

Design and Analysis of Content Caching Networks



Jim Kurose
Department of Computer Science
University of Massachusetts
Amherst MA USA

Visiting Scientist, Technicolor Paris/Palo Alto Labs
Professeur Invite, LINCS

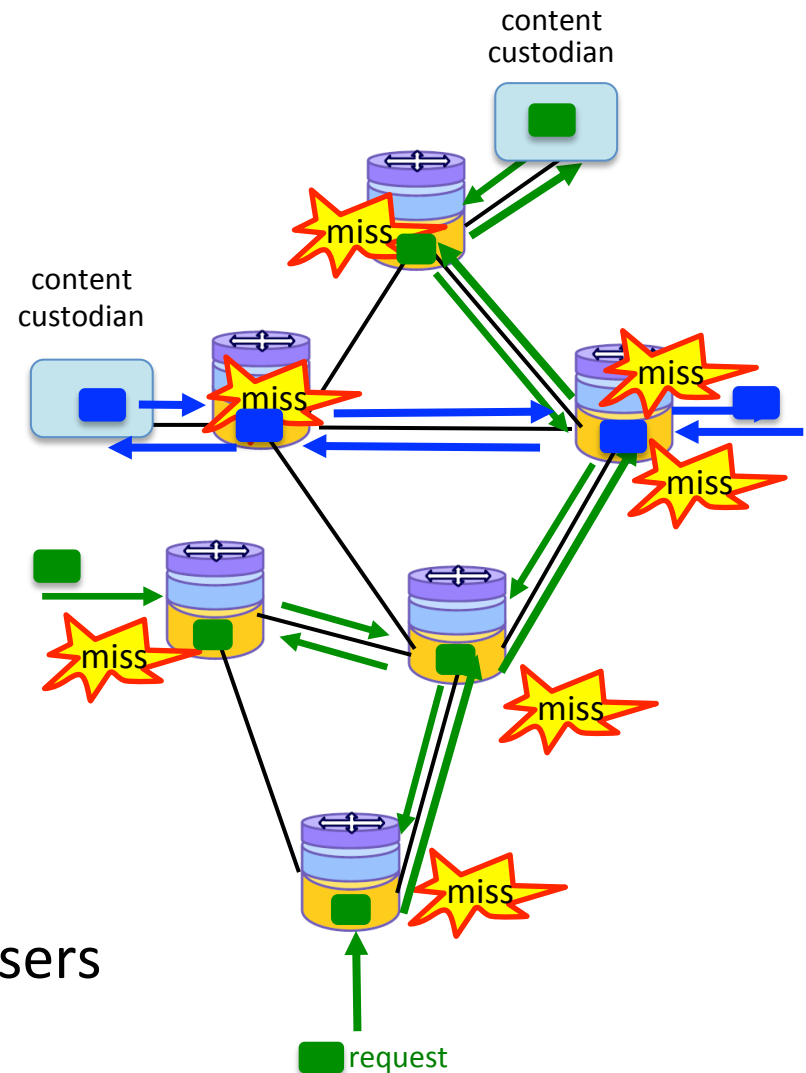
Workshop on Storage, Nov. 2012

Overview

- ❖ networks of caches
- ❖ breadcrumbs: content location/caching
- ❖ analysis of *networks* of caches
 - approximation algorithms
 - network calculus for cache networks
 - ergodicity

Networks of caches: ICN scenario

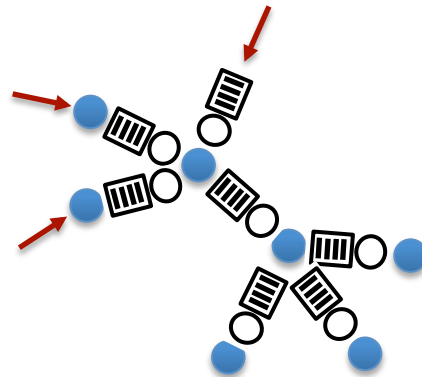
- ❖ consumer *requests content*
- ❖ request *routed* (e.g., shortest path) to known *content custodian*
- ❖ en-route to custodian, *caches* inspect request
 - *hit*: return local copy
 - *miss*: forward request towards custodian
- ❖ during download, content *stored* at caches along path
- ❖ content requests from different users interact: *cache replacement*



Note: ICN thinks wide area, but CamCube fits this model

Cache *networks*

- ❖ *network effect: interaction* among content request/reply flows from different users:
- ❖ *content replacement*: requested content by one user replaces content previously requested by others



Packet-switching:
queueing networks
(Kleinrock, 1963)

Networks of caches: challenges!

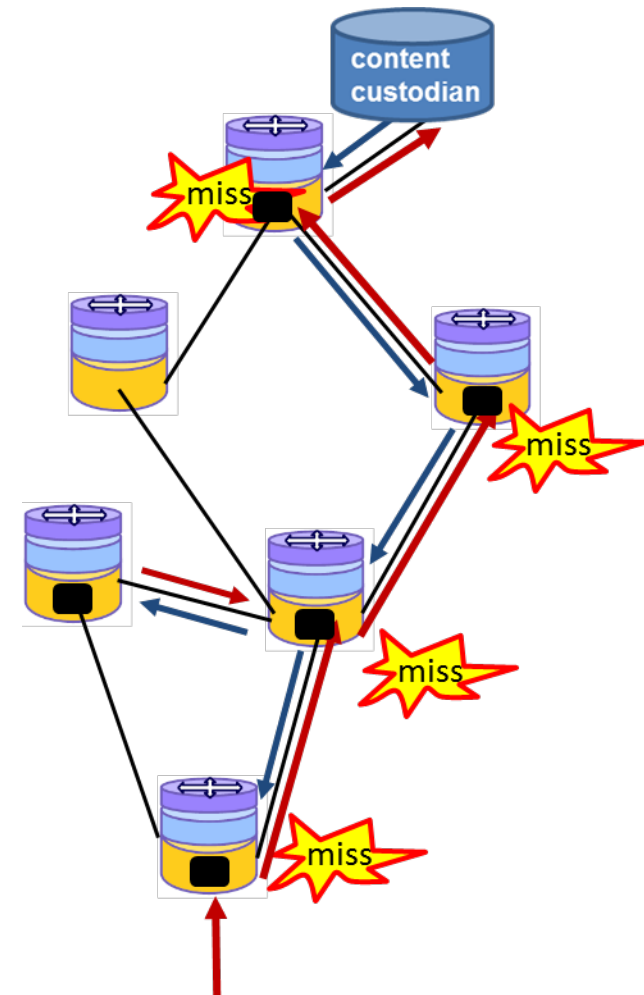
- ❖ **locating content:** known content location; opportunistic caching within network
 - ❖ *implicit:* search for content on shortest path to custodian
 - ❖ *explicit:* adaptive, state-dependent, and/or randomized search
 - ❖ similarities/difference from P2P search
- ❖ **managing content:** cache replacement, cache loading/prefetching
- ❖ **analyzing networks** of caches

