Linear dependent types for quantum circuit programming

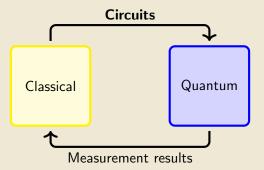
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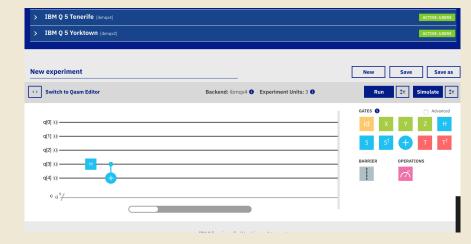
Joint work with: Kohei Kishida, Neil J. Ross, Peter Selinger

Feb 17, 2020, PL reading group, Maryland

QRAM model



Quantum circuits



Quantum circuit programming languages

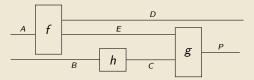


- QISKit, Q#, Cirq, ProjectQ.
- Quipper.

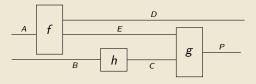
Linear types for quantum circuit programming

- No-cloning: $|\phi\rangle \mapsto |\phi\rangle \otimes |\phi\rangle$.
- Qwire.
- Proto-Quipper.

Linear types: $A \multimap B$, $A \otimes B$

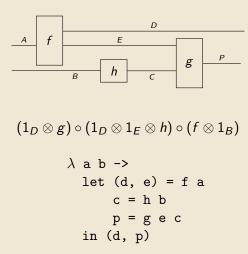


Linear types: $A \multimap B$, $A \otimes B$



$$(1_D \otimes g) \circ (1_D \otimes 1_E \otimes h) \circ (f \otimes 1_B)$$

Linear types: $A \multimap B$, $A \otimes B$



Parameters and states

- Values that are known at circuit generation time (parameters).
- Values that are known at circuit execution time (states).
- For example, a family of circuits indexed by parameters.
- Qwire: parameters and states live in different name spaces.
- Proto-Quipper: parameters and states live in the same name space.

Boxing and simple types

- We use simple types to describe (quantum) states. E.g.,
 Unit, Qubit, Qubit ⊗ Qubit.
- Box/unbox isomorphism.

box :
$$(a \multimap b) \to \mathbf{Circ}(a, b)$$

unbox : $\mathbf{Circ}(a, b) \to (a \multimap b)$

Note that a, b must be simple types.

• List of qubits is not a simple type because

 $\textbf{List Qubit} \cong \textbf{Unit} + \textbf{Qubit} + \textbf{Qubit} \otimes \textbf{Qubit} + ...$

A problem with quantum data types

- Can't box *qft* : **List Qubit** → **List Qubit**.
- Can't reverse *qft*: **List Qubit** → **List Qubit**.

A problem with quantum data types

- Can't box qft: List Qubit → List Qubit.
- Can't reverse *qft* : **List Qubit** → **List Qubit**.
- Dependent types to the rescue.
 qft: (n: Nat) → (Vec Qubit n → Vec Qubit n).
- Note that **Vec Qubit** *n* is a simple type.

$$\mathbf{Vec}\,\mathbf{Qubit}\,n\cong\underbrace{\mathbf{Qubit}\otimes\mathbf{Qubit}\otimes\ldots\otimes\mathbf{Qubit}}_{n}$$

Incorporating linear and dependent types

Let $f:(x:A) \multimap B[x]$ and a:A, what is the meaning of f:a:B[a]?

- Cervesato and Pfenning 1996. The term *a* is used twice, therefore it should be a parameter.
- McBride 2016. The term *a* in the type *B*[*a*] is not considered used, it is only "comtemplated".
- Our proposal (Fu, Kishida and Selinger 2020). The term a is used once, and the type B depends only on the shape of a.

Shape operation

Linear	\supseteq	Parameters
C1 (O 11:)		
$\mathrm{Sh}(Qubit)$	=	Unit
Sh(List $A)$	=	List $Sh(A)$
$Sh(A \otimes B)$	=	$\operatorname{Sh} A \otimes \operatorname{Sh} B$
$Sh((x:A) \multimap B[x])$	=	$(x: \operatorname{Sh}(A)) \to \operatorname{Sh}(B[x])$

Types depend on shapes

• Kinding linear dependent types:

$$\frac{\Phi, x : \mathrm{Sh}(A) \vdash B[x] : *}{\Phi \vdash (x : A) \multimap B[x] : *}$$

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• Kinding linear dependent types:

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• Typing rule for application:

$$\frac{\Gamma_1 \vdash M : (x : A) \multimap B[x] \quad \Gamma_2 \vdash N : A}{\Gamma_1 + \Gamma_2 \vdash MN : B[\operatorname{Sh}(N)]}$$

For $\Gamma_2 \vdash N : A$, we have:

$$\operatorname{Sh}(\Gamma_2) \vdash \operatorname{Sh}(N) : \operatorname{Sh}(A)$$

Proto-Quipper-D

- A linear dependent type system for programming quantum circuits.
- A categorical semantics based on state-parameter fibration.
- A prototype implementation of Proto-Quipper-D.

More details can be found in our recent drafts.

- "Linear dependent type theory for quantum programming panguages", Fu, Kishida and Selinger, 2020.
- "An introduction to quantum circuit programming in dependently typed Proto-Quipper". Fu, Kishida, Ross and Selinger, 2020.

Demo . An implementation is available at	

https://gitlab.com/frank-peng-fu/dpq-remake.

