

Concurrent systems

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Concepts

Liveness - The operation eventually returns something.

Safety - The operation never returns anything incorrect (an ad-hoc rule).

Correct - Safety and Liveness

- In out context: A process that never **fails** (stops taking steps) in the middle of middle of its operation is called **correct**. We typically assume that a correct process invokes infinitely many operations, so a process is correct if it takes infinitely many steps.

Progress - Non blocked, process completes operations in a finite amount of time.

Liveness Properties

Deadlock-free (DF) - If every process is correct, some process makes progress.

Starvation-free (SF) - If every process is correct, every process makes progress.

Lock-free / non-blocking (LF) - Some correct process makes progress (in a finite number of steps).

Wait-free (WF) - Every correct process makes progress (in a finite number of steps).

Obstruction-free (OF) - Every process makes progress if it executes in isolation from some point (it is the only correct process).

Periodic table of liveness properties

| | Independent non-blocking (finite steps) | Dependent non-blocking | Dependent blocking (infinite steps) |
|------------------------------|---|------------------------|-------------------------------------|
| every process makes progress | wait-freedom | obstruction-freedom | starvation-freedom |
| some process makes progress | lock-freedom | ? | deadlock-freedom |

Independent: Independent of concurrency.

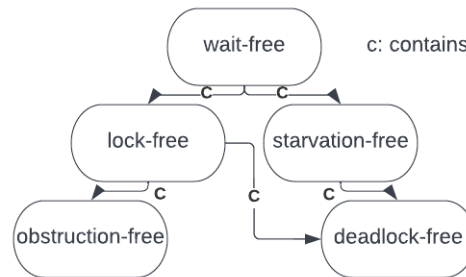
Non-blocking: At least one process make progress.

Independent non-blocking: Even in concurrency, at least one process make progress.

Dependent non-blocking: When there is no concurrency, at least one process make progress.

Dependent blocking: When there is no concurrency, it is not guaranteed that at least one process make progress.

Relations between liveness properties



Register

Dimensions

- Value ranges: The set of values that can be stored in the register.
- Access pattern: The number of processes that can read or write the register.
 - single reader: 1R
 - multi reader: MR
 - single writer: 1W
 - multi writer: MW
- Concurrent behavior: The correctness guarantees ensured when the register is accessed concurrently.
 - Atomicity: linearizable

- Safety: (single writer) If write does not overlap, return last written value, otherwise return any value in its range.

- Regularity: (single writer) If write does not overlap, return last written value, otherwise return value written or the precedent.

Python

Library: threading

Objects:

- Thread: Methods: start(), run(), join() # join wait a thread to be finished to continue.
- Lock: Methods: acquire() [block if already acquired], release(), locked()
- RLock: Reentrant lock, recursion over the same thread.
- Condition: Inherits from Lock. Work as synchronized in Java. Methods: acquire(), release(), wait(), notify(), notify_all()
- Semaphore: Inherits from Condition. Classic Dijkstra semaphore. Methods: acquire(), release()
- Event: Inherits from Condition. Communication mechanism for threads. Methods: is_set(), set(), clear(), wait()
- Timer: t = Timer(30.0, print_hello); t.start() # After 30 seconds hello will be printed.
- Barrier: b = Barrier(n) # b.wait will be blocked until n instances are waiting.

Which one use? check [this link](#)

Library: syncio