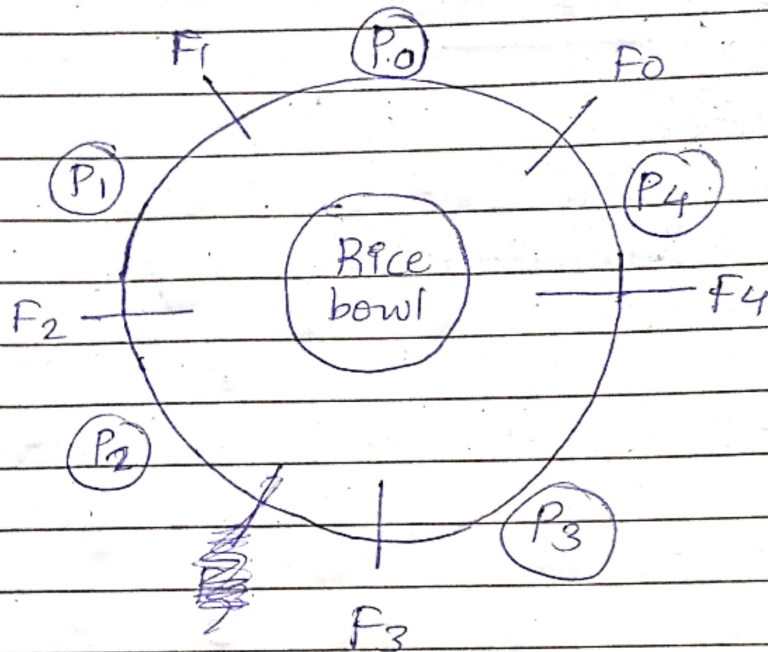


Dining Philosophers problem

Dining table \rightarrow 5 philosophers
5 forks



void Philosopher(void) {
while (true) {

Thinking();

take_fork(i); \leftarrow left fork

take_fork((i+1)%N); \leftarrow right fork

EAT();

put_fork(i);

put_fork((i+1)%N);

}

}

Philosopher

\rightarrow Think

\rightarrow Eat

Case 1: P_0 comes

$i = 0$

left-fork = 0

Right fork = $(0+1) \% N = 1 (f_1)$

Eating

f_0

f_1

end

Case 2: P_1 comes

$i = 1$

left fork = 1 (f_1)

right fork = 2 (f_2)

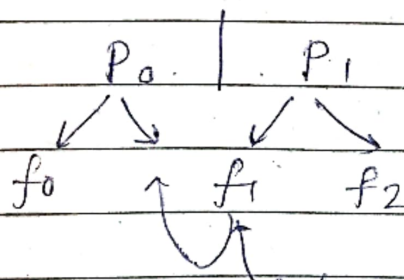
Eating

f_1

f_2

Case 3: Case 2:

P_0 picks f_0 & P_1 picks f_1



When P_1 completes eating then only P_0 would get the fork

problem of race condⁿ occurs

$S[i] \rightarrow$ use array of semaphores

	S_0	S_1	S_2	S_3	S_4	
Initially:	1	1	1	1	1	[Initialize every semaphore with 1 as when initialized with 0 it waits & gets blocked]
	1	1				
	0	0				

Code:

```

void philosopher (void) {
    while (true) {
        Thinking();
        Entry { wait(fork( $S[i]$ ))
              wait(fork( $S[(i+1) \bmod n]$ ))

        EAT();
        Signal(Put-fork( $i$ ));
        Signal(Put-fork( $(i+1) \% N$ ))
    }
}

```

?

$P_0 \rightarrow S_0 \quad S_1$

$P_1 \rightarrow S_1 \quad S_2$

$P_2 \rightarrow S_2 \quad S_3$

$P_3 \rightarrow S_3 \quad S_4$

$P_4 \rightarrow S_4 \quad S_0$

\rightarrow Case 1:

When P_0 comes S_0 & S_1 ($1 \rightarrow 0$) &
while going out ($0 \rightarrow 1$)

In DB P_0 & P_2 can come. If ϕ is a special case of mutual exclusion, as P_0 & P_2 are independent of each other.

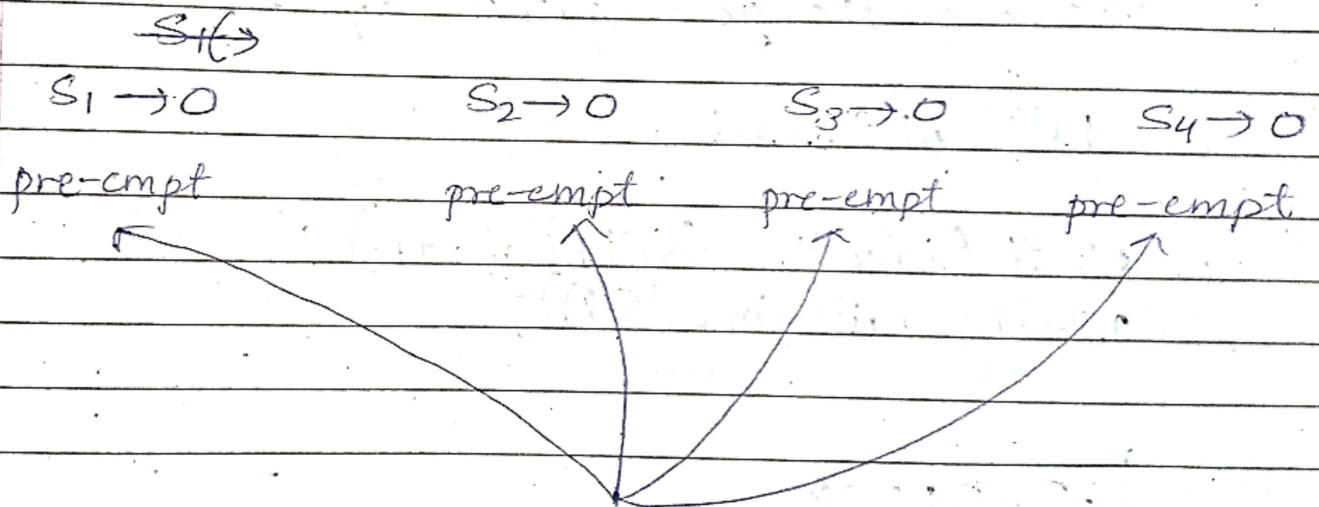
Case 3:

S_0	S_1	S_2	S_3	S_4
↓	↓	↓	↓	↓
1	1	1	1	1

P_0 comes first (~~$S_0 \rightarrow 1$~~)
 $S_0 (1 \rightarrow 0)$

& gets pre-empted

P_1 comes P_2 comes P_3 comes P_4 comes



Blocked for right fork as the right fork is acquired by next philosopher

All philosophers take the left hand side fork but gets blocked for the right one. This situation is called deadlock (All processes get blocked)

How to remove deadlock? On changing sequence of one process.

P₄ S₀ S₄

S ₀	S ₁	S ₂	S ₃	S ₄
↓	↓	↓	↓	↓
1	1	1	1	1
0	0	0	0	

blocked P₄ S₀ S₄

We can change sequence of any philosopher

Nth philosopher

wait (take fork (S(i+1) mod N))
wait (take fork (S_i));