



# Tecnológico de Monterrey

**Choo Choo Train**

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March 6, 2017

# Purpose of Project

The purpose of this project is to help children get an insight into the world of programming from a small age. Children are usually uninterested to learn how to program because they are not sure of what it actually is or what it's useful for. They think that because they don't know about computers, they will be bad at programming. The fact that programs look like huge files of unreadable text, sometimes with complicated formulas, doesn't help either. If children can relate programming to something that is familiar to them, then we might be able to help them understand what programming really is about, and possibly encourage them to pursue a Computer Science Degree in the future.

Children will use our visual language to "program" blocks, using a visual representation tool of code. Then they will be able to watch their code being executed, represented as a train going through its tracks. By mapping the elements of a language (such as loops, statements, etc.) to a specific type of train track section or element, we can help kids understand the logic and data flow of a program. This will introduce them to the world of programming while still maintaining an attractive and fun appearance that won't make them think that what they're about to do is rocket science (yet).

# Main Objective and Area

The main objective of our language is to be understandable, visual, and easy to use so that children will be motivated to learn how to program. The main characteristic is that both the input and the output are visual, in a sense that the output encompasses the step by step execution of the program, as well as any output generated. With this visual representation, complicated programming terminology and structures are abstracted and represented with objects familiar to children, in this case a train track environment. Its objective is to be an educational, fun, and pleasing programming language for children.

# Language Requirements

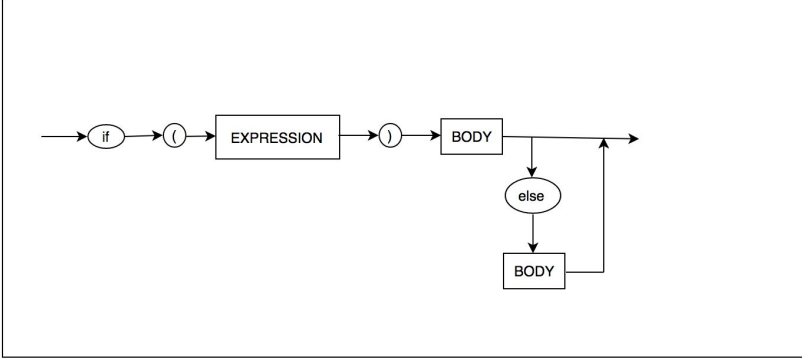
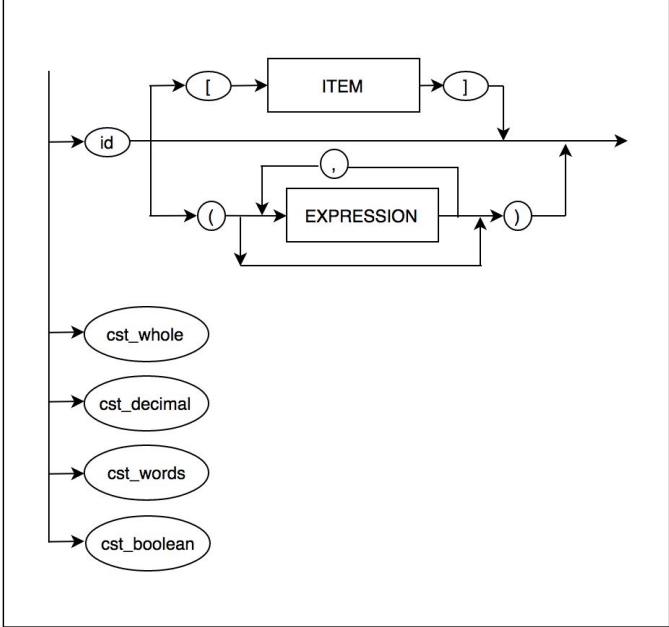
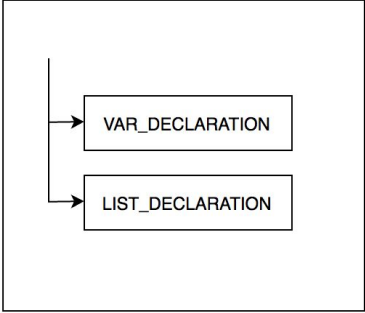
## 3.1 Tokens

