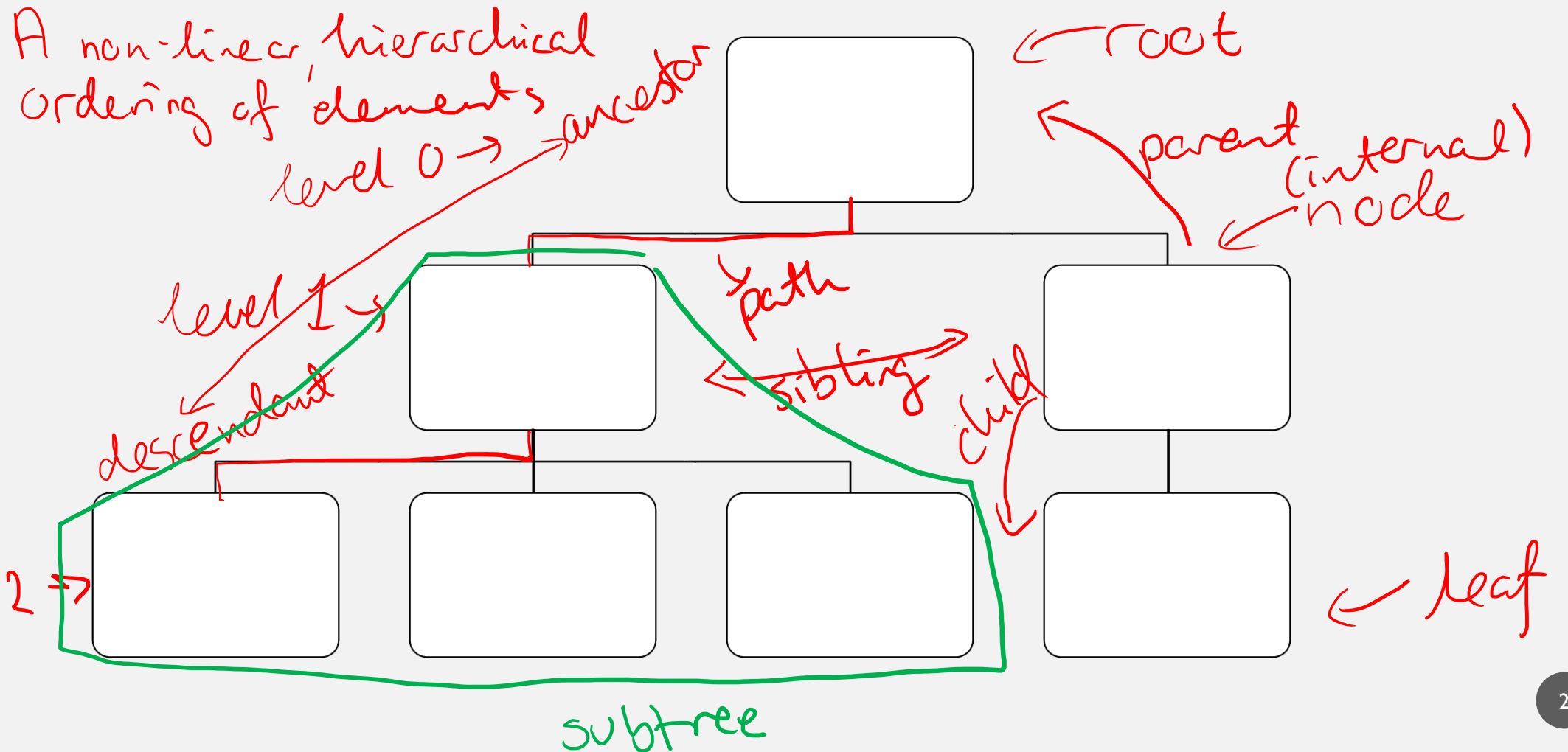


BINARY TREES

ADSI, S2023

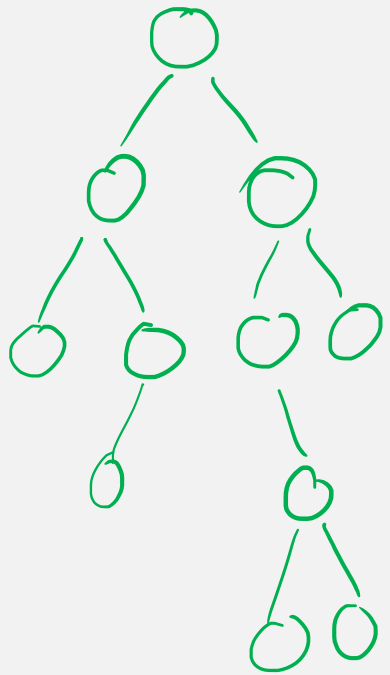
WHAT IS A TREE?

A non-linear, hierarchical
ordering of 'elements'

