



Exam 2 Review

Deterministic Systems Models / Systems Optimization

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ISE 4623/5023: Deterministic Systems Models / Systems Optimization

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Topics to be covered in Exam 2

- Module 3. Sensitivity analysis and duality
- Module 4. LP models for common structures
- Everything from previous modules (particularly in topics where the students had the lowest performance from previous exams/assignments/quizzes)



Module 3. Sensitivity analysis and duality

Including (but not limited to):

- Constructing dual problems and understand their connection with the primal
 - Meaning of dual variables and their application
 - Weak and strong duality theorems
- Sensitivity and post-optimal analysis
 - Add or remove activities/variables
 - Add or remove constraints
 - Change parameters (costs, resources, etc.)

Module 4. LP models for common structures

Including (but not limited to):

- Understanding common problems and their LP models/structures
 - Data Envelopment Analysis (DEA)
 - DMUs, inputs, outputs
 - Efficiency, reference set, etc.
 - Generalized network models
 - Transportation problem
 - Assignment problem
 - Minimum cost flow problem
 - » Connection with transportation problem, assignment problem, etc.
 - » Shortest path problem
 - » Max flow problem
- Knowing how to adapt/modify/extend common structures to fit particular contexts
 - Handle unbalanced networks (total demand is different from total supply capacity)
 - Handle possible **unmet demand** or excess production
 - Adding capacities to nodes (e.g., maximum flow that can go through some nodes)
 - Adding/adapting flow constraints for different types of capacities (number of units, weight, size, etc.)

Data Envelopment Analysis (DEA) Efficiency Model

- Sets:

N = set of DMUs

I = set of *inputs*

O = set of *outputs*

- Parameters:

y_{oj} = amount of output $o \in O$ produced by DMU $j \in N$

x_{ij} = amount of input $i \in I$ produced by DMU $j \in N$

p = Individual DMU under study

- Decision variables:

λ_j = fraction of the j th DMU ($j \in N$) used to achieve efficiency (i.e., get to the efficiency frontier)

ϕ_p = Proportional increase (growth rate) of DMU under study (p)

- Objective function:

$$\max \phi_p$$

- Constraints:

$$\sum_{j \in N} y_{oj} \lambda_j \geq y_{op} \phi_p, \forall o \in O$$

Sum of output * weights must be more than the person of comparison

$$\sum_{j \in N} x_{ij} \lambda_j \geq x_{io}, \forall i \in I$$

Sum of inputs * weights for others in sample must be less than the person of comparison?
Check notes

$$\lambda_j \geq 0, \forall j \in N$$

$$\phi_p \text{ unrestricted}$$

Transportation Problem

- Sets:

S = set of *sources* $\{1, 2, \dots, m\}$

D = set of *destinations* $\{1, 2, \dots, n\}$

- Parameters:

a_i = supply of source node $i \in S$

b_j = demand of destination node $j \in D$

c_{ij} = unit cost of flow through arc (i, j) ,
where $i \in S$ and $j \in D$

- Decision variables:

x_{ij} = flow through arc (i, j) , where $i \in S$
and $j \in D$

- Objective function:

$$\min z = \sum_{i \in S} \sum_{j \in D} c_{ij} x_{ij}$$

Flows meet the
supply or
demand

Change to \leq if
unmet demand

- Constraints:

$$\sum_{i \in S} x_{ij} = b_j, \forall j \in D$$

flow equals demand

$$\sum_{j \in D} x_{ij} = a_i, \forall i \in S$$

flow equals supply

$$x_{ij} \geq 0, \forall i \in S, \forall j \in D$$

Assignment Problem

- Sets:

W : Set of workers $\{1, 2, 3, \dots, n\}$

J : Set of jobs $\{1, 2, 3, \dots, n\}$

- Parameters:

c_{ij} : cost of assigning worker $i \in W$ to job $j \in J$

- Variables:

x_{ij} : binary variable that is 1 if worker $i \in W$ is assigned to job $j \in J$, and is 0 otherwise.

- Objective function:

$$\min z = \sum_{i \in W} \sum_{j \in J} c_{ij} x_{ij}$$

- Constraints:

$$\sum_{i \in W} x_{ij} = 1 \quad , \forall j \in J$$

$$\sum_{j \in J} x_{ij} = 1 \quad , \forall i \in W$$

$$x_{ij} \in \{0, 1\} \quad , \forall i \in W, \forall j \in J$$

Each worker must be assigned a job

Minimum Cost Flow Problem (MCFP)

- Sets:

N : Set of nodes $\{1, 2, 3, \dots, n\}$

A : Set of arcs

- Parameters:

c_{ij} : unit cost of sending a commodity through arc $(i, j) \in A$

b_i : demand/supply of commodities in node $i \in N$

u_{ij} : maximum flow through arc $(i, j) \in A$

l_{ij} : minimum flow through arc $(i, j) \in A$

- Variables:

x_{ij} : flow through arc $(i, j) \in A$

- Objective function:

$$\min z = \sum_{(i,j) \in A} c_{ij} x_{ij}$$

- Constraints:

$$\sum_{j:(i,j) \in A} x_{ij} - \sum_{j:(j,i) \in A} x_{ji} = b_i, \forall i \in N$$

$$x_{ij} \leq u_{ij}$$

$$x_{ij} \geq l_{ij}$$

$, \forall (i, j) \in A$ Within upper bound

$, \forall (i, j) \in A$ Within Lower bound

Flows meet the supply or demand

Change to \leq if unmet demand

Multicommodity Minimum Cost Flow Problem (MMCFP)

- Sets:

N : Set of nodes $\{1, 2, 3, \dots, n\}$

A : Set of arcs

K : Set of commodities

- Parameters:

c_{ijk} : unit cost of sending a commodity $k \in K$ through arc $(i, j) \in A$

b_{ik} : demand/supply of commodity $k \in K$ in node $i \in N$

u_{ij} : maximum total flow through arc $(i, j) \in A$

l_{ij} : minimum total flow through arc $(i, j) \in A$

- Variables:

x_{ijl} : flow of commodity $l \in L$ through arc $(i, j) \in A$

- Objective function:

$$\min z = \sum_{k \in K} \sum_{(i, j) \in A} c_{ijk} x_{ijk}$$

- Constraints:

$$\sum_{j: (i, j) \in A} x_{ijk} - \sum_{j: (j, i) \in A} x_{jik} = b_{ik} \quad , \forall i \in N, \forall k \in K$$

Flows meet the supply / Demand (change \leq for unmet demand)

$$\sum_{k \in K} x_{ijk} \leq u_{ij} \quad , \forall (i, j) \in A \quad \text{Within upper bound}$$

$$\sum_{k \in K} x_{ijk} \geq l_{ij} \quad , \forall (i, j) \in A \quad \text{Within Lower bound}$$



THANK YOU

QUESTIONS?

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