Token Name	Regular Expression
block	"block"
starting	"starting"
receives	"receives"
block_returns	"returns"
call	"call"
variable	"variable"
list	"list"
of	"of"
type	"type"
return_statement	"return"
do	"do"
until	"until"
if	"if"
else	"else"
colon	"."
semicolon	"."
comma	"."
input	"input"
print	"print"
curlybraces_open	"{"
curlybraces_close	"}"
parenthesis_open	"("
parenthesis_close	")"

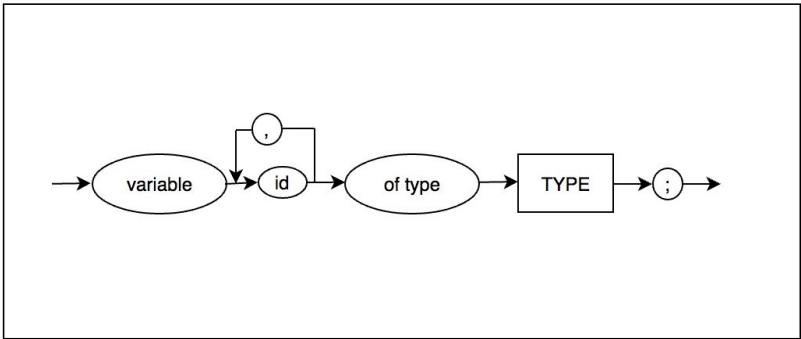
squarebracket_open	"["
squarebracket_close	"]"
op_assign	"="
op_less	"<"
op_less_equal	"<="
op_greater	">"
op_greater_equal	">="
op_equal	"=="
op_not_equal	"!="
op_and	"and"
op_or	"or"
op_negation	"not"
op_addition	"+"
op_subtraction	"_"
op_multiplication	"*"
op_division	"/"
whole	"whole"
decimal	"decimal"
words	"words"
boolean	"boolean"
cst_whole	[0-9]+
cst_decimal	[0-9]+\.[0-9]+([Ee][\+-]?[0-9]+)?
cst_words	\("[^"]"\)
cst_boolean	"true"   "false"
id	[A-Za-z]([A-Za-z0-9])*

## 3.2 Syntax Diagrams

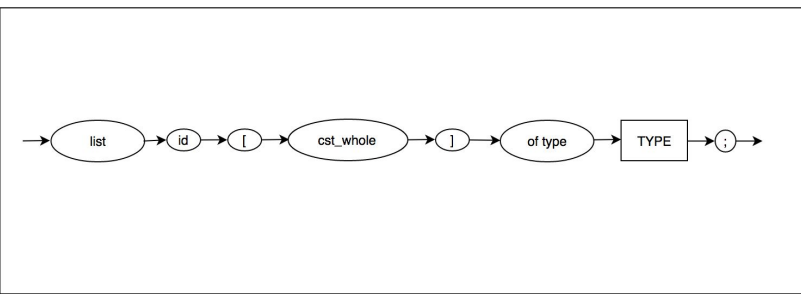
Syntactic Rule	Syntax Diagram
PROGRAM	
BLOCK	
BLOCK_BODY	
TYPE	
BODY	

CONDITION	 <pre>graph LR; Start(( )) --&gt; if((if)); if --&gt; LP("("); LP --&gt; EXPRESSION[EXPRESSION]; EXPRESSION --&gt; RP(")"); RP --&gt; BODY1[BODY]; BODY1 --&gt; Exit1(( )); BODY1 --&gt; else((else)); else --&gt; BODY2[BODY]; BODY2 --&gt; Exit2(( )); Exit1 --&gt; Join(( )); Exit2 --&gt; Join; Join --&gt; End(( ));</pre> <p>The flowchart illustrates the grammar rule for an 'if' statement. It begins with an entry arrow pointing to an oval labeled 'if'. This is followed by an oval containing '(', then a rectangular box labeled 'EXPRESSION', and another oval containing ')'. After this, the flow splits: one path goes to a rectangular box labeled 'BODY', and the other goes to an oval labeled 'else', which then leads to another rectangular box labeled 'BODY'. Both 'BODY' boxes have arrows that merge into a single path, which then leads to the final exit arrow.</p>
CONSTANT	 <pre>graph LR; Entry(( )) --&gt; id((id)); id --&gt; LP("("); LP --&gt; EXPRESSION[EXPRESSION]; EXPRESSION --&gt; RP(")"); RP --&gt; Exit1(( )); id --&gt; L "["; L --&gt; ITEM[ITEM]; ITEM --&gt; R "]" ; R --&gt; Exit2(( )); id --&gt; COMMA(","); COMMA --&gt; EXPRESSION; EXPRESSION --&gt; RP; RP --&gt; Exit3(( )); id --&gt; Exit4(( )); Exit1 --&gt; Join1(( )); Exit2 --&gt; Join1; Exit3 --&gt; Join1; Exit4 --&gt; Join1; Join1 --&gt; End(( ));</pre> <p>The flowchart for the 'id' grammar rule shows a central entry point leading to an oval labeled 'id'. From 'id', the flow branches into four main paths: 1) through an oval '[' to a box 'ITEM' and then an oval ']' to an exit; 2) through an oval '(' to a box 'EXPRESSION' and then an oval ')' to an exit; 3) through an oval ',' to a box 'EXPRESSION' and then an oval ')' to an exit; and 4) a direct exit. The exits from the first three paths are merged into a single final exit arrow. Below the main flowchart, four ovals are listed vertically: 'cst_whole', 'cst_decimal', 'cst_words', and 'cst_boolean', each with an arrow pointing to the main entry line before the 'id' oval, indicating they are constant values for this non-terminal.</p>
DECLARATIONS	 <pre>graph LR; Entry(( )) --&gt; VAR_DECLARATION[VAR_DECLARATION]; Entry --&gt; LIST_DECLARATION[LIST_DECLARATION];</pre> <p>The flowchart for declarations is simple, showing a single entry point that branches into two rectangular boxes: 'VAR_DECLARATION' and 'LIST_DECLARATION'.</p>