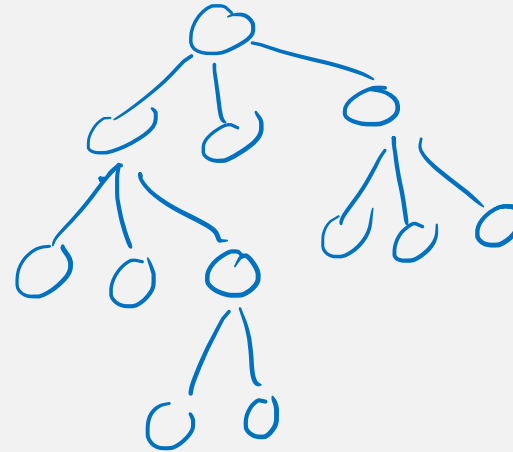


n -ARY TREES

Each node has at most n children



binary tree

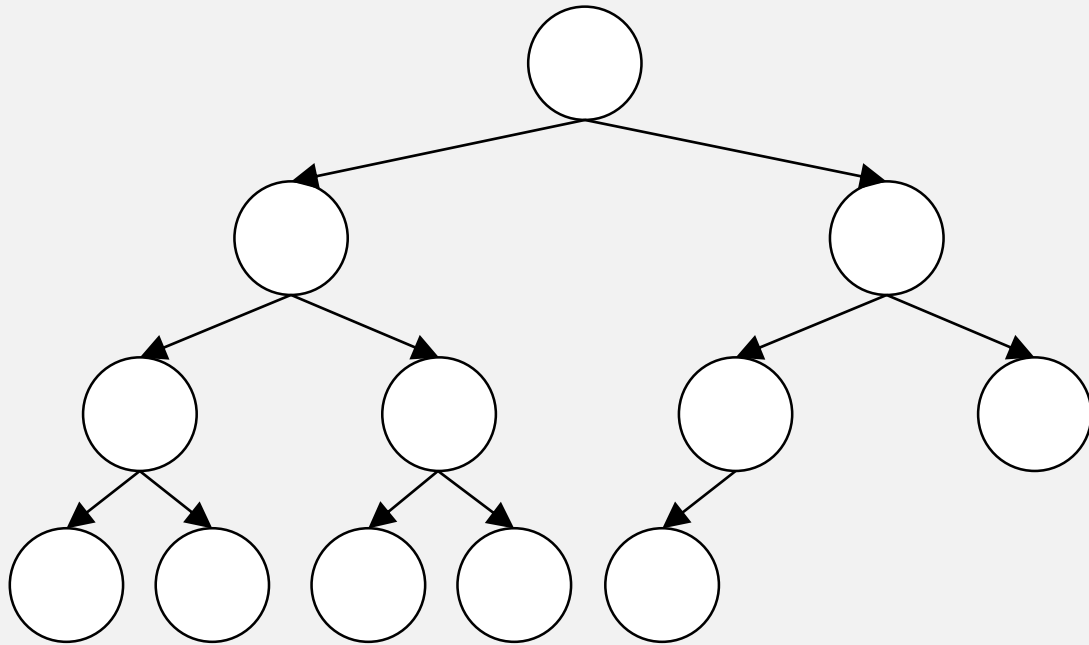


ternary tree

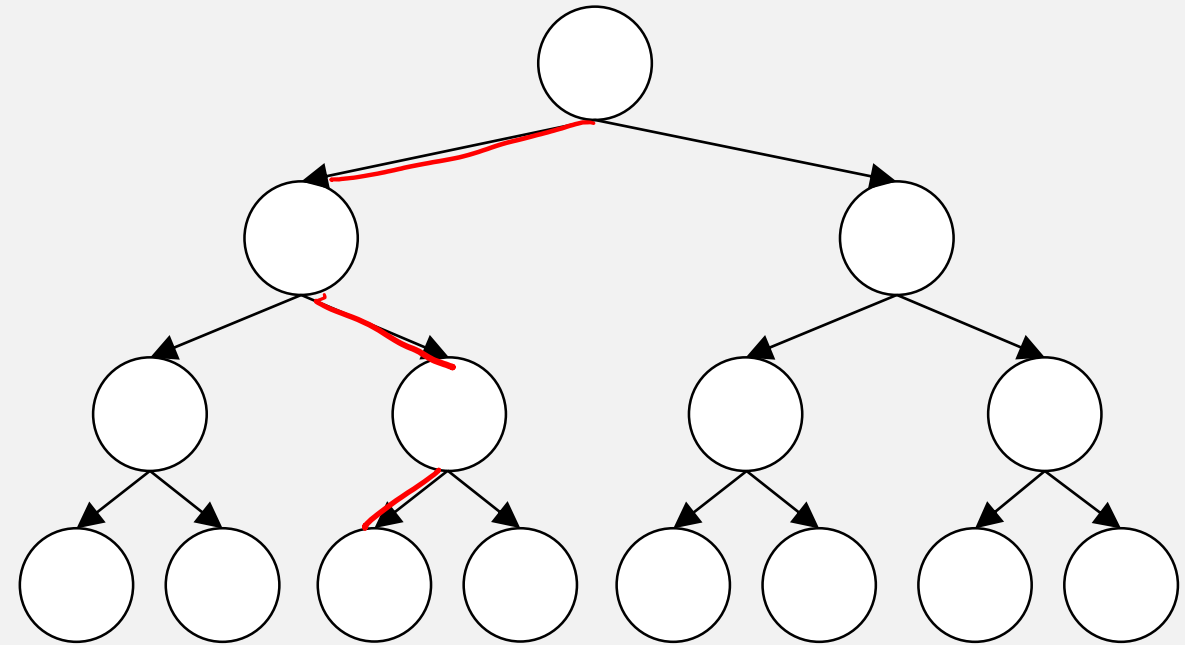
and so
on...

COMPLETE AND FULL TREES

BINARY



COMPLETE



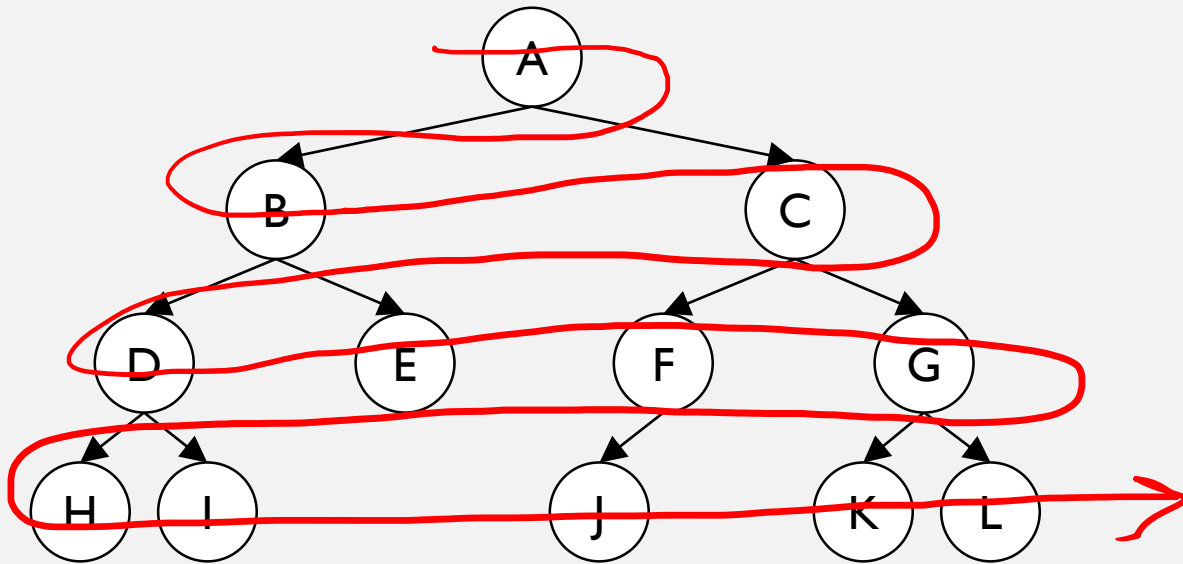
FULL

here, height
is "path length"

$$\#nodes = 2^{h+1} - 1$$

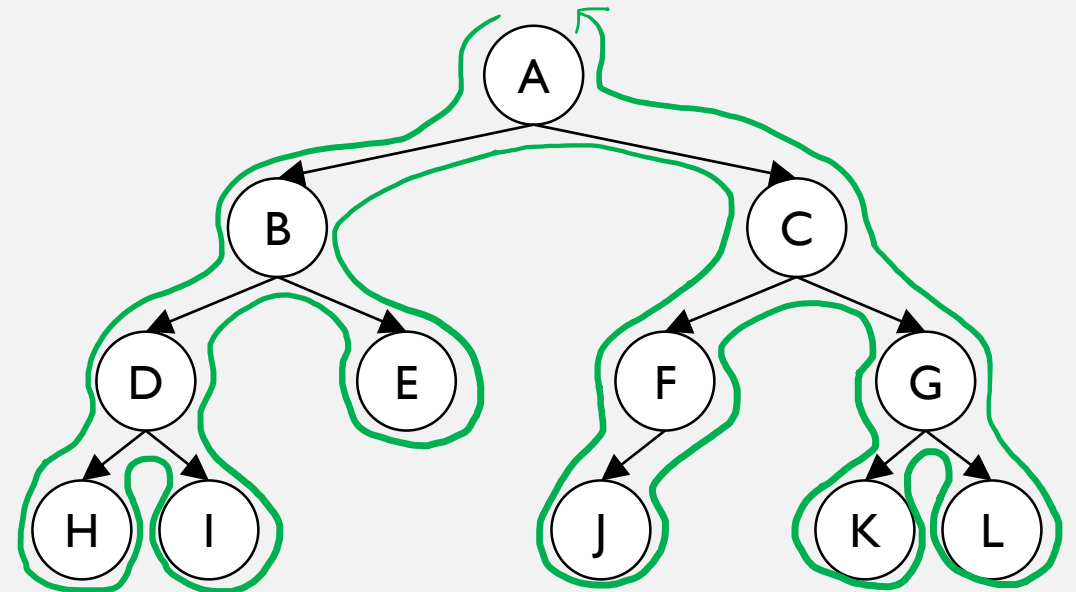
TREE TRAVERSAL

Breadth first traversal (BFT)



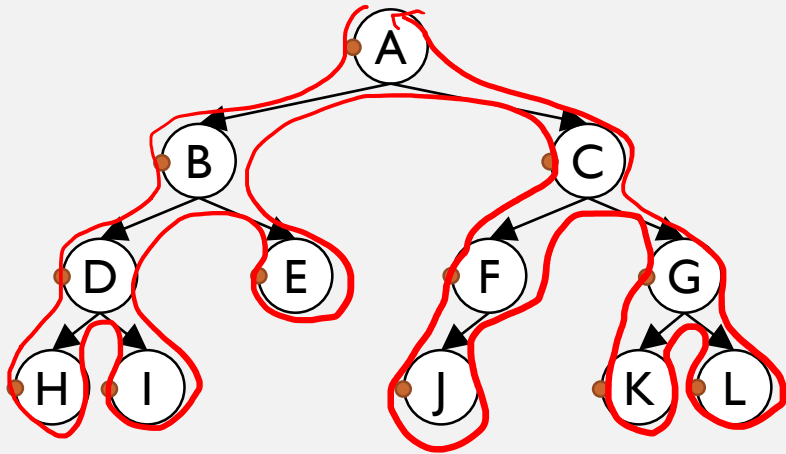
Traversal order: ABCDEFGHIJKL

Depth first traversal (DFT)



