

模板进阶

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Advanced Template

# 现代C++基础 Modern C++ Basics

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- **Supplementary**

- Template parameter that's not a type

- Type Deduction

- Friend in class template

- Laziness

- **Variadic Template**

- **SFINAE**

- **Common Techniques**

- CRTP

- Type Erasure

# Advanced Template

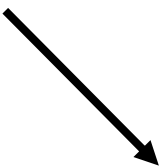
Supplementary Miscellaneous Knowledge

# Supplementary Miscellaneous Knowledge

- Supplementary
  - Template parameter that's not a type
  - Type Deduction
  - Friend in class template
  - Laziness

# Template parameter that's not a type

- First and foremost, very strangely...
  - Template parameter that's not a type  $\neq$  non-type template parameter (NTTP)
- There are three types of template parameters:
  - Type template parameter, i.e. `template<typename T>` or `template<class T>`;
  - Non-type template parameter, e.g. `template<int N>`;
  - Template template parameter.



Easier, so we talk about it first.

# Template template parameter

- First observe how a template is defined:

```
template<typename T, int N>  
class A
```

- Similar to function parameter, name can be neglected if it's not actually used:

```
template<typename, int>  
class A
```

- And that's how a template template parameter should be defined:

```
template<template<typename, int> class SomeTemplate>  
class B
```

Here we can also use `typename` since C++17.

- And we can use template in this form to fill in the parameter!
  - E.g. `B<A>`.

# Template template parameter

- A more complex example:

```
template<template<typename T, T *> class SomeTemplate>
class B
```

- Require a template with a type parameter and an NTTP, and the NTTP depends on the type parameter.
- Notice that this T cannot be used outside:

```
template<template<typename T, T *> class SomeTemplate, typename U = T>
class B
{
    T a;
};
```

- This is very like scope of variable; outer scope cannot use names in inner scope.

# Template template parameter

- Exercise: explain this template.

```
template<typename T, template<T> typename U>
class B
```

- It accepts a type parameter **T**, and a template template parameter **U**...
  - Where **U** has a NTTP, whose type is **T**.
- For example:

```
template<typename T, template<T> typename U>
class B
{
    U<T{}> a;
};
```

```
template<int Size>
class A
{
};
```

```
B<int, A> a;
```

Have **A<0>** **a**; as data member in **B<int, A>**.



# Template template parameter

- A practical example: write a stack class!
  - Users may use any container to store elements as long as it supports push/pop/..., e.g. `std::vector/deque`, `std::inplace_vector` in C++26.

```
template<typename T, template<typename> class Container = std::deque>
class Stack
{
    Container<T> cont_;

public:
    void Push(const T &elem);
};

template<typename T, template<typename> class Container>
void Stack<T, Container>::Push(const T &elem)
{
    cont_.push_back(elem);
}
```

We notice that `std::stack` uses two type parameters, i.e. `template<typename T, typename Cont = std::deque<T>>`.

# Template template parameter

- Example in standard library: `std::ranges::to`.

```
auto vec = r | std::ranges::to<std::vector>();
```

- Note 1: before C++17, template template parameter requires **exact match**.

- After C++17, default parameters are considered.

- E.g. `std::deque` actually has two parameters:
- ```
template<
    class T,
    class Allocator = std::allocator<T>
> class deque;
```

- You have to use:
- ```
template<typename T, template<typename, typename> class Container = std::deque>
class Stack
{
    Container<T, std::allocator<T>> cont_;
```

```
template<typename T, template<typename> class Container>
```

Cannot fill `std::deque`  
before C++17.

# Template template parameter

- Note 2: you can also use default parameter inside template template parameter.

```
template<typename T, template<typename = T, typename = std::allocator<T>>  
      |      |      |      |      |      class Container = std::deque>  
class Stack  
{  
    Container<> cont_;  
  
public:  
    void Push(const T &elem);  
};
```

<> Is needed even no parameter is manually filled.

# NTPP

- Previously, our NTPPs are always integers...
  - E.g. `std::array<T, N>`, `IsPrime<N>`.
- While there are more possible types, including:
  - Enumerations (special kind of integer);
  - Pointers (and `nullptr_t`), pointer to member;
  - Lvalue reference;
  - Floating points;
  - Some “simple” classes;
  - `constexpr` Lambda.

} Since C++20
- These types are referred as “structural types”.

# NTTP

- For example:

```
struct RGB
{
    int r, g, b;
};

template<decltype(&RGB::r) Channel>
void Test(RGB &color)
{
    color.*Channel = 0;
}
```

```
RGB color{ 1, 2, 3 };
Test<&RGB::g>(color);
std::cout << color.g;
```

```
void Func(int p)
{
    std::cout << p;
}

template<decltype(&Func) Pointer>
void Test2(int param)
{
    Func(param);
}
```

```
Test2<&Func>(1);
```

For a non-type *template-parameter* of reference or pointer type, or for each non-static data member of reference or pointer type in a non-type *template-parameter* of class type or subobject thereof, the reference or pointer value shall not refer or point to (respectively):

- a temporary object ([class.temporary]),
- a string literal object ([lex.string]),
- the result of a typeid expression ([expr.typeid]),
- a predefined \_\_func\_\_ variable ([dcl.fct.def.general]), or
- a subobject ([intro.object]) of one of the above.

There are actually many restrictions for passed pointer / reference, see [\[temp.arg.nontype\]](http://temp.arg.nontype).

# NTTP

- For pointer and reference, passed argument has restrictions.
  1. Its address should be determined in compile time, so only those with static storage duration are allowed.
  2. Linkage requirement: C++98 external, C++11 includes internal, C++17 includes no linkage (i.e. static variable in function).

```
template<int *Ptr>
void Test3()
{
}
template<int &Ref>
void Test4()
{
}
static int m = 0;
```

```
Test3<&m>();
Test4<m>();
```

Whether two templates are the same instantiation depends on the address.

# NTTP

```
MyClass<"hello"> x; //ERROR: string literal "hello" not allowed
```

- A special kind of “pointer” is string literals.
  - It's not allowed to use string literals to initialize **const char\***.

- Solution: 

```
extern char const s03[] = "hi"; // external linkage
char const s11[] = "hi"; // internal linkage
```

```
int main()
{
    Message<s03> m03; // OK (all versions)
    Message<s11> m11; // OK since C++11
    static char const s17[] = "hi"; // no linkage
    Message<s17> m17; // OK since C++17
}
```

- That's inconvenient since we have to introduce additional named variable...
- And normally our understanding of equivalent template should be “have the same string content”, instead of same address.
  - Since C++20 you can use class-type NTTP!

