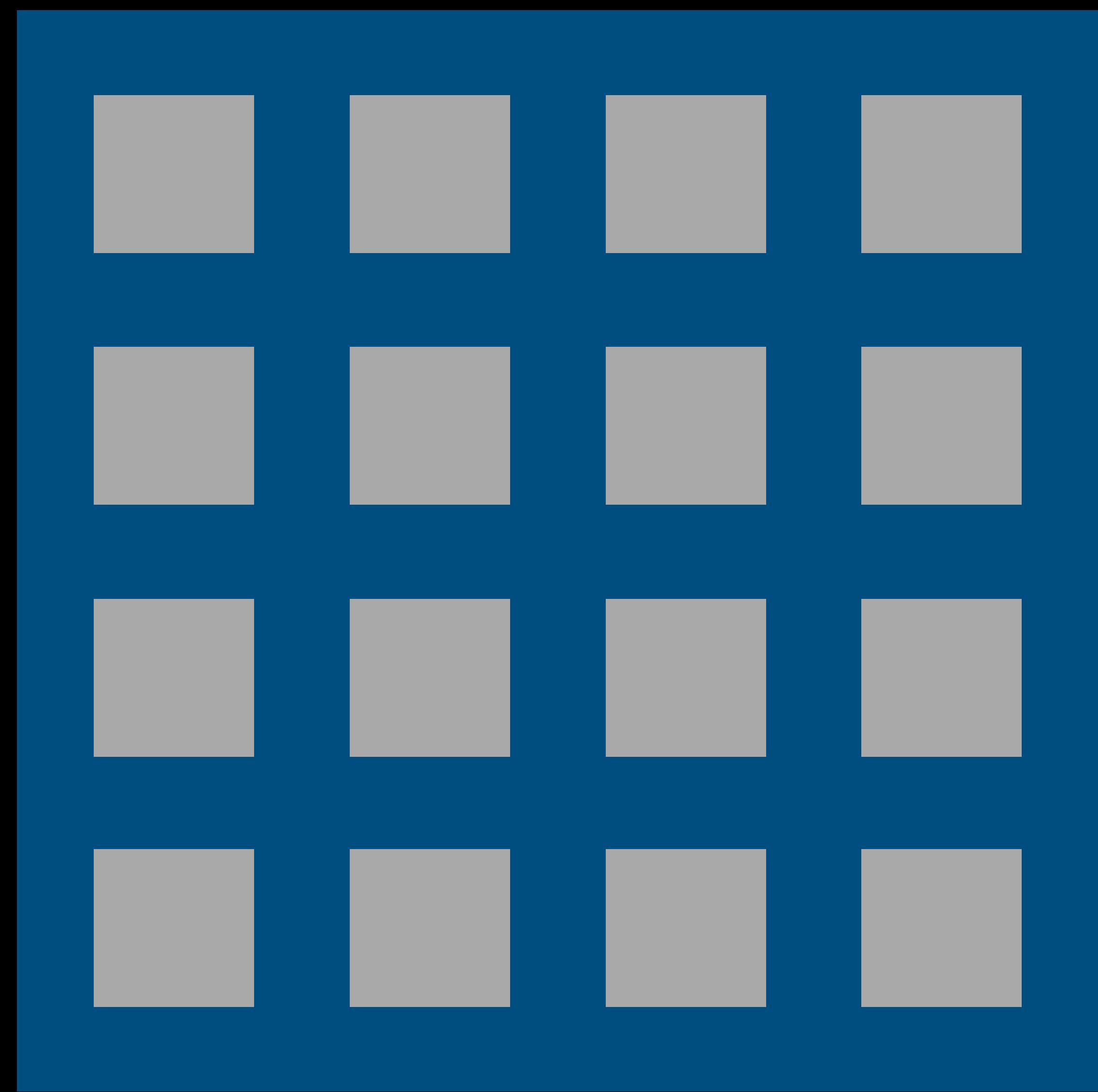
low occupancy,
non-meshable
pages



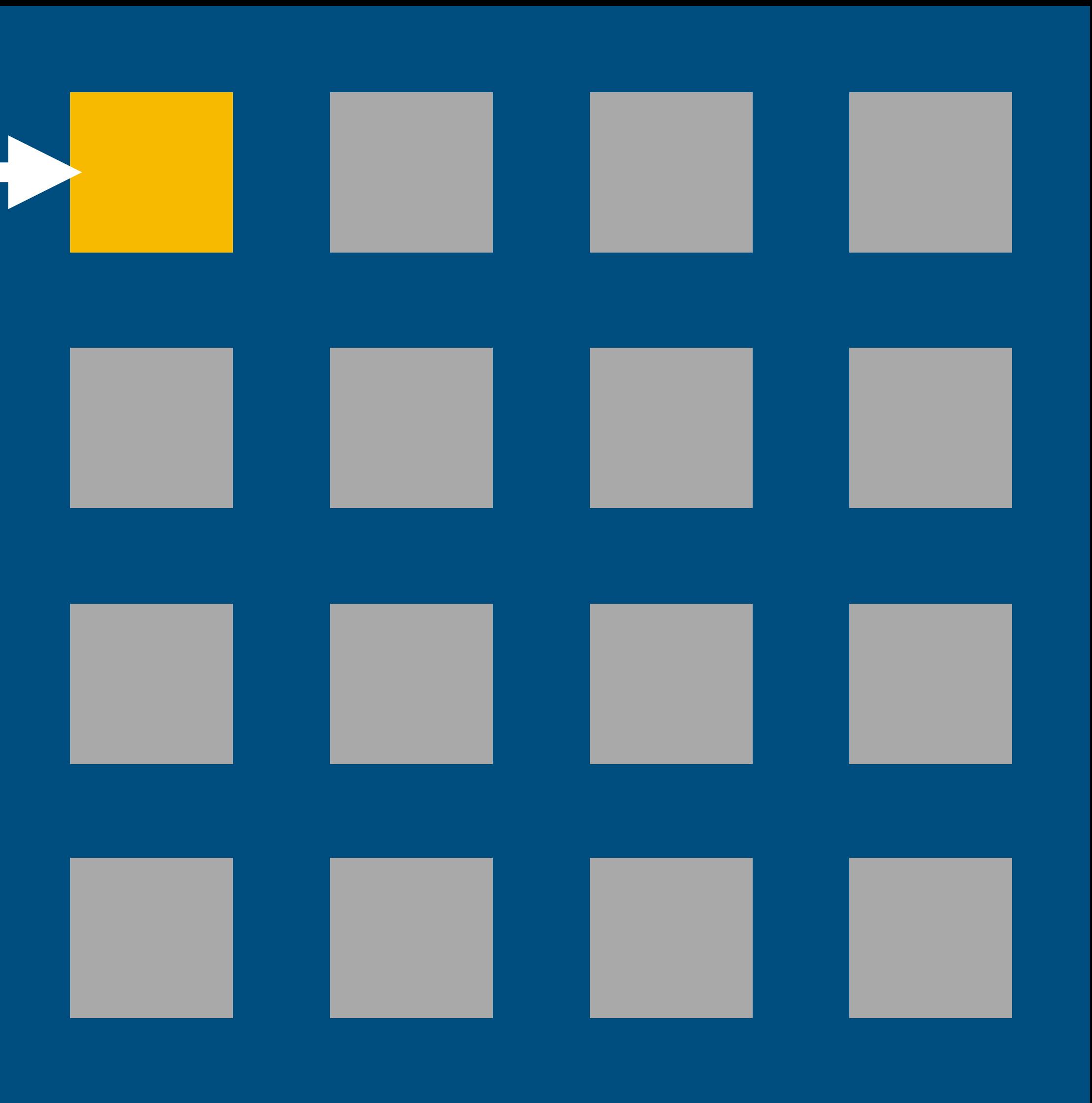
Worst Case:
many
low occupancy,
non-meshable
pages



Standard allocators



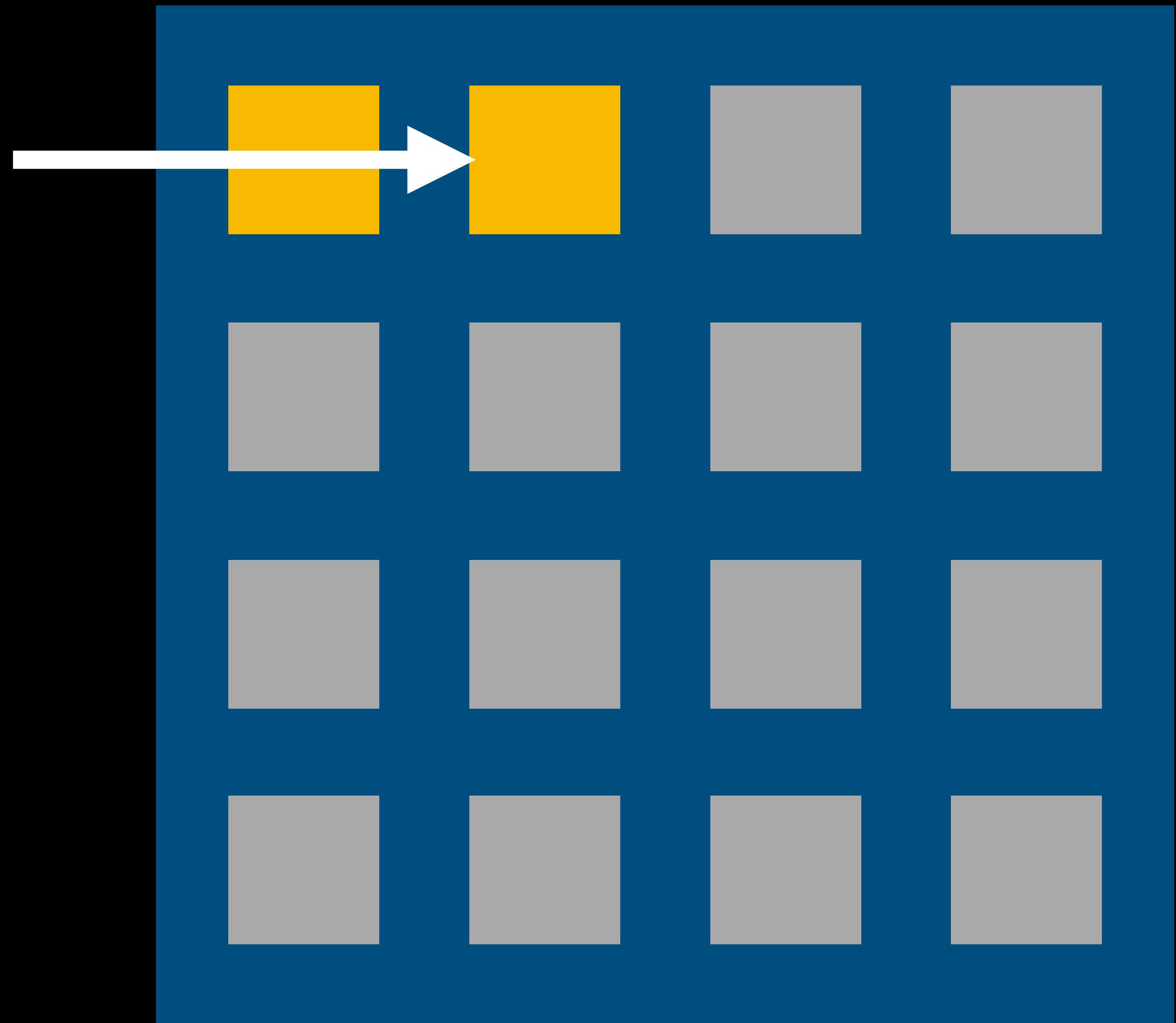
`malloc(256)`



**Standard
allocators**

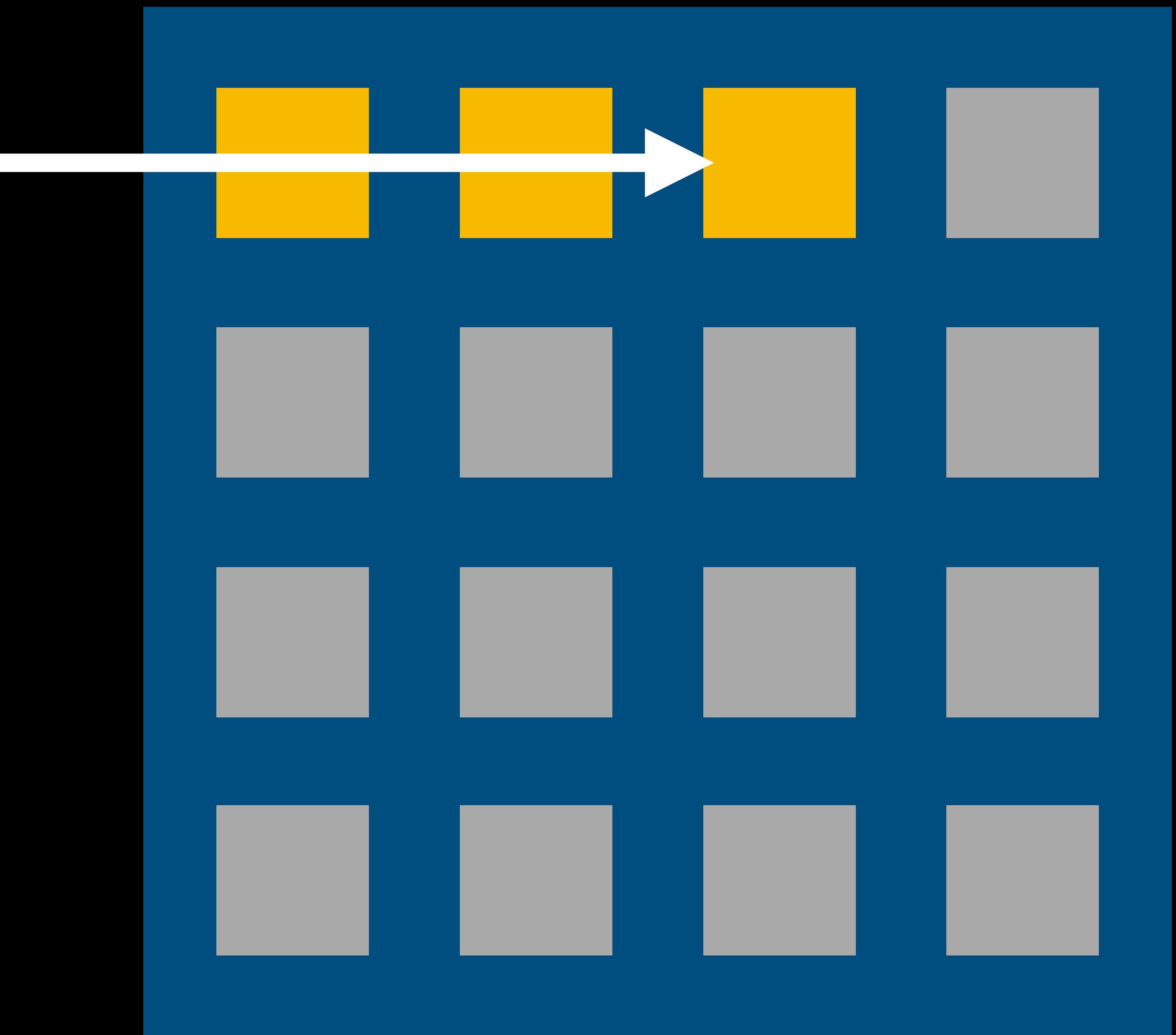
malloc(256)

**Standard
allocators**



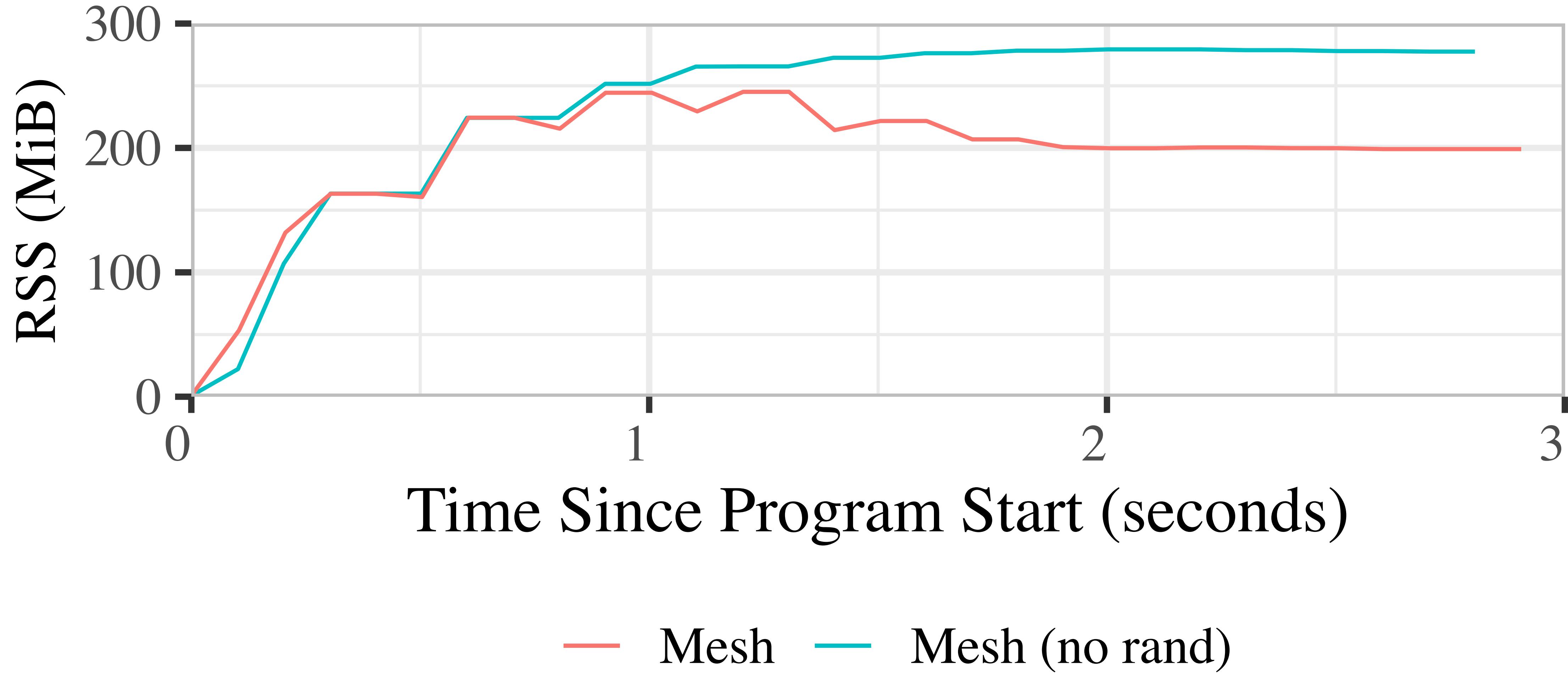
`malloc(256)`

**Standard
allocators**



Mesh uses **randomization** to ensure live objects are uniformly distributed

Regular allocation patterns are real



How to **randomize** allocation?

