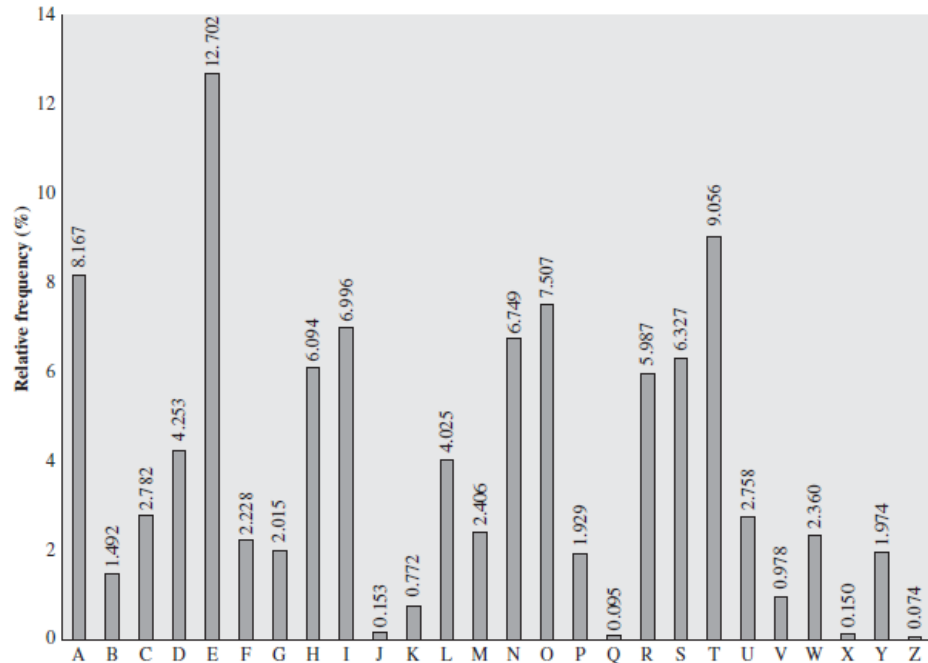


Cryptography Exam 1

Name _____ Date 2/8/16 (Closed note, computer, book)

Problem 1. A ciphertext has been generated with an affine cipher. The most frequent letter of the ciphertext is 'B', and the second most frequent letter of the ciphertext is 'U'. Break this code.



E(4) --> B(1)
T(19) --> U(20)

(P,C)=(4,1) (P,C)=(19,20)
Encryption equation: $C=(aP+b)\bmod 26$
C=Ciphertext
P=Plaintext

$$1=(4a+b)\bmod 26$$
$$20=(19a+b)\bmod 26$$

Solving two simultaneous equations, by subtracting the equations:

$a = (1/15) \cdot 19 \bmod 26 = 7 \cdot 19 \bmod 26 = 133 \bmod 26 = 3 \rightarrow a=3$
where (1/15) is found using:
 $(1/15) \cdot 15 = 1 \bmod 26$ By trial and error: $(1/15)=7$

To find b:
 $1=(4 \cdot 3+b)\bmod 26 \rightarrow b=-11 \bmod 26 = 15$

Key: (a,b) = (3,15) --> Code is broken!!

Problem 2. What is multiplicative inverse of 8 in the rings \mathbb{Z}_{12} , \mathbb{Z}_{13} , and \mathbb{Z}_{14} .

Answer:

Ring \mathbb{Z}_{12} : $\gcd(12,8)$ is not 1, so $(1/8)$ does not exist.

Ring \mathbb{Z}_{13} : $\gcd(13,8)=1$, so $(1/8)$ exists: $(1/8) \cdot 8 = 1 \pmod{13}$
Using trial and error between numbers 1 through 7: $1/8=5$

Ring \mathbb{Z}_{14} : $\gcd(14,8)$ is not 1, so $(1/8)$ does not exist.

Problem 3. Compute x in the following without the use of a calculator:

(a) $x = 3^{100} \pmod{13}$

Answer:

$$\begin{aligned} x &= (3^4)^{25} \pmod{13} = (3)^{25} \pmod{13} = (3^4)^6 \cdot 3 \pmod{13} = (3^6 \cdot 3) \pmod{13} \\ &= (3^4 \cdot 3^3) \pmod{13} = (3 \cdot 1) \pmod{13} = 3 \end{aligned}$$