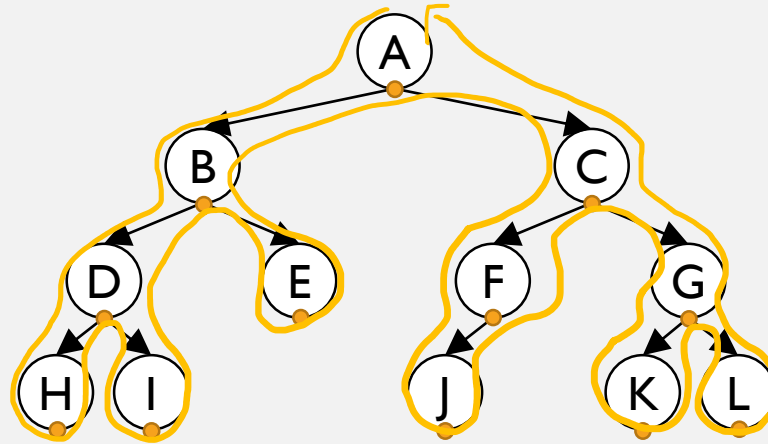
TREE TRAVERSAL

Pre-order DFT



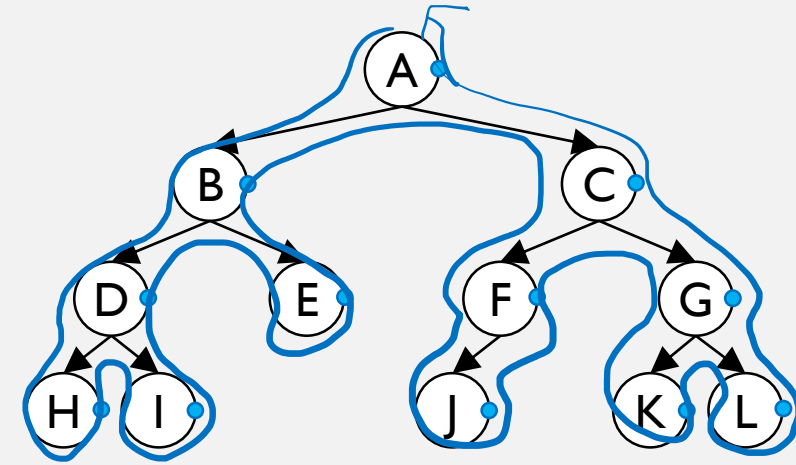
ABDHEICFJGKL

In-order DFT



HDIBEAFCKGL

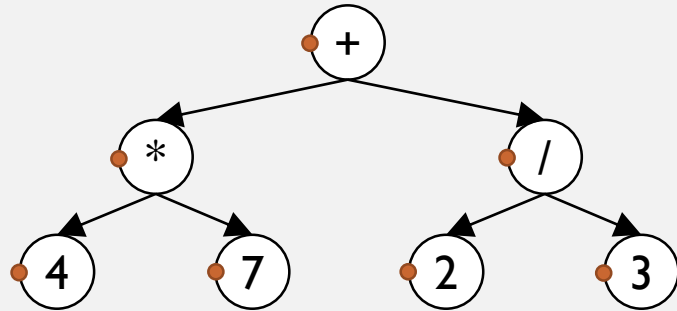
Post-order DFT



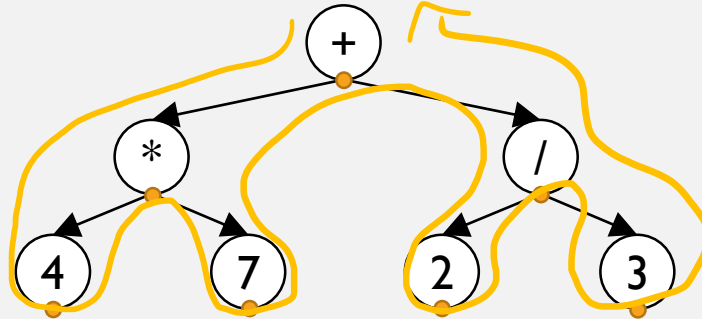
HI DEBJFKLGCA

EXPRESSION TREES

Pre-order DFT

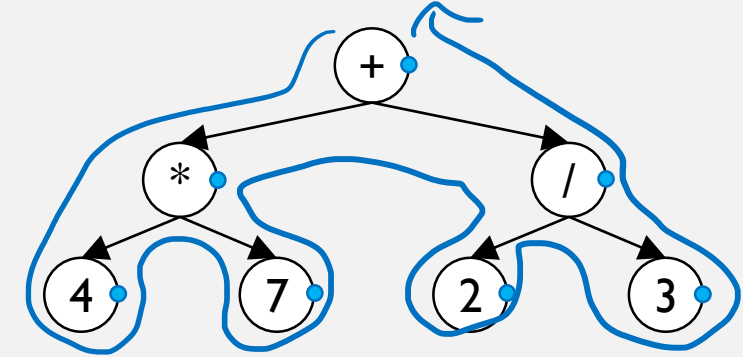


In-order DFT



$(4 * 7) + (2 / 3)$

Post-order DFT



$47 * 23 / +$

postfix
notation

THE (BINARY) TREE ADT

- getRoot
- isEmpty
- size
- contains
- inOrder
- preOrder
- postOrder
- levelOrder → BFT
- height

THE NODE ADT

- setElement
- getElement
- addLeftChild
- addRightChild
- getLeftChild
- getRightChild

REPRESENTING A BINARY TREE

"linked" binary tree

Expand the idea
of a linked list

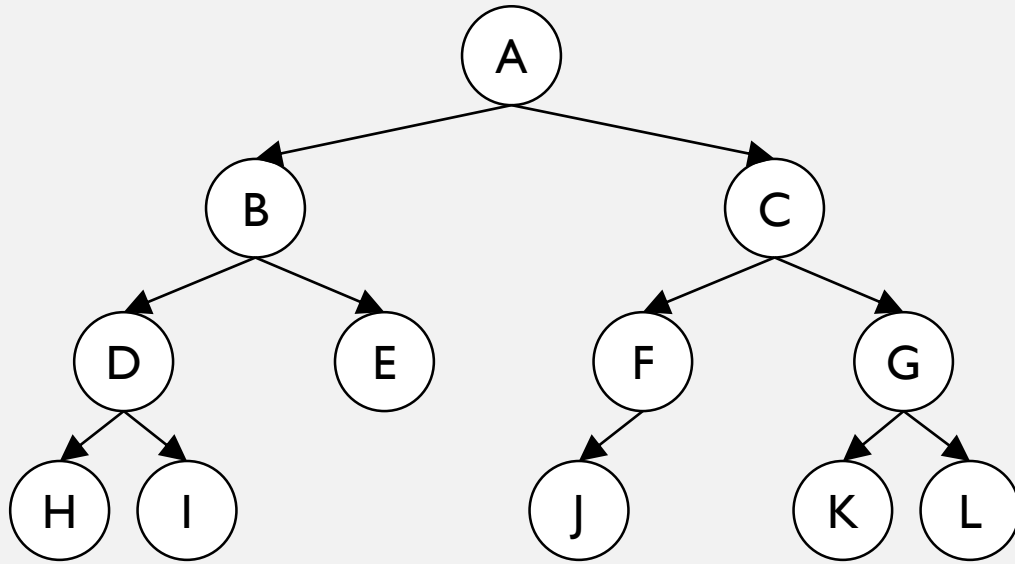
~~next~~ right parent
 left

"Array-like" binary tree
use an arraylist
→ not obvious

if we don't know if $c=l$ or r

"ARRAY-LIKE" REPRESENTATION

$$p = \left\lfloor \frac{c-1}{2} \right\rfloor$$



danger:
wasting
space

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
A	B	C	D	E	F	G	H	I	•	•	J	•	K	L