# NTPP

- For class NTPP, we first introduce literal types:

A literal type is any of the following:

Not class,  
ignored.

- possibly cv-qualified `void` (so that `constexpr` functions can return void); (since C++14)

- scalar type;
- reference type;
- an array of literal type;
- possibly cv-qualified class type that has all of the following properties:
  - has a `trivial(until C++20)constexpr(since C++20) destructor`,
  - all of its non-static non-variant data members and base classes are of non-volatile literal types, and
  - is one of

- a lambda type, (since C++17)

- an aggregate union type that
  - has no variant members, or
  - has at least one variant member of non-volatile literal type,
- a non-union aggregate type, and each of its anonymous union members
  - has no variant members, or
  - has at least one variant member of non-volatile literal type,
- a type with at least one `constexpr` (possibly template) constructor that is not a copy or move constructor.



# NTTP

For a non-type *template-parameter* of reference or pointer type, or for each non-static data member of reference or pointer type in a non-type *template-parameter* of class type or subobject thereof, the reference or pointer value shall not refer or point to (respectively):

- a temporary object ([class.temporary]),
- a string literal object ([lex.string]),
- the result of a typeid expression ([expr.typeid]),
- a predefined \_\_func\_\_ variable ([dcl.fct.def.general]), or
- a subobject ([intro.object]) of one of the above.

- And class NTTP just requires:
  - Be a literal type;
  - All base classes and non-static data members are public, non-mutable and structural types (or array of structural types).
  - Particularly, for pointer and reference member, it has the same restrictions as NTTP.
- Finally, template class is also allowed to write at NTTP, and the concrete type will be deduced automatically.

```
template<std::array arr>
void f();

f<std::array<double, 8>{}>();
```

- So we can easily write a **FixedString** class...

# NTPP

Cannot be `const char* str` as:

1. It's non-owning, cannot do complex operations like concatenation;
2. It's forbidden to point to string literal as the last page shows.

```
template<std::size_t N>
struct FixedString
{
    char str[N];
    constexpr FixedString(const char (&input)[N])
    {
        for (std::size_t i = 0; i < N; i++)
            str[i] = input[i];
    }
};

template<FixedString InputStr>
void Test() {}

int main()
{
    Test<"123">{};
}
```

We'll provide an exercise in homework to write a more reasonable class.

# NTTP

- Floating points can be seen as instantiated by “underlying binary representation”.
  - E.g. if the platform uses IEEE 754, `Test<1.0f>` can be hypothetically viewed as `Test<3f800000>`.
  - But all `NaN` will be seen as equivalent.
- And we notice that floating point rounding error is still important.

```
template<double Val>
class MyClass {
};

int main()
{
    std::cout << std::boolalpha;
    std::cout << std::is_same_v<MyClass<42.0>, MyClass<17.7>>           // always false
               << '\n';
    std::cout << std::is_same_v<MyClass<42.0>, MyClass<126.0 / 3>>       // true or false
               << '\n';
    std::cout << std::is_same_v<MyClass<42.7>, MyClass<128.1/ 3>>       // true or false
               << "\n\n";

    std::cout << std::is_same_v<MyClass<0.1 + 0.3 + 0.00001>,
                               MyClass<0.3 + 0.1 + 0.00001>>           // true or false
               << '\n';

    std::cout << std::is_same_v<MyClass<0.1 + 0.3 + 0.00001>,
                               MyClass<0.00001 + 0.3 + 0.1>>           // true or false
               << "\n\n";

    constexpr double NaN = std::numeric_limits<double>::quiet_NaN();
    std::cout << std::is_same_v<MyClass<NaN>, MyClass<NaN>>               // always true
               << '\n';
}
```

# NTPP

- Since C++17, it's also allowed to accept NTPP of any type.

```
template<auto Param> // or auto& / decltype(auto)  
class A
```

```
A<1> a;  
A<nullptr> b;
```

- Before that, you need to add an additional type parameter for assistance.
  - But you have to specify that type manually, very inconvenient...

```
template<typename T, T Param>  
class A
```

```
A<int, 1> a;  
A<std::nullptr_t, nullptr> b;
```

- It's then easy to accept lambda as NTPP:

# NTTP

`std::invocable<F, Args...>` requires `F` to be callable with arguments `Args...`; here it means `Callable` should use 0 parameter.

```
template<std::invocable auto Callable>
class A
{
public:
    constexpr decltype(auto) operator()() { return Callable(); }
};

int main()
{
    A<[]() { return 0; }> a;
    A<[]() { std::cout << "Hello, world"; }> b;
}
```

Call `Callable`.

Lambda should be `constexpr` (thus no capture is allowed), but this is implicitly added since C++17 as long as all operations are allowed in `constexpr` function.

- Notice that each lambda has its unique type even if they have same closure body, so the instantiation is unique theoretically.

# Supplementary Miscellaneous Knowledge

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  - Type Deduction
  - Friend in class template
  - Laziness

# Type Deduction

- We don't need to fill all template parameters due to type deduction...
- Each function parameter will deduce template parameter **independently**.
  - Assuming template parameter is P and passed argument type is A...
  - For each deduction, there are generally two rules:
    - For non-reference parameter P, **decayed** A is deduced.
      - P will ignore its top-level cv-qualifier.
    - For reference parameter (e.g. P&), the original A is deduced.
    - \***auto** has basically the same rule.
  - And finally, if deduction leads to conflict types, fail to match.

# Type Deduction

- For example:

Top-level cv is ignored, equiv. to use 1 to deduce T.

```
template<typename T>
void Func3(const T)
{
}
```

```
Func3(1);
```

```
template<typename T>
void Func(T a, T b);

template<typename T>
void Func2(T &a, T &b);

int main()
{
    const int a = 1, b = 1;
    Func(a, b); // a & b: T -> int
    Func2(a, b); // a & b: T -> const int

    int c[8], d[10];
    Func(c, d); // a & b: T -> int*
    Func2(c, d); // a: T -> int[8]; b: int[10]
}
```

Not decayed,  
so conflict.



# Type Deduction

- In most cases, **conversion is forbidden** in deduction.

```
template<typename T>
void Func(T a, T b);

class A
{
public:
    operator int() { return 0; }
};

int main()
{
    Func(1, A{});
}
```

Deduction failure due to conflict, rather than  $T = \text{int} + A \rightarrow \text{int}$  conversion.