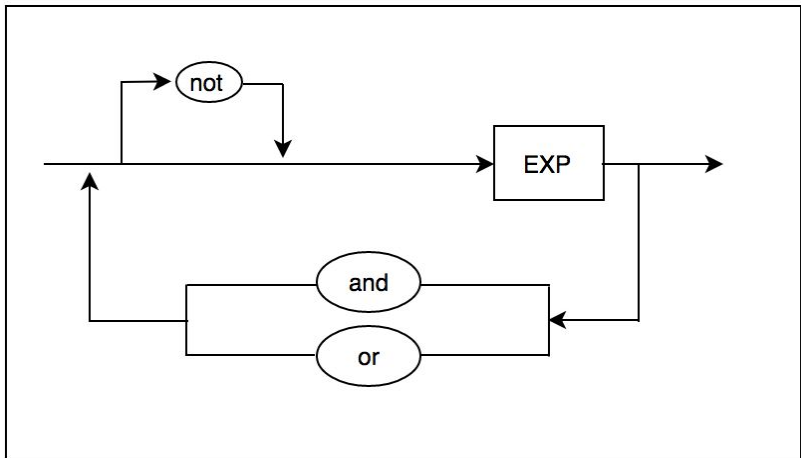
VAR\_DECLARATION



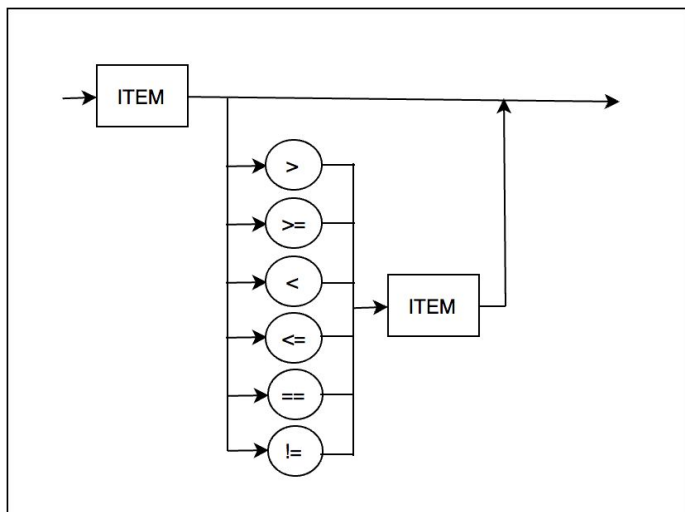
LIST\_DECLARATION

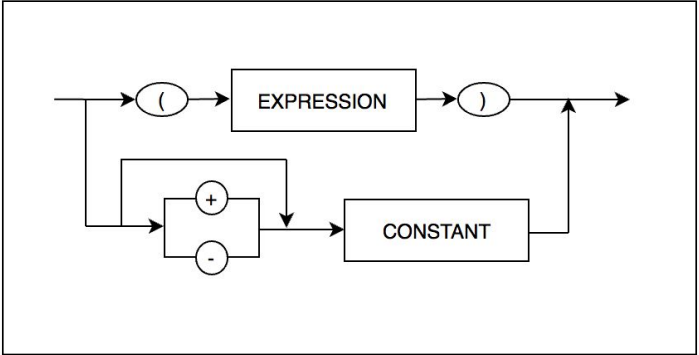
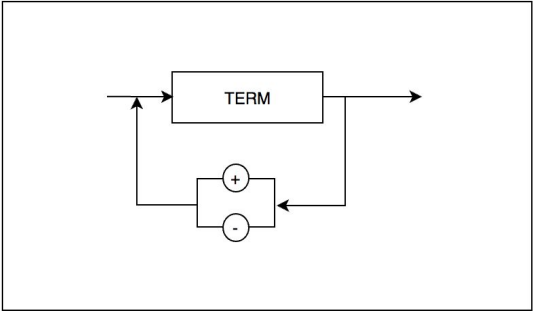
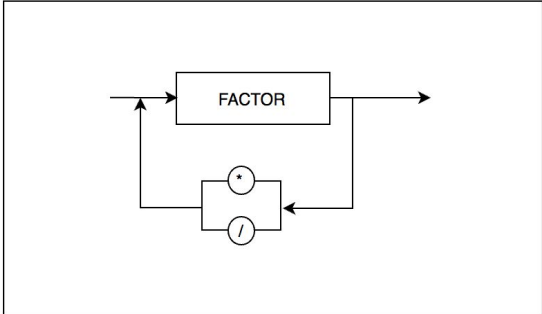
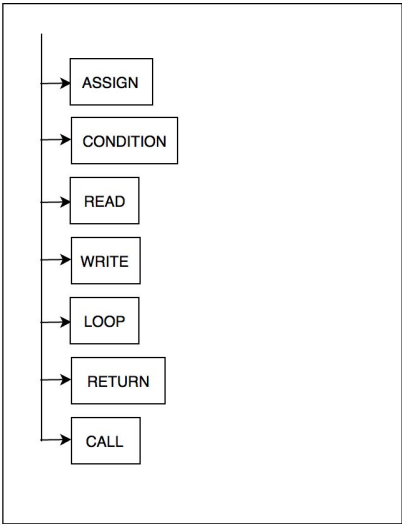


EXPRESSION



EXP

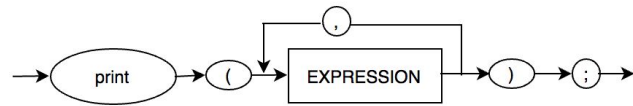


FACTOR	
ITEM	
TERM	
STATEMENT	



CALL	<pre> graph LR     start(( )) --&gt; call([call])     call --&gt; id([id])     id --&gt; LP("(")     LP --&gt; comma((,))     comma --&gt; EXPRESSION[EXPRESSION]     EXPRESSION --&gt; RP(")")     LP --&gt; RP     RP --&gt; semicolon([;])     semicolon --&gt; end(( )) </pre>
LOOP	<pre> graph LR     start(( )) --&gt; do([do])     do --&gt; BODY[BODY]     BODY --&gt; until([until])     until --&gt; LP("(")     LP --&gt; EXPRESSION[EXPRESSION]     EXPRESSION --&gt; RP(")")     RP --&gt; end(( )) </pre>
ASSIGN	<pre> graph LR     start(( )) --&gt; id([id])     id --&gt; LP("(")     LP --&gt; ITEM[ITEM]     ITEM --&gt; RP(")")     RP --&gt; equals([=])     equals --&gt; EXPRESSION[EXPRESSION]     EXPRESSION --&gt; semicolon([;])     semicolon --&gt; end(( )) </pre>
RETURN	<pre> graph LR     start(( )) --&gt; return([return])     return --&gt; EXPRESSION[EXPRESSION]     EXPRESSION --&gt; semicolon([;])     semicolon --&gt; end(( )) </pre>
READ	<pre> graph LR     start(( )) --&gt; input([input])     input --&gt; LP1("(")     LP1 --&gt; id([id])     id --&gt; LP2("(")     LP2 --&gt; ITEM[ITEM]     ITEM --&gt; RP2(")")     LP2 --&gt; RP2     RP2 --&gt; RP1(")")     LP1 --&gt; RP1     RP1 --&gt; semicolon([;])     semicolon --&gt; end(( )) </pre>