Overview

- ❖ introduction: content-centric networks
- ❖ caching networks
- ❖ *breadcrumbs*: best-effort content location/caching
- ❖ analysis of *networks* of caches
 - approximation algorithms*
 - network calculus for cache networks*
 - ergodicity*

Breadcrumbs: motivation

- ❖ *goal*: access content, given known custodian location
- ❖ explicit or implicit cache search policy
 - *implicit*: search for content on shortest path to custodian
 - *explicit*: adaptive, state-dependent, and/or randomized search
- ❖ *Q*: how much state to maintain about network cache contents?



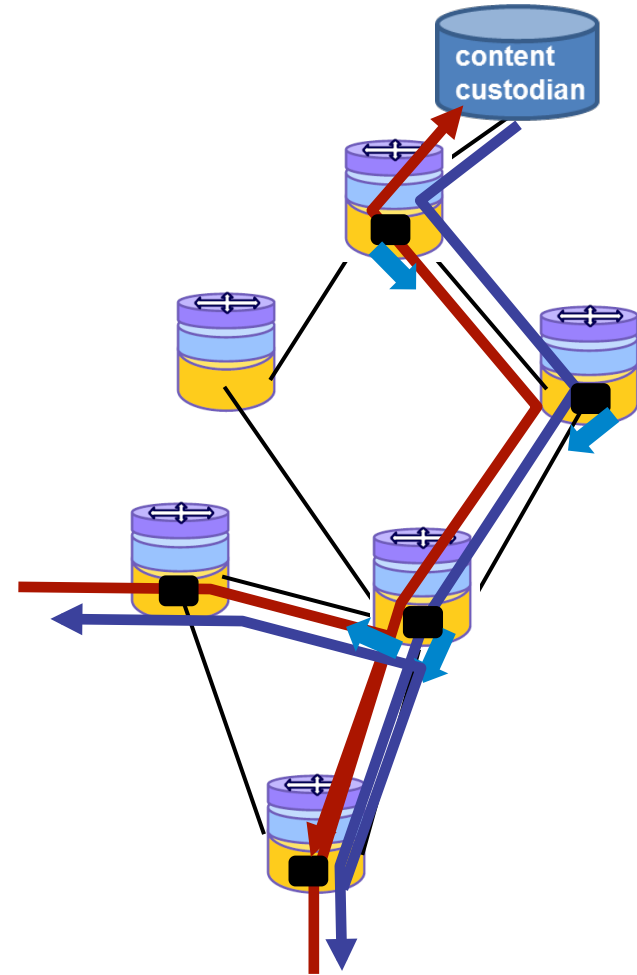
Breadcrumbs



Breadcrumb: per-file soft state - next-hop neighbor where file was last forwarded

Breadcrumbs search

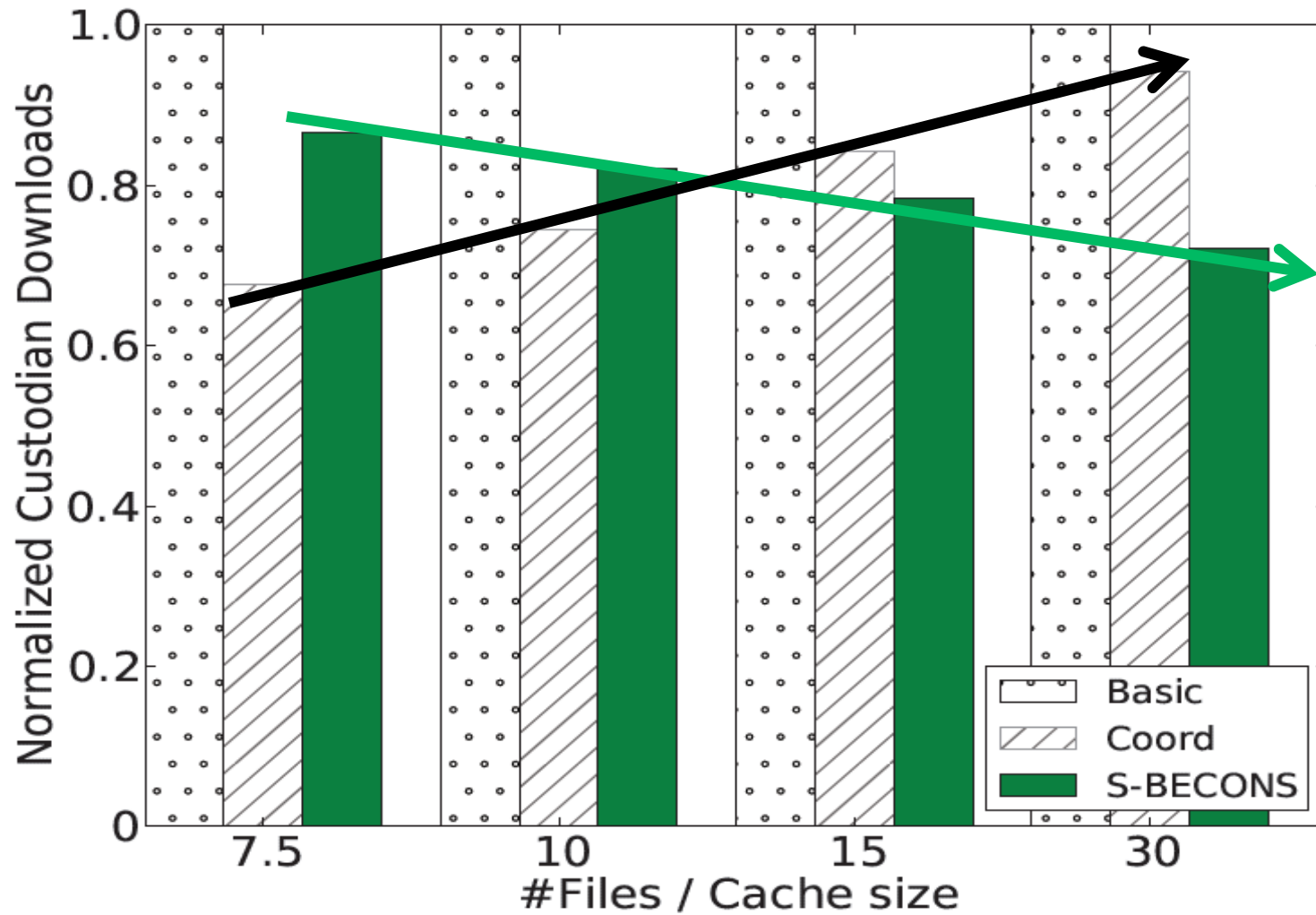
- ❖ request starts on route (shortest path) to custodian
- ❖ enroute cache holding file sends file to requestor
- ❖ cache with recent (breadcrumb can route request following breadcrumb) or to custodian
 - if **dead-end** reached: reroute to custodian
- ❖ soft-state timeout
- ❖ breadcrumb path properties: stability, invalidation