Due to separate deduction, **this is inevitable.**

Deduction failure instead of  $T = \text{int} + B<\text{int}> \rightarrow A<\text{int}>$  conversion.

This may be evitable?

```
template<typename T>
class A
{
};

template<typename T>
class B
{
public:
    operator A<T>() { return {}; }
};

template<typename T>
void Func(A<T> a);

int main()
{
    Func(B<int>());
}
```

# Type Deduction

- There are three special cases to allow conversion:
  1. If P is reference, the deduced A can be more cv-qualified than A.
  2. If P is pointer, the deduced pointer can have qualification conversion.

```
template<typename T> void Func(const T &);  
template<typename T> void Func2(const T *);  
  
int main()  
{  
    int a = 1;  
    Func(a); // conversion: int -> const int, then T = int  
    Func2(&a); // conversion: int* -> const int*, then T = int  
}
```

1 and 2 are different,  
since more cv-qualified  
**int\*** is **int\* const**, not  
**const int\***.

3. If P is base class (pointer), A is derived class (pointer), derived-to-base conversion is allowed.

```
template<typename T> class A {};  
  
template<typename T> class B : public A<T> {};  
  
template<typename T> void Func(A<T> a);  
  
int main()  
{  
    Func(B<int>{}); // conversion: B<int> -> A<int>  
}
```

# Type Deduction

- Sometimes deduction needs further match...

```
template<typename T, typename U>
void Func(T (*funcPtr)(U &))
{
}

void Test(int &);

int main()
{
    Func(Test);
}
```

If we regard it as a whole, it's not reference type (i.e. `V funcPtr`).

Then we know `V = void (*)(int&)`, the decayed function type.

And we match `T (*)(U&)`, getting `T = void` and `U = int`.

- Such match is exact, no conversion is allowed.

```
void Test(int);
```



# Type Deduction

- Lots of pattern can be used to match, not described in detail.
  - Sometimes recursive match is needed.

- `cV(optional) T;`
- `T*;`
- `T&;`

- `T&&;` (since C++11)

- `T(optional) [I(optional)];`

- `T(optional) (U(optional));` (until C++17)

- `T(optional) (U(optional)) noexcept(I(optional));` (since C++17)

- `T(optional) U(optional)::*;`
- `TT(optional)<T>;`
- `TT(optional)<I>;`
- `TT(optional)<TU>;`
- `TT(optional)<>;`

In the above forms,

- `T(optional)` or `U(optional)` represents a type or *parameter-type-list* that either satisfies these rules recursively, is a non-deduced context in P or A, or is the same non-dependent type in P and A.
- `TT(optional)` or `TU(optional)` represents either a class template or a template template parameter.
- `I(optional)` represents an expression that either is an I, is value-dependent in P or A, or has the same constant value in P and A.

- `noexcept(I(optional))` represents an **exception specification** in which the possibly-implicit `noexcept` specifier's operand satisfies the rules for an `I(optional)` above. (since C++17)

```
template<typename T>
void f1(T*);
```

```
template<typename E, int N>
void f2(E(&)[N]);
```

```
template<typename T1, typename T2, typename T3>
void f3(T1 (T2::*)(T3*));
```

```
class S {
public:
    void f(double*);
};
```

```
void g (int*** ppp)
{
    bool b[42];
    f1(ppp);      // deduces T to be int**
    f2(b);        // deduces E to be bool and N to be 42
    f3(&S::f);    // deduces T1 = void, T2 = S, and T3 = double
}
```

# Type Deduction

- There also exists non-deduced context ([\[temp.deduct.type\]](#)), i.e. where parameters cannot be deduced.
  - `std::initializer_list` cannot be deduced, as we stated before.
    - `auto` in definition can deduce it, e.g. `auto a = {1,2,3}`.
  - The first/major array bound, if P is not reference / pointer type.
    - Reference deduction won't decay; pointer has already decayed the first bound.

```
template<int i>
void f1(int a[10][i]);

template<int i>
void f2(int a[i][20]);    // P = int[i][20], array type

template<int i>
void f3(int (&a)[i][20]); // P = int(&)[i][20], reference to array

void g()
{
    int a[10][20];
    f1(a);    // OK: deduced i = 20
    f1<20>(a); // OK
    f2(a);    // error: i is non-deduced context
    f2<10>(a); // OK
    f3(a);    // OK: deduced i = 10
    f3<10>(a); // OK
}
```

# Type Deduction

- NTTP cannot be deduced from expression:

- Cannot deduce  $N = 5$  here.

- Qualified type names cannot be used for deduction:

- Cannot deduce  $T = \text{int}$  here.

- Default parameter cannot be used for deduction:

- Though `Func()` should be equiv. to `Func(1)`,  $T$  cannot be deduced as `int`.

- `decltype` type.

- Cannot deduce  $T = \text{int}$  here.

```
template<typename T> T value = {};  
  
template<typename T> void Func(decltype(value<T>));  
  
int main()  
{  
    Func(1);  
}
```

```
template<std::size_t N>  
void f(std::array<int, 2 * N> a);  
  
std::array<int, 10> a;  
f(a); // P = std::array<int, 2 * N>
```

```
template<typename T>  
void Func(T = 1);  
  
int main()  
{  
    Func();  
}
```

```
template<typename T>  
class A  
{  
public:  
    using Type = T;  
};  
  
template<typename T>  
void Func(typename A<T>::Type);  
  
int main()  
{  
    Func(1);  
}
```

# Type Deduction

- There are some other deduction contexts, not cover it in detail.
  - Address of an overload set, where return type will also be used for deduction.
    - If there are multiple best match after overload resolution, ambiguous.
  - Conversion function template.
  - ...
- Since C++17, class template argument deduction (CTAD) is introduced.
  - It deduces class template argument **from constructor**, so rules are similar to deducing from function call.

# CTAD

- Just transform template parameter of class to ctors!

```
template<typename T>
struct UniquePtr
{
    UniquePtr(T *t);
};
```



```
struct X
{
    template<typename T>
    X(T *t);
};
```

- And then use rules before...
- Wait, where is copy/move ctor?
  - As deduction will strip reference, they're combined together as a single hypothetical function:

```
template<typename T> X(UniquePtr<T>);
```

- Besides, C++ allows **user-defined deduction guides** in CTAD.



