

# Assignment problem for MAS

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# Introduction

**Starting conditions:** The MAS is located at airbases and receives multiple requests for service including:

- Location to visit;
- Number of air-vehicles required;
- Earliest time of 1-st visit;
- Latest time of 1-st visit;
- Minimum duration per visit;
- Maximum interval between visits.

## Mission Objective:

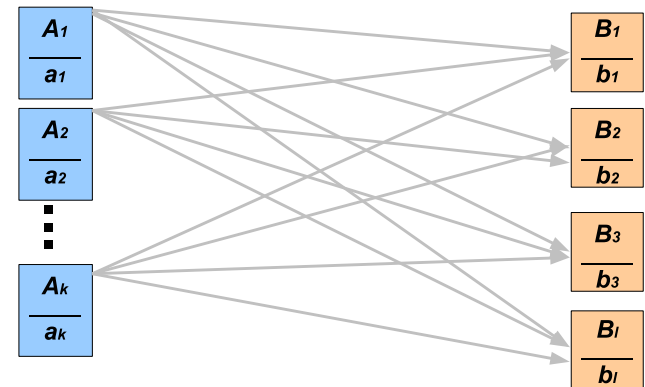
Air-vehicle(s) to be assigned to request and the corresponding paths to take to the service request; (variations to requests with minimal change if it cannot be met.)

## That:

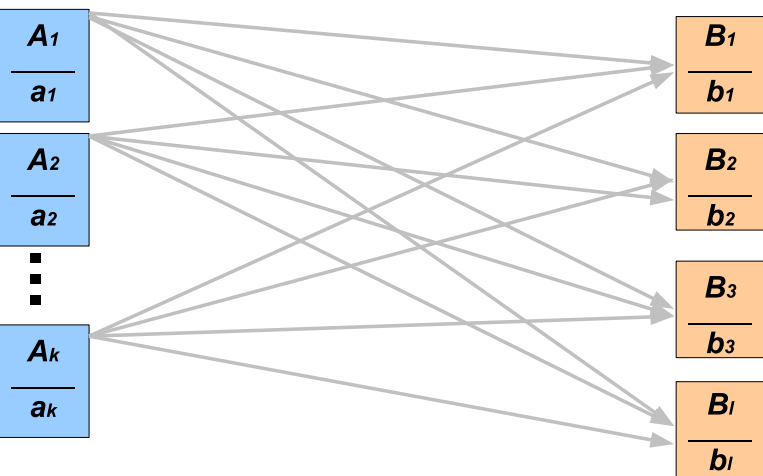
Maximize the number of service requests that can be serviced.

## Constraints

- Air-vehicle performance and dynamics;
- LOS occlusion in area of operations.



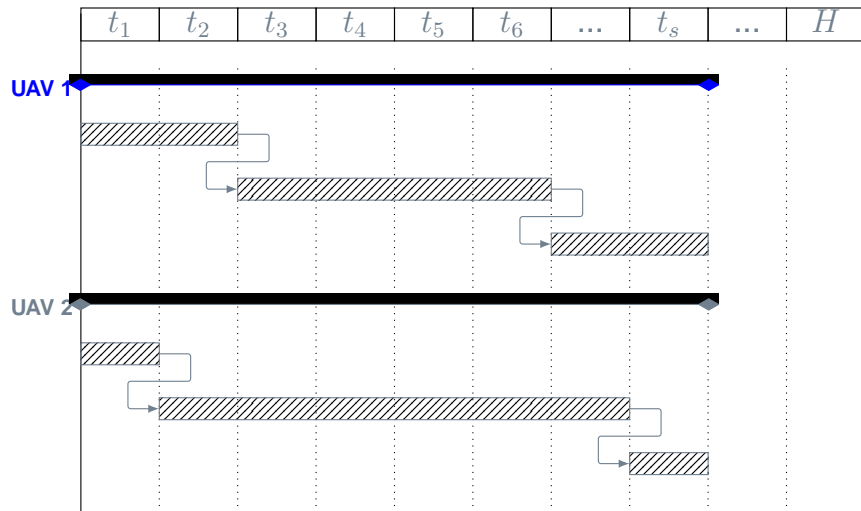
# Variables and constraints



Divide the interval  $[0, H]$  by the moments  $t_s = s\Delta$ ,  $s = 1, 2, \dots, \nu$  where  $\nu = \left\lceil \frac{H}{\Delta} \right\rceil$