Random probing:

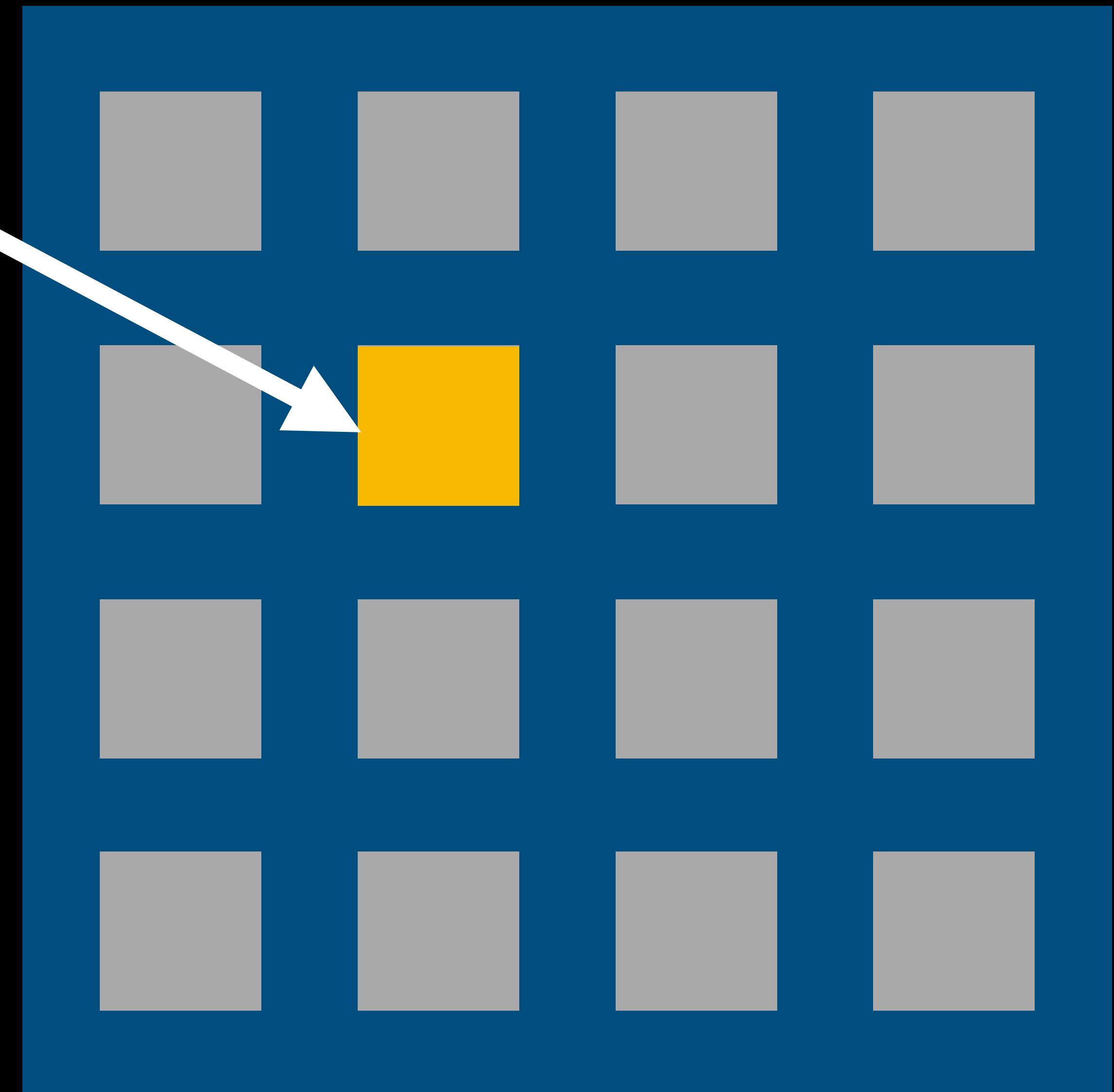
```
while true:  
    if rand_off().is_free:  
        return rand_off
```



malloc(256)

Random probing:

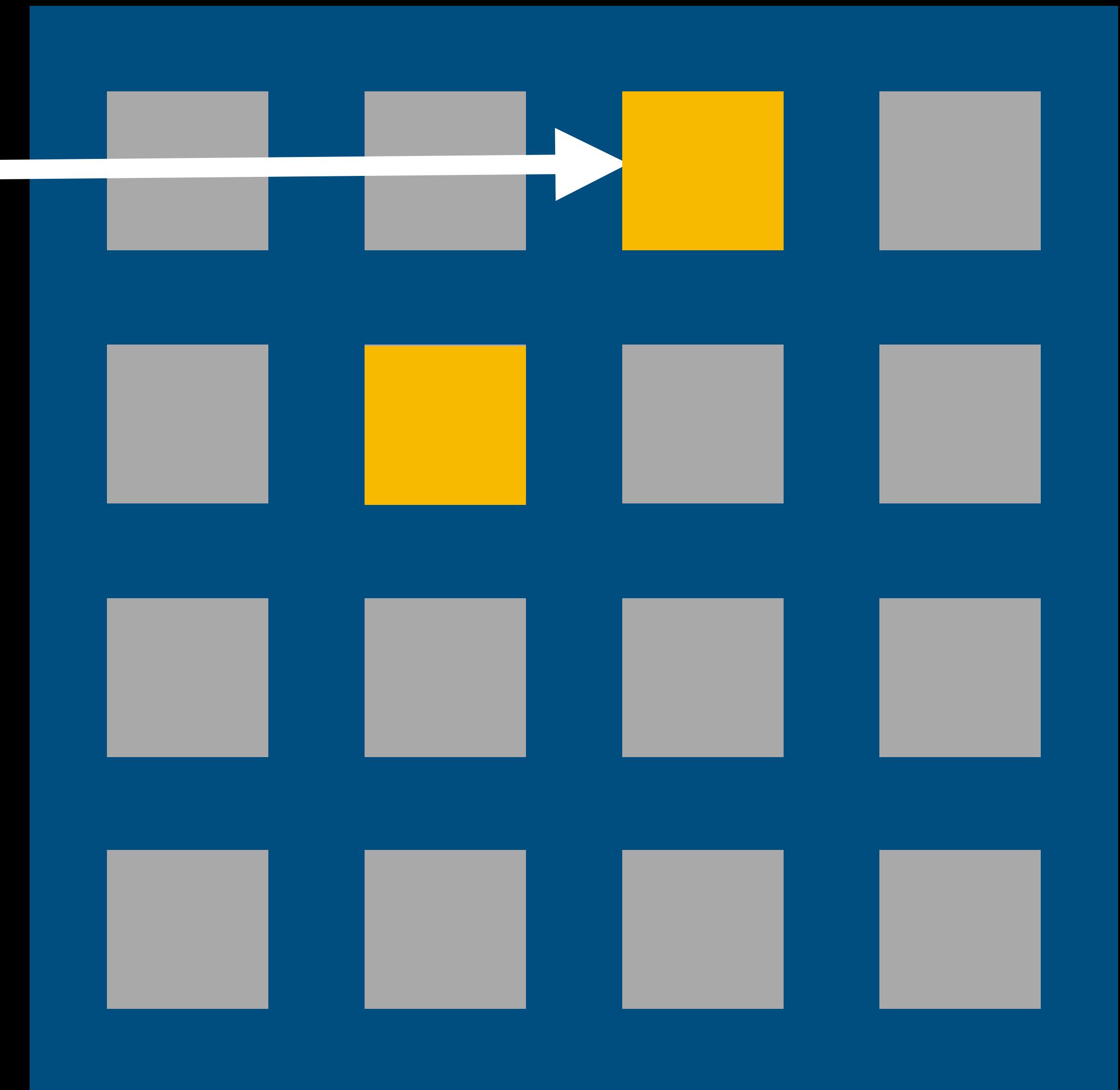
```
while true:  
    if rand_off().is_free:  
        return rand_off
```



malloc(256)

Random probing:

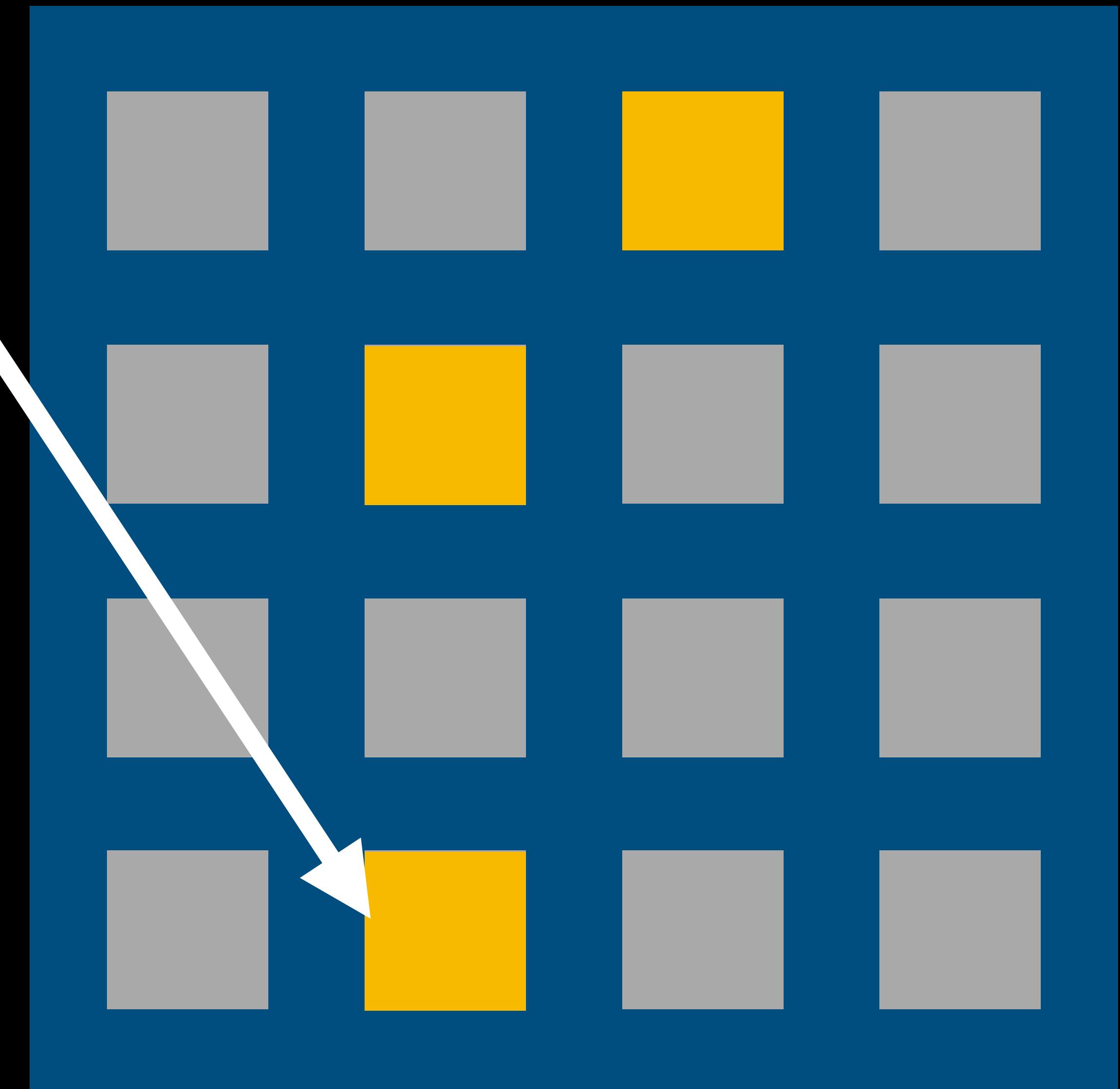
```
while true:  
    if rand_off().is_free:  
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```



```
malloc(256)
```

Random probing:

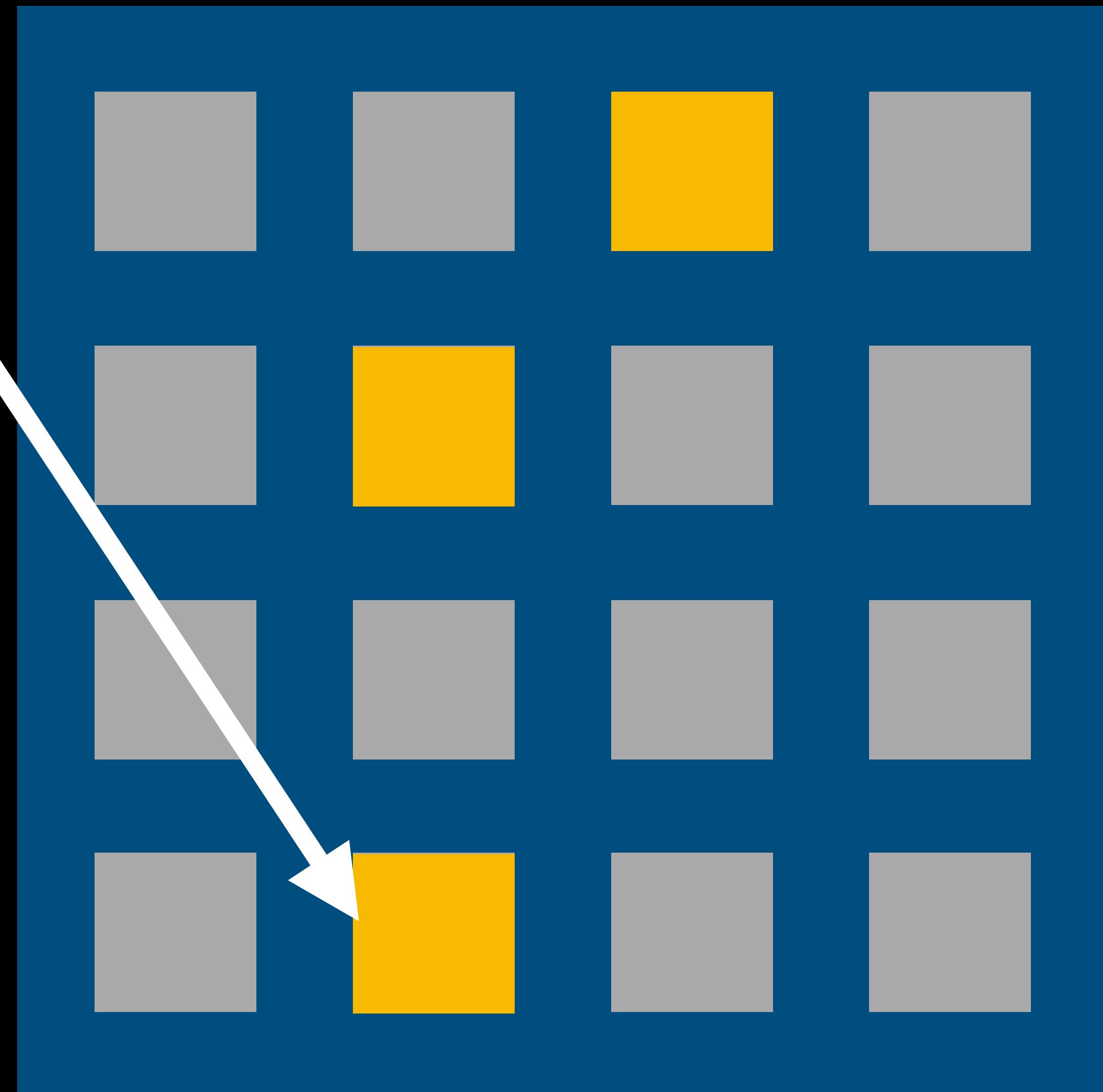
```
while true:  
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```



malloc(256)