# CTAD

- For example:

```
std::list l{ 1, 2, 3 };
std::vector v{ l.begin(), l.end() };
```

  - CTAD of `std::list` is quite simple; it comes from `list(std::initializer_list<T>)`.
  - But vector from iterator is `template<typename Iter> vector(Iter, Iter)`.
    - There is no T at all...How is class template parameter deduced?
- User-defined deduction guide is very similar to ctor, just add trailing return type to specify “what should be deduced”.

```
template<class InputIt>
vector(InputIt, InputIt)
    -> vector<typename std::iterator_traits<InputIt>::value_type>;
```

Use `std::iter_value_t<InputIt>` since C++20.

---

*explicit(optional) template-name ( parameter-list ) requires-clause(optional) -> simple-template-id ;* (1)

---

**template** <template-parameter-list> requires-clause(optional)  
*explicit(optional) template-name ( parameter-list ) requires-clause(optional) -> simple-template-id ;* (2)

# CTAD

- **explicit**: if written, disable copy initialization deduction.

```
template<class InputIt>
explicit A(InputIt, InputIt)
    -> A<typename std::iterator_traits<InputIt>::value_type>;
```

```
A v{ l.begin(), l.end() };
// A v = { l.begin(), l.end() }; // deduction fails
```

- requires clause: specify concept constraints for when deduction happens.

```
template<class InputIt>
requires std::random_access_iterator<InputIt>
explicit A(InputIt, InputIt)
    -> A<typename std::iterator_traits<InputIt>::value_type>;
```

```
// A v{ l.begin(), l.end() }; // deduction fails
```

# CTAD

- Note 1: CTAD cannot do partial deduction. All or nothing!

```
template<typename T1, typename T2> class A{ public: A(T1, T2); };  
A a{1,2}; // Okay;  
A<int> a{1,2}; // Nope, though 2 can be used to deduce int.
```

- Note 2: CTAD doesn't have to be same as some ctor; deduction and overload resolution are separate.

```
template<class InputIt>  
A(InputIt) -> A<typename std::iterator_traits<InputIt>::value_type>;
```

```
A v{ l.begin() }; // deduction succeeds, but A<int> doesn't have  
    // proper ctor (no viable candidate).
```

# CTAD

- Note 3: unlike method definition, the class name **cannot** use qualified name in deduction guide.
  - Assuming **A** is in namespace **Test**.

```
template<class InputIt>
Test::A(InputIt)->Test::A<typename std::iterator_traits<InputIt>::value_type>;
```

- So, usual practice is to write deduction guide immediately after class definition.

```
namespace Test
{
template<typename T>
class A
{
public:
    template<class InputIt>
    A(InputIt a, InputIt b)
    {
    }
};

template<class InputIt>
A(InputIt) -> A<typename std::iterator_traits<InputIt>::value_type>;

} // namespace Test
```

# CTAD

- Note 4: in class context, injected template name is preferred over CTAD.
  - Injected class name means that **A** refers to current instantiation, no need to write **A<T>** explicitly.
  - But injected class name is disabled when using qualified name, then CTAD happens.

```
A &operator=(const A &another)
{
    ::A a{ another }; // CTAD
    swap(a);
    return *this;
}
```

```
template<typename T>
class A
{
public:
    A &operator=(const A &another)
    {
        A a{ another };
        swap(a);
        return *this;
    }
}
```

- In this example CTAD or not doesn't differ...

# CTAD

- A better example:

Otherwise it's `X<T>(b, e)`, but we actually want `X<U>(b, e)`.

```
template<class T>
struct X
{
    X(T) {}

    template<class Iter>
    X(Iter b, Iter e) {}

    template<class Iter>
    auto foo(Iter b, Iter e)
    {
        return X(b, e); // no deduction: X is the current X<T>
    }

    template<class Iter>
    auto bar(Iter b, Iter e)
    {
        return X<typename Iter::value_type>(b, e); // must specify what we want
    }

    auto baz()
    {
        return ::X(0); // not the injected-class-name; deduced to be X<int>
    }
};
```

# CTAD

- Note 5: in deduction guide, **T&&** is still universal reference, not rvalue reference.
  - Here due to reference collapsing, ctor is then transformed to **A(int&)**.
  - \*So we say **A(T&&)** isn't universal reference because it cannot deduce lvalue reference without deduction guide.
- To prevent that astonishment, normally deduction guides use value type.
  - Here **A<int>** rather than **A<int&>** is deduced, compilation error as expected..




```
template<typename T> A(T) -> A<T>:
no instance of constructor "A<T>::A [with T=int]" matches the argument list
int ma
{
  in  main.cpp(69, 8): argument types are: (int)
    View Problem (Alt+F8) Quick Fix... (Ctrl+.)
    A b{ a }; // A<int> is deduced.
```

```
template<typename T>
class A
{
public:
    A(T &&) {} // rvalue reference
};

template<typename T>
A(T &&) -> A<T>; // universal reference

int main()
{
    int a = 1;
    A b{ a }; // A<int&> is deduced.
```

# CTAD

class template argument deduction for alias templates (FTM)*	P1814R0 	10	19	19.27*
class template argument deduction for aggregates (FTM)*	P1816R0  P2082R1 	10* 11*	17	19.27*

- Note 6: only the class itself can be deduced; adding reference / pointer would lead to failure.

- In most cases just use **auto**.

- Note 7: implicit deduction guide for aggregate is added since C++20.

- Example: user-defined guide must be added in C++17.

```
A b{ a }; //
A &&c{ a };
A &d = b;
A *p = &b;
```

- Note 8: C++20 also introduces deduction for alias template.

```
template<typename T> struct A { T val; };
template<typename T> A(T) -> A<T>;
// 如果没有deduction guide, 则A a{1};不正确
```

- To put it simply, for every deduction guide, use alias to deduce parameters by trailing type as much as possible. Remove deduced ones, add alias constraints and use only non-deducible ones to form new guides.



# Alias Template Deduction\*

- This is very complex so **optional**.