↑ parent p
1

↑ children
3 4

↑ parent
6

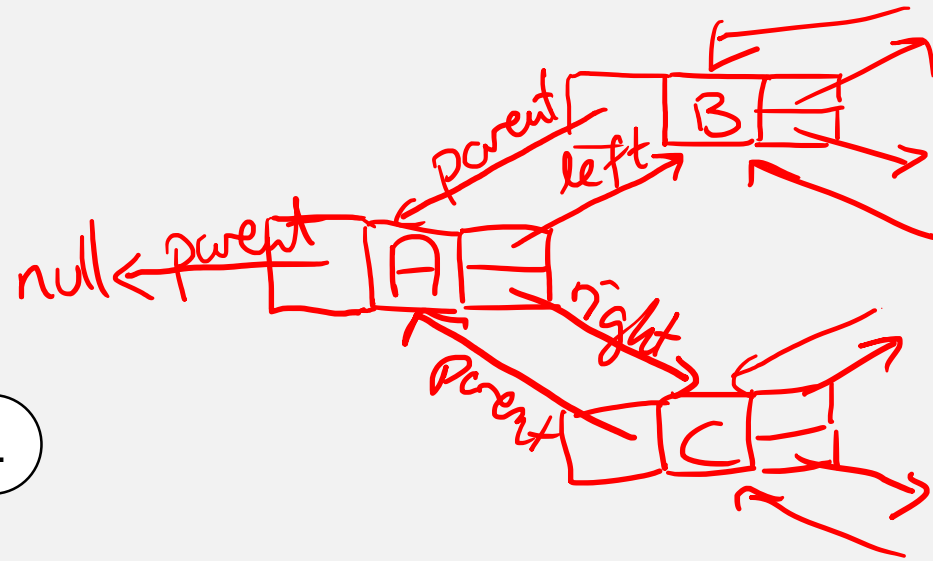
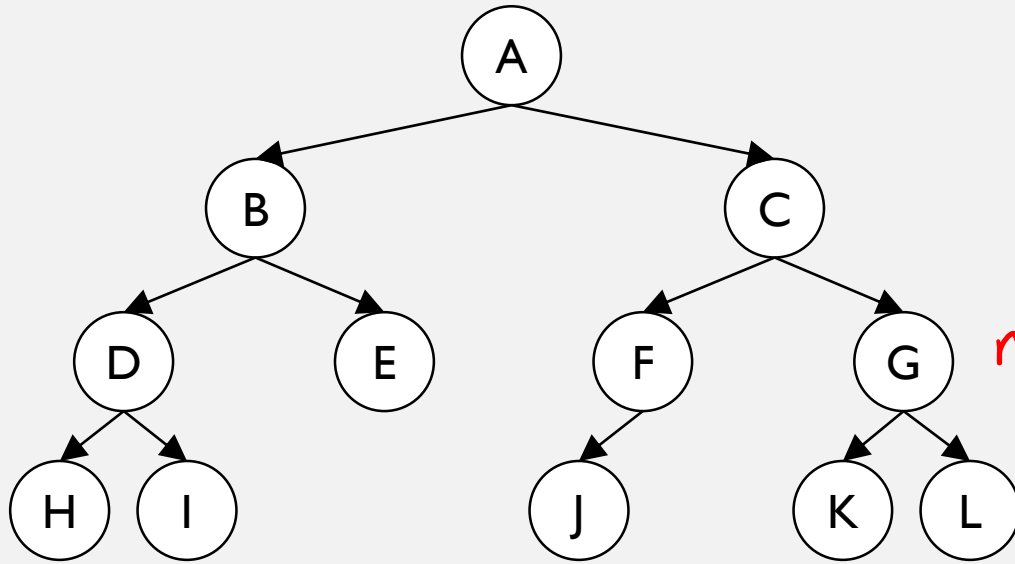
↑ children
13 14

In general, parent p :

$$\left. \begin{array}{l} \text{left child } l = 2p + 1 \\ \text{right child } r = 2p + 2 \end{array} \right\} p = \frac{c}{2} - 1 = \frac{l-1}{2}$$

"computed child links"

"LINKED" REPRESENTATION

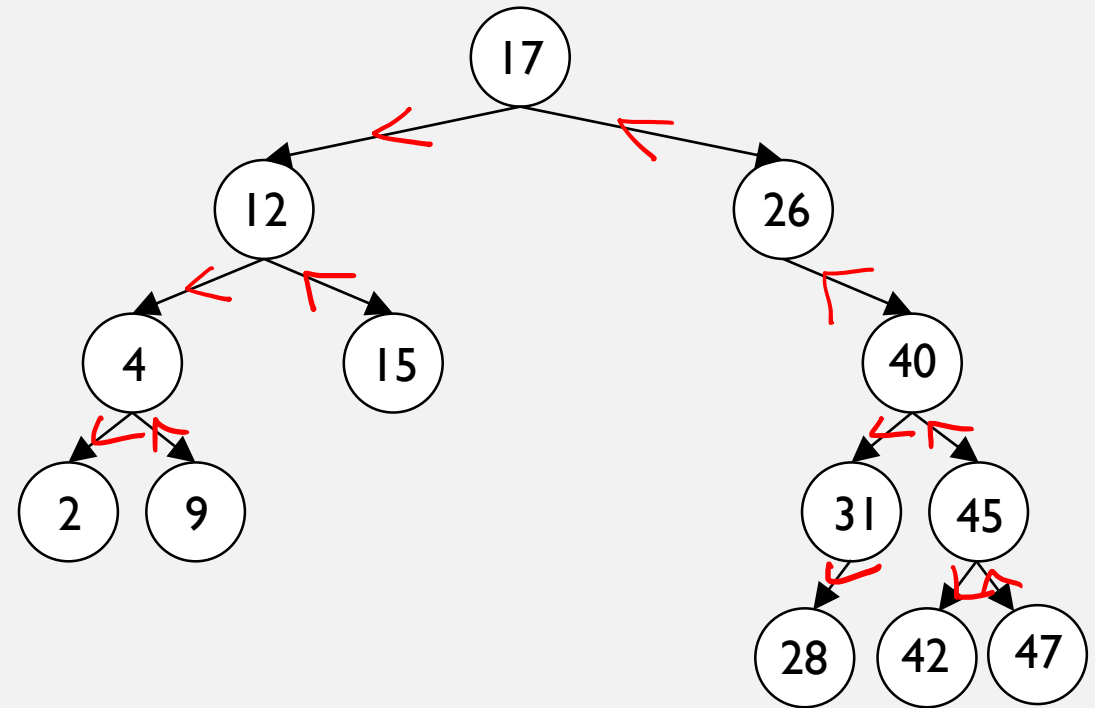


and so on...