Random probing:

```
while true:  
    if rand_off().is_free:  
        return rand_off
```



(DieHard [Berger & Zorn 2006])

Random probing fast in
expectation *iff page*
occupancy is low

Random probing fast in
expectation *iff page*
occupancy is low

**but this is at odds with
minimizing heap size!**

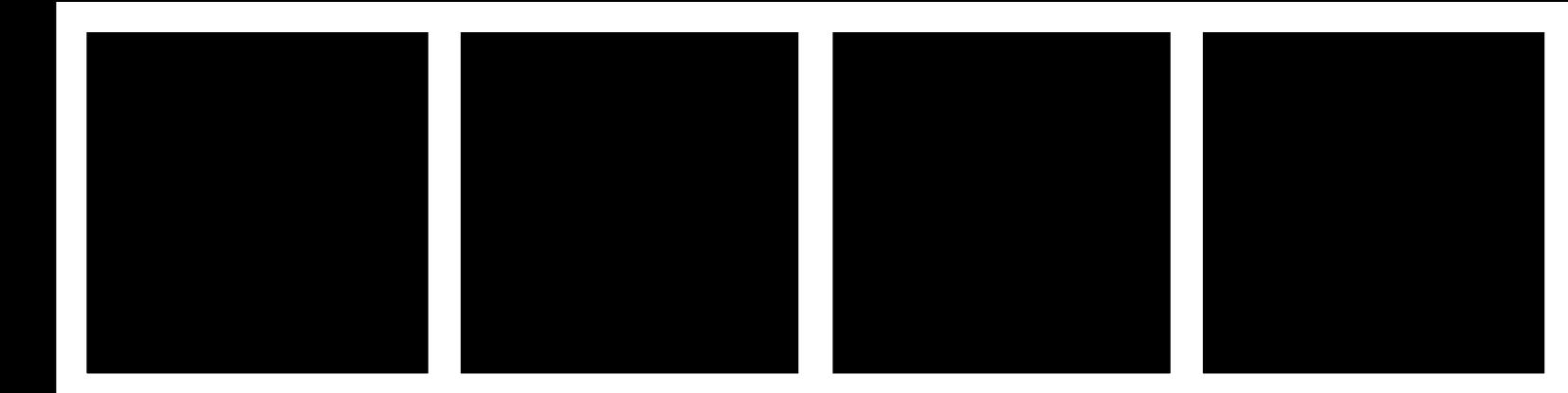
Shuffle Vector:
Fast randomized
allocation + full page utilization

Shuffle Vector: Fast randomized allocation

Page

0 1 2 3

load



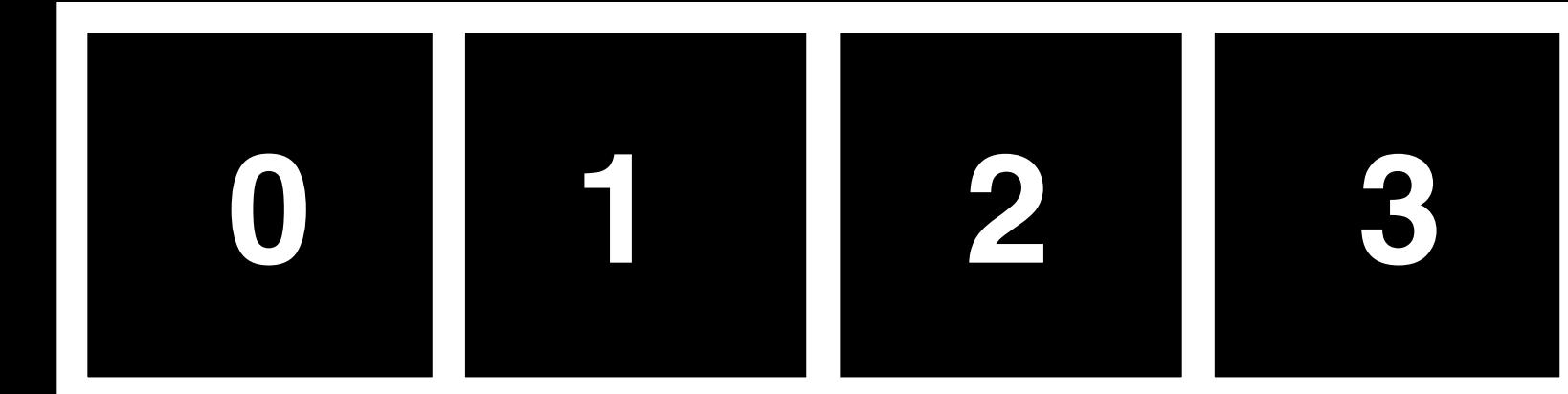
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



load



Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



shuffle (

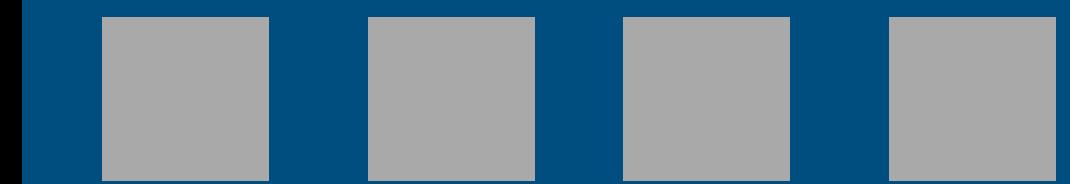
0	1	2	3
---	---	---	---

)

Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



shuffle (

2	3	1	0
---	---	---	---

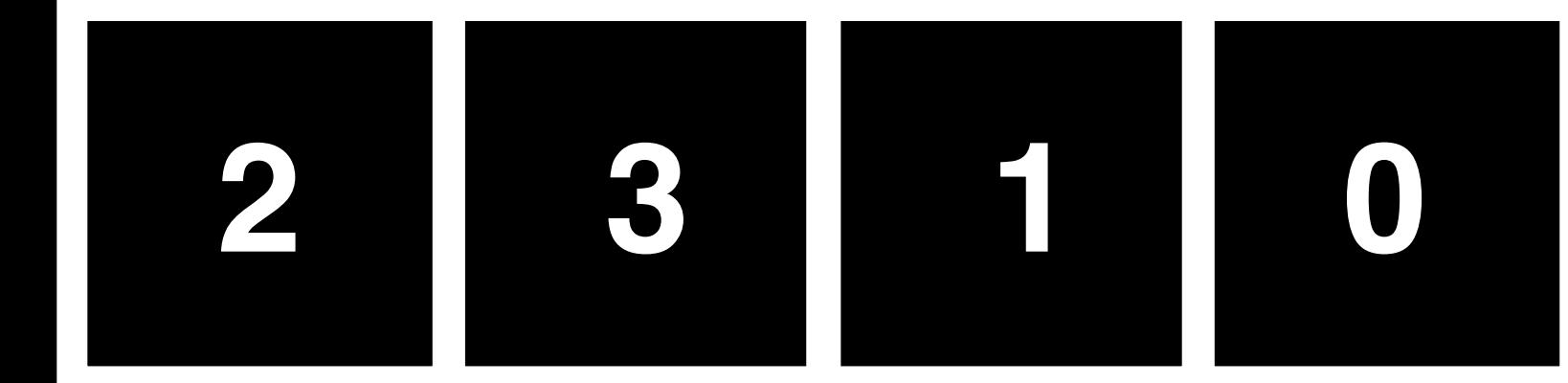
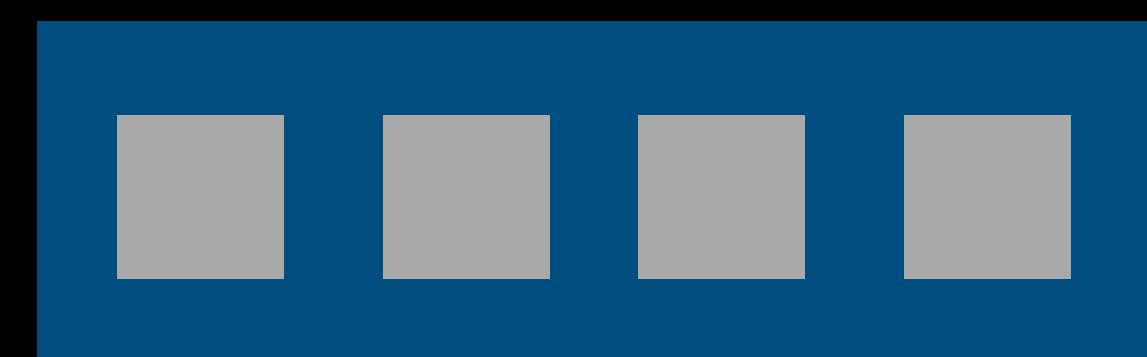
)

Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

`malloc()`

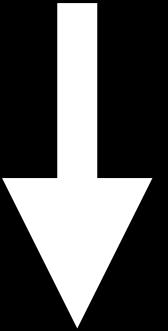
Page



Thread-local shuffle vector

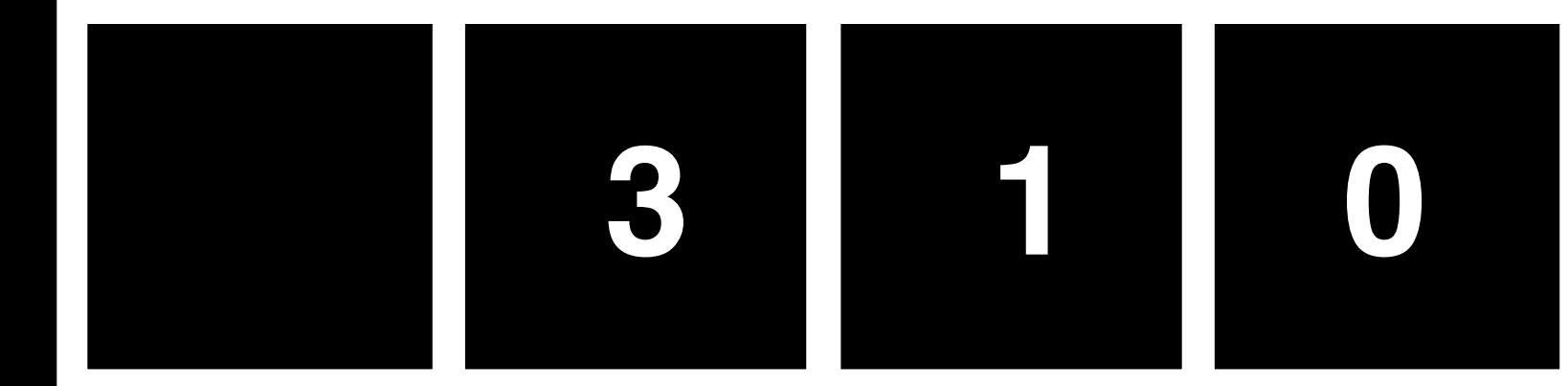
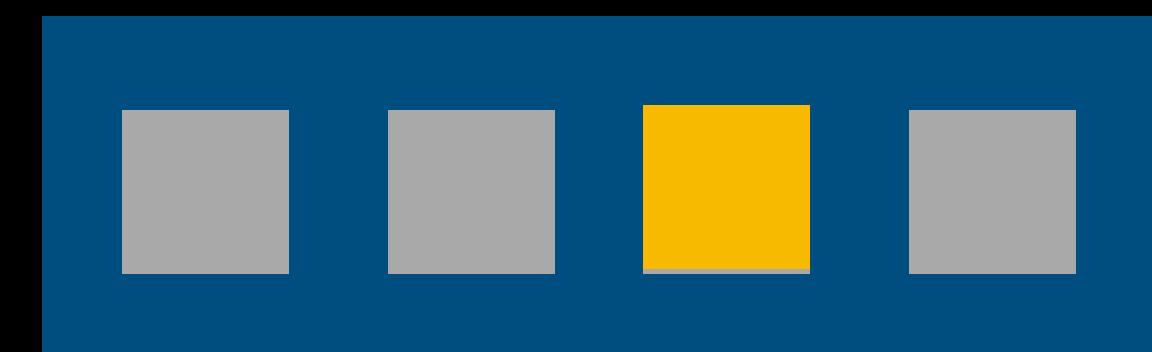
Shuffle Vector: Fast randomized allocation

malloc()



page_start +
2 * object_size

Page



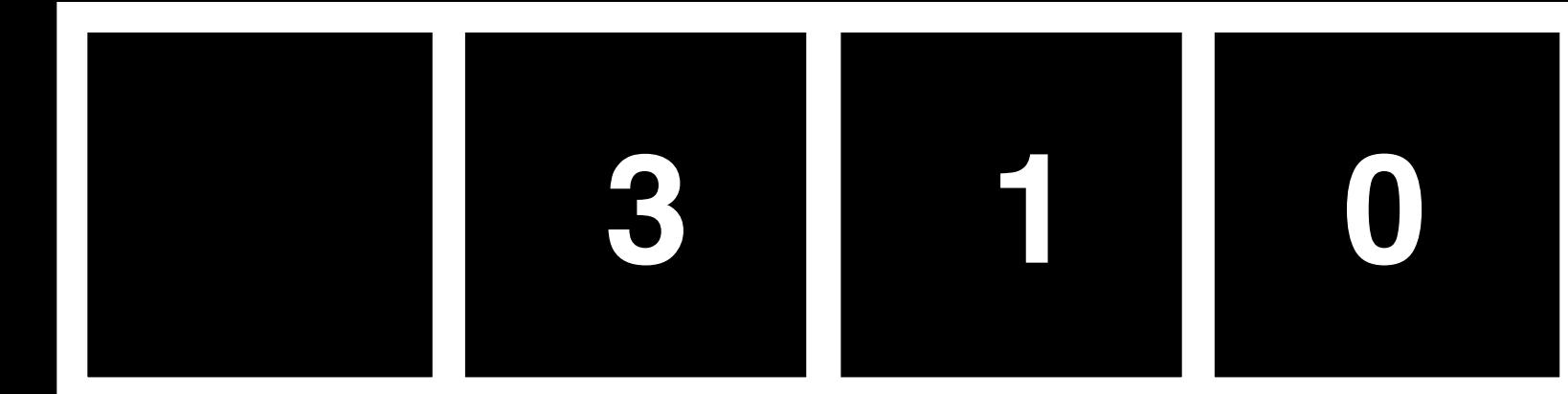
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



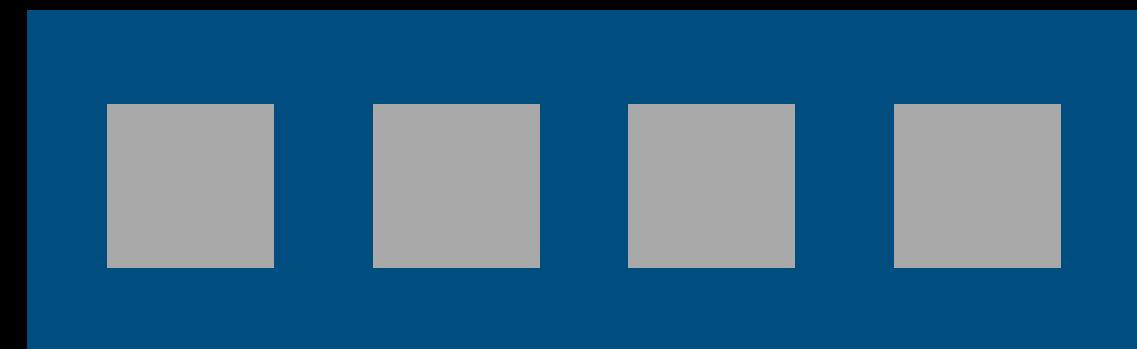
free (2)



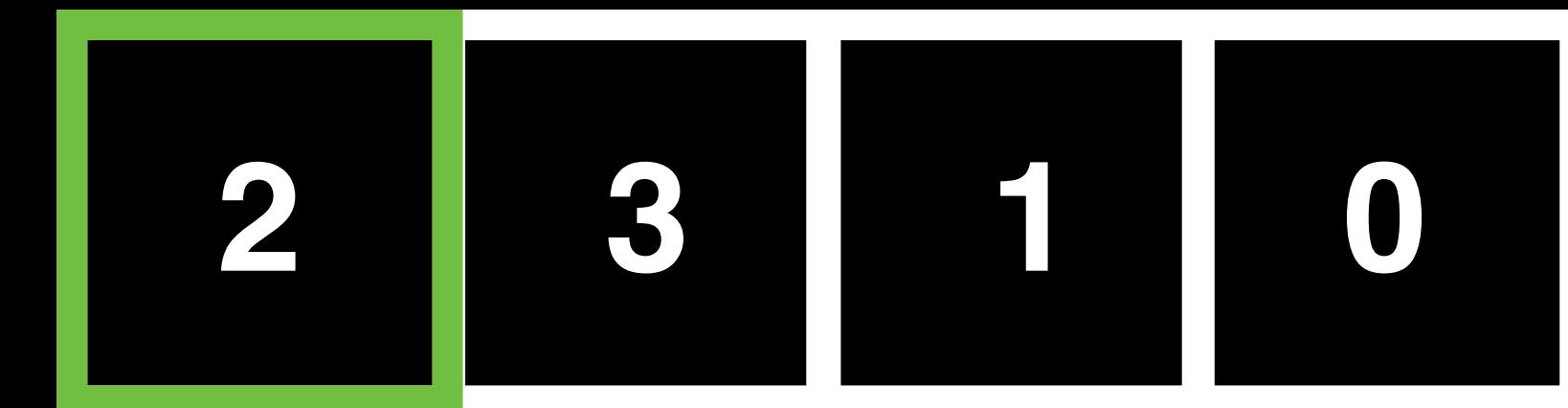
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



free()