- We use A as example here; #1 and #2 form two deduction guides.
- From #1, use `C<V*, V*>` to deduce `C<T, U>`, forming

```
template<class T, class U, class V>  
C(V *, V *) -> C<V *, V *>;
```

- From #2, use `C<V*, V*>` to deduce `C<T, std::identity_t<U>>`, forming

```
template<class T, class U, class V>  
C(V *, U) -> C<V *, std::identity_t<U>>;
```

- U is not deducible, as it's an alias of qualified type (non-deduced context).
- Strip useless parameters and add constraints (here A doesn't have constraints; if it's B, then add `std::integral` to W).

```
template <class T, class U> struct C {  
    C(T, U);  
};  
// #1  
  
template<class T, class U>  
    C(T, U) -> C<T, std::identity_t<U>>;  
// #2  
  
template<class V> using A = C<V *, V *>;  
template<std::integral W> using B = A<W>;
```

# Alias Template Deduction\*

- Now we get two guides for alias:

```
template<class V>
C(V *, V *) -> C<V *, V *>;

template<class U, class V>
C(V *, U) -> C<V *, std::type_identity_t<U>>;
```


- Finally, when it's deduced, we need to check again whether it satisfies alias.

```
int i{};
double d{};
A a1(&i, &i);    // deduces A<int> from #1
A a2(i, i);      // error: cannot deduce V * from i both #1 and #2 have V*
A a3(&i, &d);    // error: #1: cannot deduce (V*, V*) from (int *, double *)
                  // #2: cannot deduce A<V> from C<int *, double *> #2 deduces C<int*, double*>,
                  // but it doesn't match A<V>.
```

- See more practice in [C++ standard](#) and [cppreference](#).  
• gcc has a bug on this feature, see [c++ - How to write deduction guidelines for aliases of aggregate templates? - Stack Overflow](#) (and you can also do practice by this analysis).

# CTAD

- Note 9: since C++23, CTAD can happen through inherited ctor.
  - For template parameters of derived class, if they're part of parameters of base class, they'll be deduced by guides of base class.
  - Let's see examples in the standard directly:

```
template <typename T> struct B {  
    B(T);  
};  
template <typename T> struct C : public B<T> {  
    using B<T>::B;  Inherited ctor.  
};  
template <typename T> struct D : public B<T> {};  
  
C c(42);           // OK, deduces C<int>  
D d(42);           // error: deduction failed, no inherited deduction guides  
B(int) -> B<char>;  
C c2(42);          // OK, deduces C<char>
```

```
template <typename T> struct E : public B<int> {  
    using B<int>::B;  
};  
  
E e(42);           // error: deduction failed, arguments of E cannot be deduced from introduced guides  
  
template <typename T, typename U, typename V> struct F {  
    F(T, U, V);  
};  
template <typename T, typename U> struct G : F<U, T, int> {  
    using G::F::F;  
}  
  
G g(true, 'a', 1); // OK, deduces G<char, bool>
```

**T** isn't part of base class.

**T** and **U** are part of base class.

# Supplementary Miscellaneous Knowledge

- Supplementary
  - Template parameter that's not a type
  - Type Deduction
  - Friend in class template
  - Laziness

# Friend in class template

- To define a friend method for a class:
  - We can also define friend function in the class, and ADL will help us to find it.
    - Like defining `operator<<` as friend.

```
class A
{
    friend void Func(A &a);
};

void Func(A &) {}
```

- For a class template, splitting definition is harder...
  - Is it correct?

```
template<typename T>
class A
{
    friend void Func(A &a);
};

template<typename T>
void Func(A<T> &) {}
```

No!

Friend of `A<T>` is a normal function `func(A<T>)`, not a template function.

\*But unlike normal functions, their definition is instantiated only when used.

# Friend in class template

- If we call & compile:

```
A<int> a;  
Func(a);
```

```
error: main.cpp.obj : error LNK2019: 无法解析的外部符号 "v  
oid __cdecl Func(class A<int> &)" (?Func@@YAXAEAV?$A@H@@@Z  
) , 函数 main 中引用了该符号  
build\windows\x64\release\main.exe : fatal error LNK1120:  
1 个无法解析的外部命令
```

- For `A<int>`, the friend method is `Func(A<int>)`, so we may try:
  - Succeed to compile!
  - But we can never define `Func` for **any** `T`.
- Solution 1: make template method as friend instead of normal method.

```
template<typename T>  
class A  
{  
    int member;  
  
    template<typename U> friend void Func(A<U> &a);  
};  
  
template<typename U>  
void Func(A<U> &a)
```

```
template<typename T>  
class A  
{  
    int member;  
    friend void Func(A &a);  
};  
  
void Func(A<int> &a)  
{  
    a.member;  
}
```

# Friend in class template\*

- Limitation: not exactly same as our previous intention. **Optional** as the difference is subtle.
  - Since all instantiation regard template Func as friend, instead of its own normal Func.

```
template<typename T>
class A
{
    int member;

    template<typename U, typename V>
    friend void Func(A<U> &, A<V> &);

    template<typename U>
    friend void Func2(A<T> &, A<U> &);
};

template<typename U, typename V>
void Func(A<U> &a, A<V> &b)
{
    a.member;
    b.member;
}
```

Both **A<int>**  
and **A<float>**  
regard it as  
friend, so okay.

```
template<typename U>
void Func2(A<int> &a, A<U> &b)
{
    a.member;
    b.member;
}

int main()
{
    A<int> a;
    A<float> b;
    Func(a, b); Okay
    Func2(a, b); Error: cannot
                  access b.member.
}
```

**A<int>** friend:  
**void Func2(A<int>&, A<U>&);**

**A<float>** friend:  
**void Func2(A<float>&, A<U>&);**

Different friends, so **A<int>**  
friend cannot access **A<float>**  
member.

# Friend in class template

- Solution 2: specify template instantiation as friend.
  - Limitation: hard to code; we need two additional forward declarations.

```
template<typename T> class A; // 声明A
template<typename T> void Func(A<T>&); // 声明Func
template<typename T> // 定义A
class A
{
    friend void Func<T>(A<T>& a); // 声明友元函数是一个模板实例
};
template<typename T> void Func(A<T>&){} // 定义Func
```

We notice that no **template** keyword is needed here, unlike explicit instantiation.

- Solution 3: define friend methods inside the class directly, and use ADL to call it.
  - Limitation: ADL is quite obscure.



# Friend in class template

- For friend class:

```
template<typename T>
class B
{
};

template<typename T>
class A
{
    template<typename U> friend class B; // All instantiations of B are friend.
    friend class B<T>; // Only one instantiation, B<T>, is friend.
    friend T; // Template parameter as friend, never add "class" keyword.
```

- For **friend T**, if **T** isn't a class, it will be ignored.