(b) $13^x = 1 \pmod{18}$

Answer:

$$13^1 \pmod{18} = 13$$

$$13^2 \pmod{18} = 169 \pmod{18} = 7$$

$$13^3 \pmod{18} = (13^2 \times 13) \pmod{18} = (7 \times 13) \pmod{18} = (91) \pmod{18} = 1$$

$$13^4 \pmod{18} = (13^3 \times 13) \pmod{18} = (1 \times 13) \pmod{18} = 13$$

$$13^5 \pmod{18} = (13^4 \times 13) \pmod{18} = (13 \times 13) \pmod{18} = (169) \pmod{18} = 7$$

$$13^6 \pmod{18} = (13^5 \times 13) \pmod{18} = (7 \times 13) \pmod{18} = (91) \pmod{18} = 1$$

13, 7, and 1 will repeat. The answer is $x = 0, 3, 6, 9, \dots$

$$(c) x = 7^{50} \bmod 19$$

Answer:

$$\begin{aligned} x &= 7^{50} \bmod 19 = (7^2)^{25} \bmod 19 = 11^{25} \bmod 19 = (11^2)^{12} \cdot 11 \bmod 19 \\ &= 7^{12} \cdot 11 \bmod 19 = (7^2)^6 \cdot 11 \bmod 19 = 11^6 \cdot 11 \bmod 19 \\ &= (11^2)^3 \cdot 11 \bmod 19 = 7^3 \cdot 11 \bmod 19 = 7^2 \cdot 7 \cdot 11 \bmod 19 \\ &= 11 \cdot 7 \cdot 11 \bmod 19 = 7 \cdot 7 \bmod 19 = 11 \end{aligned}$$

Problem 4. Prove that double encryption with the affine cipher is only as secure as single encryption.

Answer:

$$\begin{aligned} C1 &= (aP + b) \bmod 26 \\ C2 &= (aC1 + b) \bmod 26 \end{aligned}$$

P= Original Plaintext

We have to figure out how C2 is related directly to P.

$$\begin{aligned} C2 &= (aC1 + b) \bmod 26 = [a((aP + b)) + b] \bmod 26 \\ &= (a \cdot a \cdot P + a \cdot b + b) \bmod 26 \end{aligned}$$

With choosing of: $A = a \cdot a$; $B = a \cdot b + b$

$$C2 = (A \cdot P + B) \bmod 26$$

This will be another affine cipher with the key (A=a.a , B=a.b+b)

So it is as secure as only a single affine cipher.

Problem 5. Find the plaintext in the affine cipher with the key parameter **a = 7** and **b = 22**: falszztysyzyjkywjrztjztyynaryjkyswarztyegyyj

Answer:

$$(1/a) \cdot 7 = 1 \pmod{26} \rightarrow 1/a = 15$$

$$\text{Decryption equation: } P = [(1/a) \cdot (C - b)] \pmod{26} = [15 \cdot (C - 22)] \pmod{26}$$

a: $15 \cdot (0 - 22) \pmod{26} = 15 \cdot 4 \pmod{26} = 8 \rightarrow i$
e: $15 \cdot (4 - 22) \pmod{26} = 15 \cdot 8 \pmod{26} = 16 \rightarrow q$
f: $15 \cdot (5 - 22) \pmod{26} = 15 \cdot 9 \pmod{26} = 5 \rightarrow f$
g: $15 \cdot (6 - 22) \pmod{26} = 15 \cdot 10 \pmod{26} = 20 \rightarrow u$
j: $15 \cdot (9 - 22) \pmod{26} = 15 \cdot 13 \pmod{26} = 13 \rightarrow n$
k: $15 \cdot (10 - 22) \pmod{26} = 15 \cdot 14 \pmod{26} = 2 \rightarrow c$
l: $15 \cdot (11 - 22) \pmod{26} = 15 \cdot 15 \pmod{26} = 17 \rightarrow r$
n: $15 \cdot (13 - 22) \pmod{26} = 15 \cdot 17 \pmod{26} = 21 \rightarrow v$
r: $15 \cdot (17 - 22) \pmod{26} = 15 \cdot 21 \pmod{26} = 3 \rightarrow d$
s: $15 \cdot (18 - 22) \pmod{26} = 15 \cdot 22 \pmod{26} = 18 \rightarrow s$
t: $15 \cdot (19 - 22) \pmod{26} = 15 \cdot 23 \pmod{26} = 7 \rightarrow h$
w: $15 \cdot (22 - 22) \pmod{26} = 15 \cdot 0 \pmod{26} = 0 \rightarrow a$
y: $15 \cdot (24 - 22) \pmod{26} = 15 \cdot 2 \pmod{26} = 4 \rightarrow e$
z: $15 \cdot (25 - 22) \pmod{26} = 15 \cdot 3 \pmod{26} = 19 \rightarrow t$
falszztysyzyjkywjrztjztyynaryjkyswarztyegyyj
first the sentence en...