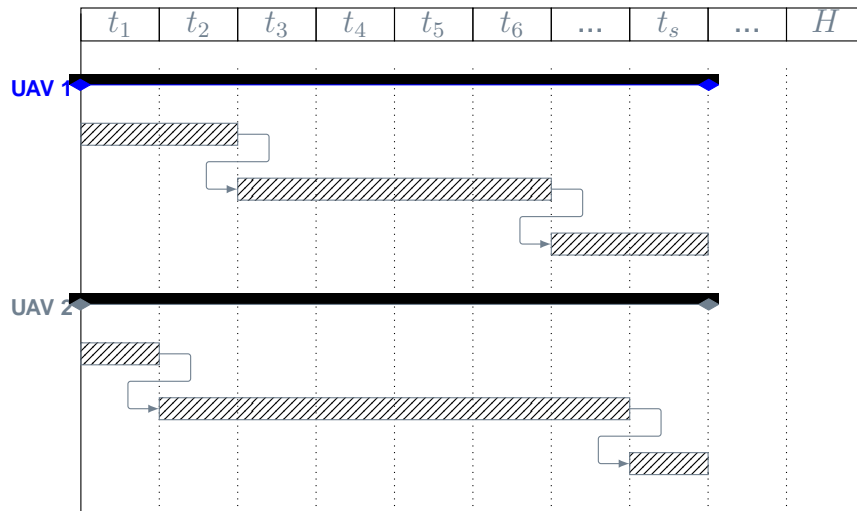
- $B_j, j = 1, \dots, l$  is the locations to visit;
- $x_{ij}(t_s)$  is the number of UAVs from  $i$ -th aerobase send to  $j$ -th zone at the moment  $t_s$ ;
- $b_j(t_s)$  is the number of UAVs that are serving the  $j$ -th zone at the moment  $t_s$ ;
- $A_i, i = 1, \dots, k$  is the locations of aerobases;
- $a_i(t_s)$  is the number of UAVs at the  $i$ -th aerobase at the moment  $t_s$ ;
- $k$  and  $l$  are the number of aerobases and zones for service, respectively;

# Variables and constraints



- $h_i$  is the flight endurance for UAVs from  $i$  - th aerobase.
- $t_{ij}$  is the flight time from  $i$  - th aerobase to  $j$  - th zone;
- $h_i - 2t_{ij}$  - maximal possible service time of the zone  $j$ ;

# Variables and constraints



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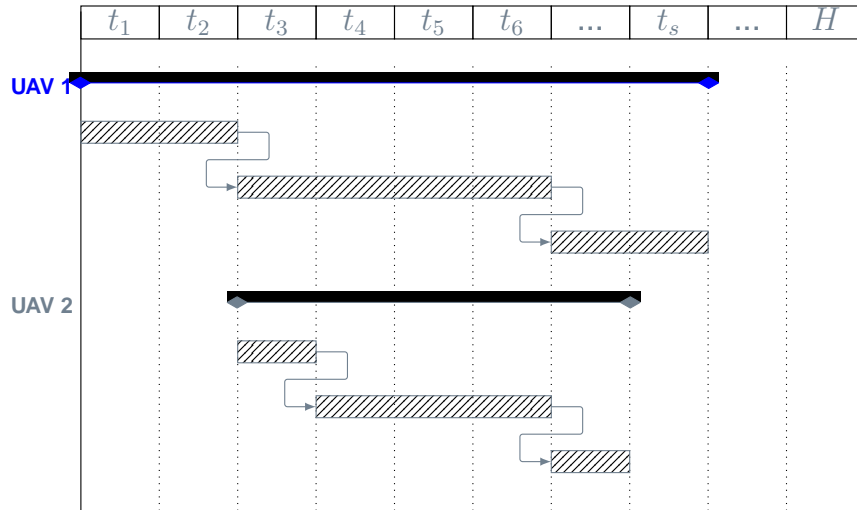
$$a_i(t_s + \Delta) = a_i(t_s) - \sum_{j=1}^l x_{ij}(t_s) + \sum_{j=1}^l x_{ij}(t_s + \Delta - h_i)$$