# Supplementary Miscellaneous Knowledge

- Supplementary
  - Template parameter that's not a type
  - Type Deduction
  - Friend in class template
  - Laziness

# Lazy Instantiation

- When a class template is instantiated, not all of its members are immediately fully instantiated.
  - Some of them will only be fully instantiated **when they're actually used**.

<sup>3</sup> The implicit instantiation of a class template specialization causes

- (3.1) — the implicit instantiation of the declarations, but not of the definitions, of the non-deleted class member functions, member classes, scoped member enumerations, static data members, member templates, and friends; and
- (3.2) — the implicit instantiation of the definitions of deleted member functions, unscoped member enumerations, and member anonymous unions.

The implicit instantiation of a class template specialization does not cause the implicit instantiation of default arguments or *noexcept-specifiers* of the class member functions.

C++ standard, section [\[temp.inst\]](#).

# Lazy Instantiation

- Example:

```
template<int N> // This should be std::size_t at best; we just want to
                // illustrate lazy instantiation, so here we need N <= 0
                // to make Array<N> ill-formed.
class Array
{
public:
    int arr[N];
};

template<typename T, int N>
class A
{
public:
    // Successful ones
    void Test() { Array<N> arr; }
    struct Test2
    {
        Array<N> arr;
    };

    // Failed ones
    void Test3(Array<N> arr) {}
    union {
        Array<N> arr;
        int m;
    };
};
```

E.g. we define  
`A<int, -1> a;`

As long as we don't call / use them  
to force their full instantiation.

Declaration is not valid.  
Only definition can be  
lazy instantiated.

# Lazy Instantiation

- Similarly, there exists “laziness” in some other cases.
  - Value template initialization: 

```
template<typename T>  
T v = T::default_value();
```

    - When value of `v` is not used, `default_value()` is allowed to not exist in `T`.
      - For example, `decltype(v<int>)` will compile successfully.
  - Default value: 

```
void Func(T a = T{}) {}
```

    - When default value isn't used, this assignment is allowed to make no sense.
    - E.g. when `Func` is called without filling default value (i.e. only calling `Func(xxx)`, never calling `Func()`), `T` is allowed to be not default constructible.
  - Pointer definition: 

```
void Test4(Array<N> *arr) {}
```

    - Definition of `Array` isn't checked so `N = -1` won't make compilation fail.
  - And concept (check our homework in the last lecture!).

5. 补充一个知识，模板类成员函数可以对其模板参数施加约束，当约束不满足时不会导致编译错误，只会去掉这个成员函数。判断以下代码中哪些语句不能够编译通过：

# Lazy Instantiation

- Some laziness depends on compiler.
  - Virtual function: some compilers always instantiate virtual function to construct full vtable.
  - Semantic error: instantiation check is after syntactic and semantic check theoretically, while some compilers will drop semantic check if it's not instantiated.

## Phase 7: Compiling

Compilation takes place: each preprocessing token is converted to a **token**. The tokens are syntactically and semantically analyzed and translated as a **translation unit**.

## Phase 8: Instantiating templates

Each translation unit is examined to produce a list of required template instantiations, including the ones requested by **explicit instantiations**. The definitions of the templates are located, and the required instantiations are performed to produce *instantiation units*.

```
void error() { Array<-1> boom; }
```

Defined in **class A**;  
Not all compilers detect this  
as error if it's not called.

# Advanced Template

Variadic Template

# Variadic Template

- Variadic Template
  - Basics (Pack expansion & ...)
  - Fold expression



# Variadic Template

- Many functions can accept unbounded number of parameters...
  - `std::vector::emplace_back`, `std::format`, `std::invoke`, etc..
- This is enabled by variadic template!
  - Let's use printing a bunch of parameters as example.

```
#include <iostream>
```

```
void print ()  
{  
}
```

```
template<typename T, typename... Types>  
void print (T firstArg, Types... args)  
{  
    std::cout << firstArg << '\n';  
    print(args...);  
}
```

Called *pack expansion*,  
meaning unpack `args`  
to `arg1`, `arg2`, ....

```
std::string s("world");  
print (7.5, "hello", s);
```

`firstArg` is 7.5  
and `T` is `double`.

`args` are "hello", `s`  
and `Types` are `const char*`, `std::string`.

Called *template parameter pack*,  
meaning "any number of template parameter".

Called *function parameter pack*,  
meaning "any number of parameter".

// call print() for remaining arguments

`print(7.5, "hello", s) -> print("hello", s) -> print(s) -> print()`

# Variadic Template

- Every function call passes value type, which incurs overhead.
  - They're read-only, make them `const&!`

```
template<typename T, typename... Args>  
void print(T firstArg, const Args &...args)
```

- And we may want to eliminate empty function...
  - That is, when number of args is 0, do nothing.
  - You can use operator `sizeof...(pack)`.

```
template<typename T, typename... Args>  
void print(T firstArg, const Args &...args)  
{  
    std::cout << firstArg << "\n";  
    if constexpr (sizeof...(args) != 0)  
    {  
        print(args...);  
    }  
}
```

Here you can also use `sizeof...(Args)`; any pack is ok.

# Variadic Template

- Exercise 1: write declaration of `std::invoke` (use universal ref, no need to add concept).

```
template<typename F, typename... Args>  
decltype(auto) Invoke(F func, Args &&...args)
```

- And we need to forward `args` to `func`.
  - Assuming that `func` is called by `operator()`. `return func(std::forward<Args>(args)...);`
- So generally, pack expansion is:
  - Write a pattern, as if operating on a single normal parameter.
    - E.g. here the pattern is `std::forward<Args>(args)`.
  - And add `...` (sometimes need additional space before `...`), meaning that applying pattern on every element and concatenating with `,`.
    - So it's like `std::forward<Arg1>(arg1), std::forward<Arg2>(arg2)...`
    - Notice that it's semantic substitution, **not a comma expression**.

