

This module provides an implementation of the constructive algorithm for x in the Chinese remainder theorem.

```
module CRT ( crt ) where
```

Helpers

The extended Euclidean algorithm finds the coefficients of Bézout's identity. For two inputs a and b it finds x and y s.t.

$$ax + by = \gcd(a, b) \tag{1}$$

We can implement this as follows:

```
extendedEu :: Integer -> Integer -> (Integer, Integer)
extendedEu _ 0 = (1, 0)
extendedEu a b =
  let
    (q, r) = quotRem a b
    (s, t) = extendedEu b r
  in
    (t, s - q * t)
```

As such, we then trivially have:

```
gcd :: Integer -> Integer -> Integer
gcd a b = a*x + b*y
  where (x, y) = extendedEu a b
```

CRT

Description

Chinese remainder theorem tells us that for a list of integers n_1, \dots, n_k and a second list a_1, \dots, a_k , there exists soem x that solves the following:

$$\begin{aligned} x &= a_1 \pmod{n_1} \\ x &= a_2 \pmod{n_2} \\ &\dots \\ x &= a_k \pmod{n_k} \end{aligned}$$

And that moreover, all solutions x are congruent modulo $N = \prod_i n_i$.

Algorithm

We can generate that x using the following algorithm. First, calculate N . Then, for each i we can find r_i and s_i s.t.

$$r_i n_i + \frac{s_i N}{n_i} = 1 \quad (2)$$

we do this by using the extended Euclidean algorithm on n_i and N/n_i . In which case s_i is the second element of the output of `extendedEu`.

We then have a solution x given by

$$x = \sum_{i=1}^k \frac{a_i s_i N}{n_i} \quad (3)$$

And we can find the minimal solution by taking mod N .

Implementation

We can implement that as follows.

First, a helper to calculate each ith element of the above sum.

```
getElement :: Integer -> (Integer, Integer) -> Integer
getElement bigN (a, n) = (a*s*bigN) `div` n
  where s = snd $ extendedEu n (bigN `div` n)
```

Then the real thing:

```
crt :: [(Integer, Integer)] -> Integer
crt cs = sum (map (getElement n) cs) `mod` n
  where n = product $ map snd cs
```