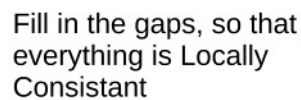
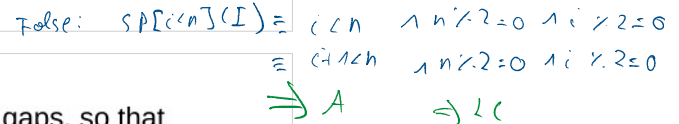


$$h = i + 2.R$$

$$n - c = k \cdot 2$$

$$n \equiv i \pmod{2}$$

1.



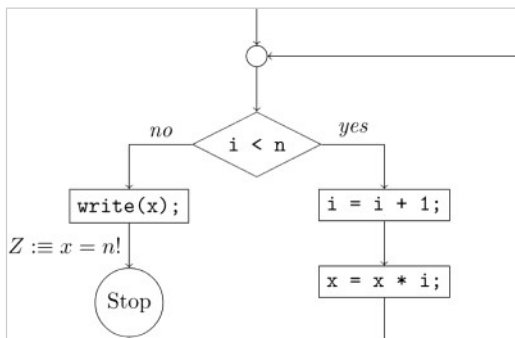
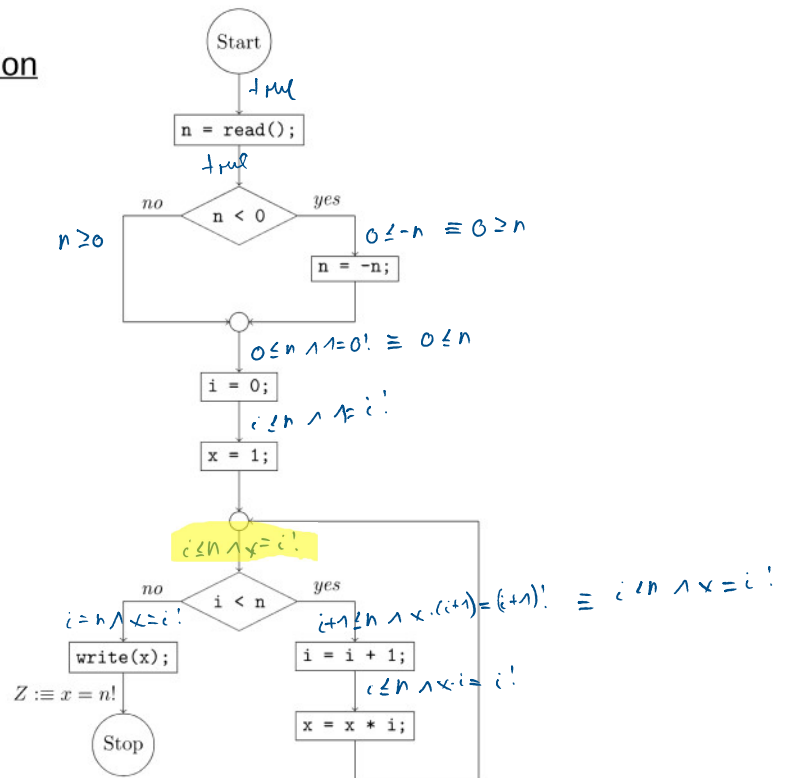
factorial