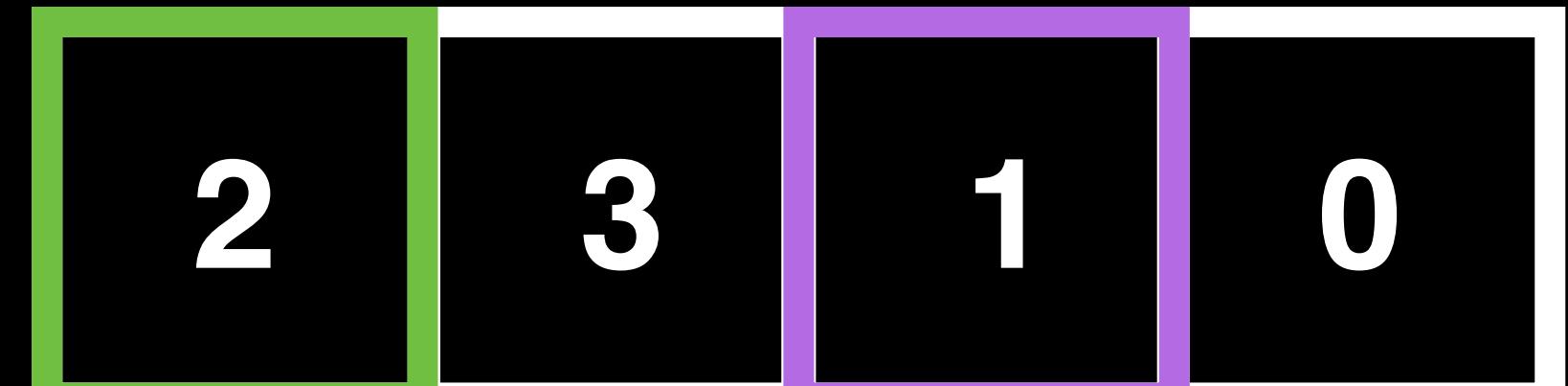
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



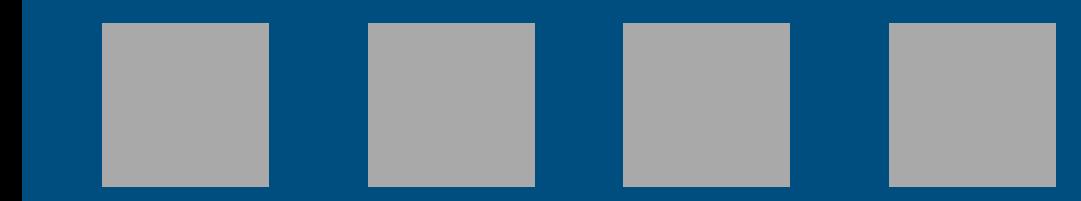
free()

shuffle_one( 0)

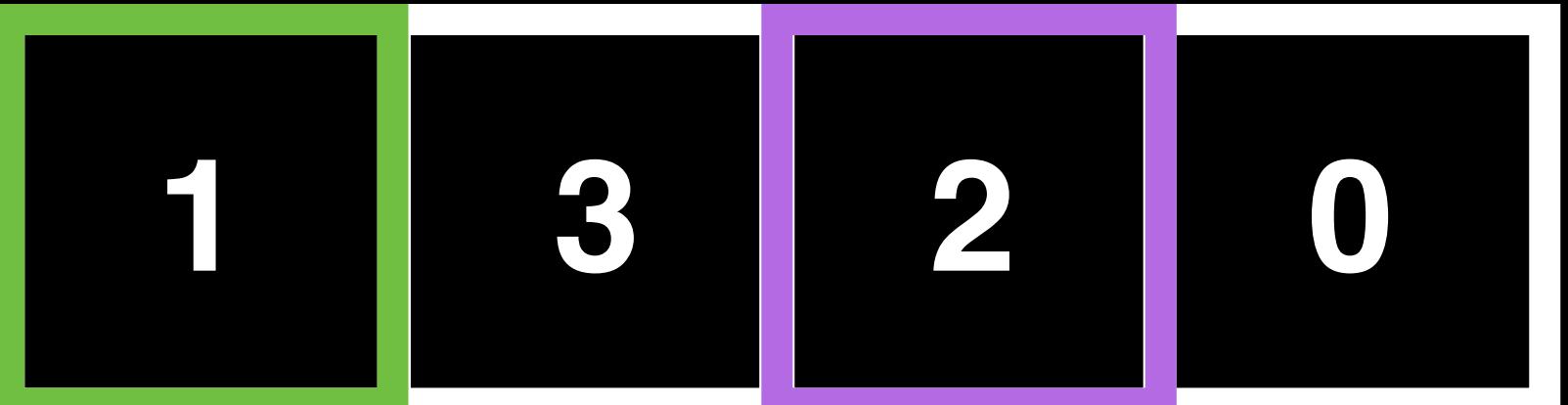
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



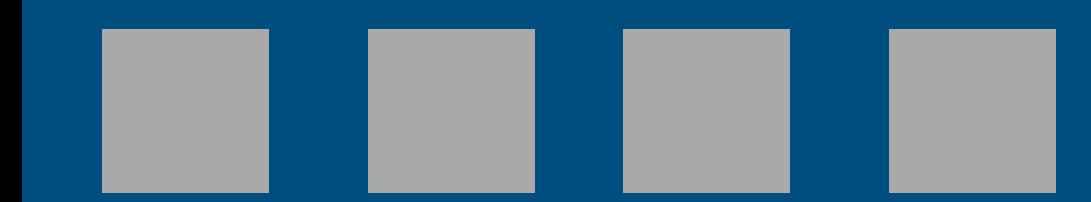
free()

shuffle_one( 0)

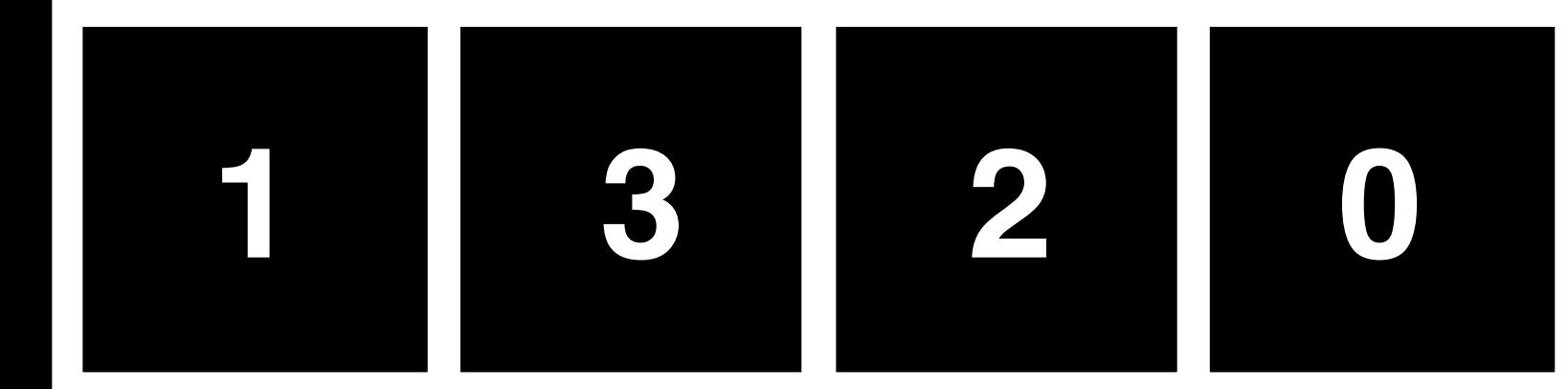
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page

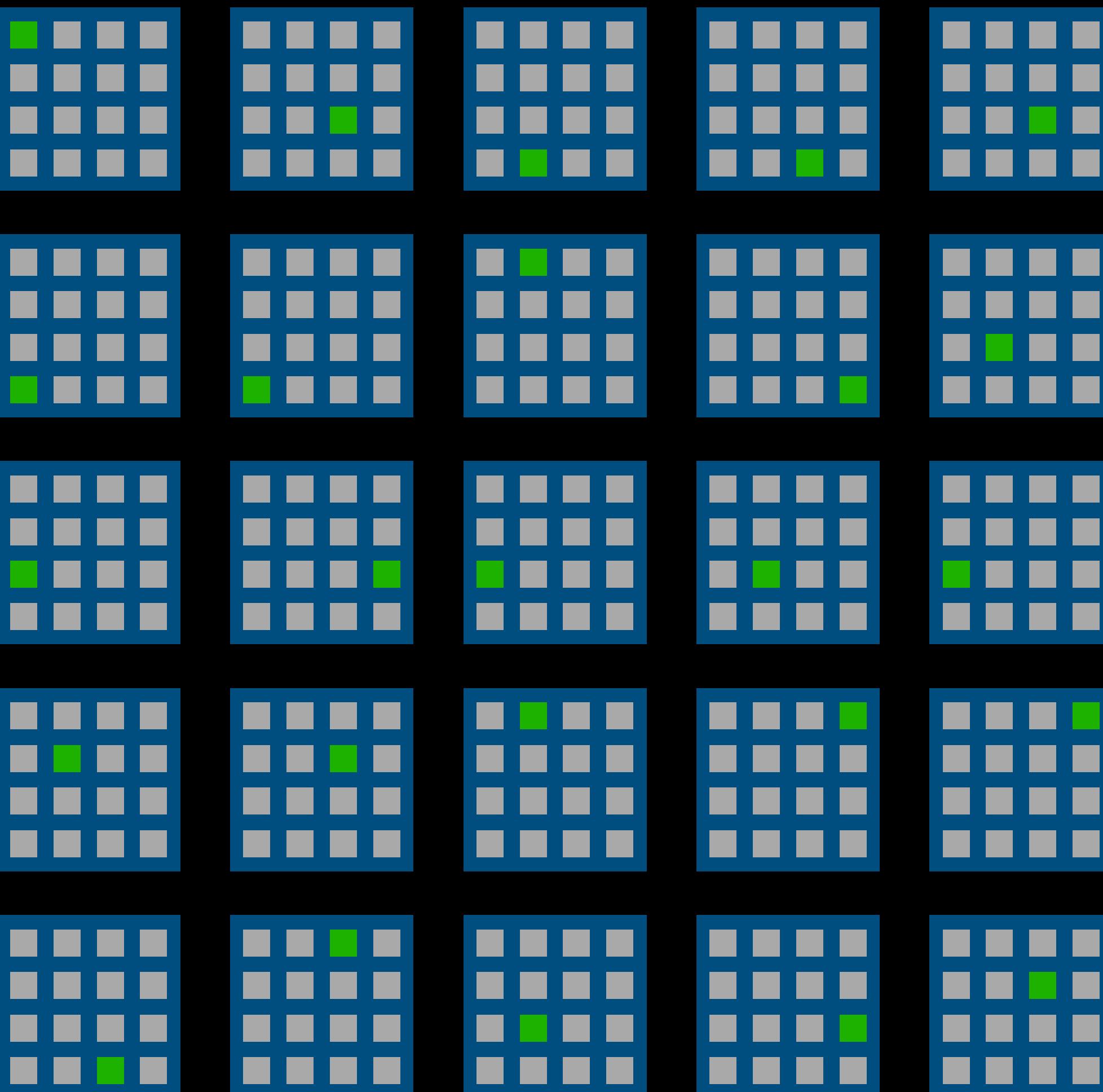


free()