$$b_j(t_s + \Delta) = b_j(t_s) - \sum_{i=1}^k x_{ij}(t_s - h_i + t_{ij}) + \sum_{i=1}^k x_{ij}(t_s - t_{ij})$$

$$a_i(t_s) + \sum_{j=1}^l x_{ij}(t_s) = a_i; \quad b_j(t_s) + \sum_{i=1}^k x_{ij}(t_s - t_{ij}) = b_j,$$

$$(t_s - h_i + t_{ij} > 0); (t_s - t_{ij} > 0); i = 1, \dots, k; j = 1, \dots, l.$$

# Types of objective function



the total service time

$$J_1(x) = \sum_{i=1}^k \sum_{j=1}^l \sum_{s=0}^{\nu} x_{ij}(t_s)(h_i - 2t_{ij})$$

the total number of UAVs "circles"

$$J_2(x) = \sum_{i=1}^k \sum_{j=1}^l \sum_{s=0}^{\nu} x_{ij}(t_s)$$

the total unobservable time for multiple zones  $J_3(x) = H - J_1(x)$

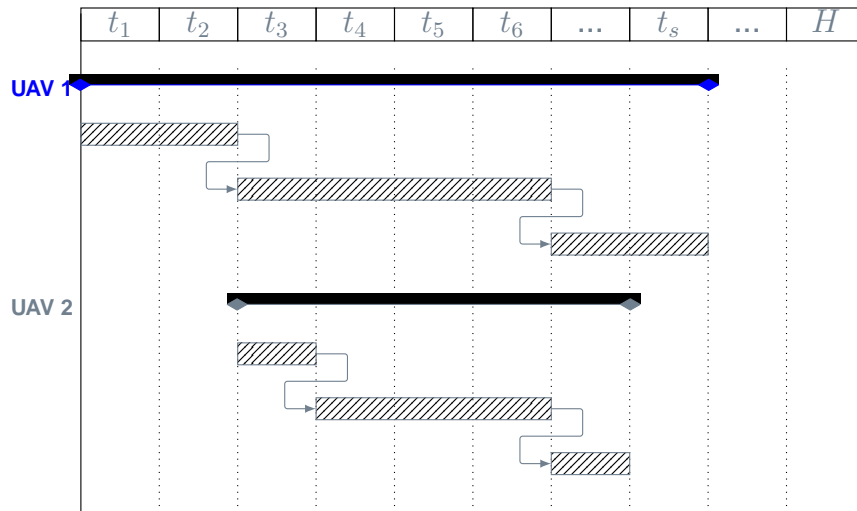
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$$(t_s - h_i + t_{ij} > 0); (t_s - t_{ij} > 0); i = 1, \dots, k; j = 1, \dots, l.$$