# Variadic Template

- Exercise 2: write a function that accepts a container and many indices, and passing corresponding elements to `print`.

```
template<typename T, typename... IdxTypes>
void print_elems(const T &container, IdxTypes... idx)
{
    print(container[idx]...);
}
```

- Exercise 3: what's the type of `result`?

```
template<typename... T1>
class A
{
public:
    template<typename... T2>
    static auto func()
    {
        return std::tuple<std::pair<T1, T2>...>();
    }
};
```

```
auto result = A<char, int>::func<float, double>();
```

Answer: `std::tuple<  
 std::pair<char, float>,  
 std::pair<int, double>  
>`

# Variadic Template

- [Hard] Exercise 4: explain what **g** accepts.

```
template<typename... OuterTypes>
class Nested {
    template<typename... InnerTypes>
    void f(InnerTypes const&... innerValues) {
        g(OuterTypes(InnerTypes(innerValues)...)...);
    }
};
```

Assuming

`sizeof...(OuterTypes) == N,`  
`sizeof...(InnerTypes) == M.`

- `InnerTypes(innerValues)...` is expanded to `I1{iv1}, I2{iv2}, ...`.
  - And these M parameters are used to construct some type `OuterTypes`, forming a new pattern.
- Expand this pattern, we get `O1{ I1{iv1}, I2{iv2}, ... }, O2{ I1{iv1}, I2{iv2}, ... }, ...`.
  - And these N parameters are passed to **g**.

# Variadic Template

- Exercise 5: write a class that **public** inherits all of its template parameters.

```
template<typename... T1>  
class A : public T1...
```

- Note: since C++17, it's allowed to inherit all methods by expansion:

```
template<typename... T1>  
class A : public T1...  
{  
public:  
    using T1::T1...; // inherit all ctor  
};
```

# Variadic Template

- Exercise 6: write a ctor to accept params to initialize each base class.

```
template<typename... T1>
class A : public T1...
{
public:
    // 注意不能用T1&&, 因为不是universal reference.
    template<typename... T2>
    A(T2 &&...args) : T1{ std::forward<T2>(args) }...
    {
    }
};
```

# Variadic Template

- Exercise 7: explain code below.

```
template<typename... Funcs>
struct Overloaded : Funcs...
{
public:
    using Funcs::operator()...;
};

int main()
{
    auto overloads = Overloaded{ [](int i) { return std::to_string(i); },
                                  [](const std::string &s) { return s; },
                                  [](auto &&val) { return std::string{}; } };
    std::variant<int, std::string> v{ 1 };
    std::visit(overloads, v);
}
```



# Variadic Template

- This is basically equivalent to generate:

```
struct Overloaded0 : A, B, C
{
public:
    std::string operator()(int i) const { return std::to_string(i); }
    std::string operator()(const std::string &s) const { return s; }
    std::string operator()(auto &&val) const { return std::string{}; }
};
```

- Here A, B, C are compiler-generated types for lambda, and each of them generates an overload for `operator()`.
- Notice that aggregates can have base class and implicit CTAD guides only since C++20; in C++17 we need additional guides.

# Variadic Template

- Exercise 8: write deduction guide for **Overloaded** in C++17.

```
template<typename... Args>  
Overloaded(Args...) -> Overloaded<Args...>;
```

- Exercise 9: what does **g(1, 2)** print?

```
template<class... Args>  
int h(Args... args)  
{  
    return sizeof...(Args);  
}  
  
template<class... Args>  
void g(Args... args)  
{  
    print(h(args) + args...);  
    print(h(args...) + args...);  
}
```

Expansion pattern will  
examine all previous  
code as a whole!

Answer: 2 3 3 4

$h(args) + args... \Leftrightarrow$   
 $h(arg1) + arg1,$   
 $h(arg2) + arg2, \dots$

$h(args...) + args... \Leftrightarrow$   
 $h(arg1, arg2, ...) + arg1,$   
 $h(arg1, arg2, ...) + arg2, \dots$

# Variadic Template

- Note 1: semantic substitution is different from plain text substitution. It has two steps:
  - Syntactic parse: examine what meaning the code wants to explain; this is unrelated to size of pack.
  - Substitution: insert types in AST.
- This explains why it works when `sizeof...(T1) == 0`:
  - If it's just text substitution, then it's `class A: public {}` (wrong syntax!).
  - If it's parsed to AST and then types are substituted, then it means "A inherits 0 type", which is correct.
- Note 2: variadic pack can only be the last parameters.

```
template<typename... T1>  
class A : public T1...
```

```
template<typename... Ts, typename T>  
void Error(Ts... args, T arg);
```

```
Error(1, 2);
```

Intuitive reason: pack will match eagerly, so `args` matches `1, 2` and `arg` matches nothing.

# Variadic Template

- But this is Okay:

```
template<typename... Ts> class Y {};  
  
template<typename... Ts, typename U>  
class Y<U, Ts...> {};
```
- Reason: **Ts...** is written at last so it matches last. This is unrelated to position of template parameter declaration (as it can be deduced)!
- Note 3: variadic template is less preferred than normal template in overload resolution.

```
template<typename T> void print(T); // preferred.  
template<typename... Ts> void print(Ts...);
```
- But when non-variadic one has more parameters filled with default, it's still ambiguous.\*

```
// ambiguous when calling print(1), #1 preferred when print(1,2)  
template<typename T> void print(T, T = T());  
template<typename T, typename... Ts> void print(T, Ts... args);
```

\*: Though this is regulated in C++ standard, only clang implements it.

# Variadic Template

- Note 4: NTPP can also use variadic template.

```
template<std::size_t... Idx, typename T>
void Access(const T &container)
{
    print(container[Idx]...);
}
```

```
std::vector<int> v{ 1, 2, 3 };
Access<1, 2>(v);
```

- But there is no way to hybrid two variadic templates.
  - i.e. impossible to define “any number of template parameter, either type or non-type”.
- Note 5: pack can be indexed at compile time since C++26.
  - But only id pack; expression pack cannot be indexed.

```
template<typename... Ts>
void Test(Ts... args)
{
    auto &arg0 = args...[0];
    using LastType = Ts...[sizeof...(args) - 1];
}
```

# Variadic Template

- And template template parameter pack isn't allowed to index.
- Note 6: friend can also be expanded pack since C++26.

```
struct C {};  
struct E { struct Nested; };
```

```
template<class... Ts>  
class R  
{  
    friend Ts...;  
};
```

```
template<class... Ts, class... Us>  
class R<R<Ts...>, R<Us...>>  
{  
    friend Ts::Nested..., Us...;  
};
```

```
R<C, E> rce;           // classes C and E are friends of R<C, E>
```

```
R<R<E>, R<C, int>> rr; // E::Nested and C are friends of R<R<E>, R<C, int>>
```

# Variadic Template

- Note 7: to capture variadic arguments in lambda, you need to:

```
void func(auto... args) {  
    int a = 0;  
    [&, args...]() { }; // 对a进行引用, args按值拷贝  
}
```

- However, you cannot transform it to a named parameter like normal capture:

```
void func(std::unique_ptr<int> ptr0) {  
    [ptr = std::move(ptr0)]() { };  
}
```

- Since C++20, this is improved so that you can write:

```
[...args_ = std::move(args)]() {  
    return foo(args_...);  
};
```

# Variadic Argument

- We notice that C also has variadic argument (e.g. `printf`), but it's very obscure to use.