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factorial

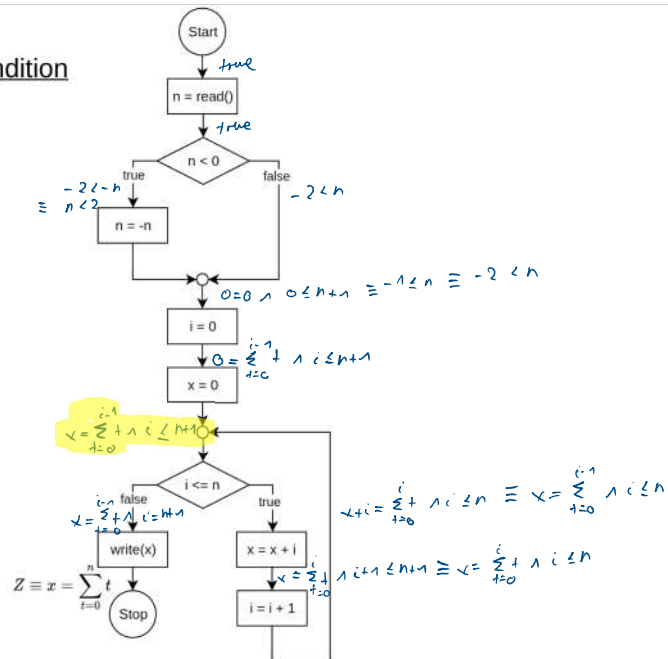
Prove Z using Weakest Precondition



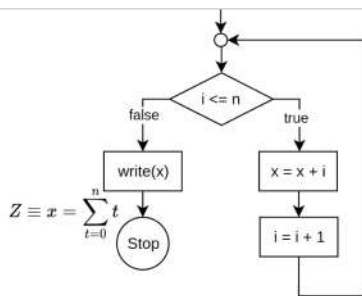


simple_sum

Prove Z using Weakest Precondition



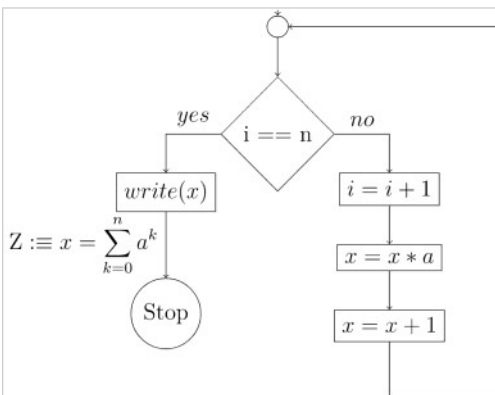
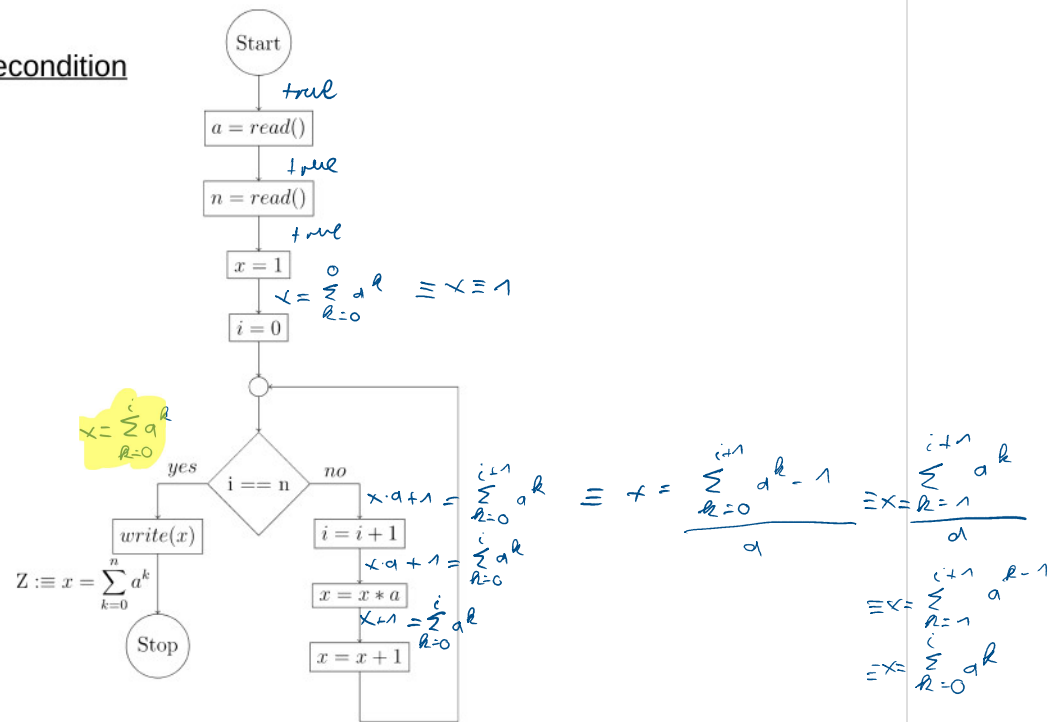
False Branch
 $\equiv i \leq n+1 \wedge i > n$
 $\equiv i-1 \leq n+1 \wedge i > n$
 $\equiv i < n+2 \wedge i > n$
 $\equiv n < i < n+2$
 $\Rightarrow i = n+1$