Breadcrumbs: simulation analysis

- ❖ nxn torus, Zipf and uniform content access probability, 300 files, varying cache sizes
- ❖ compared to
 - caching w/ no Breadcrumbs (“Basic”),
 - coordinated caching with upstream neighbor
- ❖ performance metric
 - minimize load on custodians
- ❖ sample results
 - *Breadcrumbs* performance improves (relatively) as network size grows, and as the files-to-cache size ratio grows

Breadcrumbs: sample results



Breadcrumbs: future work

- ❖ tradeoff between time to locate content, cache/custodian load balancing
- ❖ estimating, exploiting cache eviction rates of neighbors



NEC implementation

“Breadcrumbs: efficient, best-effort content location in cache networks,”
Elisha J. Rosensweig and Jim Kurose, *2009 IEEE INFOCOM Mini-conference*

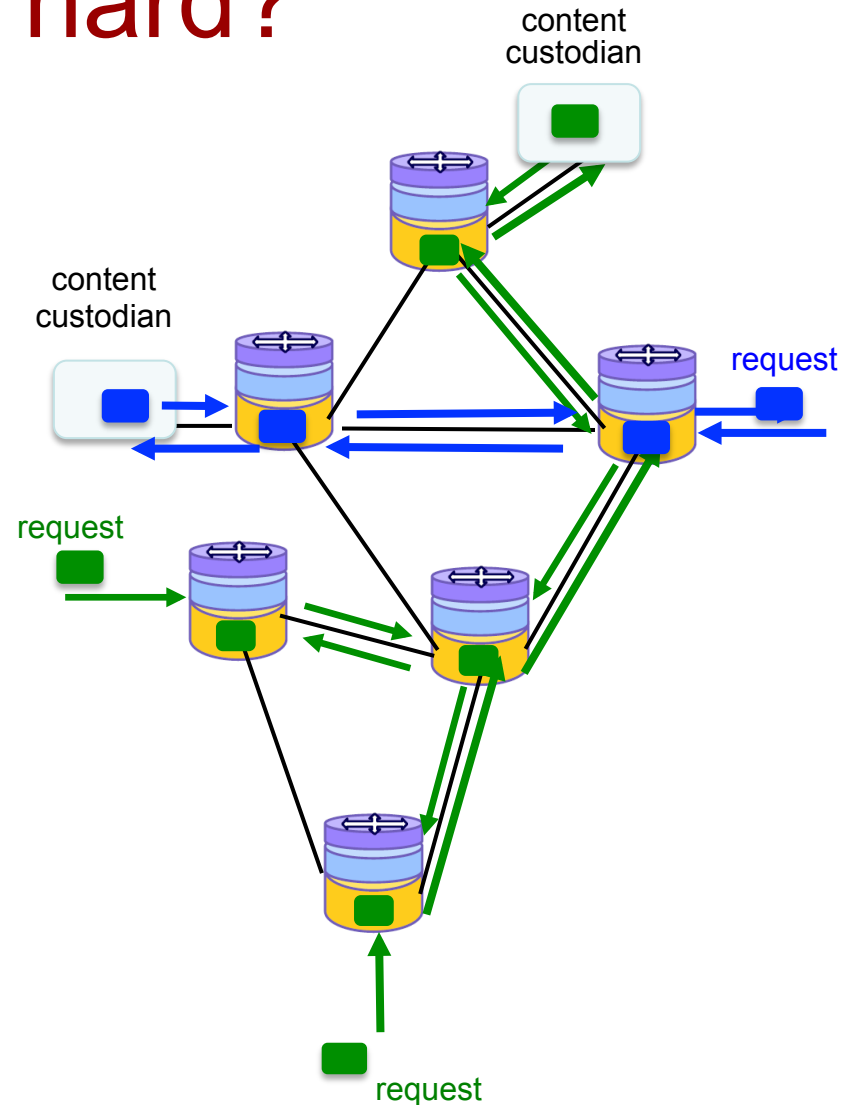
“Performance Evaluation of Partial Deployment of Breadcrumbs in
Content-Oriented Networks,” T. Tatsuhiro et al, *2012 IEEE ICC Futurenet*

Overview

- ❖ introduction: content-centric networks
- ❖ networks of caches
- ❖ content location/caching
- ❖ analyzing *networks* of caches
 - approximation algorithms
 - network calculus for cache networks
 - ergodicity

Why is the problem hard?