Thread-local shuffle vector

**All
Pages
Meshable**



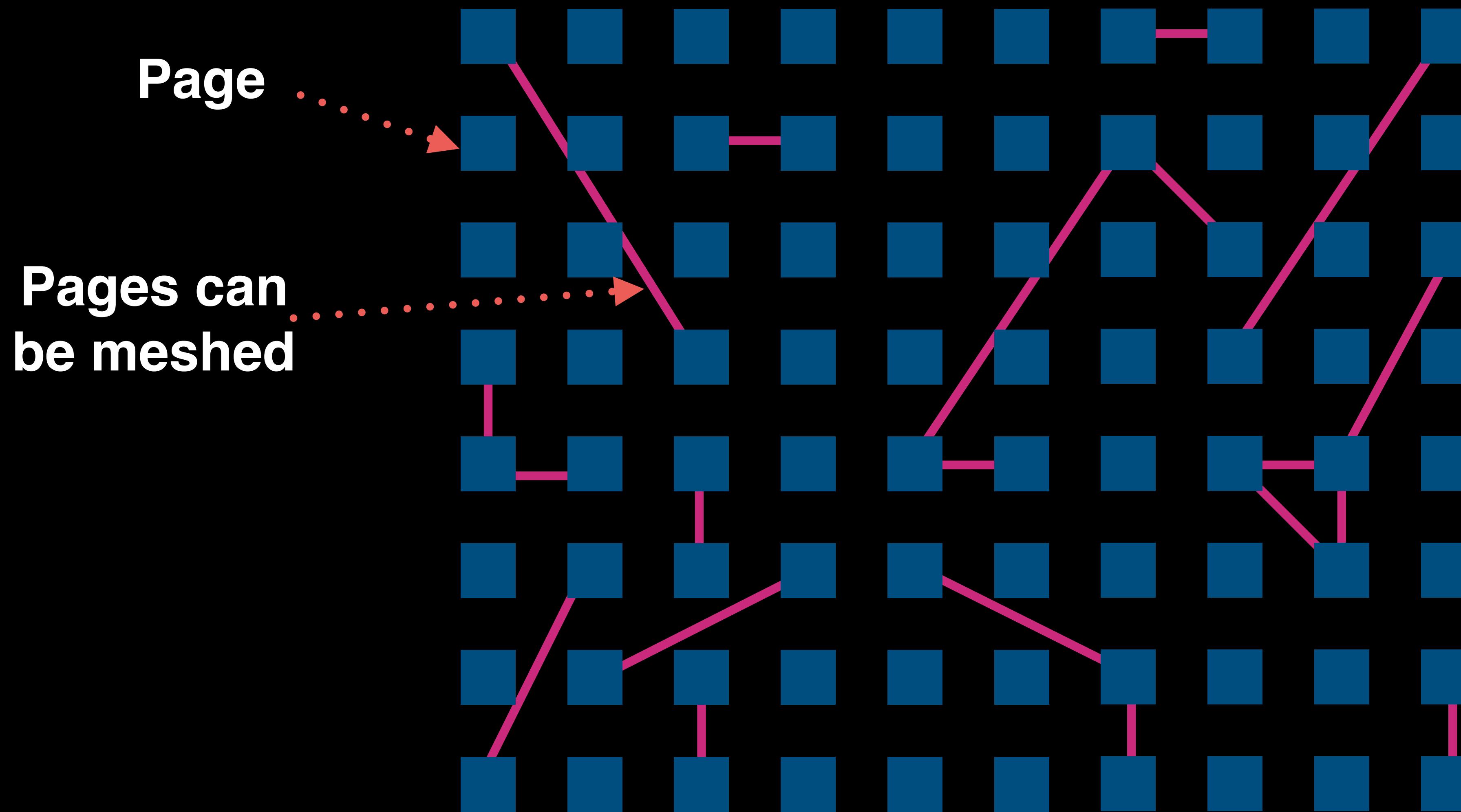
Finding pages to Mesh

Problem: Find meshing that releases maximum number of pages

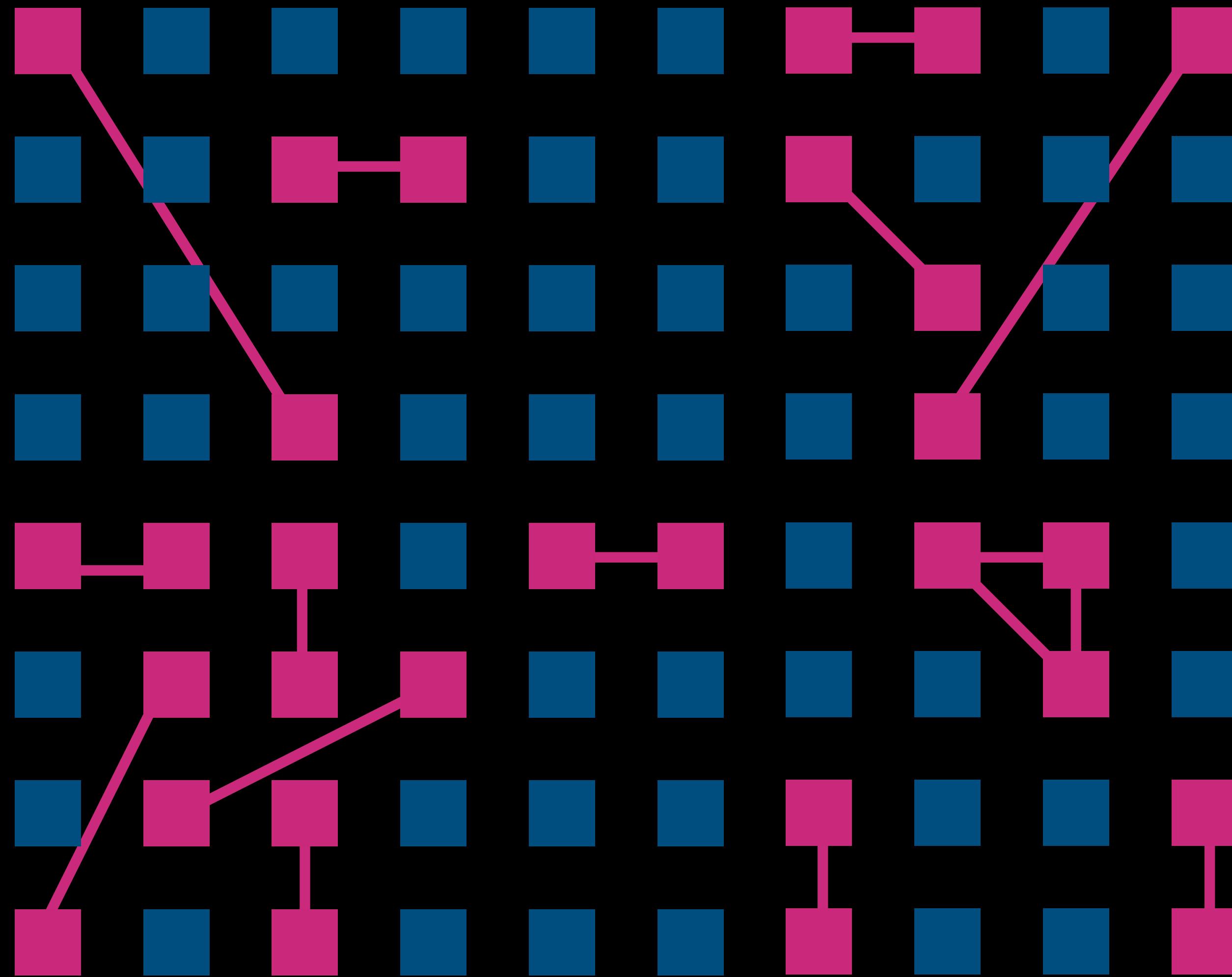
Problem: Find meshing that releases maximum number of pages

Run in the free () slowpath
At most once every 100 ms
Treat each size class independently

Problem: Find meshing that releases the maximum number of pages



MinCliqueCover



MinCliqueCover

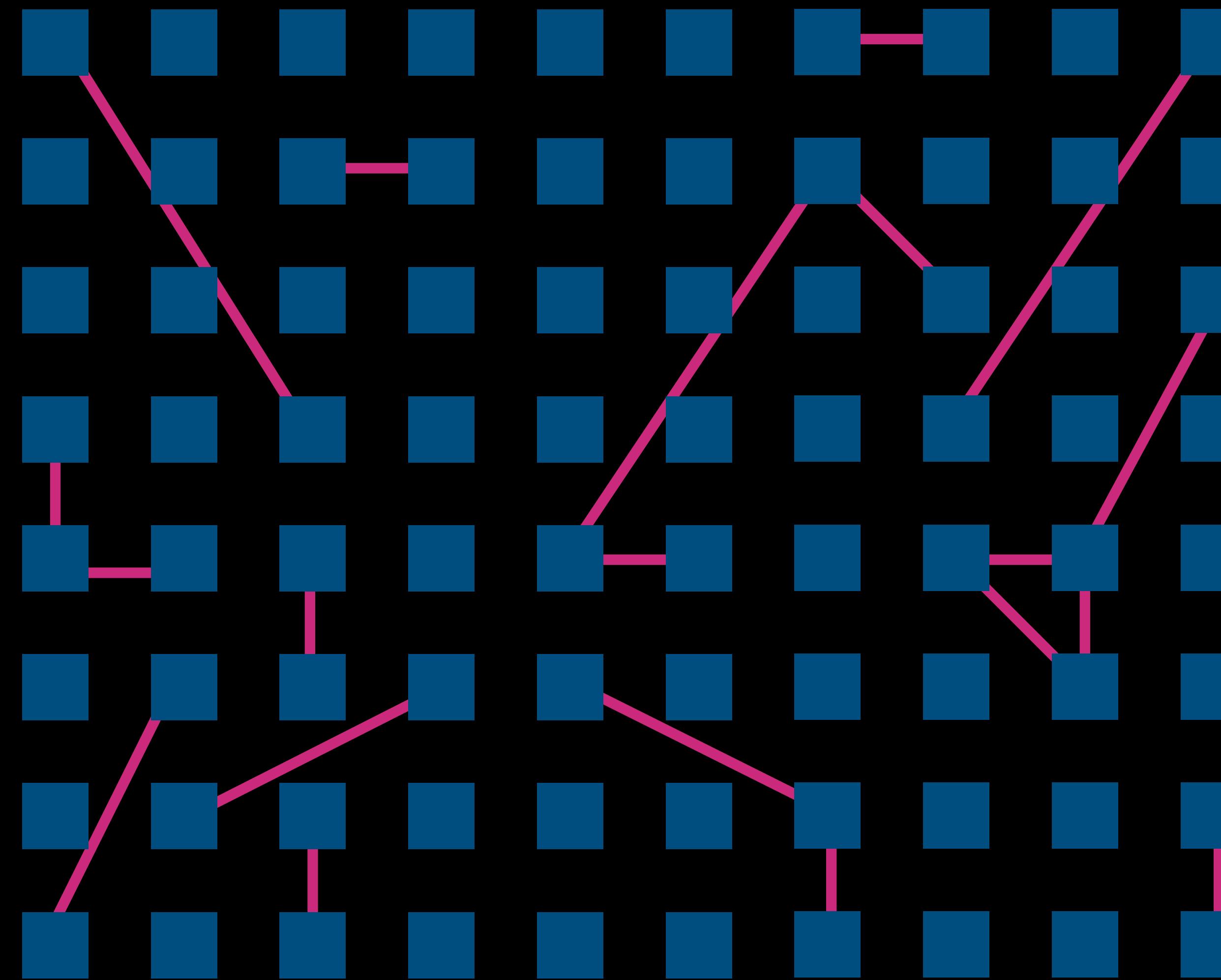
(NP-Complete)

MinCliqueCover

(NP-Complete)

BUT! Randomness ensures we can get
away with solving simpler graph problem
(Matching)

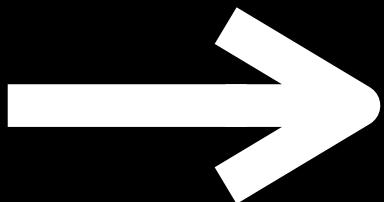
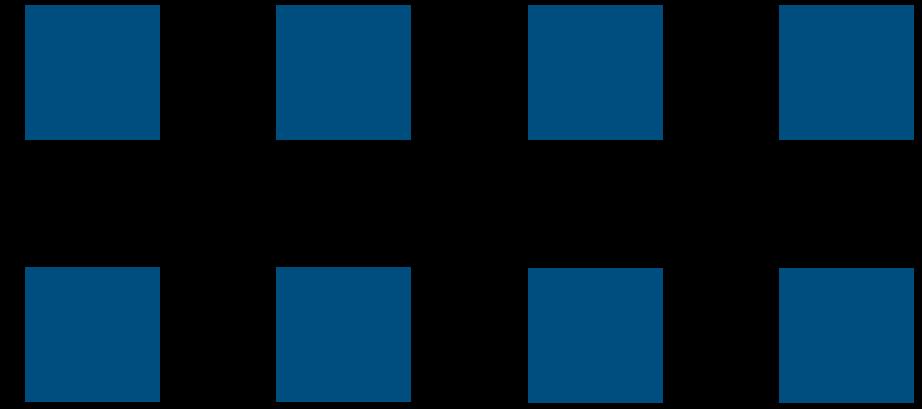
Wrinkle: building this graph would require RAM + time



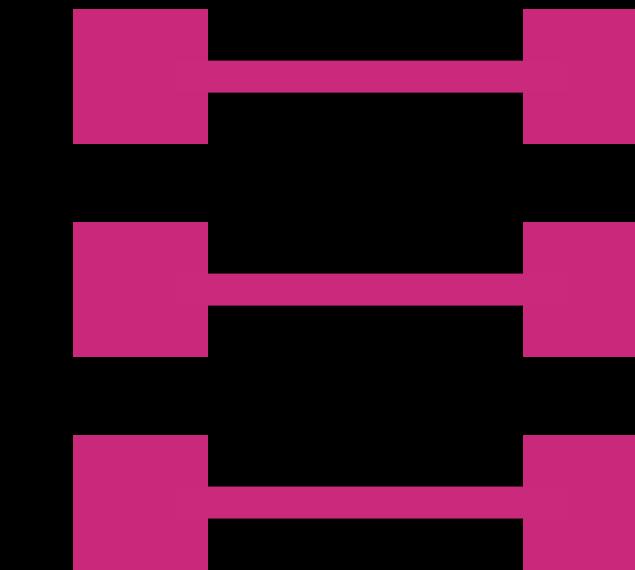
SplitMesher: approximates Matching without materializing meshing graph