Problem 6. The requirement for an encryption function **E** to be one-to-one is that for any two plaintext **p1** \neq **p2** and the key, **K**, then the **E(K,p1)** \neq **E(K,p2)**. Find and prove the condition for the affine cipher to be one-to-one.

Answer:

$$p1 \neq p2 \rightarrow E(K,p1) \neq E(K,p2)$$

This is equivalent to:

$$E(K,p1) = E(K,p2) \rightarrow p1 = p2$$

So:

$$E(K,p1) = E(K,p2) \rightarrow E((a,b),p1) = E((a,b),p2) \rightarrow (ap1+b) \bmod 26 = (ap2+b) \bmod 26 \rightarrow a.(p1-p2) \bmod 26 = 0$$

In order to have this condition satisfied, **a** should be having multiplicative inverse:

There should be such (1/a) that $(1/a).a = 1 \bmod 26$

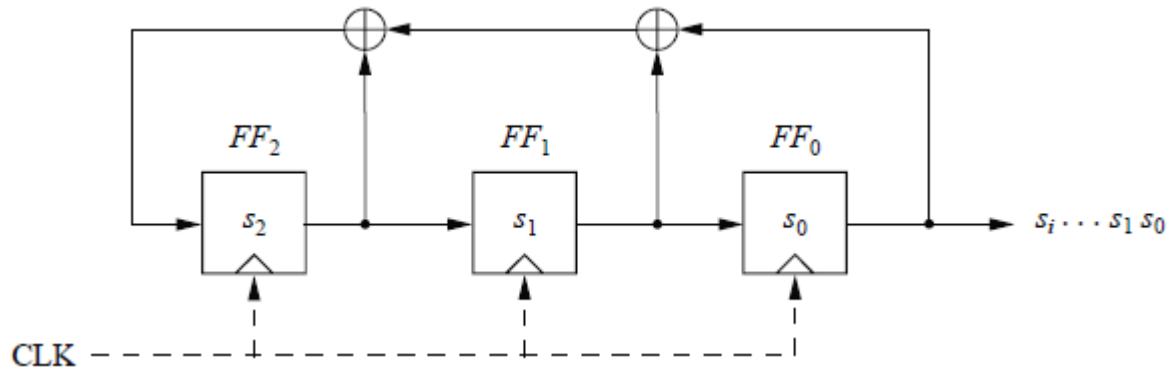
$$(1/a).a.(p1-p2) \bmod 26 = 0$$

$$\rightarrow (p1-p2) \bmod 26 = 0 \rightarrow p1 = p2$$

Condition is existence of (1/a) in mod 26:

$$\gcd(26,a)=1 \rightarrow a=\{1,3,5,7,9,11,15,17,19,21,23,25\}$$

Problem 7. In the following LFSR:



(a) Find the sequence generated from initialization vector ($S_2=1$, $S_1=0$, $S_0=1$).
What is the period?

$$S_3 = S_2 \oplus (S_1 \oplus S_0)$$

In general form:

$$S_{i+3} = S_{i+2} \oplus (S_{i+1} \oplus S_i)$$

$S_3=0$, $S_4= S_3 \oplus (S_2 \oplus S_1) = 0 \oplus (1 \oplus 0) = 1$, $S_5= S_4 \oplus (S_3 \oplus S_2) = 1 \oplus (0 \oplus 1) = 0$, $S_6= S_5 \oplus (S_4 \oplus S_3) = 0 \oplus (1 \oplus 0) = 1$,
 $S_0, S_1, S_2, S_3, S_4, S_5, S_6, \dots = 1, 0, 1, 0, 1, 0, \dots$
 Period = 2

(b) Find the sequence generated from initialization vector ($S_2=0$, $S_1=1$, $S_0=1$).
What is the period?

$$S_3 = S_2 \oplus (S_1 \oplus S_0)$$

In general form:

$$S_{i+3} = S_{i+2} \oplus (S_{i+1} \oplus S_i)$$

$S_3= S_2 \oplus (S_1 \oplus S_0) = 0$, $S_4= S_3 \oplus (S_2 \oplus S_1) = 0 \oplus (0 \oplus 1) = 1$, $S_5= S_4 \oplus (S_3 \oplus S_2) = 1 \oplus (0 \oplus 0) = 1$,
 $S_6= S_5 \oplus (S_4 \oplus S_3) = 1 \oplus (1 \oplus 0) = 0$
 $S_0, S_1, S_2, S_3, S_4, S_5, S_6, \dots = 1, 1, 0, 0, 1, 1, 0, \dots$
 Period = 4