- ❖ simultaneous allocation of resources (cache storage) from point where content is located back to requestor
- ❖ asynchronous resource release, depending on cross flows
- ❖ inter-cache flows complex:
 - ❖ miss stream from downstream caches
 - ❖ lack of independence



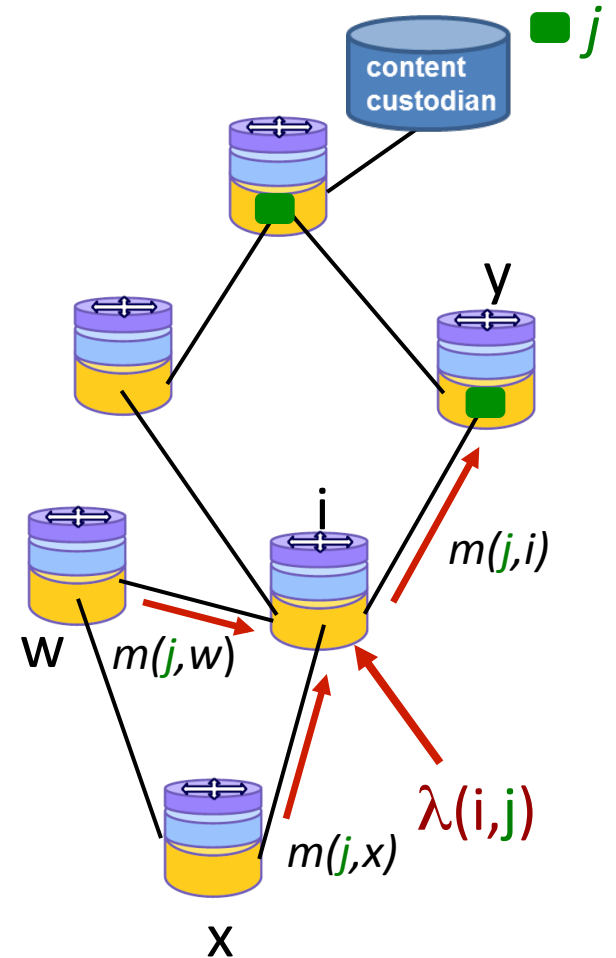
Modeling a network of caches

- ❖ node i : *exogenous* (external) arrivals for content j : $\lambda(i,j)$
- ❖ node i : *internal* arrivals (miss stream) for content j from downstream neighbors h : $m(j,h)$
- ❖ $r(i,j)$: aggregate rate of arrival requests at i for content j

$$r(i,j) = \lambda(i,j) + \sum_{\text{all downstream neighbors, } h} m(j,h)$$

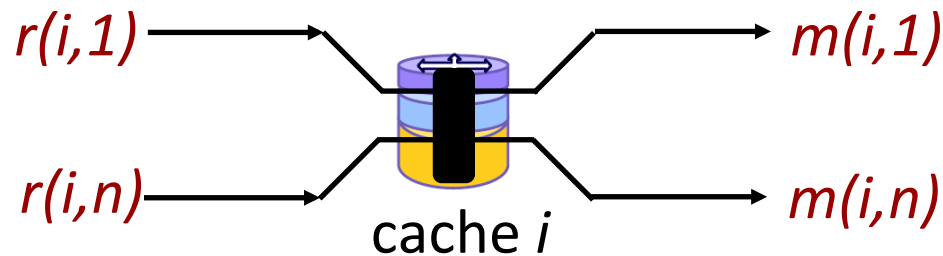
all downstream
neighbors, h

- ❖ ZDD: zero download delay assumption

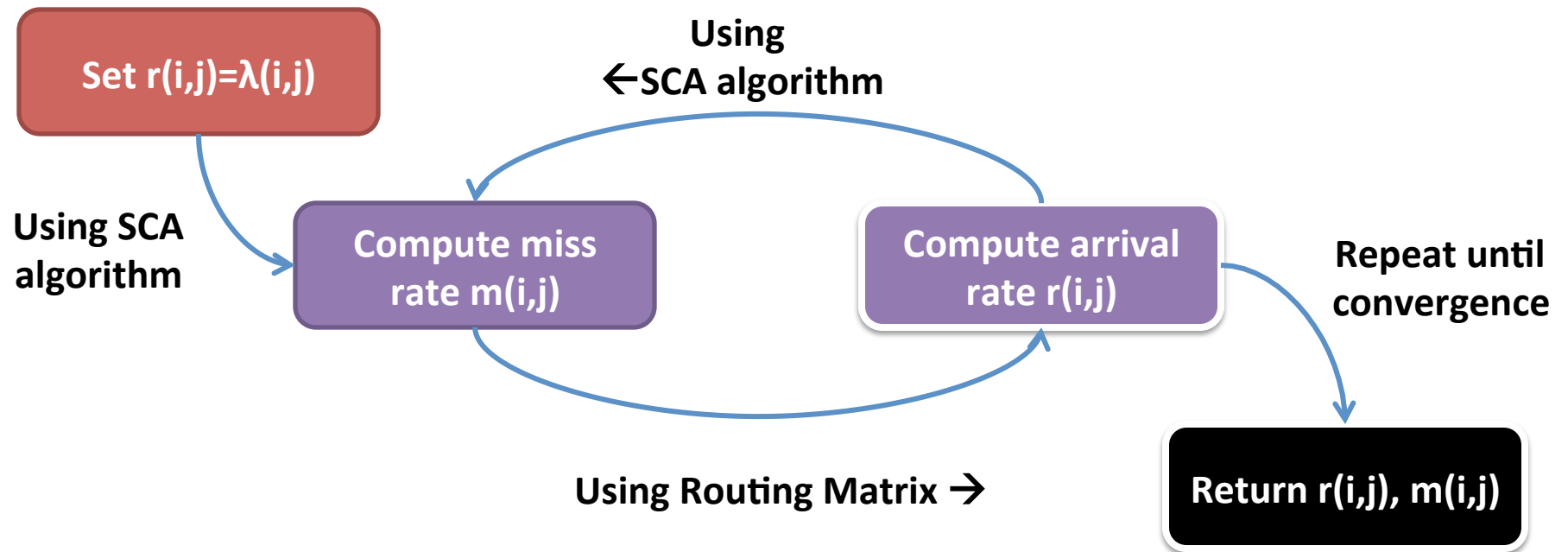


Approximating cache network performance

- ❖ **SCA:** standalone cache i approximation algorithm: given $r(i,j)$, compute miss rate for all content j
- ❖ Independence Reference Model (IRM) of incoming requests:
$$Pr(X_t = f_j \mid X_1, \dots, X_{t-1}) = Pr(X_t = f_j)$$
- ❖ SCA approximation algorithm for LRU: [Dan 1985]