SplitMesher: approximates Matching without materializing meshing graph

Set of partially full pages



Pairs of meshable pages



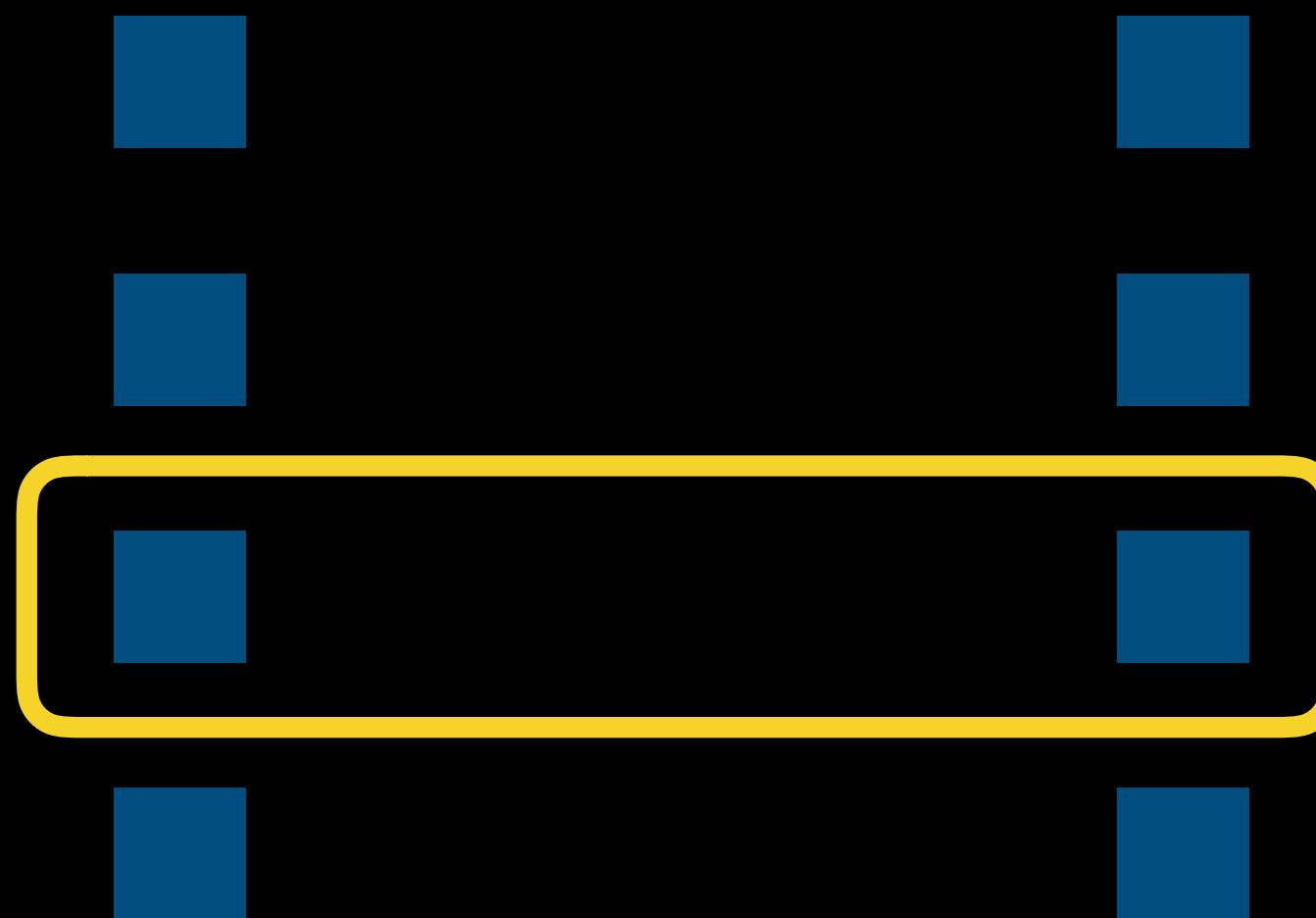
Iterate, comparing $a[i]$ to $b[i]$



Iterate, comparing $a[i]$ to $b[i]$



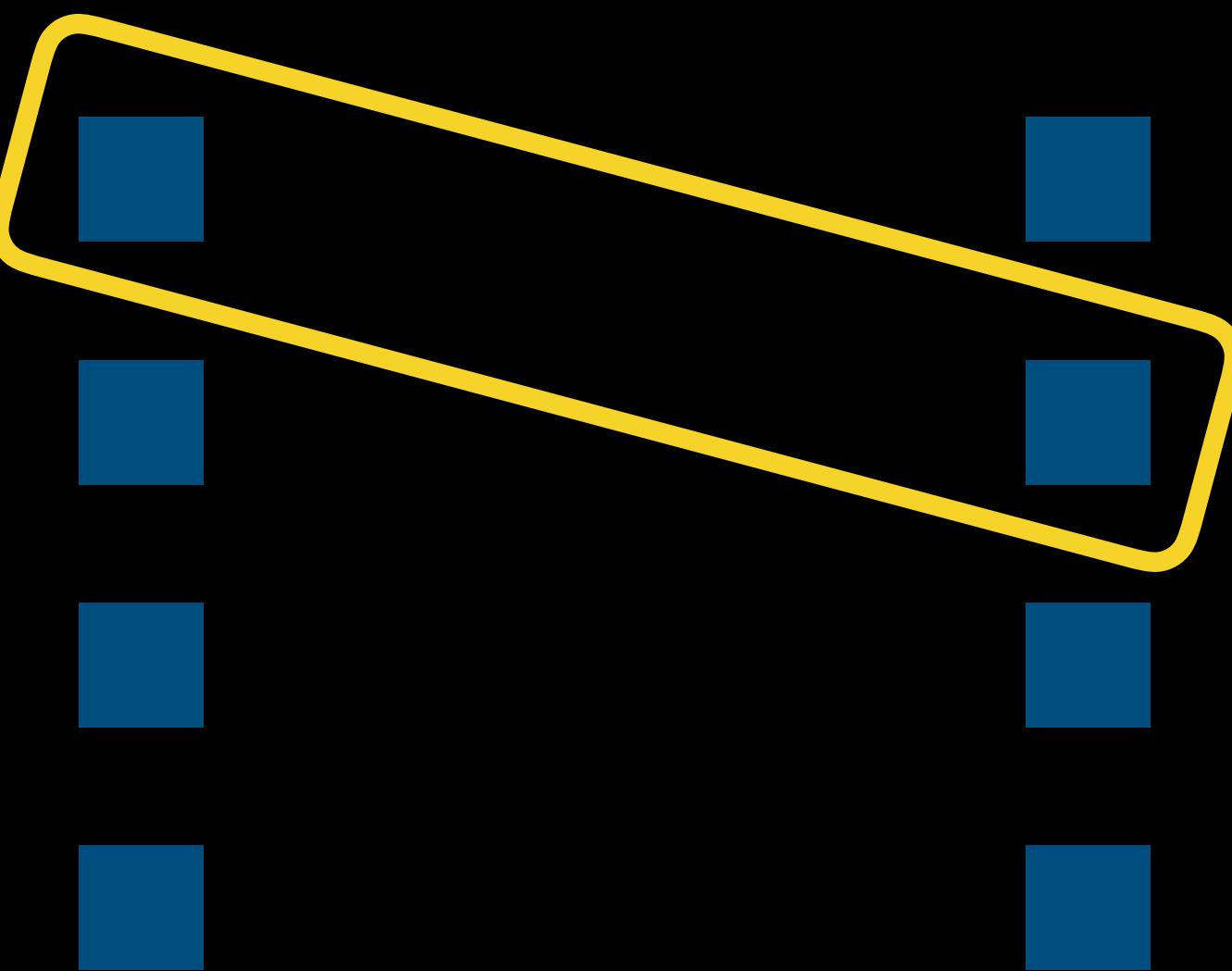
Iterate, comparing $a[i]$ to $b[i]$



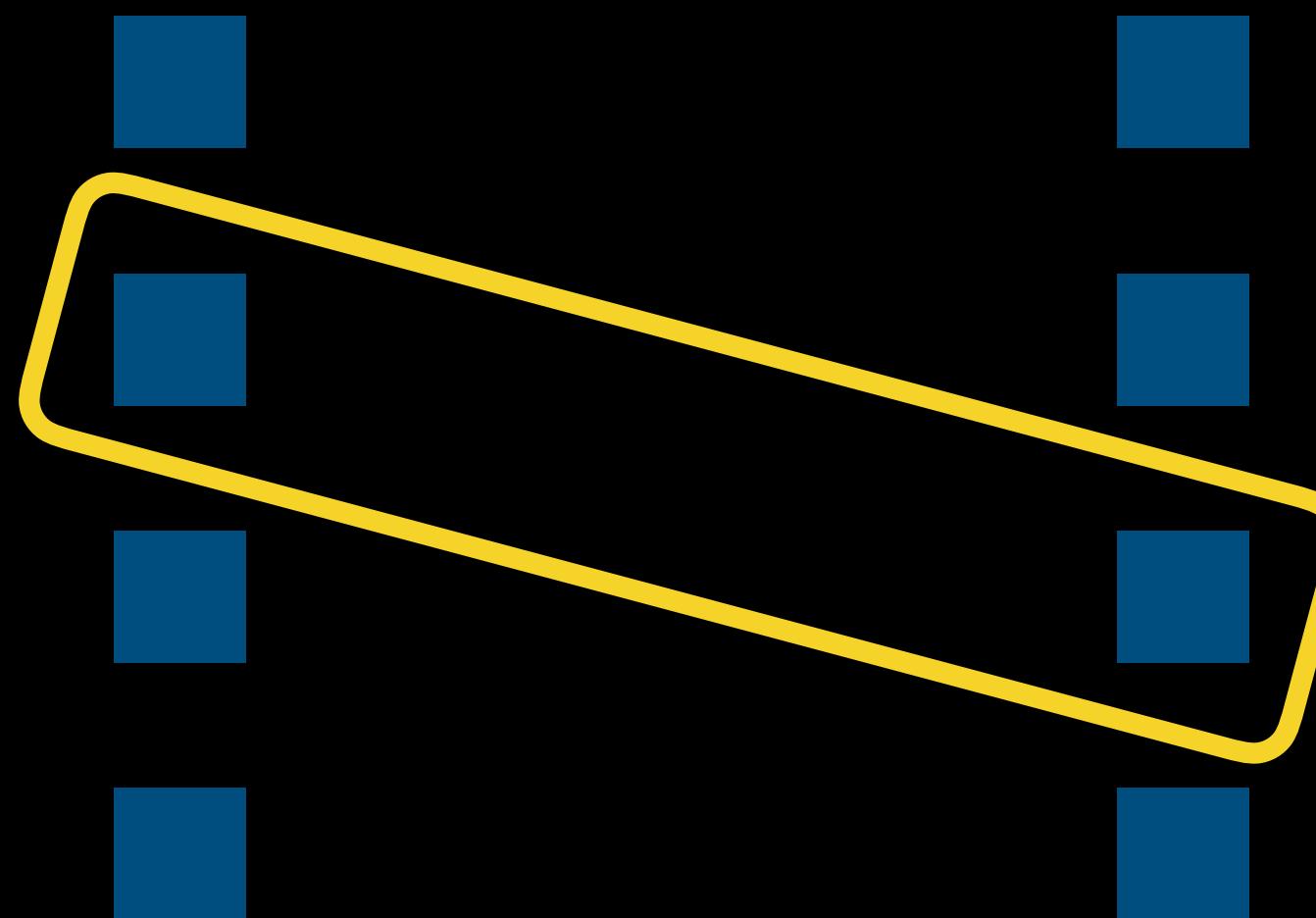
Iterate, comparing $a[i]$ to $b[i]$

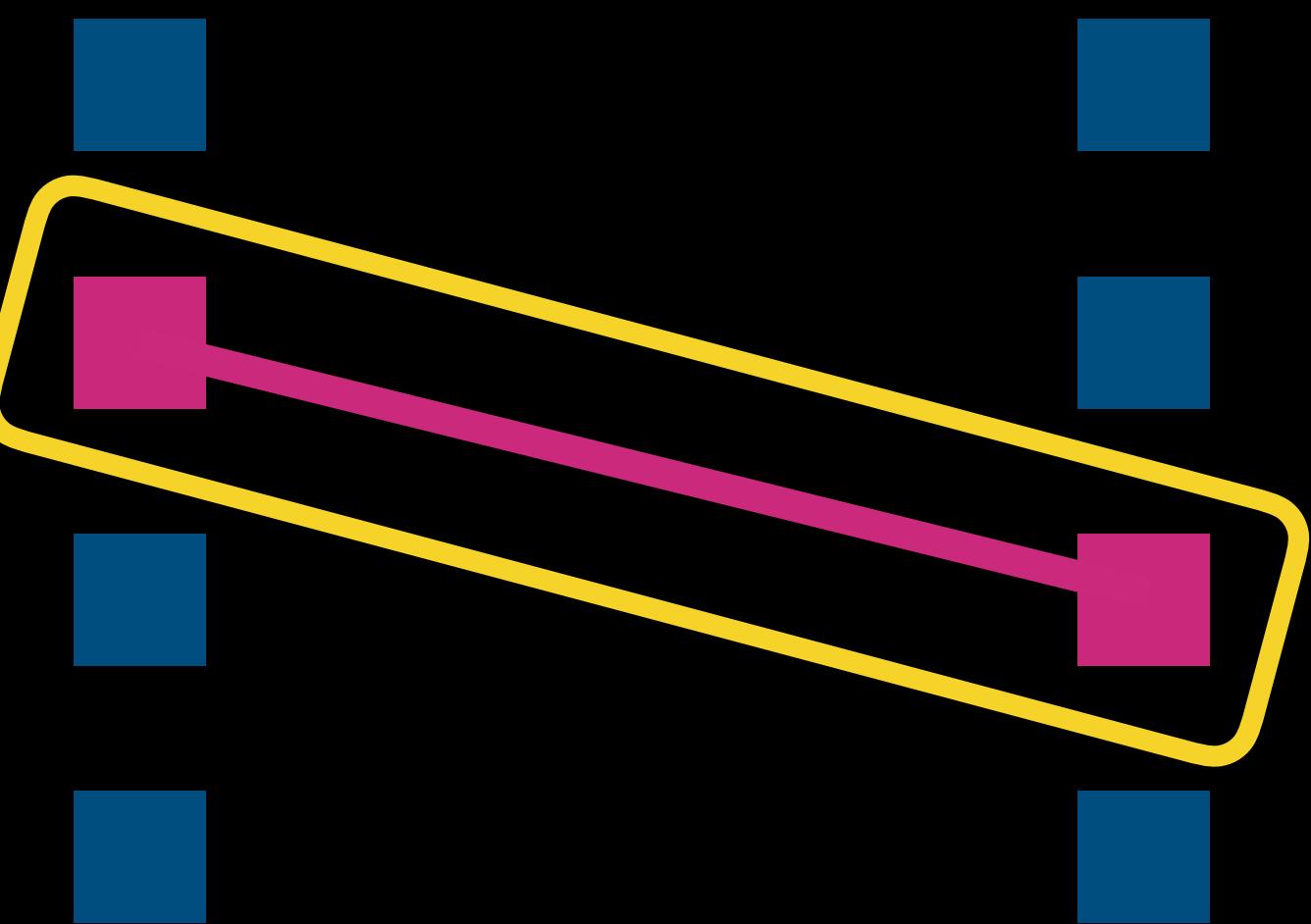


loop, comparing $a[i]$ to $b[(i+1)\%len]$

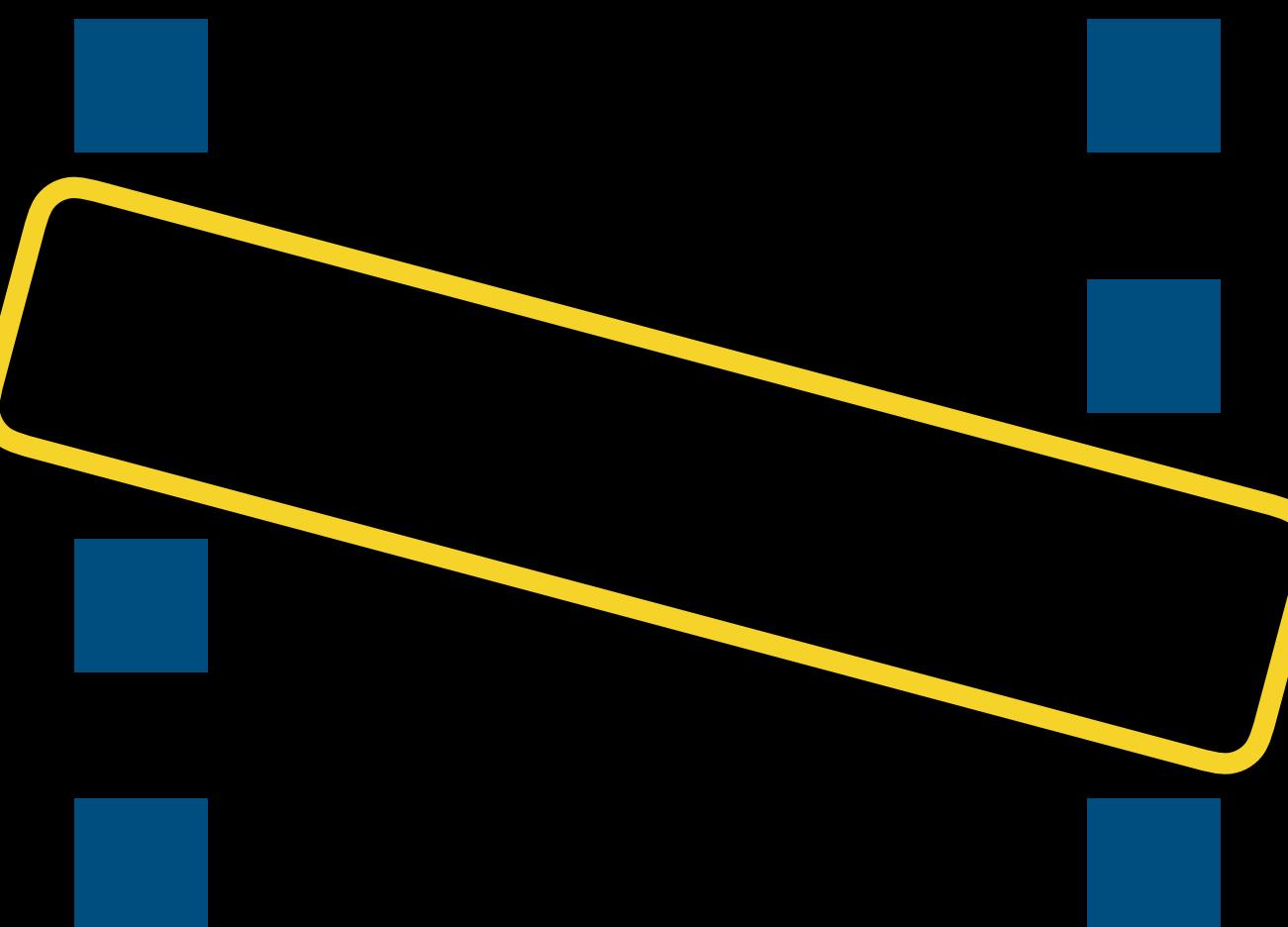


loop **fixed** number of times,
comparing $a[i]$ to $b[(i+1)\%len]$

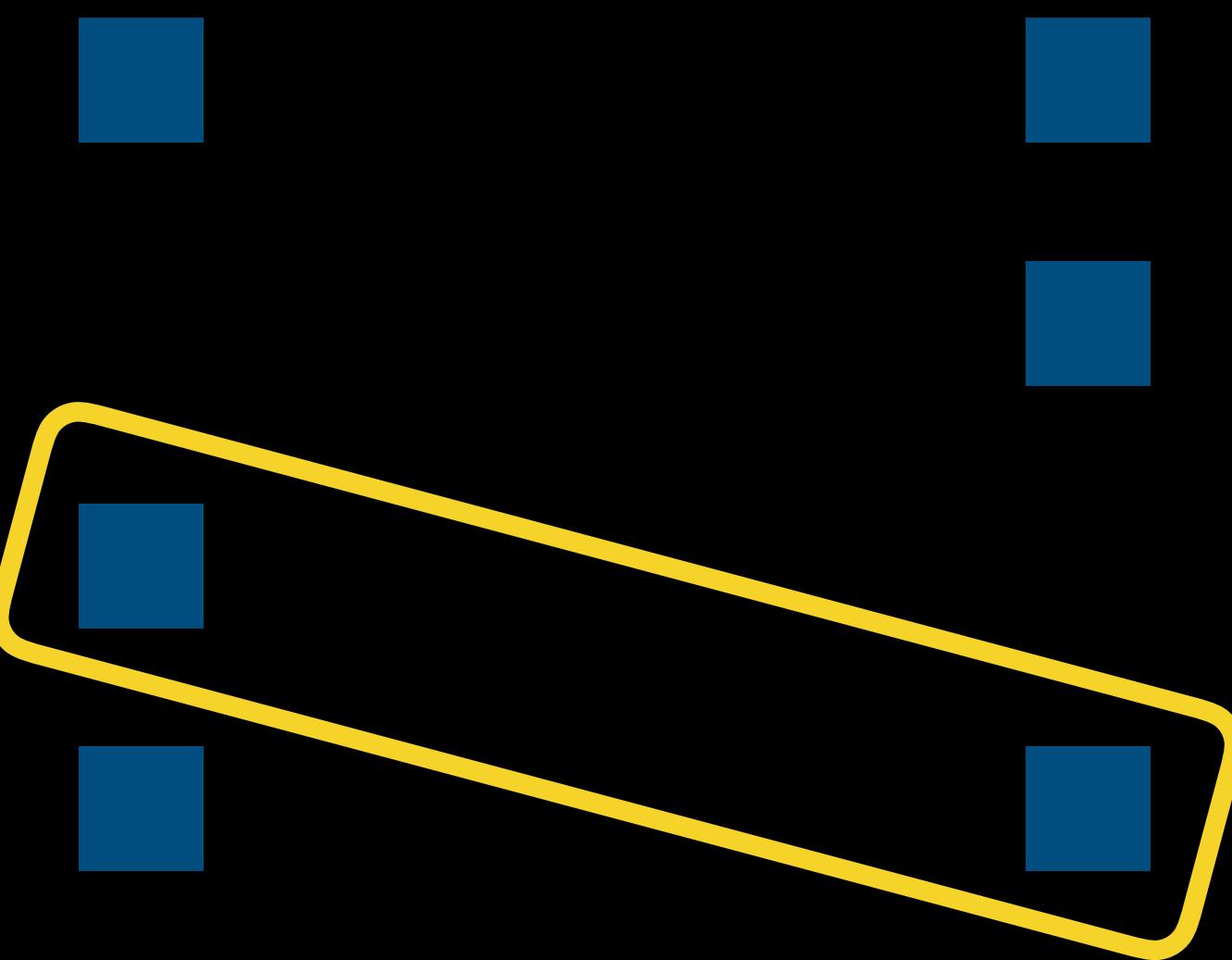




Remove found match



Continue



SplitMesher: approximates Matching without materializing meshing graph

SplitMesher: approximates Matching
without materializing meshing graph

$O(n/q)$ time

(q is the global probability of spans meshing)

SplitMesher: approximates Matching
without materializing meshing graph

$O(n/q)$ time

(q is the global probability of spans meshing)

$1/2^*$ approximation w.h.p.

MESH

*implementation
details*

MESH

*implementation
details*

built using Heap Layers

github.com/emeryberger/Heap-Layers

MESH

*implementation
details*

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(used to build Hoard, DieHard, DieHarder...)

MESH

*implementation
details*

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github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

fast framework for building custom allocators!