power_sum

Prove Z using Weakest Precondition



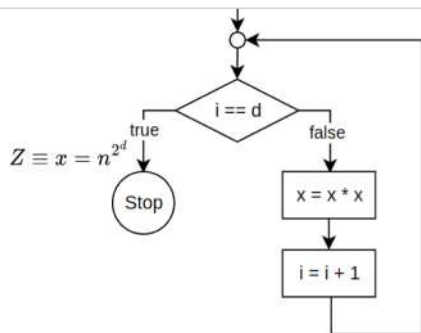
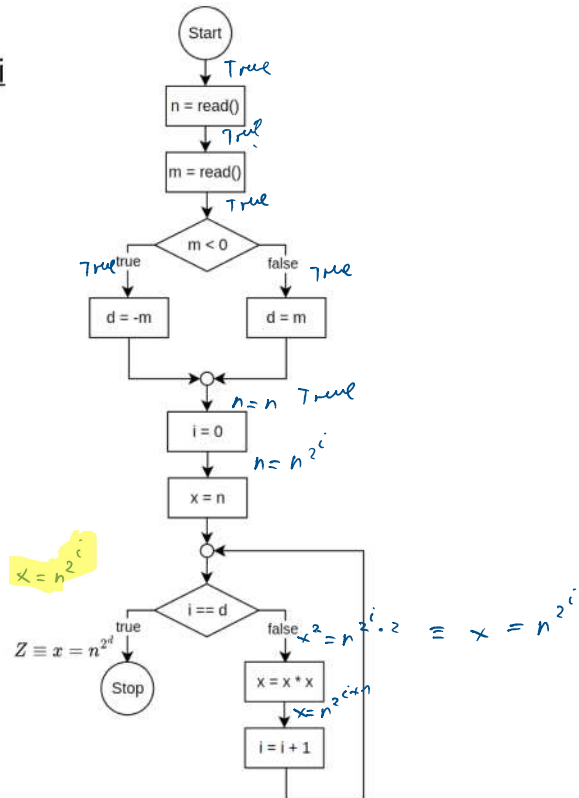
double_power

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double_po
wer

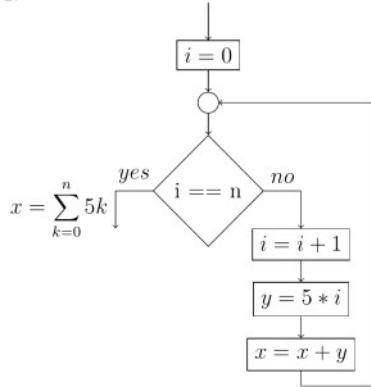
Prove Z using Weakest Precondi



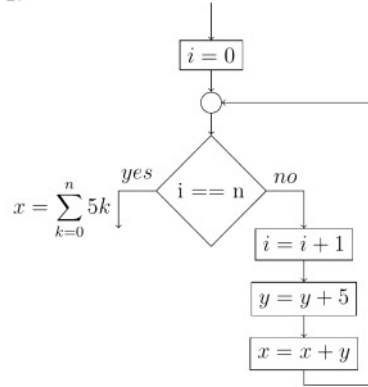


Consider these control flow graph fragments (assume x and y to be 0 initially):

1.