# Problem statement



the total service time

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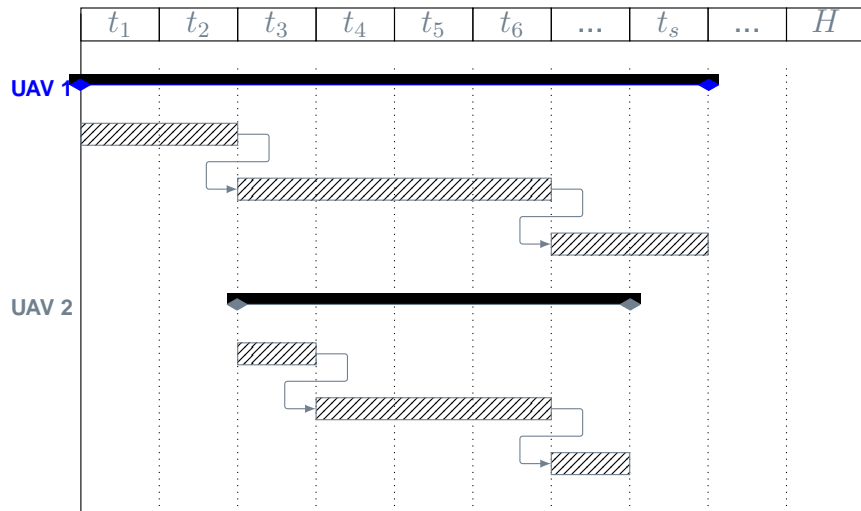
$$b_j(t_s + \Delta) = b_j(t_s) - \sum_{i=1}^k x_{ij}(t_s - h_i + t_{ij}) + \sum_{i=1}^k x_{ij}(t_s - t_{ij})$$

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$$(t_s - h_i + t_{ij} > 0); (t_s - t_{ij} > 0); i = 1, \dots, k; j = 1, \dots, l.$$

The proposed model can be extended by additional conditions (constraints) followed from the formal description.

# Problem statement



the total service time

$$J_1(x) = \sum_{i=1}^k \sum_{j=1}^l \sum_{s=0}^{\nu} x_{ij}(t_s)(h_i - 2t_{ij})$$

the total number of UAVs "circles"

$$J_2(x) = \sum_{i=1}^k \sum_{j=1}^l \sum_{s=0}^{\nu} x_{ij}(t_s)$$

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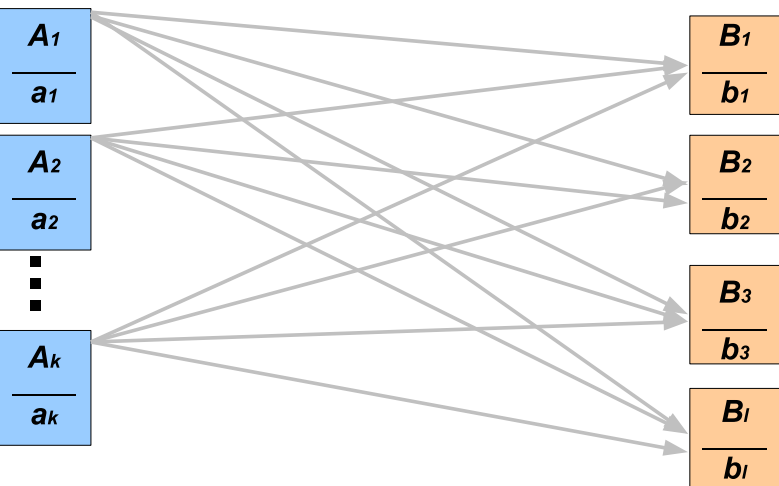
$$(t_s - h_i + t_{ij} > 0); (t_s - t_{ij} > 0); i = 1, \dots, k; j = 1, \dots, l.$$

$$\sum_{i=1}^k x_{ij}(t_s^{first}) \neq 0, s^{first} \Delta \leq t^{first} \leq (s^{first} + 1) \Delta$$

For example, if the time of earliest time of 1-st visit  $t_j^{first}$  for  $j$  zone is pre-assigned.