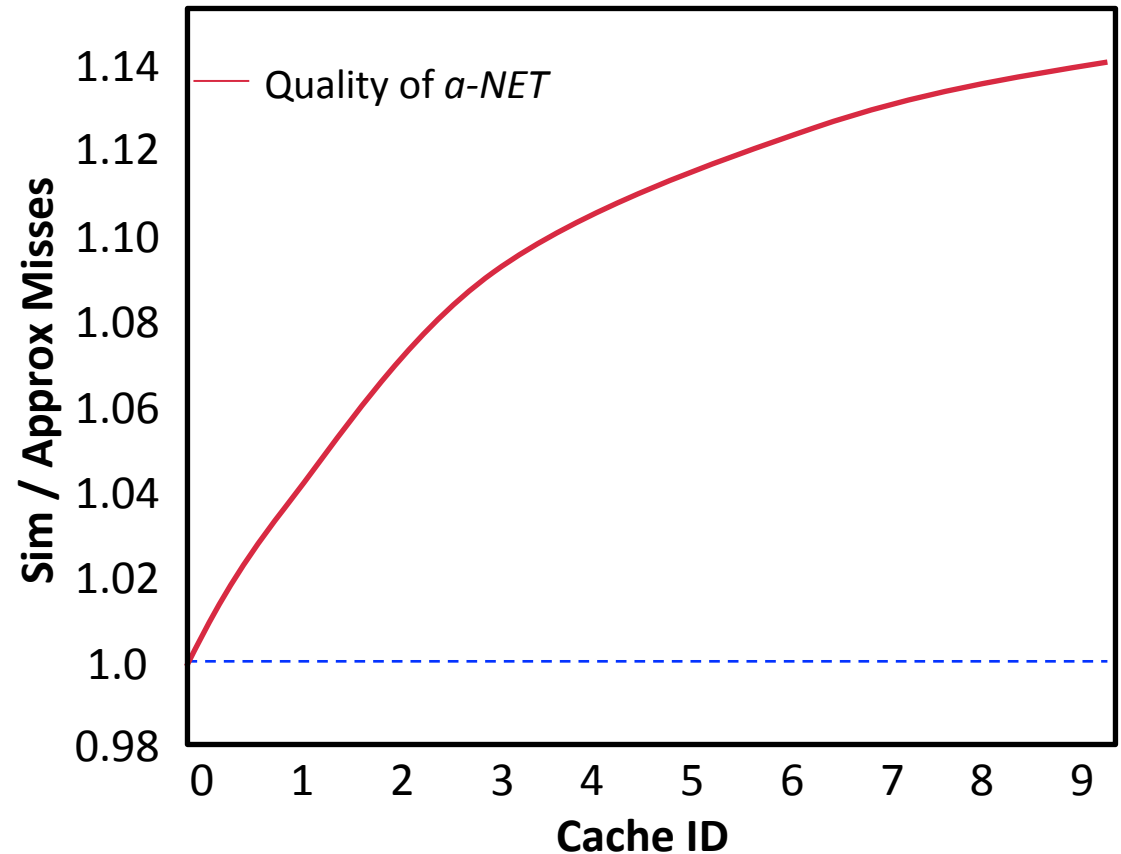
Fixed-point iteration



Note: tree-network (feed-forward) require single iteration

Fixed point approximation error

- ❖ line topology with 9 nodes
- ❖ errors decrease in networks with high node degree



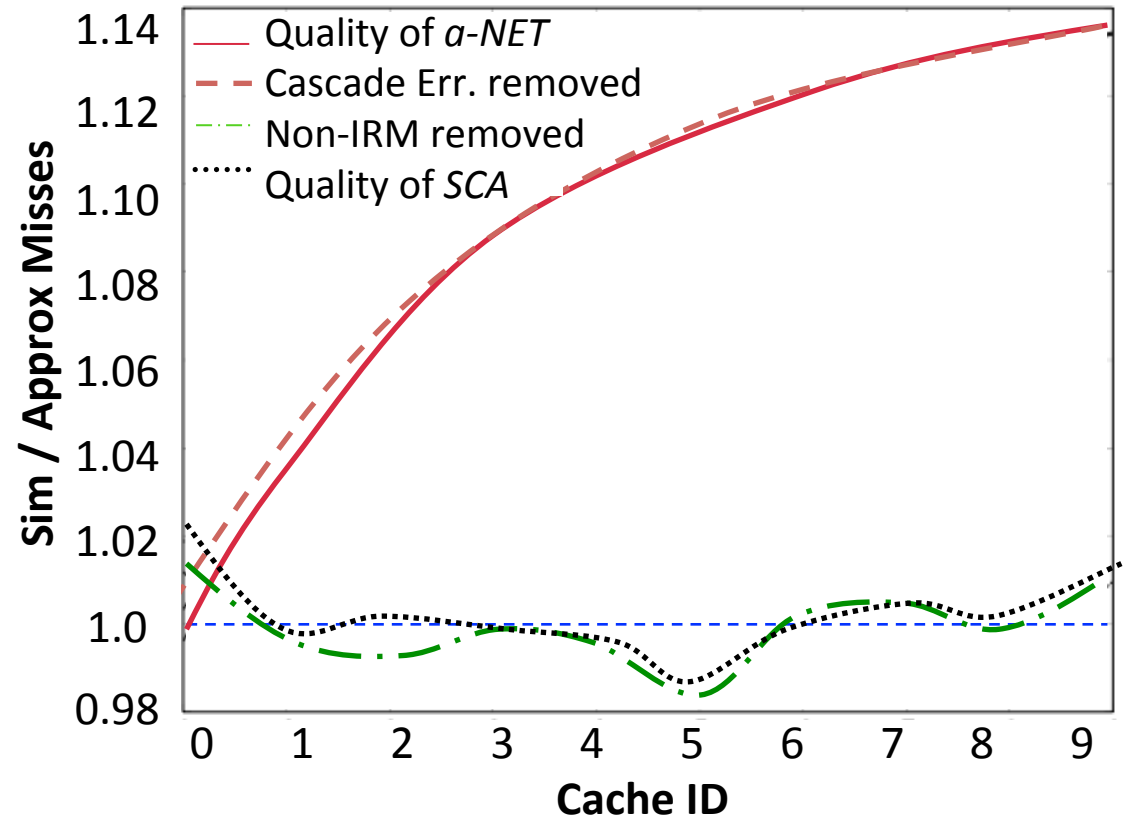
Error factor analysis

sources of approximation errors:

- ❖ SCA algorithm inaccuracies?
- ❖ cascading errors?
 - approximated output rates of one iteration is input to next iteration
- ❖ violating IRM assumed by SCA algorithm?
 - miss process for file j negatively correlated

Error factor analysis

- Factor analysis reveals that *non-IRM input* to SCA is main cause of error



“Approximate Models for General Cache Networks,”
Elisha J. Rosensweig, Jim Kurose, Don Towsley, 2010 IEEE INFOCOM

Overview

- ❖ introduction: content-centric networks
- ❖ caching networks
- ❖ *breadcrumbs*: best effort content location/caching
- ❖ analysis of *networks* of caches
 - approximation algorithms
 - network calculus for cache networks
 - ergodicity

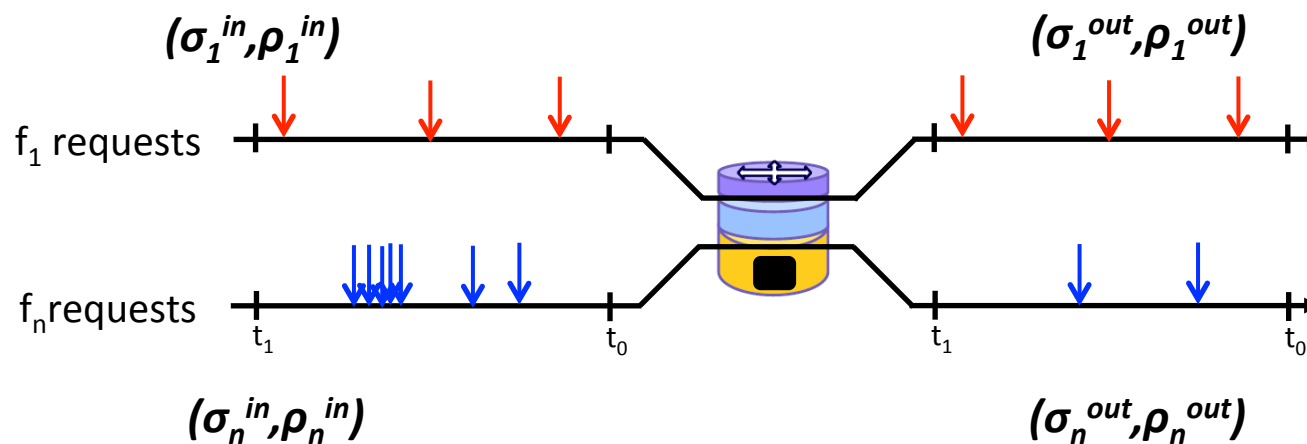
(σ_i, ρ_i) analyses of cache networks

(σ_i, ρ_i) bounds $\#$ requests for f_i over $[t_1, t_2]$:

$$\int_{t_1}^{t_2} r_i(t) dt \leq \rho_i(t_2 - t_1) + \sigma_i$$

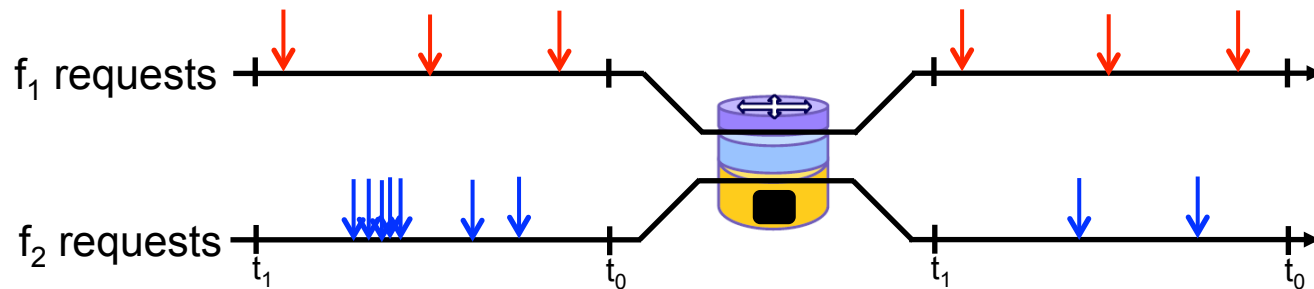
where $r_i(t)$ = request rate for f_i at time t

Goal: a network calculus for cache networks:



(σ_i, ρ_i) cache networks: observations

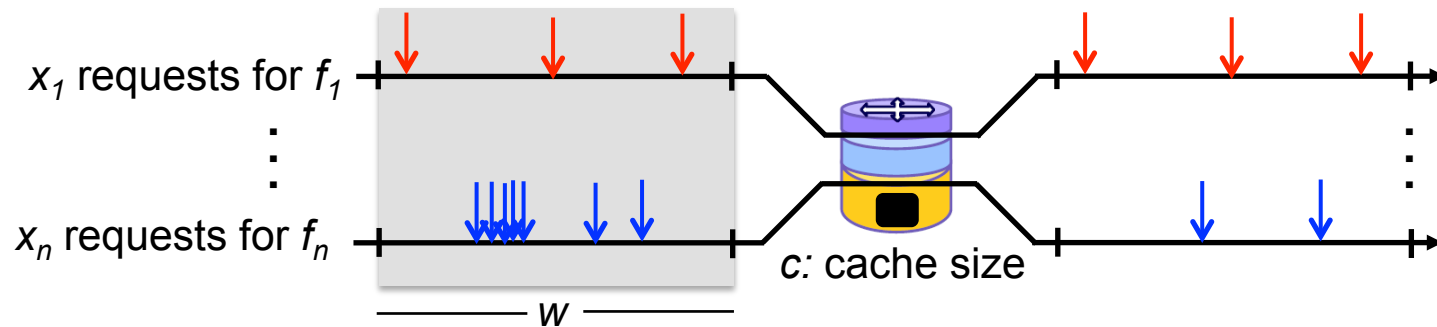
- ❖ not all requests arriving at cache will leave (unlike queue)



- ❖ stream of input requests for one file only generates no output
 - interactions among files in cache critical
- ❖ “burst” of request for same file generates one output
 - different intuition (from queues) about “performance damage” of bursts