WRITE



### 3.3 Context Free Grammar

Syntactic Rule	Context Free Grammar (LR Parsing)
<b>PROGRAM</b>	<b>PROGRAM_AUX</b>
<b>PROGRAM_AUX</b>	<b>BLOCK</b>   <b>PROGRAM_AUX BLOCK</b>
<b>BLOCK</b>	<b>BLOCK_AUX</b> block id <b>RECEIVES_AUX RETURNS_AUX</b> <b>BLOCK_BODY</b>
<b>BLOCK_AUX</b>	starting   $\epsilon$
<b>RECEIVES_AUX</b>	receives colon id of type <b>TYPE RECEIVES_AUX1</b>   $\epsilon$
<b>RECEIVES_AUX1</b>	comma id of type <b>TYPE RECEIVES_AUX1</b>   $\epsilon$
<b>RETURNS_AUX</b>	block_returns <b>TYPE</b>   $\epsilon$
<b>BLOCK_BODY</b>	curlybraces_open <b>BLOCK_BODY_AUX</b> curlybraces_close
<b>BLOCK_BODY_AUX</b>	<b>DECLARATIONS BLOCK_BODY_AUX</b>   <b>BLOCK_BODY_AUX1</b>
<b>BLOCK_BODY_AUX1</b>	<b>STATEMENT BLOCK_BODY_AUX1</b>   $\epsilon$
<b>TYPE</b>	whole   decimal   words   boolean
<b>BODY</b>	curlybraces_open <b>BODY_AUX</b> curlybraces_close
<b>BODY_AUX</b>	<b>STATEMENT BODY_AUX</b>   $\epsilon$
<b>CONDITION</b>	if parenthesis_open <b>EXPRESSION</b> parenthesis_close <b>BODY</b> <b>CONDITION_AUX</b>
<b>CONDITION_AUX</b>	else <b>BODY</b>   $\epsilon$
<b>CONSTANT</b>	id <b>CONSTANT_AUX</b>   cst_whole   cst_decimal

	cst_words   cst_boolean
<b>CONSTANT_AUX</b>	squarebracket_open <b>ITEM</b> squarebracket_close   parenthesis_open <b>CONSTANT_AUX1</b> parenthesis_close   ε
<b>CONSTANT_AUX1</b>	<b>EXPRESSION</b> <b>CONSTANT_AUX2</b>   ε
<b>CONSTANT_AUX2</b>	comma <b>EXPRESSION</b> <b>CONSTANT_AUX2</b>   ε
<b>DECLARATIONS</b>	<b>VAR_DECLARATION</b>   <b>LIST_DECLARATION</b>
<b>VAR_DECLARATION</b>	variable id <b>VAR_DECLARATION_AUX</b> of type <b>TYPE</b> semicolon
<b>VAR_DECLARATION_AUX</b>	comma id <b>VAR_DECLARATION_AUX</b>   ε
<b>LIST_DECLARATION</b>	list id squarebracket_open <b>cst_whole</b> squarebracket_close of type <b>TYPE</b> semicolon
<b>EXPRESSION</b>	<b>EXPRESSION_AUX</b> <b>EXP</b> <b>EXPRESSION_AUX1</b>
<b>EXPRESSION_AUX</b>	op_negation   ε
<b>EXPRESSION_AUX1</b>	<b>op_and</b> <b>EXPRESSION</b>   <b>op_or</b> <b>EXPRESSION</b>   ε
<b>EXP</b>	<b>ITEM</b> <b>EXP_AUX</b>
<b>EXP_AUX</b>	op_less <b>ITEM</b>   op_less_equal <b>ITEM</b>   op_greater <b>ITEM</b>   op_greater_equal <b>ITEM</b>   op_equal <b>ITEM</b>   op_not_equal <b>ITEM</b>   ε
<b>FACTOR</b>	parenthesis_open <b>EXPRESSION</b> parenthesis_close   <b>FACTOR_AUX</b>
<b>FACTOR_AUX</b>	<b>op_addition</b> <b>CONSTANT</b>   <b>op_subtraction</b> <b>CONSTANT</b>   <b>CONSTANT</b>
<b>ITEM</b>	<b>TERM</b> <b>ITEM_AUX</b>
<b>ITEM_AUX</b>	op_addition <b>ITEM</b>