# LP assignment problem



$$F = \sum_{i=1}^k \sum_{j=1}^l c_{ij} x_{ij} \rightarrow \min_{x_{ij}}$$

$$\sum_{i=1}^k x_{ij} = b_j, \quad j = 1, 2, \dots, l$$

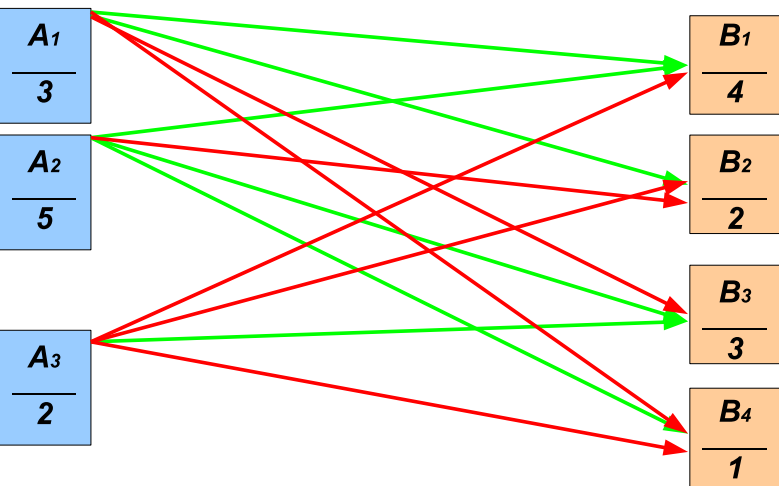
$$\sum_{j=1}^l x_{ij} = a_i, \quad i = 1, 2, \dots, k$$

$$\sum_{i=1}^k a_i = \sum_{j=1}^l b_j, \quad x_{ij} \geq 0, x_{ij} \in \mathbb{N}.$$

The most of methods include the following basic steps:

- To find initial plan  $x_{ij}$ ;
- Check optimality condition for that plan;
- Construct the improved plan in case of nonoptimality.

# Optimal solution



$$F = \sum_{i=1}^3 \sum_{j=1}^4 c_{ij} x_{ij} \rightarrow \min_{x_{ij}}$$

$$\sum_{i=1}^3 x_{ij} = b_j, \quad j = 1, 2, 3, 4$$

$$\sum_{j=1}^4 x_{ij} = a_i, \quad i = 1, 2, 3$$

$$\sum_{i=1}^3 a_i = \sum_{j=1}^4 b_j, x_{ij} \geq 0, x_{ij} \in \mathbb{N}.$$

	$B_1$	$B_2$	$B_3$	$B_4$	$a_i$	$u_i$
$A_1$	1	2	0	0	3	
$A_2$	3	0	1	1	5	
$A_3$	0	0	2	0	2	
$b_j$	4	2	3	1	$F=46$	
$\nu_j$						

Optimal solution:

$$x_{11} = 1; x_{12} = 2;$$

$$x_{21} = 3; x_{23} = 1;$$

$$x_{24} = 1; x_{33} = 2.$$

# Outline of the method

The method for LP assignment problem as well as simplex methods are:

- iterative,      • exact(satisfied all constraints),
- finite,      • relaxed(in a sense of the value of objective function).

Thus in some sense this method is analog of simplex method, but the ideas of this method is more naturally can be applied to assignments LP problems.