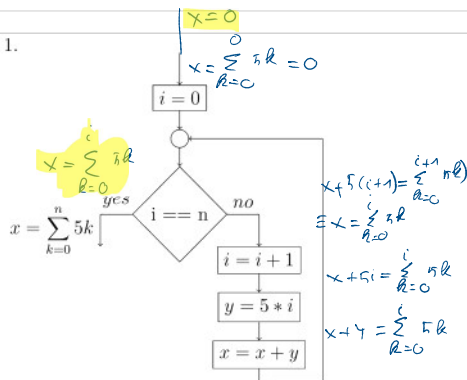


2.

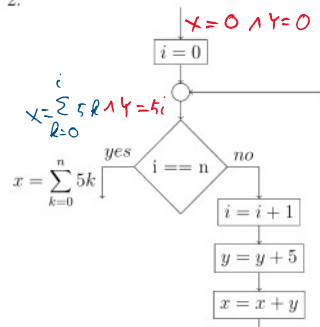


Find suitable loop invariants and prove them locally consistent. Discuss, why these invariants have to be like that.

1.



2.



1. Try the same loop invariant again
2. Why didn't it work? What information are we missing?
3. How can add that information to our proof?

$$\begin{aligned}
 x + y + 5 &= \sum_{k=0}^{i+1} 5k \equiv x = \sum_{k=0}^i 5k + 5(i+1) - y - 5 \equiv x = \sum_{k=0}^i 5k + 5i - y \\
 x + y + 5 &= \sum_{k=0}^i 5k \wedge y + 5 = 5(i+1) \\
 x + y + 5 &= \sum_{k=0}^i 5k \wedge y = 5i
 \end{aligned}$$

$$\begin{aligned}
 &\wedge y + 5 = 5(i+1) \\
 &\equiv x = \sum_{k=0}^i 5k + \frac{5i - y}{0} \wedge y = 5i \\
 &\equiv x = \sum_{k=0}^i 5k
 \end{aligned}$$

i	0	1	2	3	4	
y	0	5	10	15	20	
x	0	5	15	30	50	

$y = 5i$

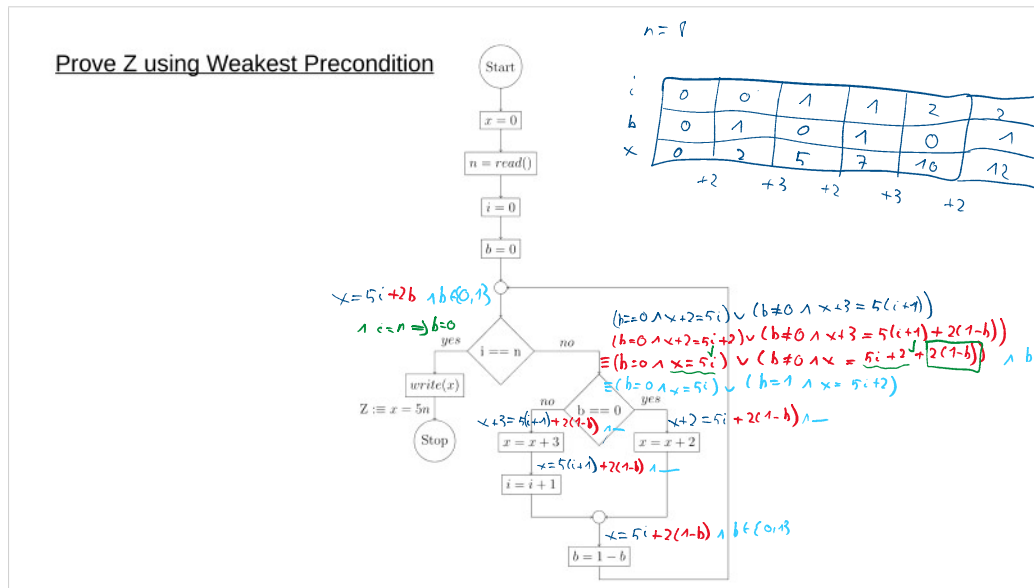
y	0	5	10	15	20	
x	0	5	15	30	50	

$$Y = 5i$$

$$x = \sum_{k=0}^i 5k$$



the_second_strenght...



1.

2.

$b = f_b(i)$ \times

$x = f_x(i, b) = 5i + 2b$

Reusing the invariant.