Building block: miss set, M_i

miss set for f_i : set of requests for c unique files, other than i

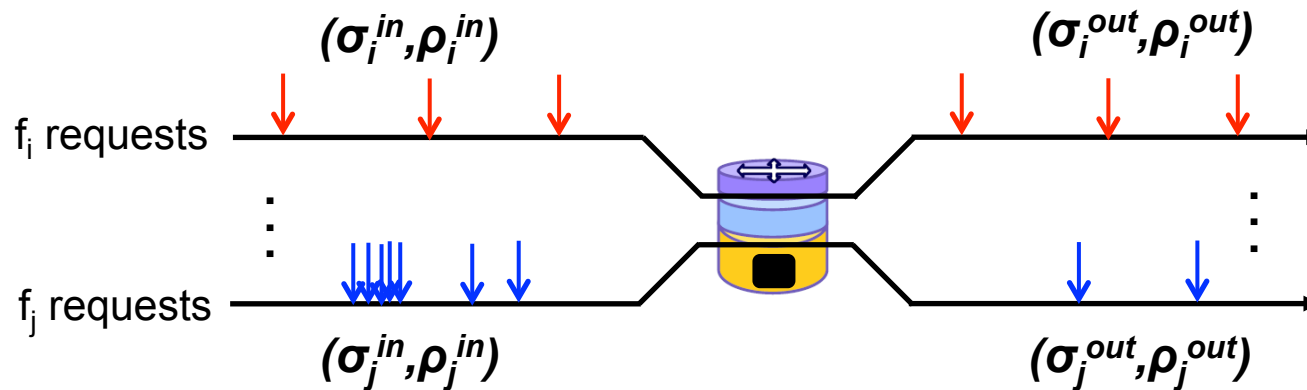


$M_i(x_1, \dots, x_n, c)$: max number miss sets for file f_i , given $\{x_i^{in}\}$ arrivals, cache of size c .

properties:

- $M_i = \min(x_i^{in}, M)$ **← !!**
- $x_i^{out} \leq M_i$, and this bound is achievable

From $(\sigma_i^{in}, \rho_i^{in})$ to $(\sigma_i^{out}, \rho_i^{out})$:

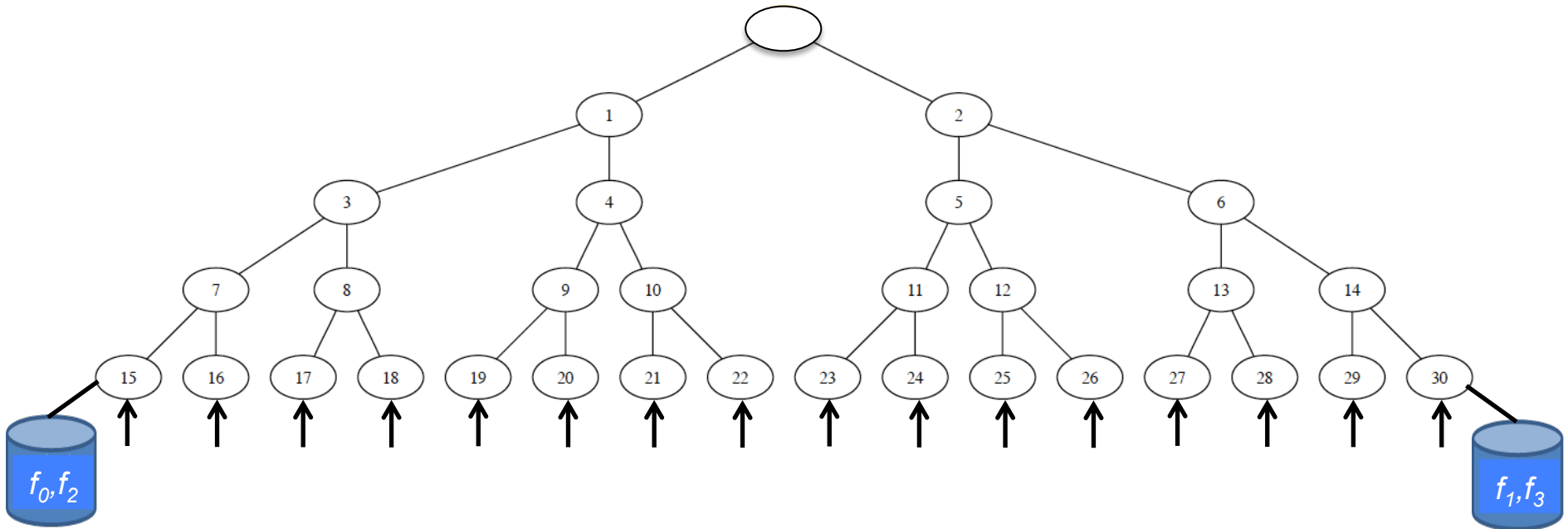


Theorem: For a cache of size c :