MESH

*implementation
details*

built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

< 5K SLOC, C++

MESH

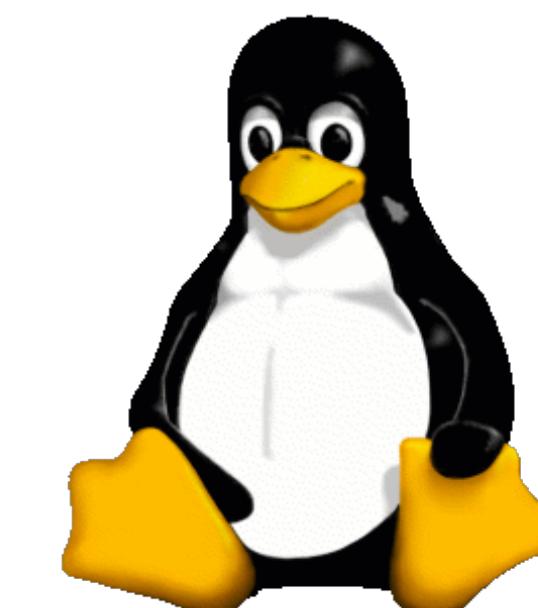
*implementation
details*

built using Heap Layers

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< 5K SLOC, C++



MESH

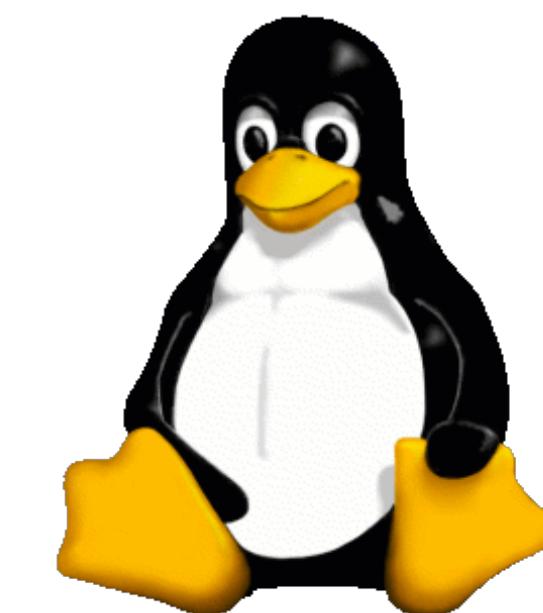
*implementation
details*

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MESH

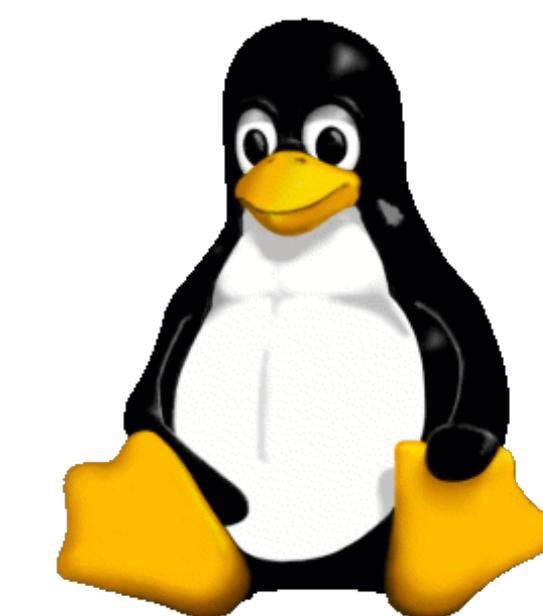
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details*

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MESH

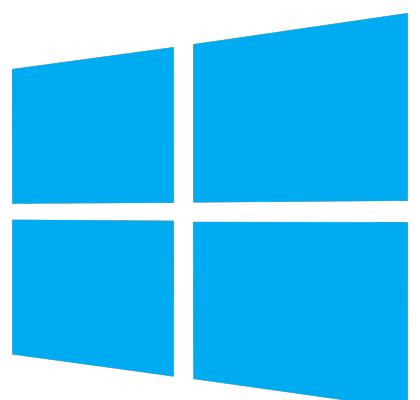
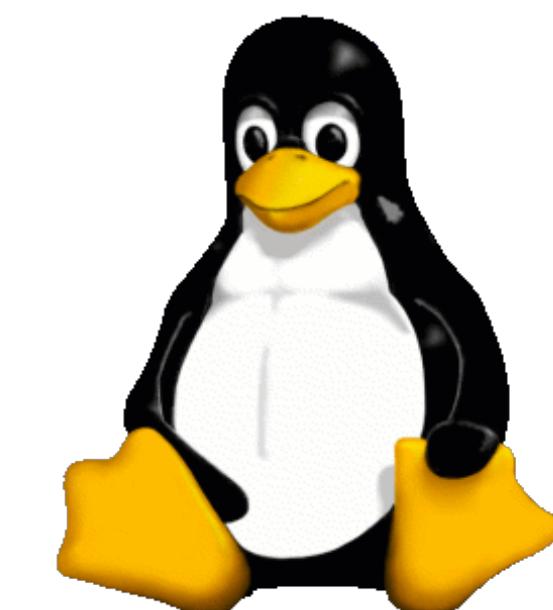
*implementation
details*

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(used to build Hoard, DieHard, DieHarder...)

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MESH

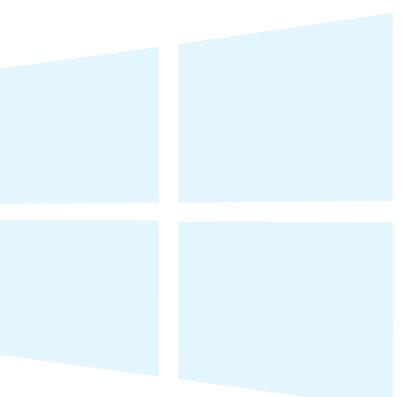
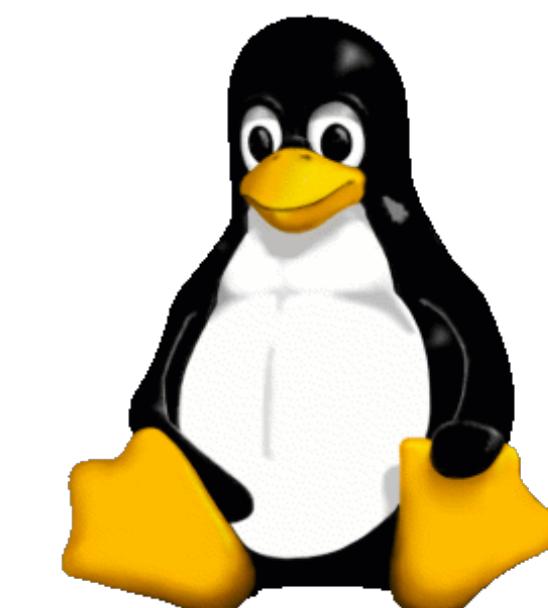
*implementation
details*

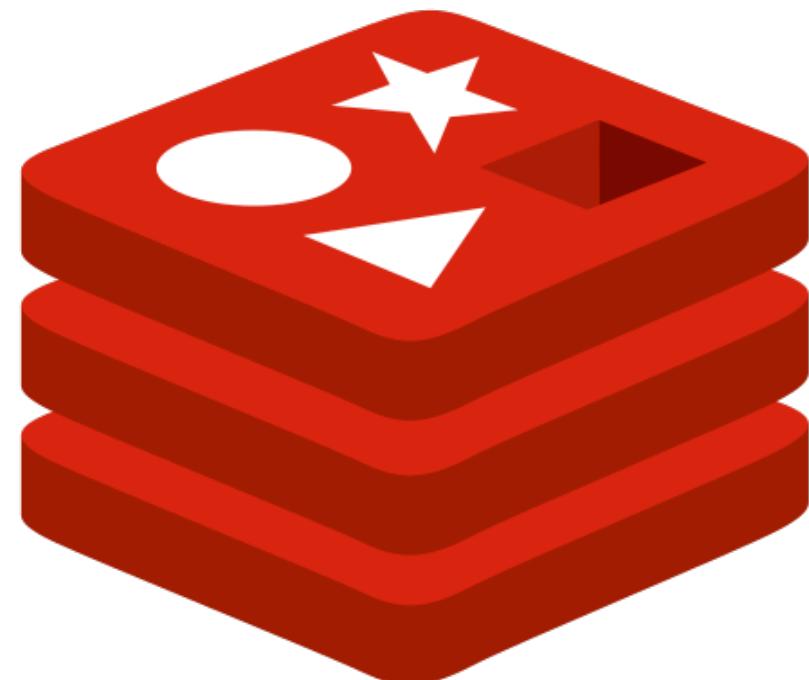
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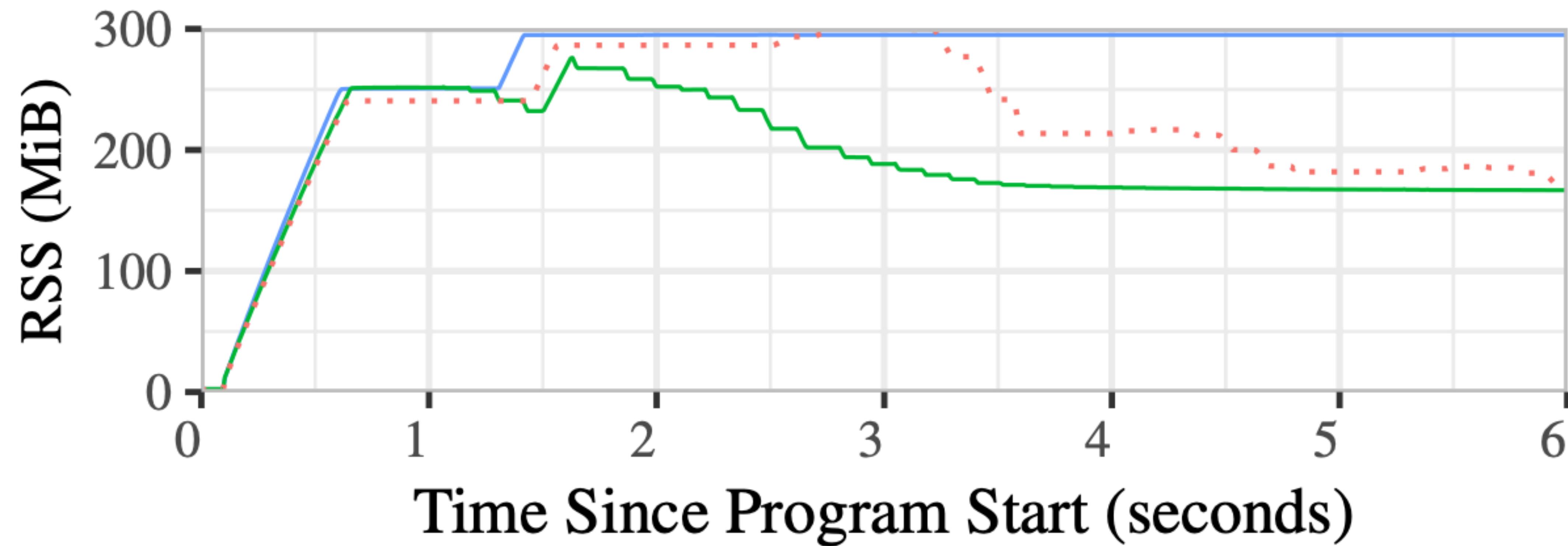
< 5K SLOC, C++