- We're not going to talk about how to use it, but several notes:

1. It still uses ellipsis, but not on a pack expansion:

```
int printx(const char* fmt, ...);
```

- Comma before ... can be omitted (deprecated since C++26), so you can write something like:

```
template<typename T> template<typename... Ts>  
void Test(T...);      void Test(Ts.....);
```

Variadic  
template

Variadic  
argument

2. It has lowest overload resolution precedence, so `void Test(...);` means "the least preferred overload".

- We'll utilize this property in SFINAE.



# Variadic Template

- Variadic Template
  - Basics (Pack expansion & ...)
  - Fold expression

# Fold expression

e.g.  $f1() +_1 f2() +_2 f3()$ , it's  $\text{root}(+_2) \rightarrow \text{lChild}(f1() + f2()) \rightarrow \text{rChild}(f3())$ , while  $\text{lChild}$  is  $\text{root}(+_1) \rightarrow \text{lChild}(f1()) \rightarrow \text{rChild}(f2())$ ;

- We can know before  $+_1$  is evaluated,  $\text{lChild}$  and  $\text{rChild}$  is first evaluated.
- However, you can evaluate in the sequence of:
  - $\text{lChild}$  evaluates  $f1()$
  - $\text{rChild}$  evaluates  $f3()$ , gets the value.
  - $\text{lChild}$  evaluates  $f2()$ , gets the value.
  - This still obeys our rules, e.g.  $f1()$  and  $f2()$  evaluated before  $\text{lChild}$ .
- So if we output  $a$  in  $f1()$ ,  $b$  in  $f2()$ ,  $c$  in  $f3()$ , any permutation of  $abc$  is possible!

- With only simple pack expansion, print is still very strange...
  - Why do we have to do recursion? Normally it should just be a loop...
- To reduce recursive operations in pack, fold expression is introduced since C++17.
  - For example, to add all elements in a pack:

```
template<typename... Args>
auto AddAll(Args... args)
{
    auto result = (args + ... + 0);
    return result;
}
```

- $(\text{Pack OP } \dots \text{ OP Init})$ , meaning  $(\text{Pack}_1 \text{ op } (\dots \text{ op } (\text{Pack}_{N-1} \text{ op } (\text{Pack}_N \text{ op Init}))))$
- We notice that it doesn't necessarily mean  $\text{Pack}_N \text{ op Init}$  evaluates first; this depends on **evaluation order**, as we reviewed in Lecture 1.

# Fold expression

- Exercise: write print by fold expression.
  - Hint: comma expression.

`()` is necessary to raise precedence; in fold expression, pack should be with precedence not lower than cast.

```
template<typename... Args>
void print(const Args &...args)
{
    ((std::cout << args << " " ), ..., 0);
}
```

Comma expression will evaluate one by one from left to right, so the output is determined.

- Besides binary right fold, there also exists binary left fold:
  - $(\text{Init OP } \dots \text{ OP Pack})$ , meaning  $(((((\text{Init op Pack}_1) \text{ op Pack}_2) \text{ op } \dots) \text{ op Pack}_N)$
  - They differ when operation is not commutative:

```
template<typename... Args>
void Test(const Args &...args)
{
    std::cout << (args - ... - 0) << " " << (0 - ... - args) << '\n';
}
```

$\text{args} == 1, 2, 3,$   
 $\text{args} - \dots - 0 \Leftrightarrow$   
 $(1 - (2 - (3 - 0))) = 1 - (-1) = 2,$   
 $0 - \dots - \text{args} ==$   
 $((0 - 1) - 2) - 3 = -6,$

# Fold expression

- And besides binary fold, there also exist unary fold:
  - Unary right fold:  $(Pack\ OP\ \dots)$ , meaning  $(Pack_1\ op\ (\dots\ op\ (Pack_{N-1}\ op\ Pack_N)))$
  - Unary left fold:  $(\dots\ OP\ Pack)$ , meaning  $((Pack_1\ op\ Pack_2)\ op\ \dots)\ op\ Pack_N$
- Without **Init**, this expression is only valid when  $sizeof...(Pack) > 0$ !

```
template<typename... Args>
auto AddAll(Args... args)
{
    return (args + ...);
}
```

**AddAll()** will lead to compilation error.

- Exception 1: comma expression with 0 size is right, which means nothing.

```
template<typename... Args>
void print(const Args &...args)
{
    ((std::cout << args << " "), ...);
}
```

**print()** is correct, nothing is done.  
(Formally result is **void**).

# Fold expression

Notice that it may be astonishing when `operator&&/||` is overloaded to return types other than `bool`. Then:  
`sizeof...(values) == 0` → return type is `bool`;  
`sizeof...(values) > 0` → return type is overloaded type.

- Exception 2: logical expression with 0 size is correct.
  - Return a value that doesn't interfere its operation result (i.e. not short-circuit).
  - For `&&`, return true; for `||`, return false.

```
template<typename... Values>
int allTrue(const Values &...values)
{
    return (... && values);
}
template<typename... Values>
int anyTrue(const Values &...values)
{
    return (... || values);
}
int main()
{
    std::cout << allTrue(1, 1, 0) << allTrue(1, 1) << allTrue() << '\n'; // 011
    std::cout << anyTrue(1, 1, 0) << anyTrue(0, 0) << anyTrue() << '\n'; // 100
    return 0;
}
```

# Fold expression

- Since C++26, it's also allowed to reasonably use fold expression in constraints.
  - Before that, it's valid but **doesn't subsume**.

```
template <class T> concept A = std::is_move_constructible_v<T>;  
template <class T> concept B = std::is_copy_constructible_v<T>;  
template <class T> concept C = A<T> && B<T>;
```

```
// in C++23, these two overloads of g() have distinct atomic constraints  
// that are not identical and so do not subsume each other: calls to g() are ambiguous  
// in C++26, the folds are expanded and constraint on overload #2 (both move and copy  
// required), subsumes constraint on overload #1 (just the move is required)
```

```
template <class... T>  
requires (A<T> && ...) void g(T...); // #1
```

```
template <class... T>  
requires (C<T> && ...) void g(T...); // #2
```