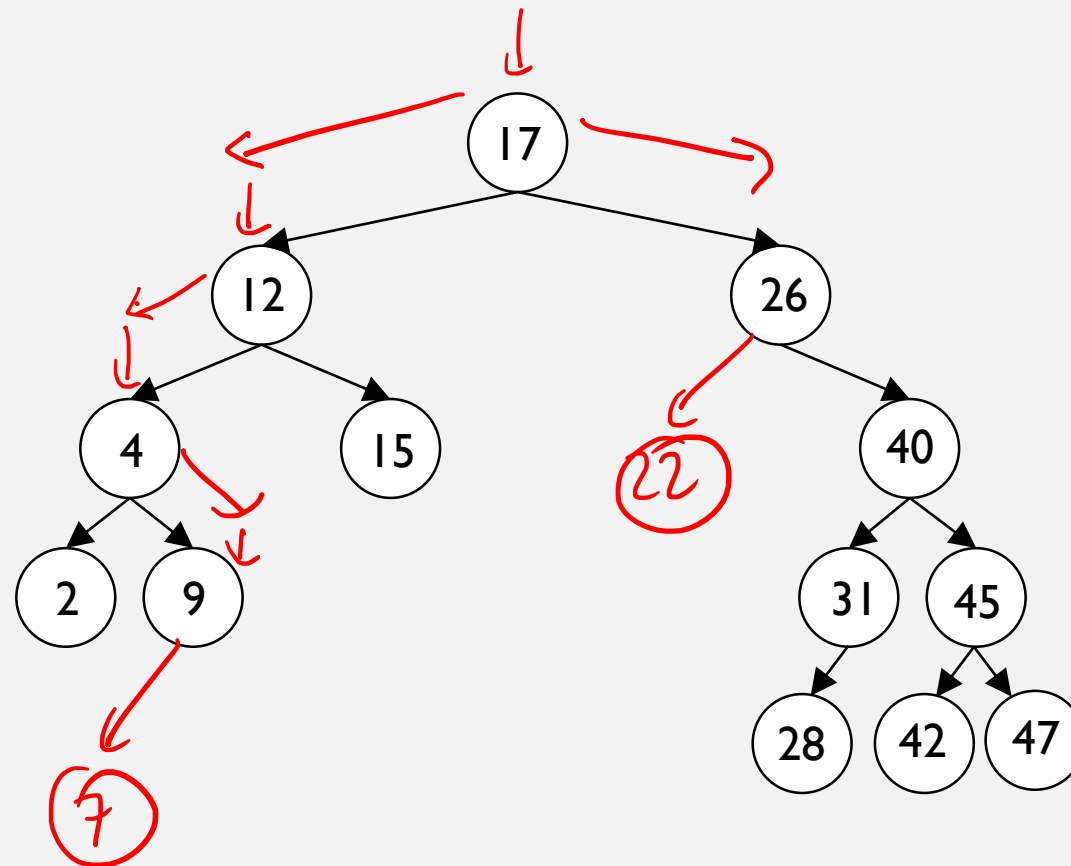
LinkedBinaryTree in java

BINARY SEARCH TREES

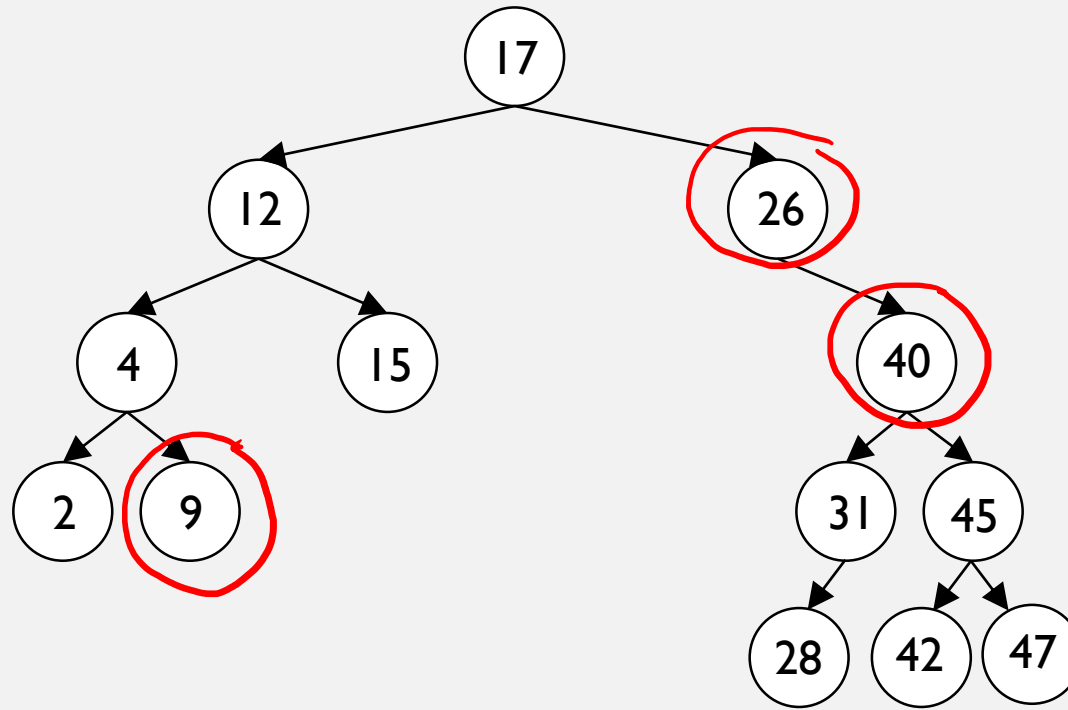
For every node n :
Each element in
left subtree $<$ element
in n and each
element in right
subtree $>$ element
in n .



