Functional Programming
IO

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Referential transparency and substitutivity

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- Every variable can be replaced by its definition
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Enables reasoning

```
1 -- sequence of function calls does not matter
2 f () + g () == g () + f ()
3 -- number of function calls does not matter
4 f () + f ( ) == 2 * f ()
```

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- By substitutivity, this expression must behave like

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  which reads two inputs!
- VERY WRONG!!!

The dilemma

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A contradiction?

No!

- Instead of performing IO operations directly, there is an abstract type of **IO instructions**, which get executed lazily by the operating system
- Some instructions (e.g., read from a file) return values, so the abstract IO type is parameterized over their type
- Keep in mind: instructions are just values like any other

Haskell IO

The main function

Top-level result of a program is an IO "instruction".

```
1 main :: IO ()
2 main = undefined
```

- an instruction describes the **effect** of the program
- effect = IO action, imperative state change, ...

Kinds of instructions

Primitive instructions

```
1 -- defined in the Prelude
2 putChar :: Char -> IO ()
3 getChar :: IO Char
4 writeFile :: FileName -> String -> IO ()
5 readFile :: FileName -> IO String
```

and many more

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No op instruction

```
1 return :: a -> IO a
```

The IO instruction `return 42` performs no IO, but yields the value 42.

Combining two instructions

The bind operator >>=

Intuition: next instruction may depend on the output of the previous one

```
1 (>>=) :: IO a -> (a -> IO b) -> IO b
```

The instruction `m >>= f`

- executes `m :: IO a` first
- gets its result `x :: a`
- applies `f :: a -> IO b` to the result
- to obtain an instruction `f x :: IO b` that returns a `b`
- and executes this instruction to return a `b`

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Example

```
1 readFiles f1 f2 =
2     readFile f1 >>= \xs1 -> readFile f2
```

More convenient: do notation

```
1 copyFile source target =
2   undefined
3
4 doTwice io =
5   undefined
6
7 doNot io =
8   undefined
```

Translating do notation into >>= operations

- `do` *lastaction*

  $\longrightarrow$

  lastaction

- `do { x <- action1; instructions }`

  $\longrightarrow$

  `action1 >>= \x -> do { instructions }`

- `do { action1; instructions }`

  $\longrightarrow$

  `action1 >> do { instructions }`

- `do { let binding; instructions }`

  $\longrightarrow$

  `let binding in do { instructions }`

Instructions vs functions

Functions

behave the same each time they called

Instructions vs functions

Functions

behave the same each time they called

Instructions

may be interpreted differently each time they are executed, depending on context

Underlying concept: **Monad**

What's a monad?

- abstract datatype for instructions that produce values
- built-in combination >>=
- abstracts over different interpretations (computations)

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IO is a special case of a monad

- one very useful application for monad
- built into Haskell
- but there's more to the concept
- many more instances to come!

Hands-on task

Define a function

```
1 sortFile :: FilePath -> FilePath -> IO ()
2
3 -- sortFile inFile outFile
4 -- reads inFile, sorts its lines, and writes the result to outFile
5
6 -- recall
7 -- sort :: Ord a => [a] -> [a]
8 -- lines :: String -> [String]
9 -- unlines :: [String] -> String
```

Utilities

```
1 sequence :: [IO a] -> IO [a]
2 sequence_ :: [IO a] -> IO ()
```

Another hands-on task

Define a function

```
1 printTable :: [String] -> IO ()
2
3 {-
4 printTable ["New York", "Rio", "Tokio"]
5 outputs
6 1: New York
7 2: Rio
8 3: Tokio
9 -}
```

Wrapup

First meeting with monads

- abstract data type of instructions returning results
- next instruction can depend on previous result
- instructions are just values
- basis for Haskell's standard IO