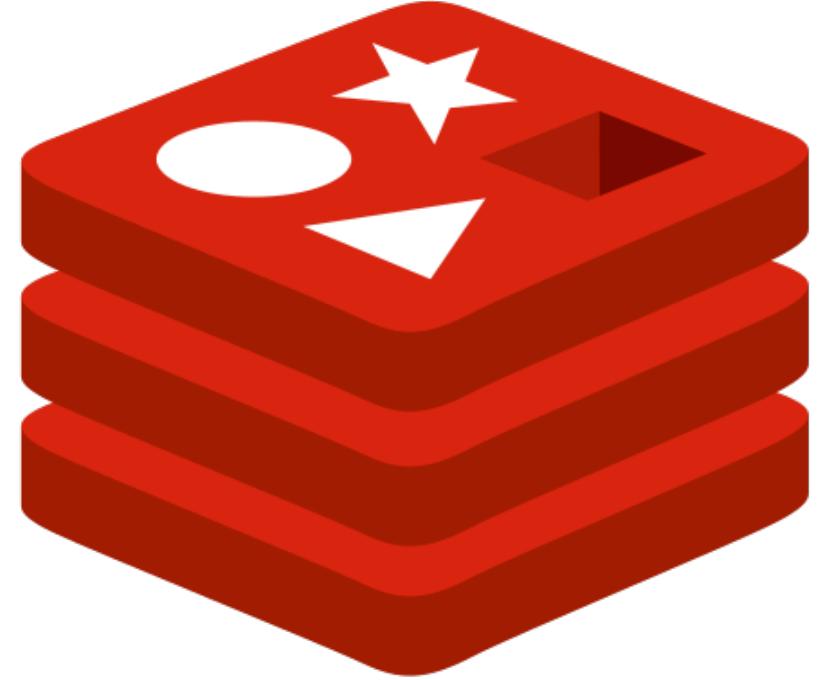




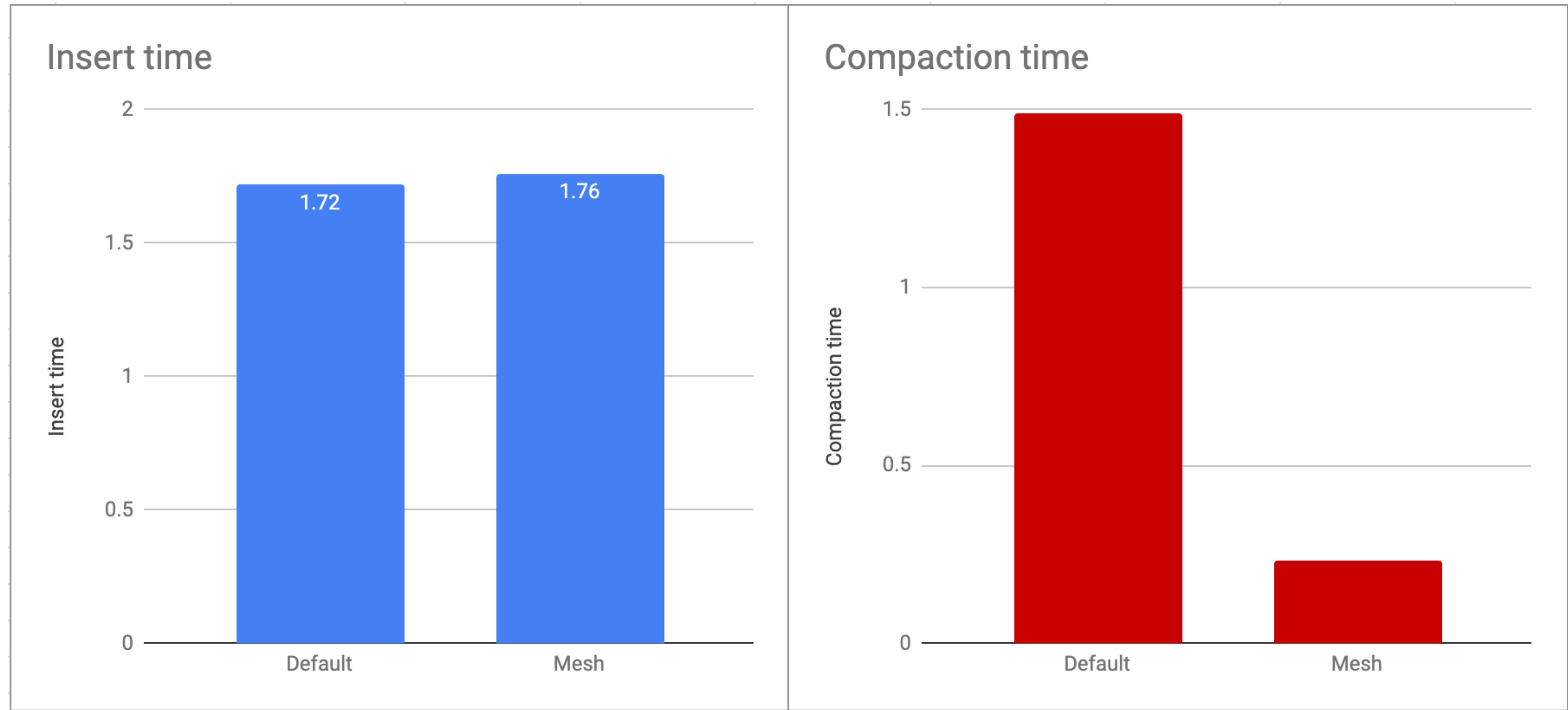
redis

- jemalloc + activatedefrag
- Mesh
- no compaction



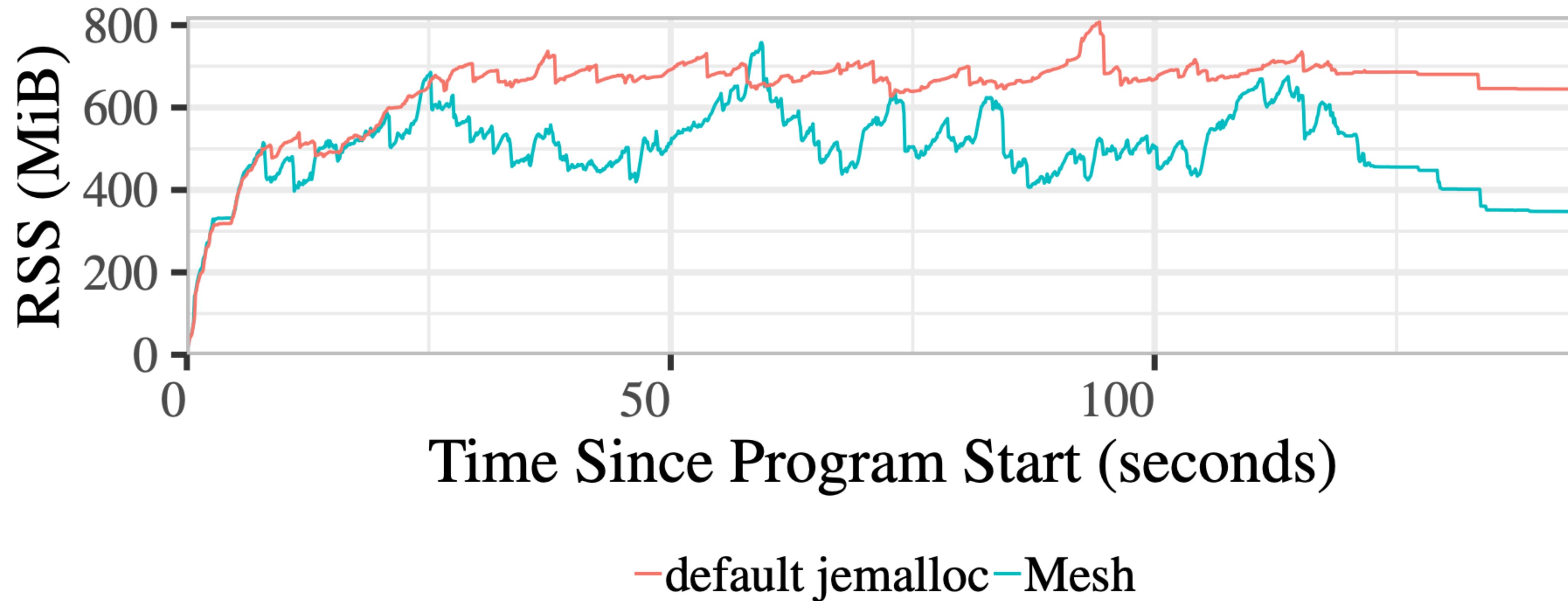


redis + MESH

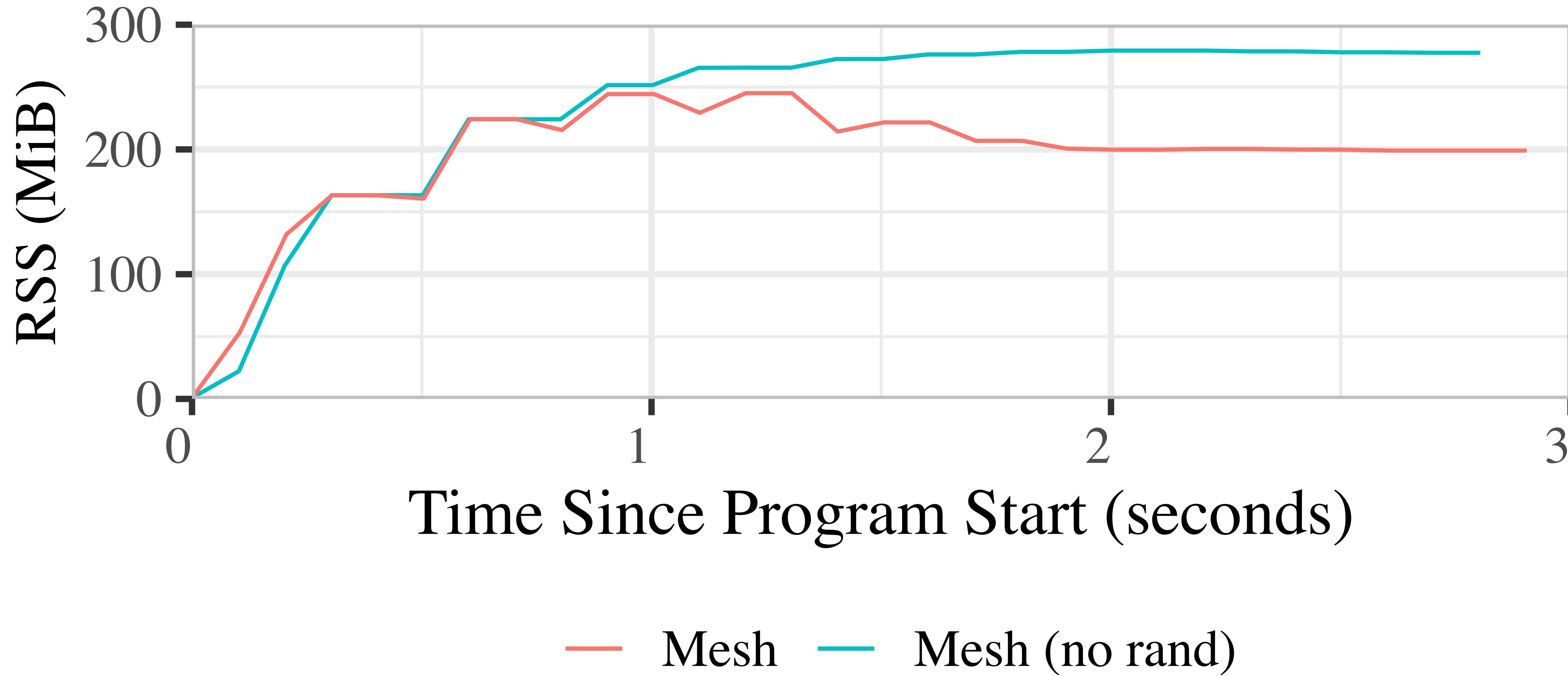


17% heap size reduction

< 1% performance overhead

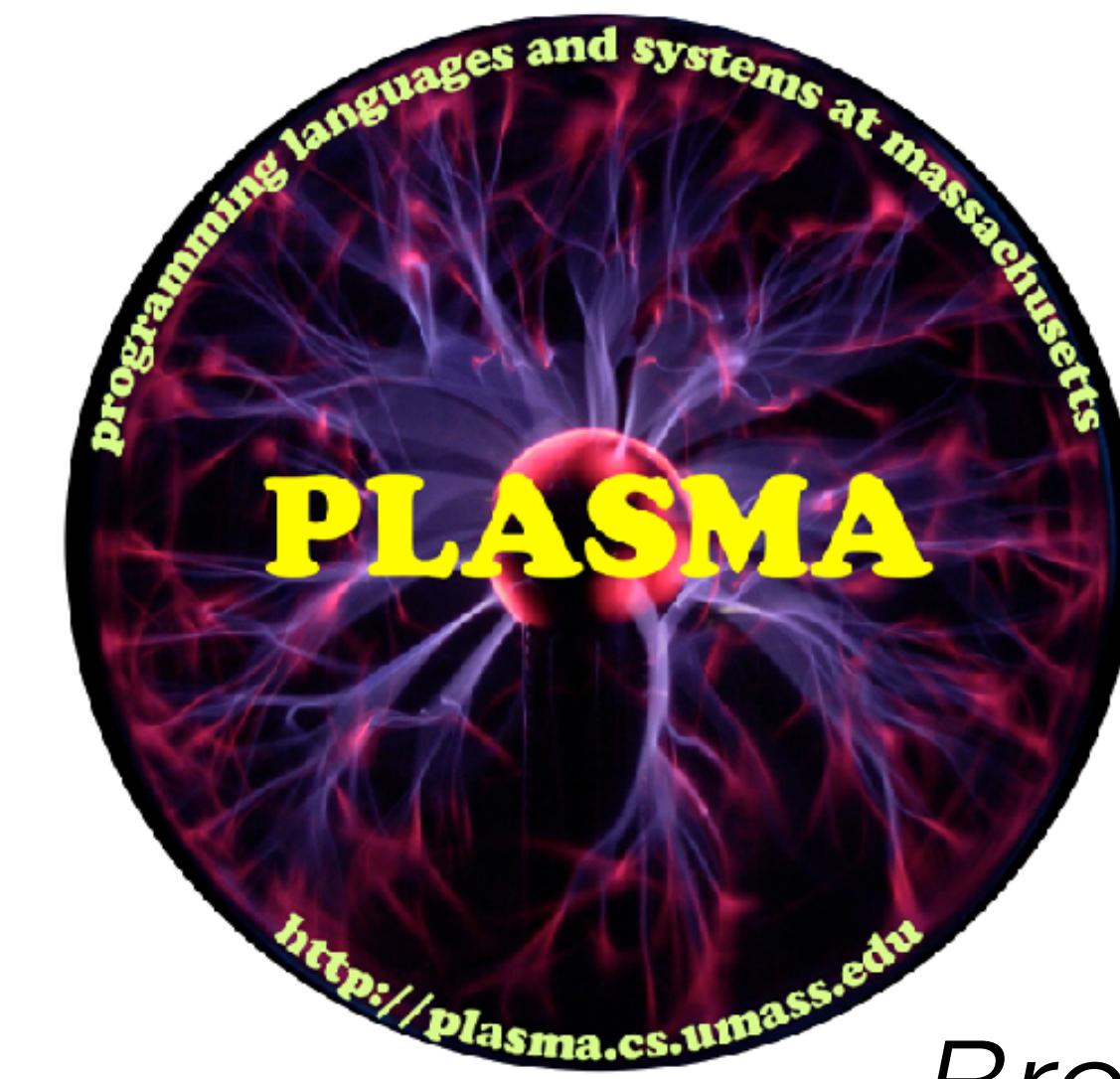
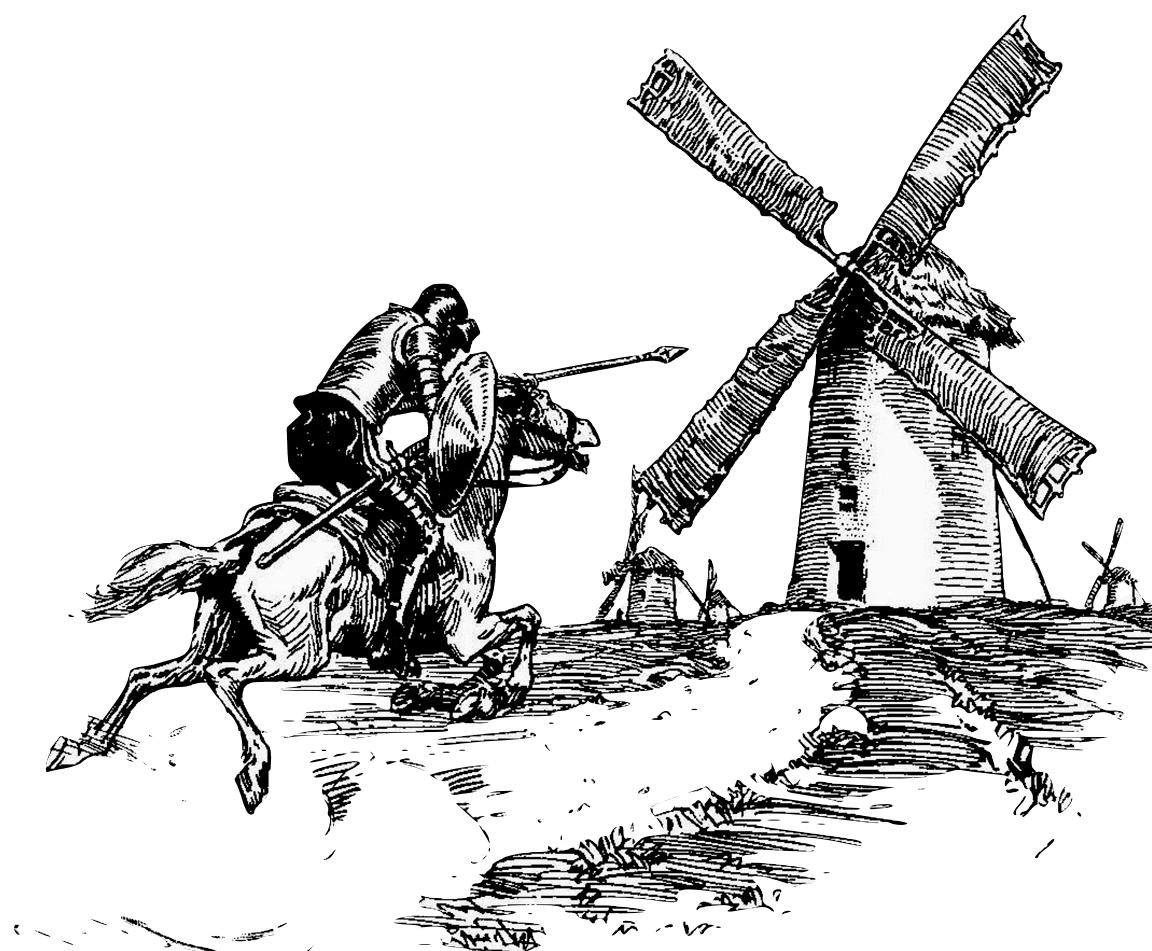


Ruby Compaction for Free



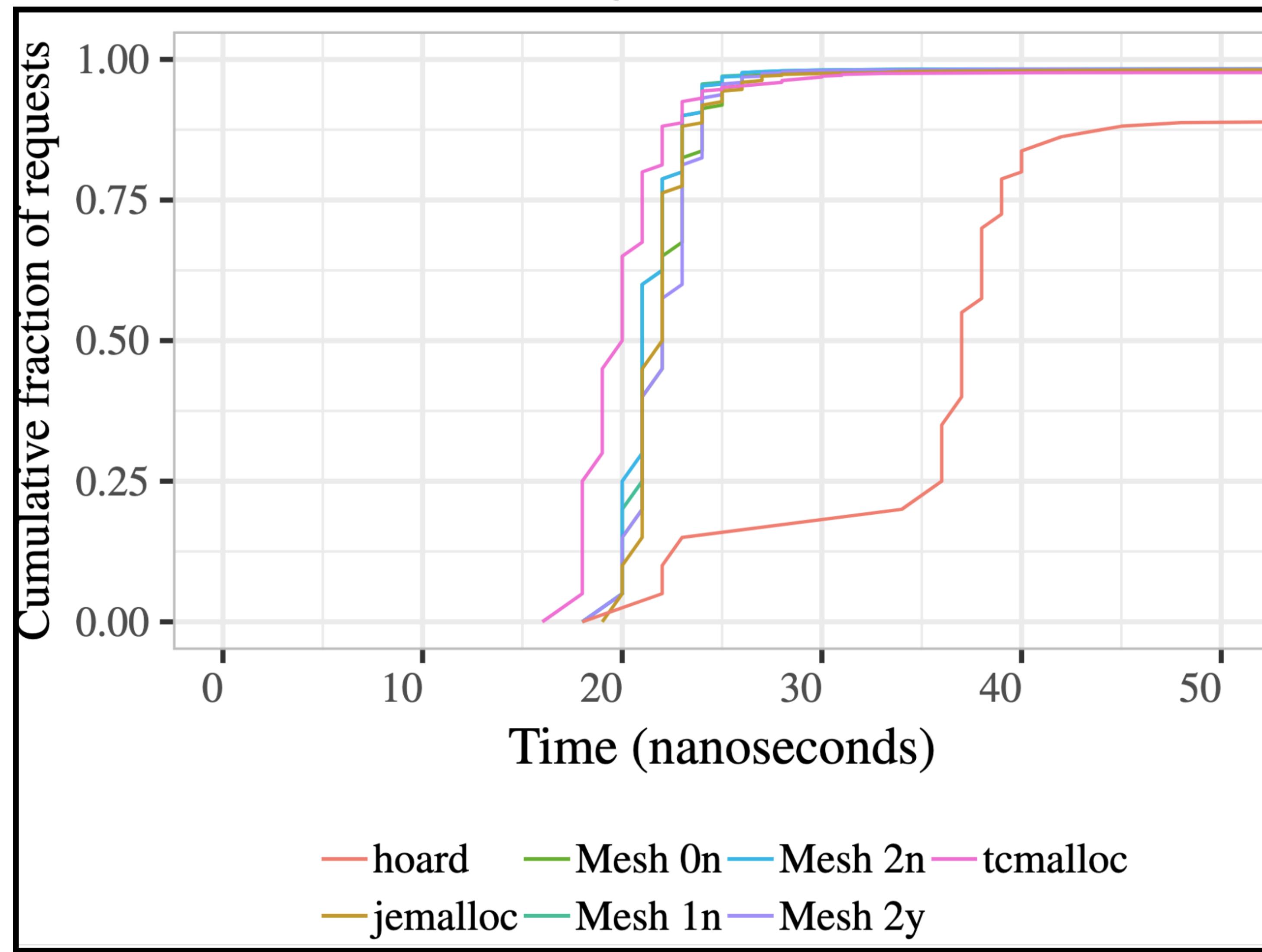
<http://LIBMESH.org>

¡Compactación sin Relocación!
(compaction without relocation)



Browsix, Coz, DieHard, Hoard...

Latency: malloc



Latency: free