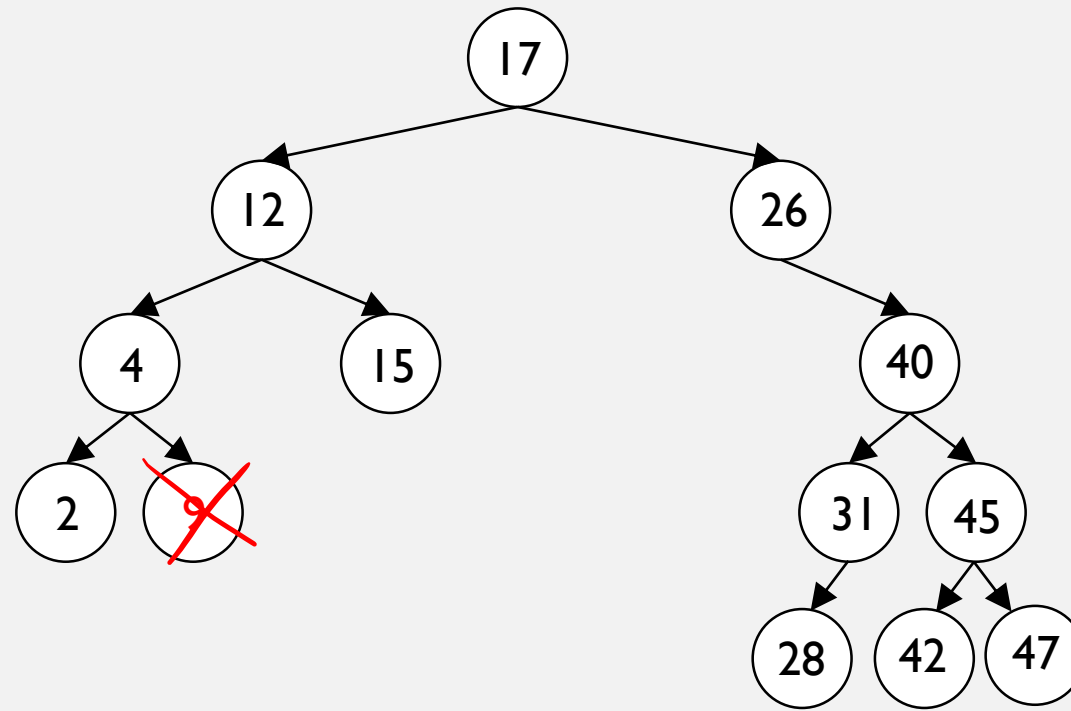
ADDING TO A BINARY SEARCH TREE



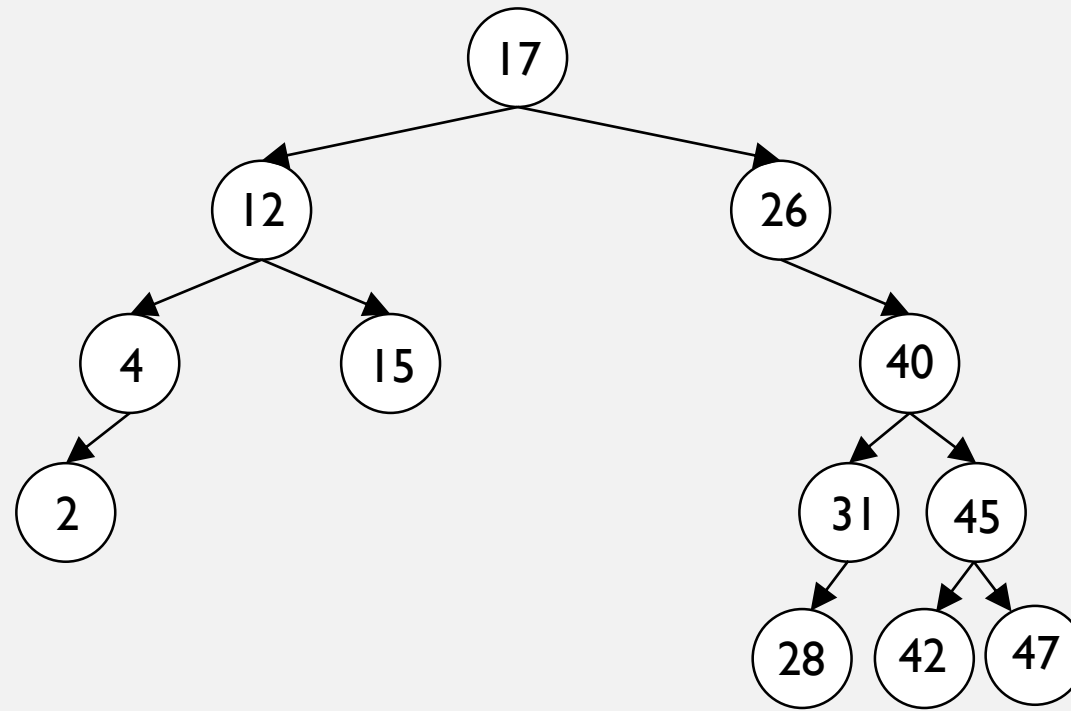
DELETING FROM A BINARY SEARCH TREE



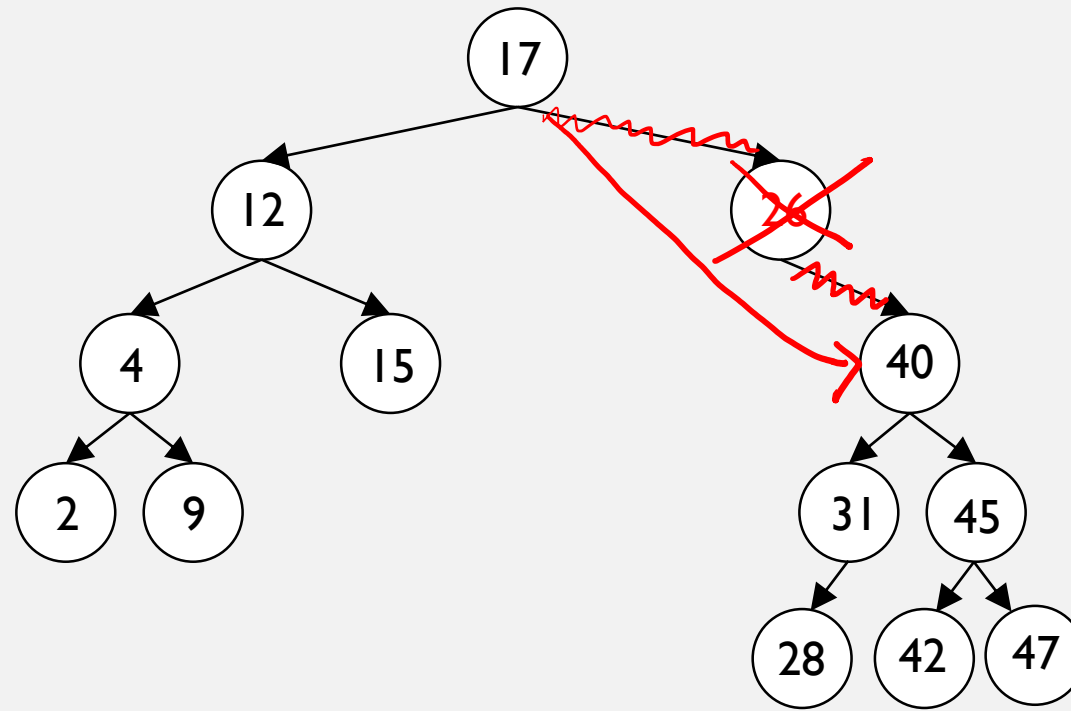
CASE I



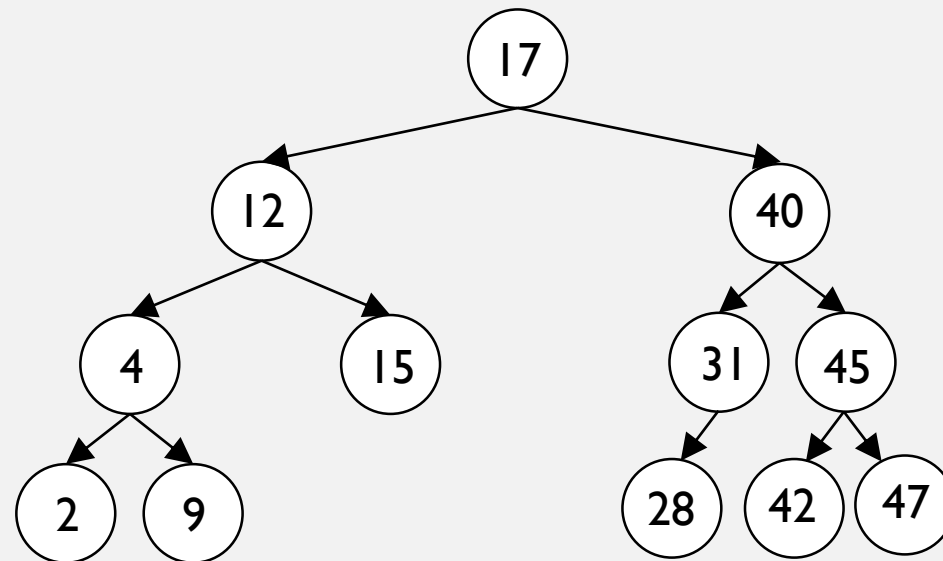
CASE I



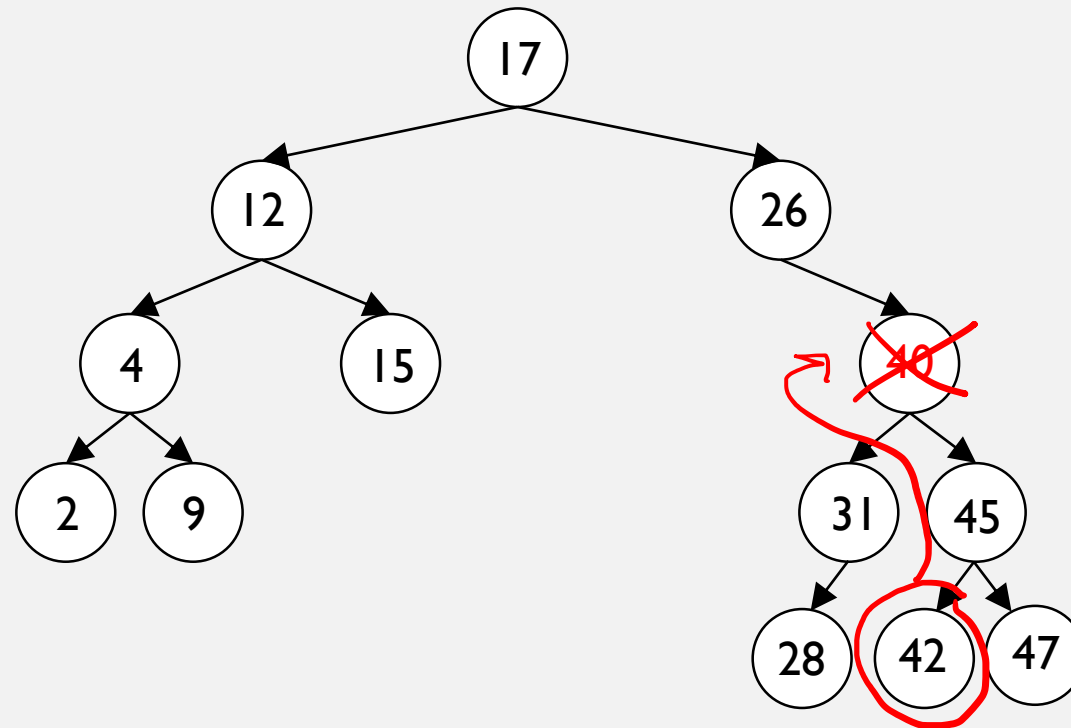
CASE II



CASE II

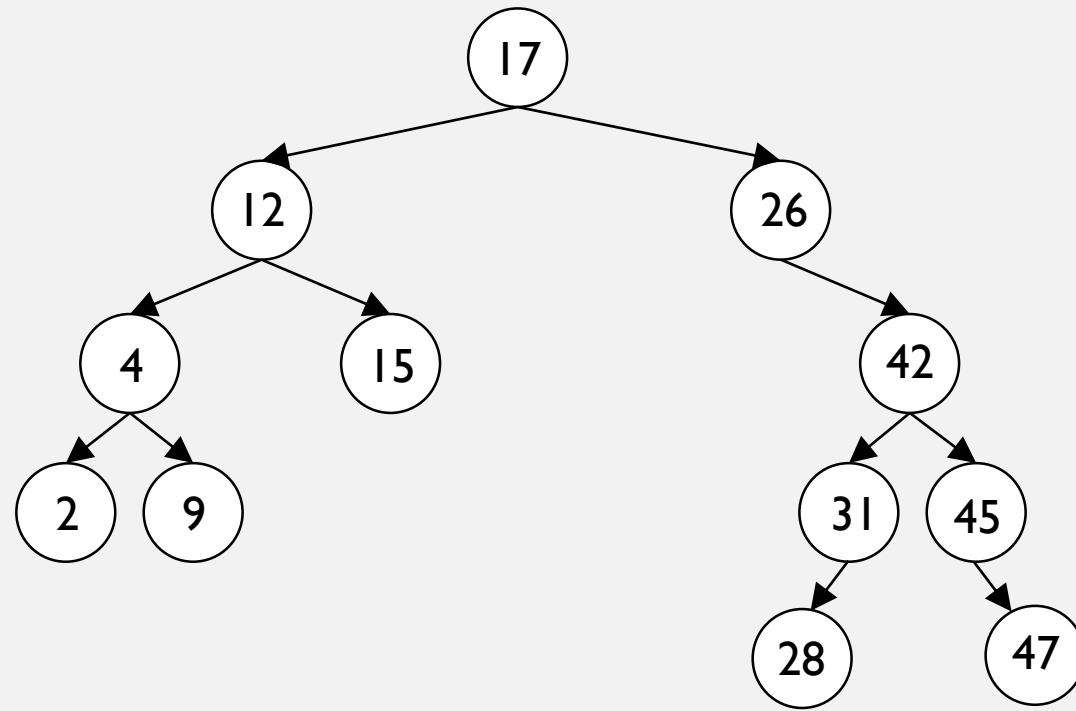


CASE III



left most
element
in right
subtree

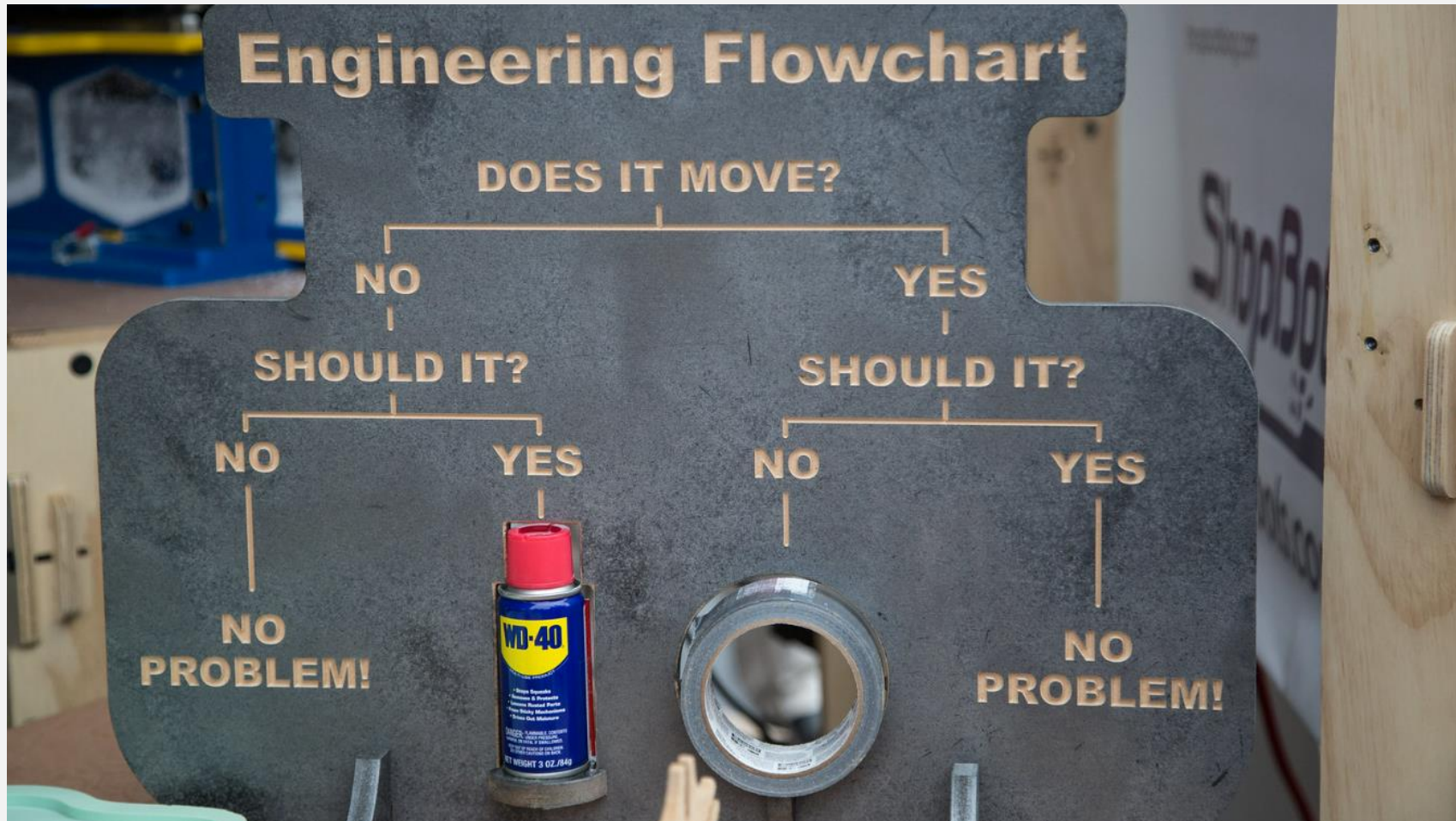
CASE III



SETS AND MAPS IMPLEMENTED AS BINARY SEARCH TREES