$$\rho_i^{out} = \min(\rho_i^{in}, M_i(\rho_1^{in}, \dots, \rho_n^{in}, c))$$

Can calculate ρ_i^{out} from $\{\rho_i^{in}\}$

Numerical example



cache size = 2 at each node

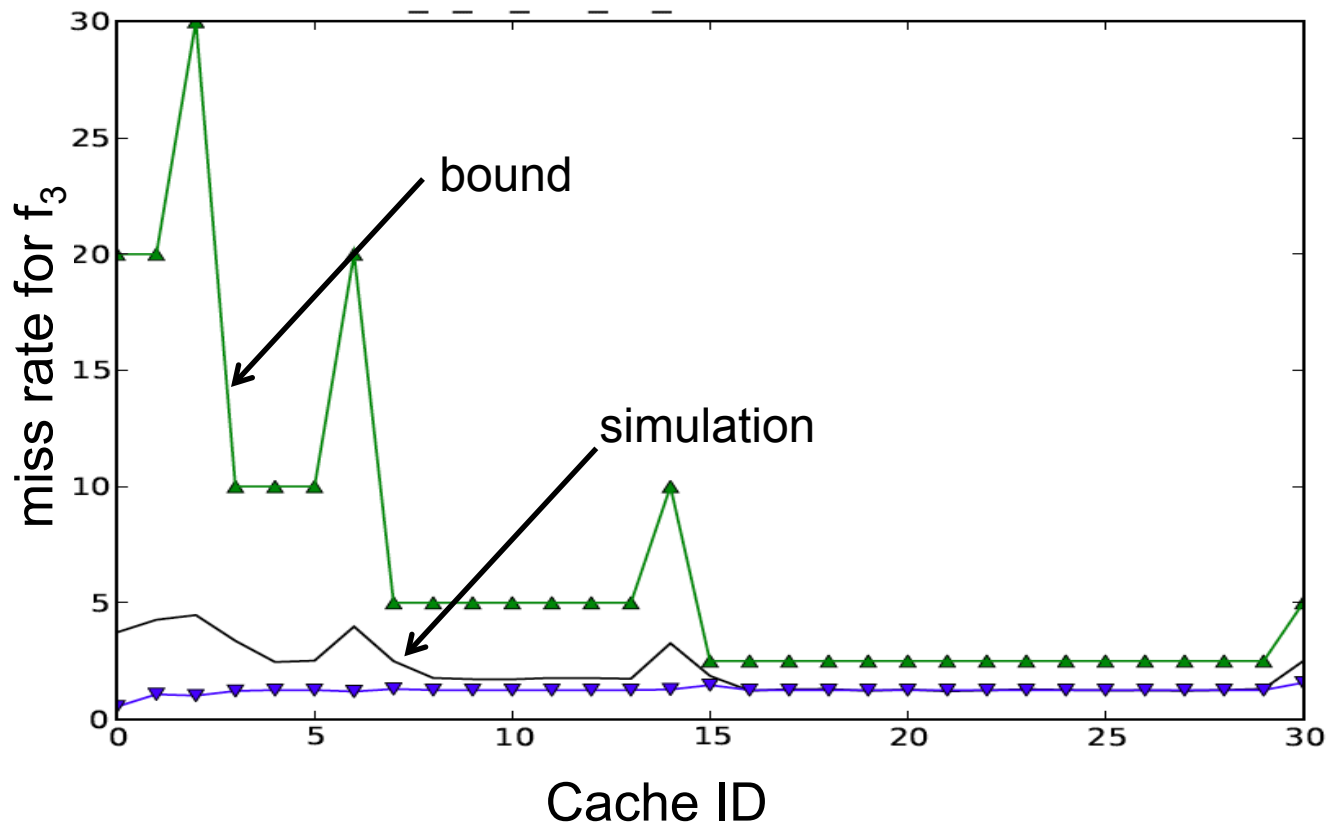


homogeneous IRM arrivals, exponential interarrival times



4 files, uniform popularity distribution

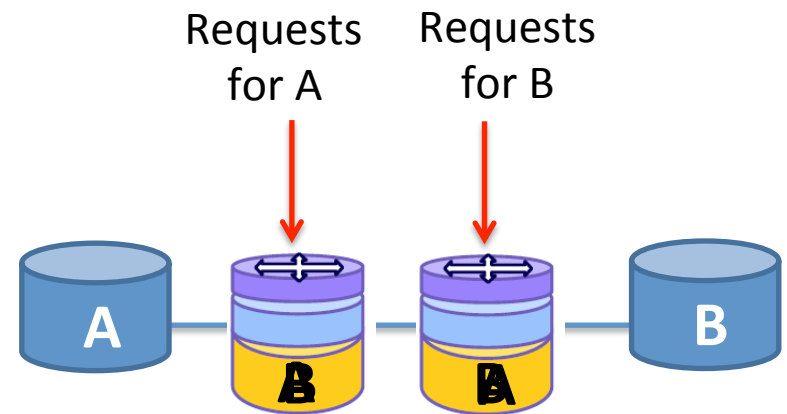
Numerical example: bounding results



E. Rosensweig, J. Kurose, "A Network Calculus for Cache Networks," to appear in *IEEE Infocom 2013*.

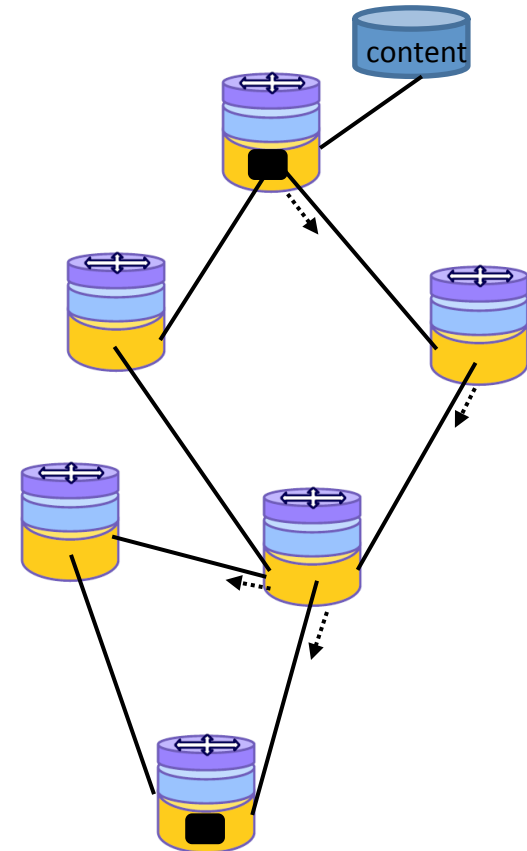
Ergodicity of cache networks

- does steady state performance depend on initial conditions (ergodicity)?
 - shown existence of non-ergodic cases (replacement policy, topology, cache size)
 - derived sufficient conditions for ergodicity
 - topology (single-custodian trees)
 - from individual ergodicity to system ergodicity



Ergodicity of cache networks

- ergodicity (continued):
 - showed random replacement: ergodic
 - defined class of non-protective policies (including LRU): all ergodic



E. Rosensweig, D. Menasche, J. Kurose, "On the Steady-State of Cache Networks," to appear in *IEEE Infocom 2013*.

Future work

- ❖ cache network protocols
 - multipath: download from multiple custodians
 - mobility
 - proactive caching, replacement in video-content distribution nets
- ❖ network calculus
 - continued development of deterministic bounds
 - stochastic bounding techniques
- ❖ approximate analysis of cache networks
 - asymptotic (Kelly, Bonald) results

The end

?? || /* */