	op_subtraction <b>ITEM</b>   $\epsilon$
<b>TERM</b>	<b>FACTOR TERM_AUX</b>
<b>TERM_AUX</b>	op_multiplication <b>TERM</b>   op_division <b>TERM</b>   $\epsilon$
<b>STATEMENT</b>	<b>ASSIGN</b>   <b>CONDITION</b>   <b>READ</b>   <b>WRITE</b>   <b>LOOP</b>   <b>RETURN</b>   <b>CALL</b>
<b>CALL</b>	call id parenthesis_open <b>CALL_AUX</b> parenthesis_close semicolon
<b>CALL_AUX</b>	<b>EXPRESSION CALL_AUX2</b>   $\epsilon$
<b>CALL_AUX2</b>	comma <b>EXPRESSION CALL_AUX2</b>   $\epsilon$
<b>LOOP</b>	do <b>BODY</b> until parenthesis_open <b>EXPRESSION</b> parenthesis_close
<b>ASSIGN</b>	id <b>ASSIGN_AUX</b> op_assign <b>EXPRESSION</b> semicolon
<b>ASSIGN_AUX</b>	squarebracket_open <b>ITEM</b> squarebracket_close   $\epsilon$
<b>RETURN</b>	return_statement <b>EXPRESSION</b> semicolon
<b>READ</b>	input parenthesis_open id <b>READ_AUX</b> parenthesis_close semicolon
<b>READ_AUX</b>	squarebracket_open <b>ITEM</b> squarebracket_close   $\epsilon$
<b>WRITE</b>	print parenthesis_open <b>EXPRESSION WRITE_AUX</b> parenthesis_close semicolon
<b>WRITE_AUX</b>	comma <b>EXPRESSION WRITE_AUX</b>   $\epsilon$

### 3.3 Semantic Characteristics

- The keyword “starting” denotes the initial function.
- Variables have to be declared before being used in any way.
- Id names can’t be repeated.
- A square bracket should only be used after an id that represents an array.
- Variables have a local scope.
- Only the following operations between types are allowed (note that the order of the types are irrelevant and every other combination is considered an error):

Valid Operations			
Operand 1	Operator	Operand 2	Result
whole	+   -   *   /	whole	result of type whole number: operation is mathematically evaluated
whole	+   -   *   /	decimal	result of type decimal number: operation is mathematically evaluated
decimal	+   -   *   /	decimal	result of type decimal number: operation is mathematically evaluated
words	+	words	result of type words: Operand 2 is added to Operand 1 (string concatenation)
whole	<   <=   >   >=   ==   !=	whole	result of type boolean: evaluates the relational operation
whole	<   <=   >   >=   ==   !=	decimal	result of type boolean: evaluates the relational operation
decimal	<   <=   >   >=   ==   !=	decimal	result of type boolean: evaluates the relational operation
words	==   !=	words	result of type boolean: evaluates the relational operation
boolean	and	boolean	result of type boolean: evaluates the logical operation
boolean	or	boolean	result of type boolean: evaluates the logical operation
-	not	boolean	result of type boolean: evaluates the unary operation

### 3.4 Special Functions

There are some functions that work similar to those used in C++ or Python. The logic is the same but keywords are changed to facilitate the reading of code for children. For example:

Conventional Name	Choo Choo Train Name
Function/Method	block
var	variable
while	do until
&&	and
	or
!	not

There are also three modifications for the language. Parameters in a block definition are defined with the keyword “receives” so that it is known that parameters are variables that the block receives. There is a keyword for “of type” that specifies any variable’s type, instead of the conventional way of defining variables as first specifying type and then the variable id. While loops are called “do until” loops, since this improves the clarity and readability of the code for little children.

### 3.5 Primitive Data Types

There are 4 primitive data types in our language. We wanted to make it simple for children to understand, so while we kept the types:

- integer
- float
- string
- boolean

We decided to change their names to: whole, decimal and words, respectively; excluding the boolean type. These are more commonly used terms in the english language, so it will help children understand the values being represented by each type. The boolean type remains untouched, and is named the same, for there is no term that would semantically represent the type boolean better than its own name. To make it easier for children, values will be restricted when declaring boolean variables only to True and False.

Lists are also part of the language, and can hold values for any of the available data types mentioned above. They are declared by following the same format as a

variable declaration, but square brackets enclosing the size of the list follow the id name.

## Language and Computers

The project will be developed using Python, Google's Blockly Library, and using two Macbook Pros.

## Bibliography

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<<https://developers.google.com/blockly/>>