Adding
Deleting } all $O(\log n)$
Access }

APPLICATIONS OF TREES



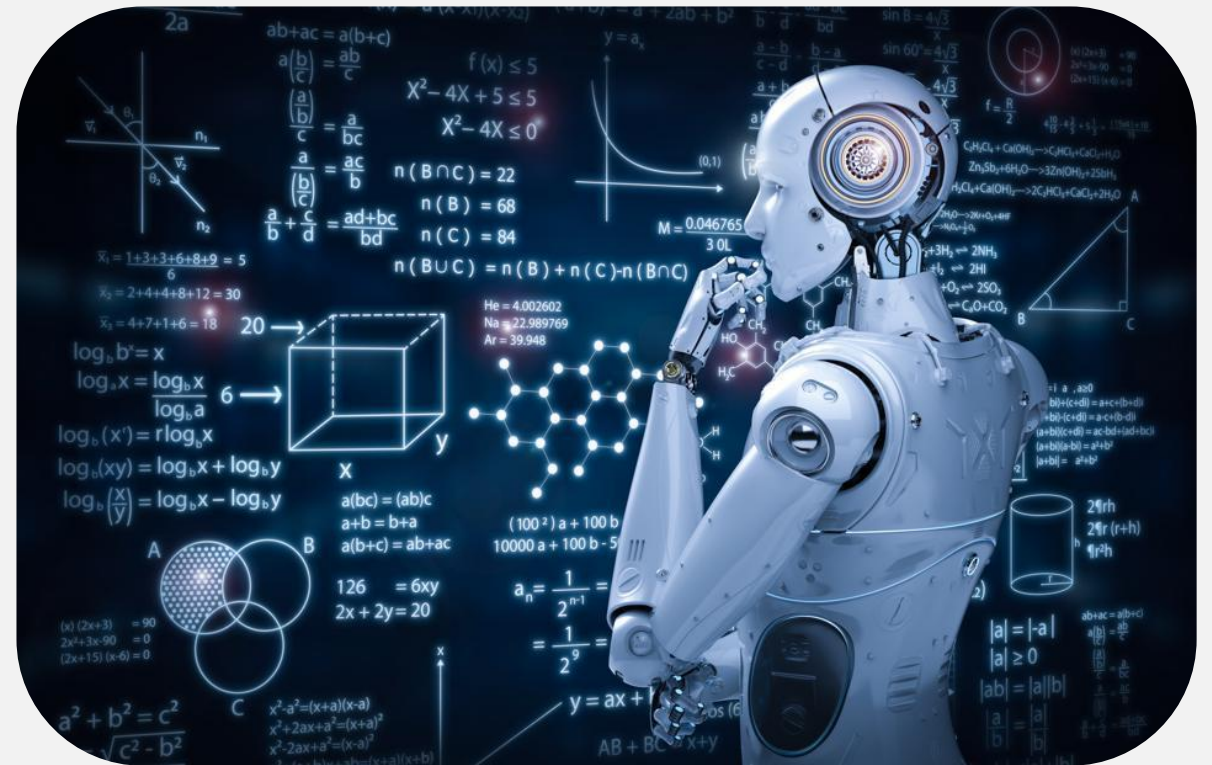
DECISION TREES

A machine learning model
→ Classify things

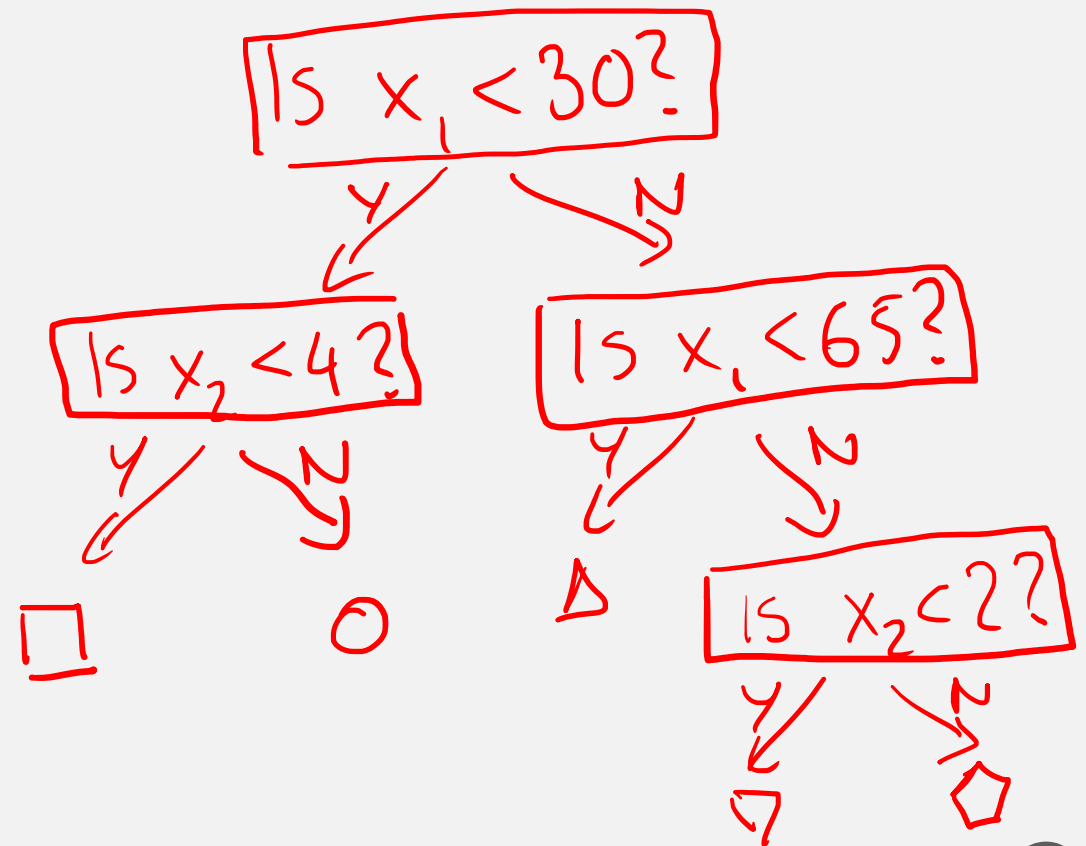
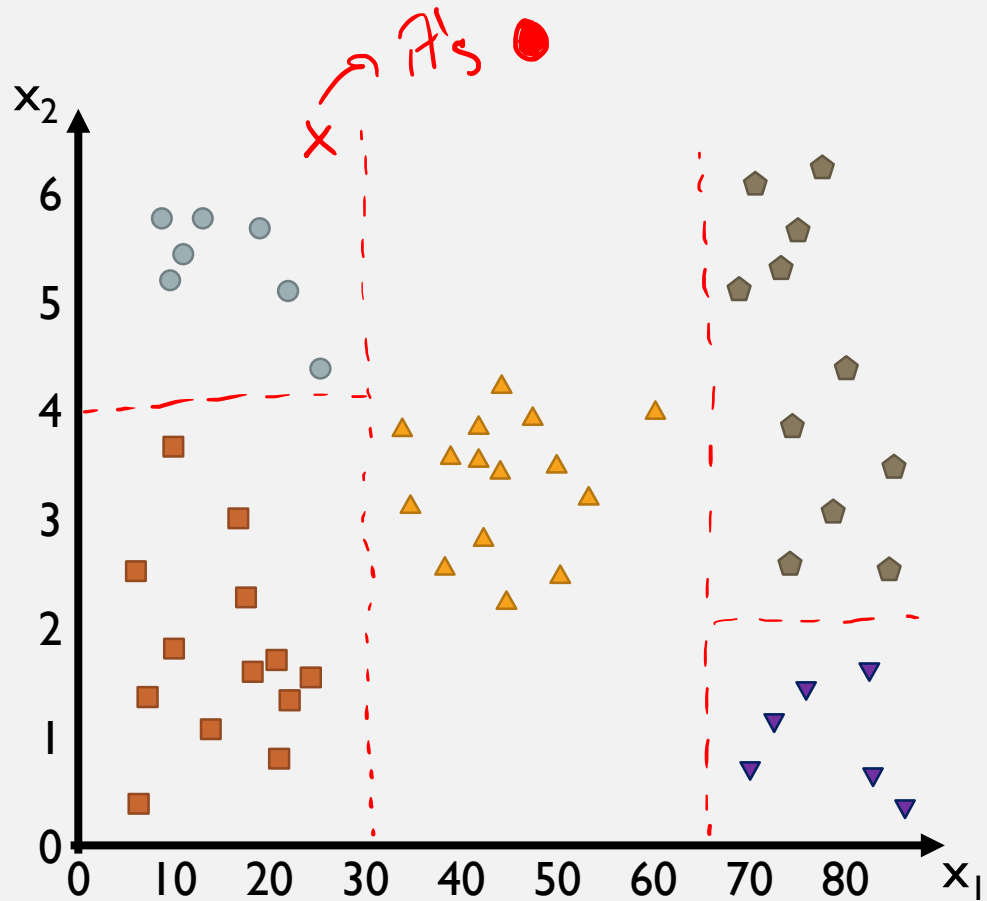
images of traffic signs
(self-driving cars)

mushroom edible?
(make an app)

recognizing diseases
(tumor malignant or benign)



DECISION TREES



HUFFMAN CODING



Efficient lossless
data compression

HUFFMAN CODING

AN_ANT_ATE_A_BANANA