## *Further work*

The major task is to solve the specific assignment problem for MAS by developing

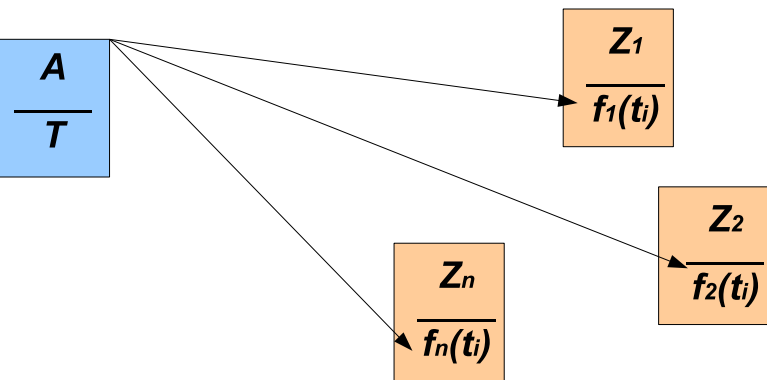
- new optimality condition,
- $\epsilon$  - optimality condition,
- realize sensitivity analysis (robustness analysis),  
*Siarhei Dymkou, Kai Yew Lum, Jian Xin Xu, Comparison of the adaptive method with classical simplex method for linear programming.(2011) (in preparation)*
- consider objectives with very general functional forms.

# The end

Three horizontal bars of varying lengths and colors (yellow, black, and grey) extending from the left side of the slide.

Thank you!

# Extra slides (DP)



- Invariant embedding of the problem into  $P(k, y), k \in [1; n]; y \in [0; T]$ , then define the Bellman function;
- Construct the Bellman equation;

$$\sum_{i=1}^n f_i(t_i) \rightarrow \max$$

$$\sum_{i=1}^n t_i \leq T, t_i \geq 0, i = 1, \dots, n$$

$\Rightarrow$

$$P(k, y) : \sum_{i=1}^k f_i(t_i) \rightarrow \max,$$

$$\sum_{i=1}^k t_i \leq y, t_i \geq 0, i = 1, \dots, k$$

Bellman equation

$$B_k(y) =$$

$$\max_{0 \leq z \leq y} \left[ f_k(z) + B_{k-1}(y - z) \right]$$

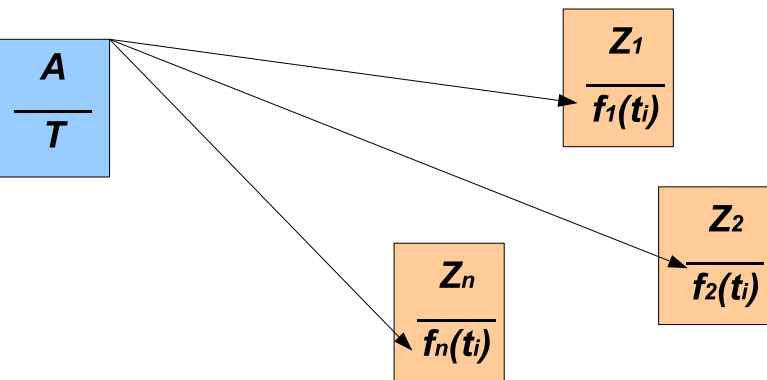
$\Rightarrow$

$$B_k(y) = \max_{t_i} \sum_{i=1}^k f_i(t_i),$$

$$\sum_{i=1}^k t_i \leq y, t_i \geq 0, i = 1, \dots, k$$

"back"

# Extra slides (DP)



- Invariant embedding of the problem into  $P(k, y), k \in [1; n]; y \in [0; T]$ , then define the Bellman function;
- Construct the Bellman equation;
- Solve the Bellman equation, and apply the solution of it to initial problem

$$t_n^0, t_{n-1}^0, \dots, t_2^0, t_1^0$$



*Continue this procedure we will find*

*the optimal solution of our problem*

Put  $k = n, y = T$  and find the value  $t_n^0 \doteq z^0(T)$  for the zone  $Z_n$

$k = n - 1, y = T - t_n^0$  and find the value  $t_{n-1}^0 \doteq z^0(T - t_n^0)$

$$f_n(z_n^0) + B_{n-1}(T - z_n^0) = \max_{0 \leq z \leq T} [f_n(z) + B_{n-1}(T - z)]$$



$$f_{n-1}(z_{n-1}^0) + B_{n-2}(T - z_{n-1}^0) = \max_{0 \leq z \leq T - t_n^0} [f_{n-1}(z) + B_{n-2}(T - t_n^0 - z)]$$