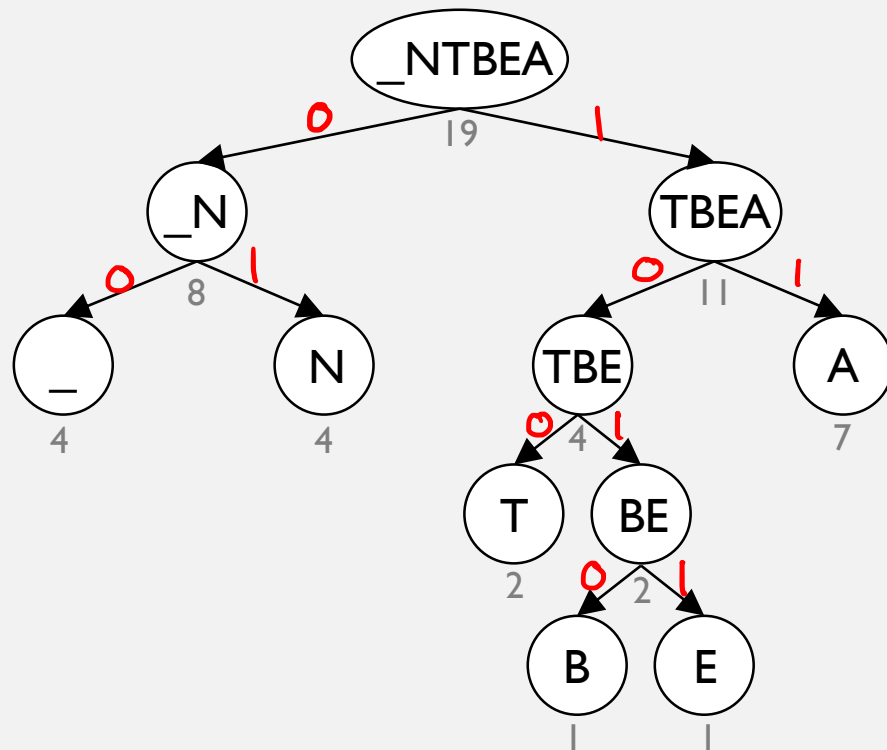
Symbol	Binary code
A	01000001
B	01000010
E	01000101
N	01001110
T	01010100
_	01011111

010000010100111001011111010000010100111001011111
010111110100000101010100010001010101111101000001
010111110100001001000001010011100100000101001110
01000001

152 bits

HUFFMAN CODING

AN_ANT_ATE_A_BANANA



Symbol	Huffman code
A	11
B	1010
E	1011
N	01
T	100
_	00

11010011011000011100101100110010101101110111

44 bits

EVERYTHING RUNS IN LOGARITHMIC
TIME ... IF ...

the tree is balanced.